

# 15-496 : A Hand-on Introduction to Wireless Networks

## Lecture 10: Wireless and the Internet

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<http://www.cs.cmu.edu/~prs/wireless08/>

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## Outline

- **Wireless and the Internet**
- **Mobility: Mobile IP**
- **Security: EAP and Radius**
- **Diversity**
  - › Loss: TCP over wireless
  - › Delay: multimedia
  - › Devices: WAP

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## Internet Architecture Assumptions

- **Host are (mostly) stationary**
  - › Address assignment, routing
- **Links in the network are fairly homogeneous**
  - › Transport protocols, applications
- **Hosts are fairly powerful**
  - › End to end principle: push functionality to end points
- **Security is an end host issue**
  - › No security inside the network (architecturally)

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## Mobility

- **Many clients today are mobile**
- **Mobility inside a subnet is supported**
  - › E.g. moving across APs that are part of a single EBSS
- **Mobility across subnets is harder because the IP address is used as address and identifier**
  - › Identifier: who you are
  - › Address: where you can be found
- **Keep IP address: network gets confused**
  - › Delivers packets to wrong "old" subnet
- **New IP address: host gets confused**
  - › Transport protocols, applications, etc.

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## Link Heterogeneity

- **Original links were basically telephone lines**
- **Today: huge diversity**
  - › Optical fiber (wavelengths) ... copper ... wireless
- **Throughput is most visible diverse features**
  - › Optical links 40 Gbps ... Wireless links of Kbps
- **But other features are different as well:**
  - › Error characteristics: higher on wireless
  - › Latency: absolute delay and variance in delay
- **Mobility adds to diversity**
  - › E.g. hand off can cause delays and sudden changes in available bandwidth

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## Device Capabilities

- **Originally: mainframes and personal computers**
- **Today: sensors ... supercomputers**
- **Note: almost any networked device today is more capable than early computers!**
  - › But our requirements and expectations have increased
  - › Anything at or above "PC class" is very capable
- **Laptop: view as mobile PC**
- **PDA: view as 5 year old PC (kind of)**
- **Cell phone: not even close to a PC**
- **Sensors: often run as private networks**

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## Security

- Access to the network enables all kinds of attacks
  - › Argues for pushing security inside the network
  - › Firewalls are a very ad hoc way of doing this
- Wireless creates unique challenges
  - › Do not need physical connection to sniff or send
- Discussed (briefly) in the 802.11 lectures
  - › WEP, 802.1x, etc.
- But wireless security needs to be linked into system wide security

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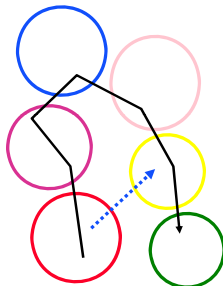
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## Mobility across IP Subnets

- When moving across IP subnets, different protocols have conflicting requirements
- Network layer wants IP address in current subnet
  - › Needed for routing of packets
- Transport layer wants IP address that was used to create connection
  - › Is used to identify the connection
- Applications often do not care but in practice they want to keep the IP address the same
  - › Tied to sockets



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## Mobile IP Goals

- Communicate with mobile hosts using their "home" IP address
  - › Allows any host to contact mobile host using its "usual" IP address
- Mobility should be transparent to applications and higher level protocols
  - › No need to modify the software
- Minimize changes to host and router software
  - › No changes to communicating host
- Security should not get worse

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## Mobile IP

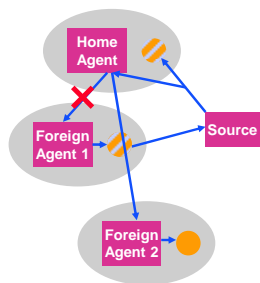
- Home network has a home agent that is responsible for intercepting packets and forwarding them to the mobile host.
  - › E.g. router at the edge of the home network
  - › Forwarding is done using tunneling
- Remote network has a foreign agent that manages communication with mobile host.
  - › Point of contact for the mobile host
- Binding ties IP address of mobile host to a "care of" address.
  - › binding = (IP address, foreign agent address)
  - › binding includes time stamp

