

XIA: eXpressive Internet Architecture - A Proposal for a Future Internet Architecture

15-441/641: Computer Networking

Lecture 27: What is Next?

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www.cs.cmu.edu/~prs/15-441-F14

Outline

- Background
- The eXpressive Internet Architecture – a proposal
 - Example and concepts
 - Research thrusts
- Research examples: AIP and APIP
- User privacy survey

NOTE: this lecture describes a research project.
The goal is to have you think outside of the box

2

Key Internet Features

What we learned about the current Internet:

- Simple core with smart endpoints
- The IP narrow waist supports evolution
- Packet based communication
- All IP hosts can exchange packets
- Non-essential functions are services
- End-to-end transport protocols
- Security is not part of the architecture

But may be there are better ways?

3

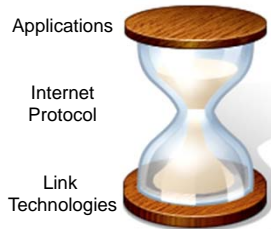
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4

“Narrow Waist” of the Internet Key to its Success

- Has allowed Internet to evolve dramatically
- But now an obstacle to addressing challenges:



- No built-in security
- New usage models a challenge – content and services, not hosts
- Hard to leverage advances in technology in network
- Limited interactions between network edge and core
- But where do we get started?

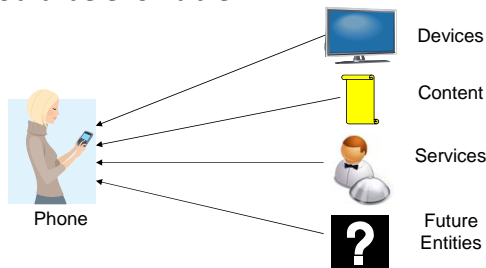
Three Simple Ideas

- Support multiple types of destinations
 - Not only hosts, but also content, services, etc.
 - Not having to force communication at a lower level (e.g., hosts) reduces complexity and overhead
- Intrinsic security guarantees security properties as a direct result of the design of the system
 - Do not rely on external configurations, data bases, ..
- Flexible addressing gives network more options for successfully completing communication operations
 - Include both “intent” and “fallback” address
 - Supports evolvability, network diversity, fault recovery, mobility, ..

6

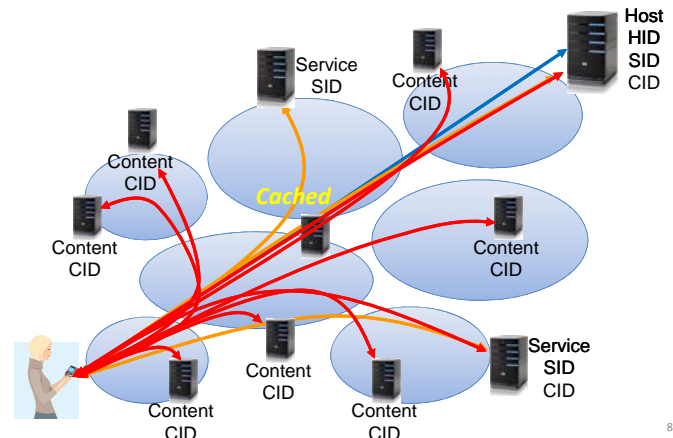
Multiple Principal Types

- Identifying the intended communicating entities reduces complexity and overhead
 - Have different forwarding semantics
- Set should be *evolvable*



7

Multiple Principal Types - Example



8

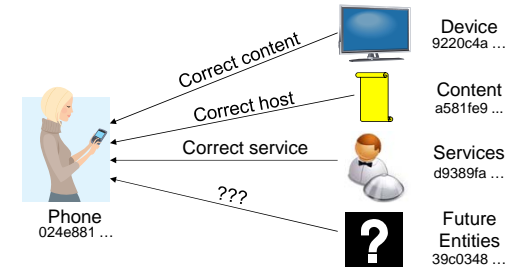
Using Principal Types

- Content and service addresses directly supports cross-application service selection and caching
 - Complex today: DNS indirection infrastructure, deep packet inspection, transparent proxies, etc.
- Routing protocols for hosts, content and services
 - Metrics driving by context, different concerns
 - Public internet: policies, business, ...
 - Intra-networks: usage models, super fast recovery, ...
- Add new (custom) functionality to the network
 - E.g., caching + service -> diverse multicast variants
 - Dealing with disruptions

9

Security as Intrinsic as Possible

- Communication security properties are a direct result of the design of the system
 - Do not rely on correctness of external configurations, actions, data bases



