15-441 Computer Networking

Lecture 21 – Queue Management and Quality of Service

Peter Steenkiste

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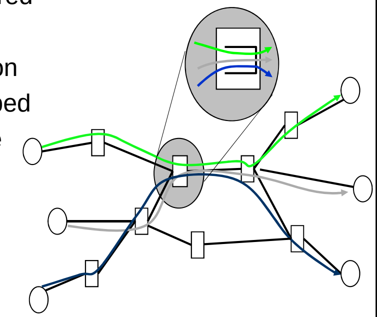
www.cs.cmu.edu/~prs/15-441-F10

Traffic and Resource Management

- Resources statistically shared
  $\sum \mathbf{Demand}_i(\mathbf{t}) > \mathbf{Resource}(\mathbf{t})$
- Overload causes congestion
  - packet delayed or dropped
  - application performance suffer
- Local vs. network wide
- Transient vs. persistent
- Challenge
  - high resource utilization
  - high application performance



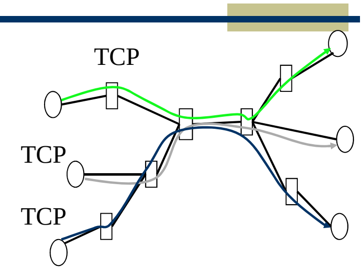
Resource Management Approaches

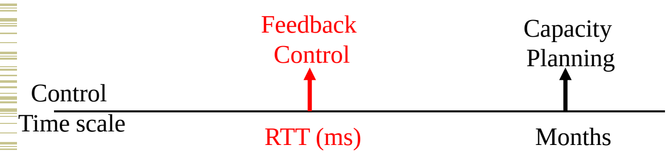
$\sum \mathbf{Demand}_i(\mathbf{t}) > \mathbf{Resource}(\mathbf{t})$

- Increase resources
  - install new links, faster routers
  - capacity planning, provisioning, traffic engineering
  - happen at longer timescale
- Reduce or delay demand
  - Reactive approach: encourage everyone to reduce or delay demand
  - Reservation approach: some requests will be rejected by the network

Congestion Control in Today's Internet

- End-system-only solution (TCP)
  - dynamically estimates network state
  - packet loss signals congestion
  - reduces transmission rate in presence of congestion
  - routers play little role



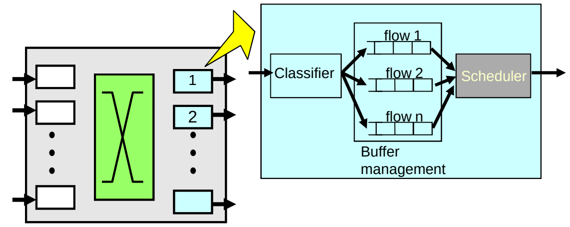


How to Improve Traffic Management

- Improve TCP
  - Stay with end-point only architecture
- Enhance routers to help TCP
  - Random Early Discard
- Enhance routers to control traffic
  - Rate limiting
  - Fair Queueing
- Provide QoS by limiting congestion

Router Mechanisms

- Buffer management: when and which packet to drop?
- Scheduling: which packet to transmit next?



Overview

- Queue management & RED
- Fair-queuing
- Why QOS?
- Integrated services

Queuing Disciplines

- Each router must implement some queuing discipline
- Queuing allocates both bandwidth and buffer space:
  - Bandwidth: which packet to serve (transmit) next
  - Buffer space: which packet to drop next (when required)
- Queuing also affects latency

Typical Internet Queuing

- FIFO + drop-tail
  - Simplest choice - used widely in the Internet
- FIFO (first-in-first-out)
  - Implies single class of traffic
  - All packets treated equally
- Drop-tail
  - Arriving packets get dropped when queue is full regardless of flow or importance
- Important distinction:
  - FIFO: scheduling discipline
  - Drop-tail: drop policy

9

FIFO + Drop-tail Problems

- Leaves responsibility of congestion control completely to the edges (e.g., TCP)
- Does not separate between different flows
- No policing: send more packets → get more service
- Synchronization: multiple end hosts may react to same events at about the same time
  - E.g. due to busts of packet loss

10

FIFO + Drop-tail Problems

- Full queues
  - Routers are forced to have have large queues to maintain high utilizations
  - TCP detects congestion from loss
    - Forces network to have long standing queues in steady-state
    - May no longer be able to absorb bursts
- Lock-out problem
  - Drop-tail routers treat bursty traffic poorly
  - Traffic gets synchronized easily → allows a few flows to monopolize the queue space

11

Active Queue Management

- Design active router queue management to aid congestion control
- Why?
- Router has unified view of queuing behavior
  - Routers see actual queue occupancy (distinguish queue delay and propagation delay)
  - Routers can decide on transient congestion, based on workload