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## Mobile IP Operation

- Agents advertise their presence.
  - › Using ICMP or mobile IP control messages
  - › Mobile host can solicit agent information
  - › Mobile host can determine where it is
- Registration process: mobile host registers with home and foreign agent.
  - › Set up binding
- Tunneling
  - › forward packets to foreign agent
  - › foreign agent forwards packets to mobile host
- Supporting mobility
  - › invalidating old caches in a lazy fashion



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## Optimizations

- **Mobile host can be its own the foreign agent.**
  - › Mobile host acquires local IP address
  - › performs tasks of the mobile agent
- **Short circuit the home location by going directly to the foreign agent.**
  - › Routers in the network store cache bindings and intercept and tunnel packets before they the mobile host's home network
  - › Need a protocol to update/invalidate caches
  - › Raises many security questions and is not in the standard

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## Security Considerations

- **Authentication of mobile host, home agent**
  - › avoid invalid interception of traffic
- **Authentication of mobile host, foreign agent**
  - › Desirable, but more difficult, so not required
- **Use encryption for sensitive data**
- **Replay of registration messages.**
  - › standard problem
  - › use standard solutions, e.g. timestamps, nonce, ..
- **Dealing with the firewalls at the foreign site.**
  - › most easily sent directly to destination
  - › but has a "strange" IP network address in source field
  - › reverse tunneling by foreign agent or mobile host
  - › Can replace tunnels by UDP for NATs

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## Discussion

- **Mobile not used in practice**
- **Was not designed for truly mobile users**
  - › I.e. for continuous operation across subnets
  - › Switching between subnets is heavy weight
- **Was designed for nomadic users**
  - › I.e. visitors are a remote site
- **Value for nomadic is limited to:**
  - › Mobile user is contacted using home IP address
  - › User maintains connections between subnets
  - › Neither of them are very likely or common
  - › Laptops are mostly used for accessing servers

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## Wireless Security

- **Security is not just about authentication and encryption**
- **Must also consider business and deployment issues**
  - › AAA: Authentication, Authorization, and Accounting
  - › Supporting users at different levels

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## Authentication in WLAN Hotspots

- **Upon association with the AP, only authentication traffic can pass through, as defined by IEEE 802.1x**

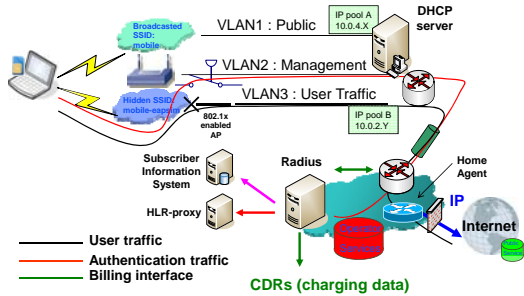


- **The protocol used to transport authentication traffic is the EAP standard (RFC3748)**

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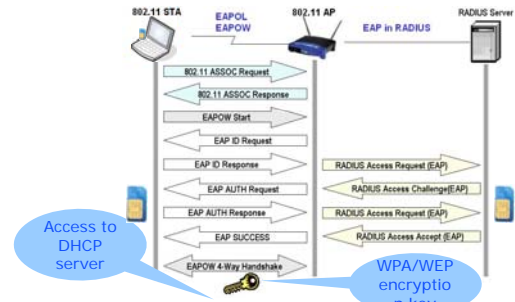
## Dual SSID Approach



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## 802.1x and EAP Protocol Execution



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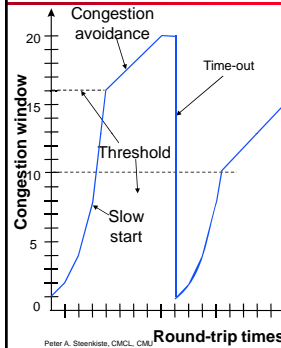
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## TCP Congestion Control