10

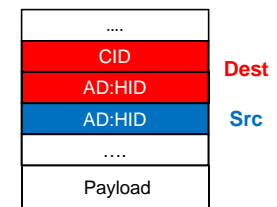
Use of Intrinsic Security

- Name-> address look automatically provides public key associated with the address
 - May not need for separate key management infrastructure
 - Can help, e.g., with network partitioning
- Changing of addresses in session in network layer
 - Sign change with private key associated with old address
- New types of intrinsic security that might
 - Variants for services, contents and hosts; new types
 - Support for existing key management processes
- Simplify comprehensive security mechanisms

11

Supporting Evolvability: Flexible Addressing

- Introduction of a new principal type will be incremental – no “flag day”!
 - Not all routers and ISPs will provide support from day one
- Creates chicken and egg problem - what comes first: network support or use in applications
- Solution: provide an *intent* and *fallback* address
 - Intent address allows in-network optimizations based on user intent
 - Fallback address is guaranteed to be reachable



12

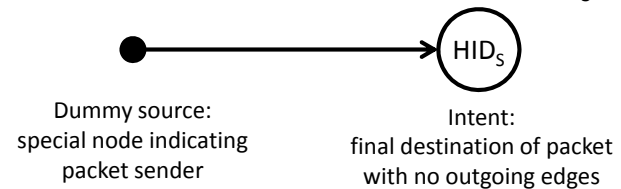
Addressing Requirements

- Fallback: intent that may not be globally understood must include a backwards compatible address
 - Incremental introduction of new XID types
- Scoping: support reachability for non-globally routable XID types or XIDs
 - Needed for scalability
 - Generalize scoping based on network identifiers
 - But we do not want to give up leveraging intent
- Iterative refinement: give each XID in the hierarchy option of using intent

13

Our Solution: DAG-Based Addressing

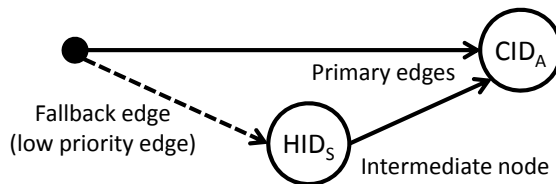
- Uses direct acyclic graph (DAG)
 - Nodes: typed IDs (XID; expressive identifier)
 - Outgoing edges: possible routing choices
- Simple example: Sending a packet to HID_S



14

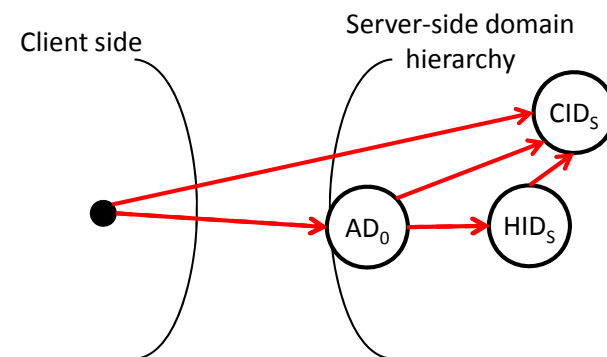
Support for Fallbacks with DAG

- A node can have **multiple outgoing edges**
 - Primary edges
 - Fallback edge (low priority edge)
- Outgoing edges have **priority** among them
 - Forwarding to HID_S is attempted if forwarding to CID_A is not possible – Realization of fallbacks



15

DAGs Support Scoping and Iterative Refinement



"XIA: Efficient Support for Evolvable Internetworking", NSDI 2012

16

Open Source XIA Release

<https://github.com/xia-project/>

- XIA Prototype released in May 2012
 - Includes full XIA protocol stack and utilities
- Being used to support evaluation, applications, services
- New functionality is being added regularly

17

Ongoing Networking Research

- Transport protocols: congestion control, error recovery
- Intrinsic security and mobility, ...
- Incremental deployment of network architectures (features)
- Verification of third party services using TPMs
- Very fast lookup of flat IDs in huge tables
- Optimize use of network features under user control
- Native Unix XIA implementation – extreme evolvability
- Design of a network control plane
- Supporting DTNs, pub-sub systems, group communication, ...
- Routing and forwarding for services, content
- Network diagnostics, centralized versus distributed control
- Video streaming as a use case for XIA
- Economic incentives and implications of cryptographic identifiers
- Balancing user accountability and privacy

18

Broader XIA Research Agenda

Trust Management (Adrian Perrig, Dave Andersen)

User Studies (Sara Kiesler)

Policy and Economics (Marvin Sirbu, Jon Peha, Jon Byers)

Core Network (Aditya Akella, John Byers, Peter Steenkiste, Dave Andersen, Srinu Seshan, Hui Zhang, Bruce Maggs)

19

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20

Examples of XIA-related Research

- The Accountable Internet Protocol
 - Accountable Internet Protocol (AIP). David Andersen, et al, *ACM SIGCOMM 2008*
 - Example of a protocol that provides accountability for host-based communication
- The Accountable and Private Internet Protocol
 - Balancing Accountability and Privacy (APIP). David Naylor, et al, *ACM SIGCOMM 2014*
 - Expands on AIP to support user privacy