12

Design Objectives

- Keep throughput high and delay low
  - High power (throughput/delay)
- Accommodate bursts
- Queue size should reflect ability to accept bursts rather than steady-state queuing
- Improve TCP performance with minimal hardware changes

Lock-out Problem

- Random drop
  - Packet arriving when queue is full causes some random packet to be dropped
- Drop front
  - On full queue, drop packet at head of queue
- Random drop and drop front solve the lock-out problem but not the full-queues problem

Full Queues Problem

- Drop packets before queue becomes full (early drop)
- Intuition: notify senders of incipient congestion
- Example: early random drop (ERD):
  - If qlen > drop level, drop new packets with fixed probability *p*
  - Does not control misbehaving users

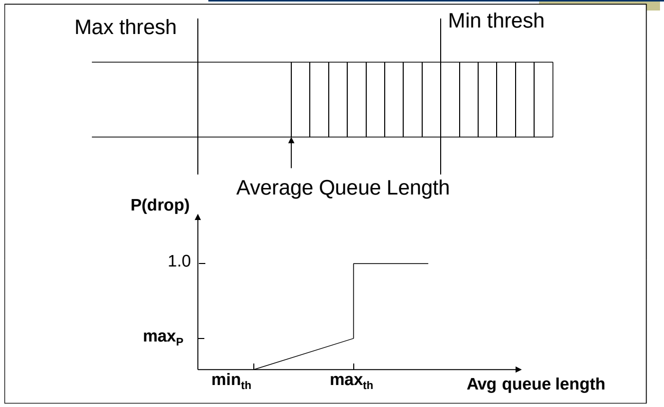
Random Early Detection (RED)

- Detect incipient congestion
- Assume hosts respond to lost packets
- Avoid window synchronization
  - Randomly mark packets
- Avoid bias against bursty traffic

RED Algorithm

- Maintain running average of queue length
- If avg < min_{th} do nothing
  - Low queuing, send packets through
- If avg > max_{th}, drop packet
  - Protection from misbehaving sources
- Else drop packet in a manner proportional to queue length
  - Notify sources of incipient congestion
  - Can instead mark the packet if ECN is used (later)

RED Operation



Explicit Congestion Notification (ECN) [Floyd and Ramakrishnan 98]

- Traditional mechanism: packet drop as implicit congestion signal to end systems
  - TCP will slow down
- Works well for bulk data transfer
- Does not work well for delay sensitive applications
  - audio, WEB, telnet
- Explicit Congestion Notification (ECN)
  - borrow ideas from DECBit
  - use two bits in IP header
    - ECN-Capable Transport (ECT) bit set by sender
    - Congestion Experienced (CE) bit set by router

Congestion Control Summary

- Basis: end system detects congestion and slow down
  - slow start/congestion avoidance
    - packet drop detected by retransmission timeout RTO as congestion signal
    - AIMD: search for right bandwidth, considering fairness
- TCP Improvement:
  - fast retransmission/fast recovery
    - packet drop detected by three duplicate acks
  - SACK: better feedback to source
- Router support
  - RED: early signaling
  - ECN: explicit signaling

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21

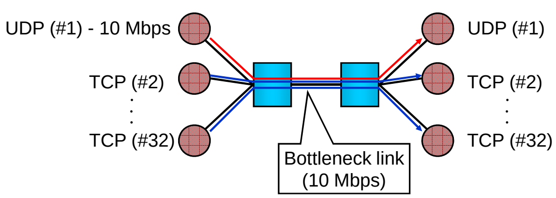
Problems to achieving fairness

- Works only if most sources implement TCP
  - most sources are cooperative
  - most sources implement homogeneous/compatible control law
    - compatible means less aggressive than TCP
- What if sources do not play by the rule?

22

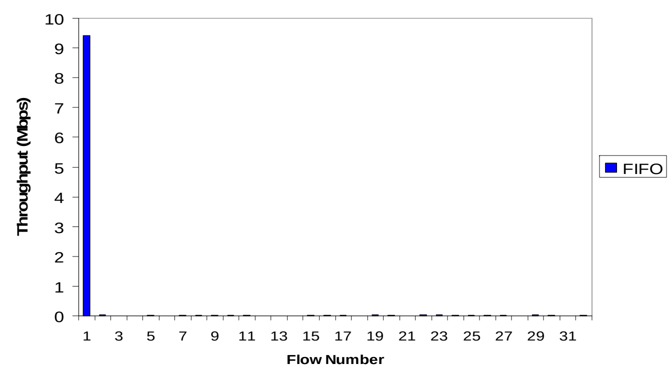
An Example

- 1 UDP (10 Mbps) and 31 TCPs sharing a 10 Mbps line



23

Throughput of UDP and TCP Flows With FIFO



24

Fairness Goals

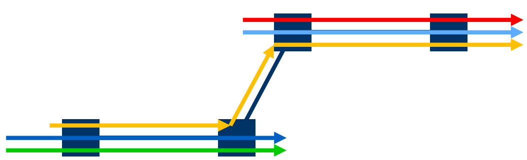
- Allocate resources fairly
- Isolate ill-behaved users
  - Router does not send explicit feedback to source
  - Still needs e2e congestion control
- Still achieve statistical multiplexing
  - One flow can fill entire pipe if no contenders
  - Work conserving → scheduler never idles link if it has a packet

What is Fairness?

