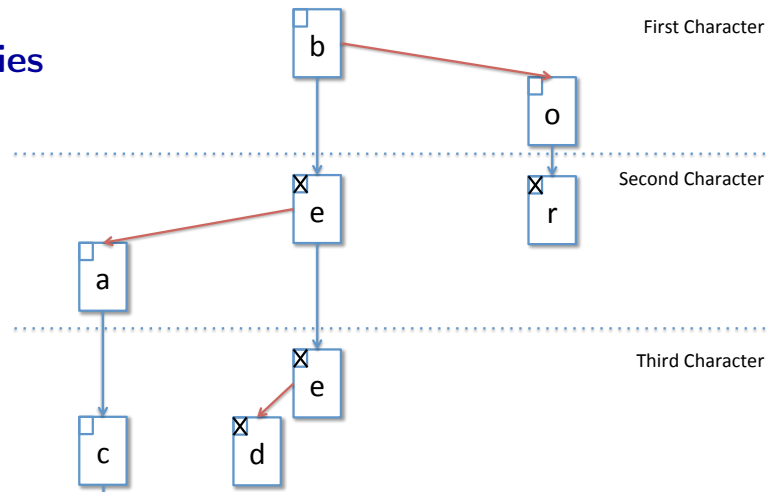


Multi-way tries



Ternary search tries have the following invariants:

- The left and right of a trie node act as a BST. In other words, for any tnode T , IGNORING middle pointers, every tnode in the $T \rightarrow \text{left}$ subtree represents a character whose ASCII value is less than $T \rightarrow c$, and every tnode in the $T \rightarrow \text{right}$ subtree represents a character whose ASCII value is greater than $T \rightarrow c$.
- For any tnode T , $T \rightarrow \text{middle}$ is a valid tnode.
- If, for some tnode T , $T \rightarrow \text{middle} == \text{NULL}$, then $T \rightarrow \text{elem}$ must not be NULL (otherwise this trie would not be a prefix for any element!).

The same `tnode_lookup` function in the implementation can be used to implement both `trie_member` (the given string is in the trie) and `trie_prefix` (this string is a proper prefix of a string in the trie).

```

1 tnode *tnode_lookup(tnode *T, char *s, size_t i) {
2   REQUIRES(is_tnode_root(T) && s != NULL && i < strlen(s));
3
4   _____
5
6   _____
7
8   _____
9
10  _____
11
12  _____
13 }
14
15 bool trie_member(trie TR, char *s) {
16   REQUIRES(is_trie(TR) && s != NULL && strlen(s) > 0);
17   tnode *T = tnode_lookup(TR->root, s, 0);
18   return T != NULL && T->is_end;
19 }
20
21 bool trie_prefix(trie TR, char *s) {
22   REQUIRES(is_trie(TR) && s != NULL && strlen(s) > 0);
23   tnode *T = tnode_lookup(TR->root, s, 0);
24   return T != NULL && T->middle != NULL;
25 }

```

Definition-as-use

In C0, all variable declarations were comprised of a *type* and a *variable name*.

```
the type (int pointer)  the variable name (myptr)
vvvv                    vvvvvv
int*                    myptr;
```

In C, you have to think about declarations differently: as a *base type* and a *pattern describing how the variable name is used*.

```
the type (int)  it is used by dereferencing (myptr)
vvv            vvvvvv
int           *myptr;
```

On one level, this is just a stylistic difference: C isn't that whitespace sensitive. But there are two places where this makes a critical difference.

Checkpoint 0

One one line, declare an integer *i*, a pointer to an integer *p*, and an array of pointers to integers *A*.

```
int _____;
```

Checkpoint 1

You can define and use a function pointer like this:

```
1 typedef int string_hasher(char *s);
2
3 int hash_string(char *s) { REQUIRES(s != NULL); ... }
4
5 int main() {
6     string_hasher *f = &hash_string;
7     printf("Hash of 'hello world': %d\n", (*f)("hello world"));
8 }
```

Using the idea of definition-as-use, declare *f* in the example above without using the intermediate typedef.

```
_____ = &hash_string;
```

Checkpoint 2

Consider the following:

```
1 typedef int compare(void *x, void *y);
2 void sort(void **A, int len, compare *F);
3 // requires A != NULL
4 // requires \length(A) == len
5 // requires compare != NULL
```

How would you use this procedure to sort an array of strings? How would you write a generic implementation of `sort` using the selection sort algorithm from earlier in the semester? Which of these preconditions could we actually check in the `sort` function, and how?