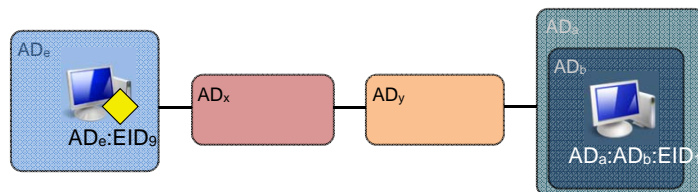
21

AIP Motivation

- Many security challenges are a result of not being able to unambiguously determine who is responsible for a specific action
 - Source spoofing, DOS attacks, untraceable spam, ..
- Add accountability to the Internet architecture
- Key idea is to use “self-certifying” addresses for both hosts and domains
- Avoid dependence on external configurations
 - E.g. global trust authority

22

Addressing and Routing



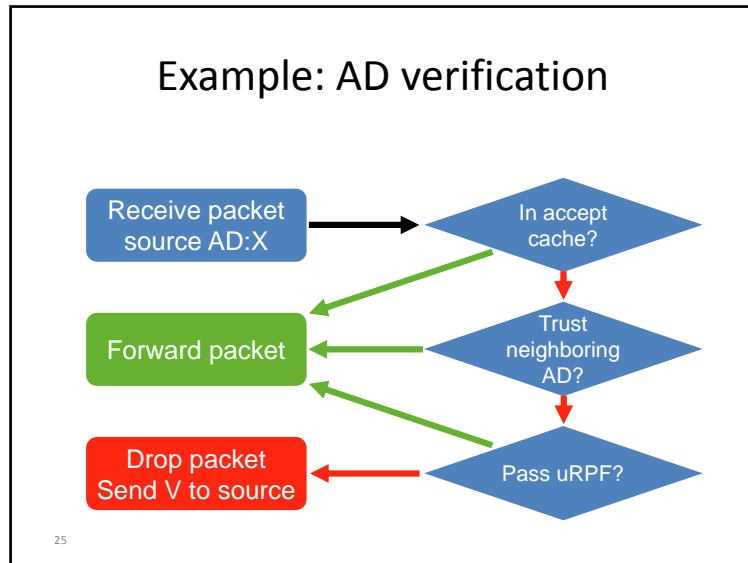
- Addresses are hierarchical, similar to today’s Internet
 - But each level has a flat address, i.e. no CIDR
- Until packet reaches destination AD, intermediate routers use only destination AD to forward packet
 - Effectively uses a pointer in a stack of domain identifiers
- Upon reaching destination AD, forward based on EID

23

Self-Certifying Identifiers

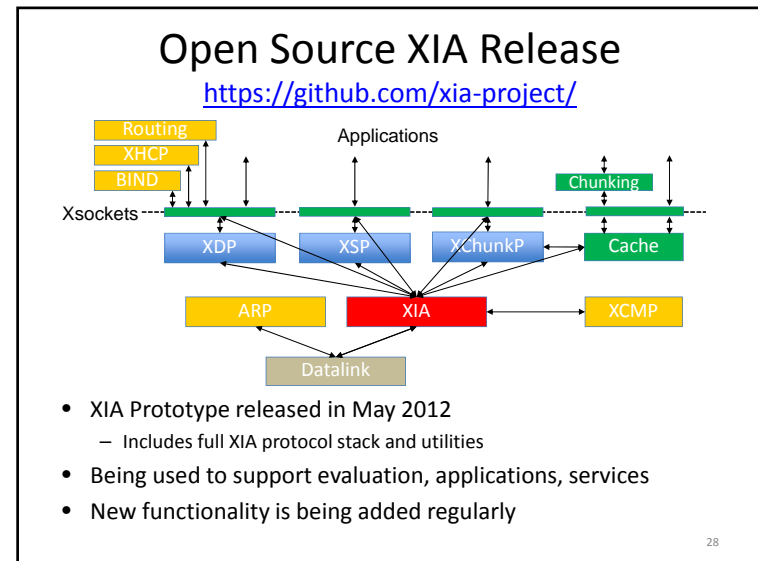
- Identifier of object is public key of object
 - Convenient to use hash of object (e.g. fixed size)
 - Need way of securely mapping user readable name into the identifier
- AD is hash of public key of domain
- EID is hash of public key of host
- Provides a means of verifying the correctness of the “source” identifiers in a packet
 - Effectively by sending a challenge to the source that it must sign with its private key

