Multi-armed Bandit Problems

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August, 2008

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Summary

- Introduction
- Classical formulation
- Properties
- Computational issues
- Extensions
- Example
- Discussion

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Definition

- Multi-armed bandit (MAB) problems
 - sequential resource allocation
 - among competing (mutually exclusive) projects
- Difficulty related to conflict between
 - allocating resources that yield good rewards
 - trying "not so promising" projects
 - but maybe with better future prospects

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Examples

- control theory problems
- allocating researchers to projects
- clinical trials
- sensor management

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Definition

- Classical definition
 - single resource
 - allocated to one of many competing projects (bandits, arms)
 - project w/ resource can change its state
 - other projects remain frozen
 - discrete time, no switching costs

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Solving

- In general this is solvable via Dynamic Programming
 - backwards induction
 - $V^*(s, N) = R_N(s), \quad \forall s$
 - ► $V^*(s, N-1) =$ max $R_{N-1}(s, a) + \gamma \sum_{s'} T(s, a, s') V^*(s', N)$
 - Bellman equations
 - $V^*(s) = \max_a R(s, a) + \gamma \sum_{s'} T(s, a, s') V^*(s')$
 - very general stochastic optimization problems
 - VI, PI, RL
 - Curse of dimensionality

Solving

- But MAB are simpler and allow for "index-type" solutions
 - for each bandit associate a dynamic allocation index (DAI)
 - depends only on that bandit
 - one k-armed bandit vs k single-armed bandits
 - at each time, choose the bandit with highest DAI
 - leads to optimal allocation policy
 - DAIs are also known as "Gittins Indices"

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- (single-armed) bandit process
 - described by pair random of sequences
 - $\{X(0), X(1), ...\}$
 - { $R(X(0)), R(X(1)), \ldots$ }
 - X(n): state after arm has been operated *n* times;
 - R(X(n)): reward obtained on the *n*-th operation
 - state evolution:
 - $X(n) = f_{n-1}(X(0), \dots, X(n-1), W(n-1))$
 - thus, arm not necessary Markov
 - W(n): independent sequence of RVs; independent also from X(0)

multi-armed bandit process

- k independent arms
- one controller
- controller operates exactly one arm at a time
- machines described by time-dependent sequences:
 - $\vdash \{X_i(N_i(t)), R_i(X_i(N_i(t)))\} \quad \forall i \forall t$
 - N_i(t): number of times machine i has been operated up to time t
 - $N_i(t)$ is the "local time" of machine *i*
- control is $U(t) = \{U_1(t), \dots, U_k(t)\}$, ie, in the form $\{00 \dots 1 \dots 000\}$

- System evolution
- $X_i(N_i(t+1)) =$ $f_{N_i(t)}(X_i(0),...,X_i(N_i(t)),W_i(N_i(t))) \quad \text{if } U_i(t) = 1$ $X_i(N_i(t)) \quad \text{if } U_i(t) = 0$
- ► $N_i(t+1) =$ ► $N_i(t) + 1$ if $U_i(t) = 1$
 - $N_i(t)$ if $U_i(t) = 0$
- $R_i(t) = R_i(X(N_i(t)), U_i(t)) =$ $R_i(X_i(N_i(t))) \quad \text{if } U_i(t) = 1$ $0 \quad \text{if } U_i(t) = 0$

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- Scheduling policy
- $\gamma = (\gamma_1, \gamma_2, \ldots)$
- such that $U(t) = \gamma_t(Z_1(t), ..., Z_k(t), U(0), ..., U(t-1))$
- and $Z_i(t) = \{X_i(0), \dots, X_i(N_i(t))\}$
- In other words, policy might depend on full history of arms' states and previous actions

Goal is to find scheduling policy γ that maximizes

$$J^{\gamma} = E\Big(\sum_{t=0}^{\infty} \beta^t \sum_{i=1}^{k} R_i(X_i(N_i(t)), U_i(t)) \mid Z(0)\Big)$$