- Congestion detection based on lost packets:
  - › Time-out or
  - › Receipt of duplicate ACKs
- Packet loss results in cutting window in half
  - › To reduce congestion in network
- Assumes all (or most) packet losses are due to queue overflow
  - › I.e. two are equivalent
- TCP also estimates roundtrip time
  - › Used to set timeout value

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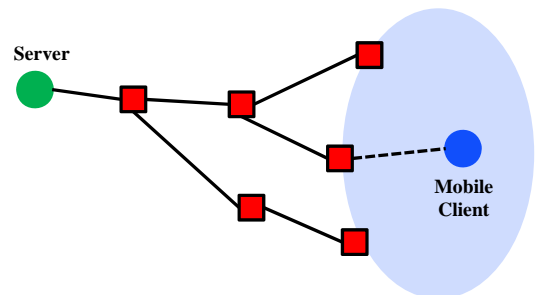
## Challenges in Wireless

- **Wireless network may drop packet due to link errors**
  - › Will be interpreted as congestion losses by TCP sender
  - › Will result in backoff and loss of throughput
  - › TCP tends to perform poorly when error rates go above 1%
- **High variability of properties wireless links could also confuse TCP**
  - › TCP likes nice predictable paths
- **Roundtrip time can be variable**
  - › Due to wireless retransmissions, handoff, ...
- **Available bandwidth can fluctuate**
  - › Rate adaptation, interference, handoff, ...

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## Solution Ideas?



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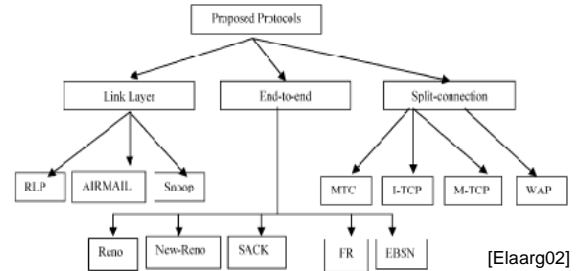
## Solution Space

- **Modify TCP for wireless paths**
  - › Would maintain status quo for wired paths
  - › What would wireless TCP look like?
  - › Difficult to do: there are many Internet hosts
  - › Traditionally, hosts have no information about path properties
- **Modify TCP for all paths**
  - › Not clear what that modification would be!
  - › Similar problems: need to modify many hosts
- **Modify TCP only on the mobile host**
  - › Interesting idea – kind of what WAP does in an extreme way
- **Keep end hosts the same but tweak things at the wireless gateway**
  - › Keep end-end TCP happy despite wireless links

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## Possible Classification of Solutions



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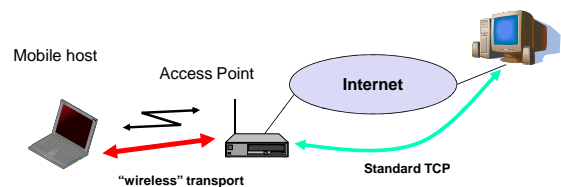
## Connection Split: Indirect TCP or I-TCP

- **Do not change TCP on the wire-line part**
- **Split the TCP connection at the wireless gateway into two parts**
  - › One optimized for the wireless link
  - › The second for the wire-line communication (TCP)
- **No real transport-layer end-to-end connection**
  - › Although host on wired network does not know this
- **Wired host should not notice the characteristics of the wireless part**
  - › This is a challenge since wireless gateway is limited in what it can send and when, e.g. cannot prematurely acknowledge data
  - › Certain things cannot be hidden: delay, dramatic throughput variations

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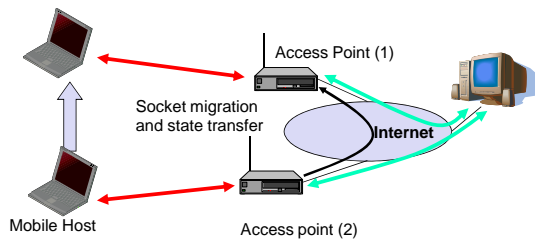
## I-TCP