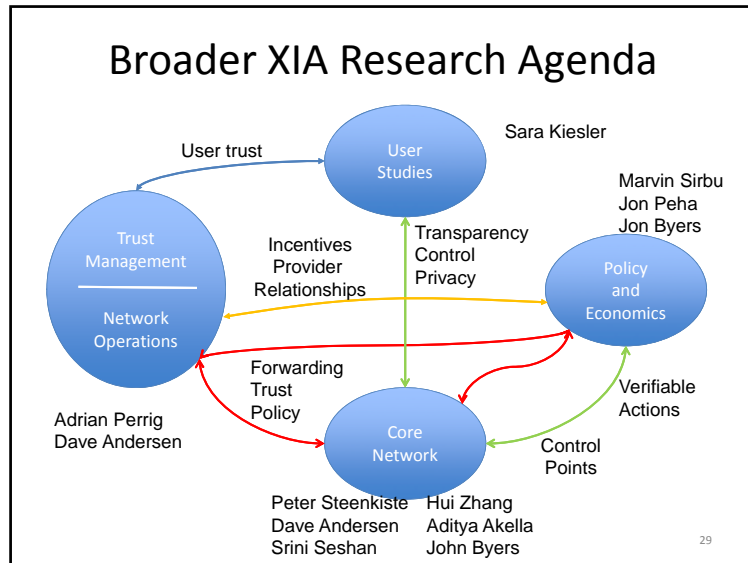
24



- ### Verification Packet
- Router sends a challenge V to Source containing:
 - Source and destination identifier
 - Hash of the packet P
 - Interface of the router
 - A secret signed by R
 - Source signs V with its private key and sends it back to R
 - But only if it recognizes the hash
 - R verifies that it was signed correctly using the public key from the source field
 - If they match, R add S to its cache
- 26

- ### AIP Discussion
- AIP adds complexity to routers ...
 - Crypto support, caches, larger forwarding tables, ..
 - ... but accountability helps address number of security challenges
 - Reduces complexity and cost in rest of networks
 - Research question
 - Fast look up in large tables of flat identifiers
 - Managing keys (revocation, minting, ...)
 - Evolving of the crypto
- 27





- ### Growing User Concern about Privacy
- Fueled by personal experience and reports, e.g., social networks, vendors, Snowden, ...
 - More privacy is always better?
 - Privacy can be expensive
 - Obvious example: strong anonymity using TOR
 - More subtle costs associated with HTTPS
 - “The Cost of “S” in HTTPS”, Naylor et. al., ACM CoNext, Dec 2014
 - Lack of accountability
 - AIP provides accountability – price is loss of privacy
 - TOR is the other way around!

Source Addresses, or Balancing Privacy and Accountability

- Source address are assumed to be essential but you can build a network without them
- What are source addresses used for?

Hard to balance Privacy and Accountability:

Tor versus AIP

“Tussle” controlled by on/off switch

~~Return address~~

~~Identify sender~~

Accountability

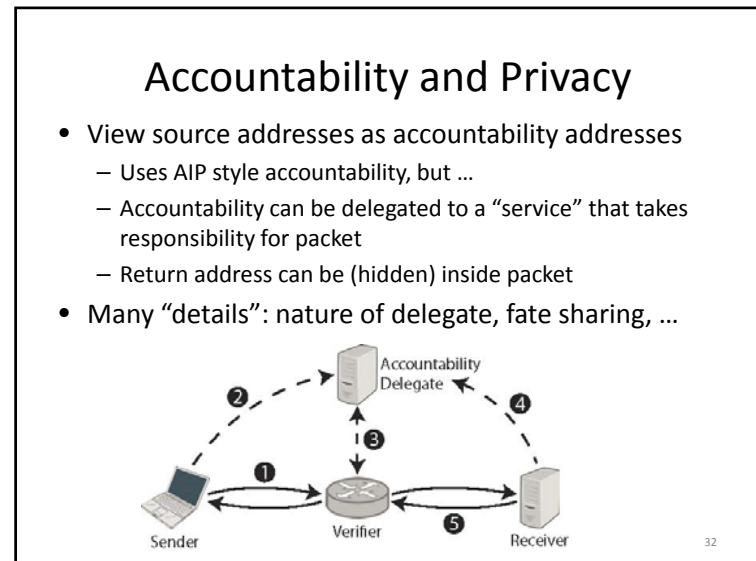
Error reporting

Flow ID

Used by:

Destination

Network



XIA Project

- More information:
 - <http://www.cs.cmu.edu/~xia>
- XIA faculty
 - Peter Steenkiste, CS/ECE, Carnegie Mellon
 - Dave Andersen, David Eckhardt, Srinu Seshan, Hui Zhang, CS, Carnegie Mellon
 - Sara Kiesler, HCII, Carnegie Mellon
 - Jon Peha, Marvin Sirbu, EPP, Carnegie Mellon
 - Adrian Perrig, ETH/Carnegie Mellon
 - Aditya Akella, CS, University of Wisconsin
 - John Byers, CS, Boston University
 - Bruce Maggs, CS, Duke





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34