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- simplest policy: myopic decisions (1 look-ahead)
- not optimal, in general
- T-steps look-ahead
 - take decisions that maximize expected reward for the next T steps
- Generalization: do not fix T
 - let τ be the number of look-ahead steps
 - τ is a RV that depends at each time on how the system evolves
 - τ is considered a *stopping time*

- in order to maximize J^{γ} , we must
 - choose a rule γ for taking a sequence of decisions
 - choose a value for \(\tau\)
- such that that rule, when used for τ steps, gives the max J^γ
- This extension of T-steps look-ahead works by
 - At t = 0, given Z(0), select γ_1 and τ_1
 - Apply γ_1 for τ_1 steps
 - ► repeat, choosing the next γ_t , τ_t , conditioned on the new information gained
 - notice: decisions based only on current states of arms
 - "forward" because keeps deciding next policies for the future

- in general this is not optimal
 - route choosing example
 - problem are irrevocable decisions
 - some alternatives available at some stage are not available later
- if any decisions made are not irrevocable, forward induction is optimal
 - every arm not used is kept frozen
 - thus can deliver the same sequence of rewards later on (up to β)

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Gittins proved that the following index is optimal

$$v_{X_i}(x_i(0)) = \max_{\tau > 0} \frac{E\left(\sum_{t=0}^{\tau-1} \beta^t R_i(X_i(t)) \mid x_i(0)\right)}{E\left(\sum_{t=0}^{\tau-1} \beta^t \mid x_i(0)\right)}$$

- suppose we are allowed to take decisions only while they're worth it,
 - then v_{X_i} gives a "retirement" value
 - ie, a value in which we are indifferent to continuing operating *i* or quitting
 - only quit *i* (and work on some *j*) if *j* has a better prospect than the retirement offered

- When in decision stage I, for each arm i,
- and considering information $x_i^{\prime}(\omega) = (x_i(0), \ldots, x_i(N_i(\tau_i(\omega))))),$

$$v_{X_i}(x_i^l(\omega)) = \max_{\tau > \tau_l(\omega)} \frac{E\left(\sum_{t=\tau_l(\omega)}^{\tau-1} \beta^t R_i(X_i(N_i(\tau_l) + t - \tau_l(\omega))) \mid x_i^l(\omega)\right)}{E\left(\sum_{t=\tau_l(\omega)}^{\tau-1} \beta^t \mid x_i^l(\omega)\right)}$$

easier if arm is Markov

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Computational issues

- Focus on Markov arms
- State space $S_i = \{1, 2, \dots, \Delta_i\}$

$$v_{X_i}(x_i(t)) = \max_{\tau > t} \frac{E\left(\sum_{t'=t}^{\tau-1} \beta^t R_i(X_i(t')) \mid x_i(t)\right)}{E\left(\sum_{t'=t}^{\tau-1} \beta^t \mid x_i(t)\right)}$$

Need to compute v for each state of each arm

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Computational issues

- Offline approach: compute indices for all states, all machines
- Online approach: only index for the last used machine
- Continuation/stopping sets
 - remember, v is retirement value
 - only quit machine *i* if reach state from which *j* would be better
 - $C_i(x_i)$: all states with index higher than x_i 's
 - $S_i(x_i)$: all states with index lower than x_i 's