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## I-TCP and Mobility



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## I-TCP Discussion

- **I-TCP Advantages**
  - › No changes in the fixed network or hosts (TCP protocol), all current TCP optimizations still work
  - › Wireless transmission errors do not "propagate" to the wire-line network
  - › Simple, effective (in the best case)
- **I-TCP Disadvantages**
  - › End-to-end semantics become less clear, e.g. what happens if the wireless gateway crashes
  - › Higher end-to-end delays due to buffering and forwarding to the gateway

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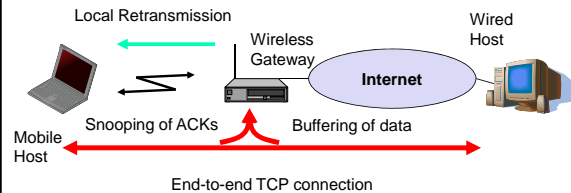
## Snooping TCP

- “Transparent” extension of TCP within the wireless gateway
  - › End hosts are not modified
- Hides wireless losses from wired host
  - › Buffer packets sent to the mobile host
  - › Local retransmission: Lost packets on the wireless link, for both directions, are retransmitted immediately by the mobile host or foreign agent
- Wireless gateway “snoops” the packet flow so it can cover up signs of packet loss
  - › E.g. recognizes acknowledgements in both directions and suppresses duplicate ACKs

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## Snooping TCP



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## Snooping TCP Discussion

- Data transfer to the mobile host
  - › FA buffers data until it receives ACK from the MH
  - › FA detects packet loss via duplicated ACKs or time-out
- Data transfer from the mobile host
  - › FA detects packet loss on the wireless link via sequence numbers
  - › FA answers directly with a NACK to the MH
  - › MH can now retransmit data with only a very short delay
- Integration of the MAC layer
  - › MAC layer often has similar mechanisms to those of TCP
- Problems
  - › Snooping TCP does not isolate the wireless (as I-TCP)
  - › Snooping might be useless if encryption is used

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## How about Mobility?

- Split connection approach
  - › Hard state at access point must be moved to new base station
- Snoop approach
  - › Soft state need not be moved
  - › While the new access point builds new state, packet losses may not be recovered locally
- Frequent handoffs remain a problem
  - › Hard state should not be lost
  - › Soft state needs to be recreated for performance

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## Mobile TCP

- Handling of long and frequent disconnections
- M-TCP splits connection as I-TCP does
- Wireless gateway monitors all packets but
  - › Does no caching, no retransmission
  - › If it detects a disconnection
    - It sets the sender’s window size to 0
    - Sender automatically goes into “persist” mode
- Advantages
  - › ETE semantics are maintained, no buffer forwarding
- Disadvantage
  - › Wireless link loss propagates to the wire-line network

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## An Internet Style Approach

- Use aggressive retransmission in the wireless network to hide retransmission losses
  - › Most deployed wireless network in fact do that already
  - › Would sell few products if they did not
- Wireless losses are translated into increased delay
  - › But TCP roundtrip time estimation is very conservative, e.g. increases if variance is high
- Also: persistent high loss rate results in reduced available bandwidth -> congestion response is appropriate and needed
- Works remarkably well!

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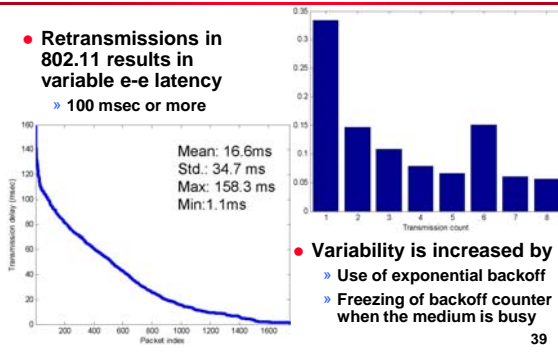
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## Variable Latency in Wireless Networks