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Offline calculation

- Computing $C_i(x_i)$ and $S_i(x_i)$
 - ordering on states: $I_1, I_2, \ldots, I_{\Delta_i}$ s.t.
 - $\mathsf{v}_{X_i}(I_1) \geq \mathsf{v}_{X_i}(I_2) \geq \ldots \geq \mathsf{v}_{X_i}(I_{\Delta_i})$
- For machine *i*, set $I_1 = \arg \max_{x_i} R_i(x_i)$
 - Now consider probabilities in Pⁱ only for transitioning to "better" states;
 - Given reward matrix R_i (reward per state);
 - For each state x_i , calculate $D_{x_i}^{i,n}$
 - expected discounted reward considering next (better) states
 - Calculate B^{i,n}
 - expected total "discounts", considering probabilities of transitions

$$V_{X_i}(x_i) = \frac{D_{x_i}^{l,n}}{B_{x_i}^{i,n}}$$

Online calculation

- Also uses the continuation/stopping sets approach
- Assume we are operating machine i in state a
 - now, we are given opportunity to switch to state x_i
 - maximize expected discounted reward over infinite horizon

$$V(a) = \max \left\{ R_i(a) + \beta \sum_{b \in \{1, ..., \Delta_i\}^i} P^i_{a, b} V(b), R(x_i) + \beta \sum_{b \in \{1, ..., \Delta_i\}} P^i_{x_i, b} V(b) \right\}$$

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Online calculation

- ► Now C_i(x_i) is the set of states with expected reward larger than V(x_i);
- $v_{X_i}(x_i) = (1-\beta)V(x_i)$
- Questions:
 - Why maximize infinite horizon is equivalent?
 - Why (1β) and not $\frac{1 \beta}{\beta}$?

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Superprocesses

- Same as before, but now each arm *i* receives control input U_i ∈ {0,..., M_i}
- U_i = 0 is a freezing action; rest are continuation actions
- If control policies are fixed, degenerates to regular MAB
- Otherwise, state evolution are rewards depend on current state and on current control input
 - Not a Markov Chain, but a Markov Process
- Scheduling policy γ controls exactly one machine

$$J^{\gamma} = E^{\gamma} \Big(\sum_{t=0}^{\infty} \beta^t \sum_{j=1}^{k} R_j(X_j(N_j(t)), U_j(t)) + Z(0) \Big)$$

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Superprocesses

- Time evolution of arm is controlled
- More complex than MAB; in general, Gittins Indices not optimal
- Unless each arm (desc. by seq. X states, rewards) has a dominating arm

$$L(X,\mu) = max_{\tau>0} \Big(\sum_{t=0}^{\tau-1} \beta^t [R(X(t)) - \mu]\Big)$$

X dominates Y iff

 $L(X,\mu) \ge L(Y,\mu) \qquad \forall \mu \in \mathbf{R}$

• μ is "retirement" value; $L(X,\mu)$ the expected gain over μ

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Superprocesses

- If there is dominance, optimal because
 - No matter how big the offered retirement is (to quit *i*), there's always a better arm *j*
- In practice, this is a quite restrictive assumption

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Example

Arm-acquiring bandits

- Regular MAB, but new arms can be created
- Gittins Indices are optimal
 - Decisions are not irrevocable
 - Decisions based on indices with K_i arms consider all of them
 - But decisions prior to this *did not* have all *K_i* arms available;
 - no way a prior decision could be "wrong"

Switching penalties

- Regular MAB, but there is a cost c for switching arms
- Gittins Indices are not optimal (example in book)
- If the index is

$$v_{X_{i}}^{s}(x_{j}(0)) = \max_{\tau > 0} \frac{E\left(\sum_{t=0}^{\tau-1} \beta^{t} R_{j}(t) - c \mid x_{j}(0)\right)}{E\left(\sum_{t=0}^{\tau-1} \beta^{t} \mid x_{j}(0)\right)}$$

- then only qualitative results are known [11]
- the general nature of the scheduling policies is not known
- solution usually requires full use of DP (backwards induction)

Multiple plays

- ▶ Regular *k*-processes MAB, but is *m* processors
- At each time allocate each processor to exactly one process
- No process being operated by more than one processor
- Only processes being processed generate reward
- Allocation according to *m* highest indices: not optimal
- Optimal if indices are sufficiently separated (C1, p.141)
 - How to guarantee this beforehand?
- For different criteria (eg: regret minimization) optimal policies are known [7,8]

Restless bandits

- k machines, m processors
- machines' states evolve over time even when not being processed
- reward of non-processed machines might be assumed to be zero
- performance criterion is

$$J^{\gamma} = E^{\gamma} \Big(\sum_{t=0}^{\infty} \beta^t \sum_{j=1}^k R_j(X_j(N_j(t)), U_j(t)) + Z(0) \Big)$$

 Goal is to find policy that maximizes infinite horizon expected discounted reward

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Restless bandits

- In general, Gittins Indices are not optimal
- But for some other optimization criterion, indices are optimal
 - eg: infinite horizon average reward-per-time-per-machine criterion

$$\frac{1}{k}\left(\lim_{T\to\infty}\frac{1}{T}E\left(\sum_{t=1}^{T}\sum_{i=1}^{k}R_{i}(X_{i}(t-1),U_{i}(t))\right)\right)$$

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Restless bandits

- Gittins indices for RB are related to "gift" values given to non-processed machines
- Argument is similar to that of the "retirement" value
 - index is a "gift" value that makes us indifferent to running or not the machine
 - it is only worth to run the machine if the expected gain is greater than the "gift" value
 - this values allow us to index all machines

Example

- find one stationary target hidden in one of k cells
- prior probability of the target in cell *i* is $p_i(0)$
- sensor can look into just one cell at a time
- sensor is imperfect
 - *P*(sensor finds target in *i* | target is in cell *j*) = $\delta_{i,j}q_j$
 - where δ is the Kronecker delta function;
 - q_j (?) is probability of false positive
- reward upon completion is β^t (ie, we want to find the target ASAP)
- which sensor to activate at each time?



- let p_i(t) be the posterior probability of target being in cell i
- $p_i(t)$ is state of cell (arm) *i* at time *t*

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Example

• For a policy γ , expected reward is

$$= \sum_{t=0}^{\infty} \beta^{\tau} P(\text{target is found at } \tau, \text{ analyse correct cell})$$
$$= \sum_{t=0}^{\infty} \beta^{\tau} \sum_{i=1}^{k} p_i(t) q_i P^{\gamma}(U(t) = e_i)$$
$$= \sum_{t=0}^{\infty} \beta^{\tau} \sum_{i=1}^{k} R_i(p_i(t), u_i(t))$$

• where reward is given for *i* iff *i* is activated at t (U(t) = i)

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Example

- Unfortunately, updates in p_i affect all other probabilities (states)
- Thus, not a regular MAB
- Easy to solve if we consider unnormalized probabilities

$$p_i(t+1)$$

= $p_i(t)$ if $u_i(t) = 0$
= $p_i(t)(1-q_i)$ if $u_i(t) = 1$

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Extensions

Example

- We try to maximize the long-term expected reward
 - ► remember, $R_i(p_i(t), u_i(t)) = p_i(t)q_i$ iff $u_i(t) = 1$, zero otherwise

$$\sum_{t=0}^{\infty}\beta^t\sum_{i=1}^{k}R_i(p_i(t),u_i(t))$$

- Gittins Index of every machine is always achieved at τ = 1 (?), so:
 - $v_{X_i}(p_i(t)) = p_i(t)q_i$
 - which is by the definition of GI, for one-step look-ahead
 - β can be ignored from the denominator because it is constant

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Example

- ► If sensor operates in *M* modes: superprocess
- If there is cost to switch targetting area: MAB w/ switching penalties
- ▶ If there are *m* sensors: MAB w/ multiple plays
- If target is moving: m sensors, restless bandit

Conclusion

- Gittins indices simplify the policy calculation for a class of sequential decision problems
- MAB are very simple problems, but might be extended
 - extensions are often related with one another
 - arm-acquiring → superprocess [240]
 - switching costs \rightarrow restless bandits [91]
 - Tax problem (minimization of cost of frozen machines) → MAB

Thanks



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