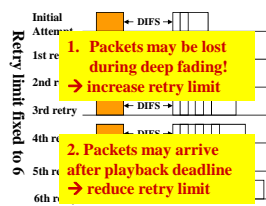
- Retransmissions in 802.11 results in variable e-e latency
  - › 100 msec or more



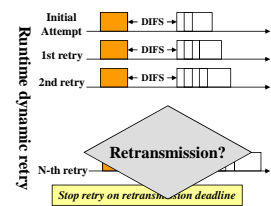
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## Time Aware Adaptive Retry

### Conventional 802.11



### TAR-aided 802.11

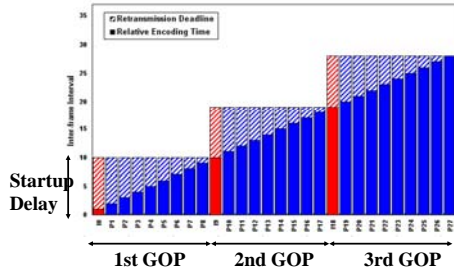


- 802.11 retransmission takes a fixed try limit for every packet (count based)
- CAR dynamically determines retry limit for each packet using its "retransmission deadline"

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## TAR: Retransmission Deadline



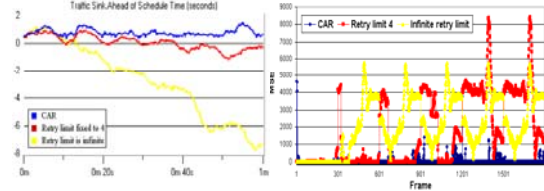
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## Comparison of Visual Quality

### Results of the Ahead of Schedule Time

### Mean Square Error of the Y Component



Packets with a value smaller than zero is behind its playback schedule

$$MSE = \frac{1}{M \times N} \sum_{i=0}^{M-1} \sum_{j=0}^{N-1} [x(i, j) - \hat{x}(i, j)]^2$$

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## What About Cell Phones?

- **Cell phones have very limited capabilities:**
  - › Low bandwidth
  - › Small display and limited “keyboard”
  - › Small memory and processing capabilities
- **Raises issues at different layers:**
  - › HTML is too rich for cell phone I/O capabilities
  - › General purpose protocol likely to be too inefficient, big, and unresponsive
- **Result: define a completely new stack and glue it to the IP stack using a gateway**
  - › Ugly but practical

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## Wireless Access Protocol (WAP)

- **Protocol stack for small low-powered devices**
  - › E.g. cell phones
- **Layered Protocol, like IP stack**
  - › Two versions of protocol stack
    - WAP1.x Protocol Stack
    - WAP2.0 Protocol Stack
  - › Based on some older, proprietary protocols
  - › Provide security
- **Connects WAP devices with a WAP gateway**
- **Gateways translate wireless protocols into Internet protocols**
  - › Located near Mobile Telephone Exchange

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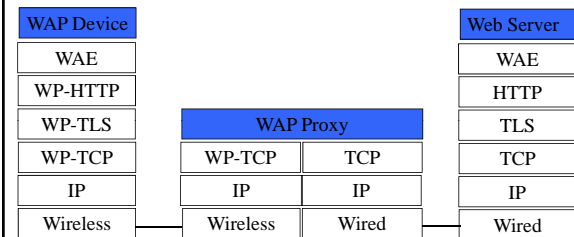
## WAP 2.x Protocol Stack

WAP Device	
WAE	Wireless Application Environment
WP-HTTP	Wireless Profiled HTTP
WP-TLS	Wireless Profiled TLS (SSL)
WP-TCP	Wireless Profiled TCP
IP	Internet Protocol
Wireless	2.5G and 3G

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## WAP 2.0 Proxy



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## Wireless Markup Language

- **Very specialized for cell phones**
  - › Needs to be compact
- **Window has limited size**
  - › Server needs to optimize for specific devices
  - › Standards ways of chaining windows and window layout
- **Keyboard is limited**
  - › Supports concise ways of using standard buttons on cells phones

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