- At what granularity?
  - Flows, connections, domains?
- What if users have different RTTs/links/etc.
  - Should it share a link fairly or be TCP fair?
- Maximize fairness index?
  - Fairness = $(\Sigma x_i)^2/n(\Sigma x_i^2)$ $0<\text{fairness}<1$
- Basically a tough question to answer!
- Good to separate the design of the mechanisms from definition of a policy
  - User = arbitrary granularity

Max-min Fairness

- Allocate user with "small" demand what it wants, evenly divide unused resources to "big" users
- Formally:
  - Resources allocated in terms of increasing demand
  - No source gets resource share larger than its demand
  - Sources with unsatisfied demands get equal share of resource



Implementing Max-min Fairness

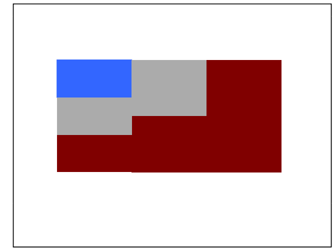
- Generalized processor sharing
  - Fluid fairness
  - Bitwise round robin among all queues
- Why not simple round robin?
  - Variable packet length → can get more service by sending bigger packets
  - Unfair instantaneous service rate
    - What if arrive just before/after packet departs?

Bit-by-bit RR

- Single flow: clock ticks when a bit is transmitted. For packet i:
  - P_i = length, A_i = arrival time, S_i = begin transmit time, F_i = finish transmit time
  - $F_i = S_i+P_i = \max(F_{i-1}, A_i) + P_i$
- Multiple flows: clock ticks when a bit from all active flows is transmitted → round number
  - Can calculate F_i for each packet if number of flows is know at all times
    - Why do we need to know flow count? → need to know A → This can be complicated

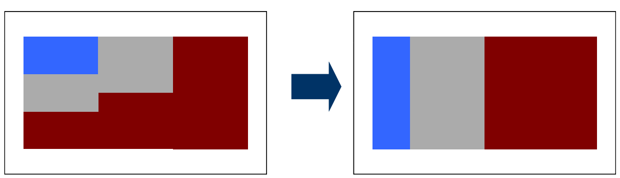
Bit-by-bit RR Illustration

- Not feasible to interleave bits on real networks
  - FQ simulates bit-by-bit RR

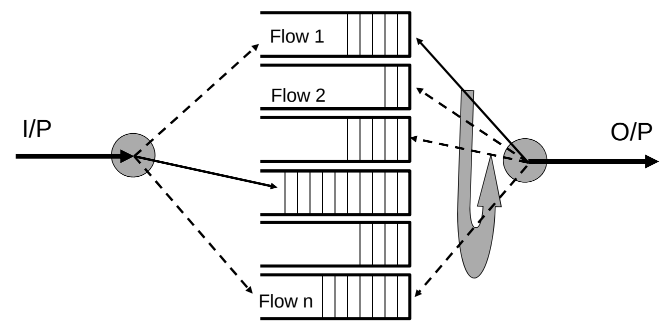


Fair Queuing

- Mapping bit-by-bit schedule onto packet transmission schedule
- Transmit packet with the lowest F_i at any given time
  - How do you compute F_i?

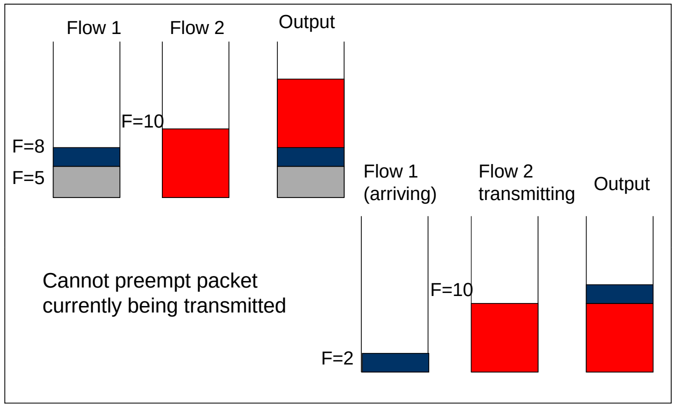


FQ Illustration



Variation: Weighted Fair Queuing (WFQ)

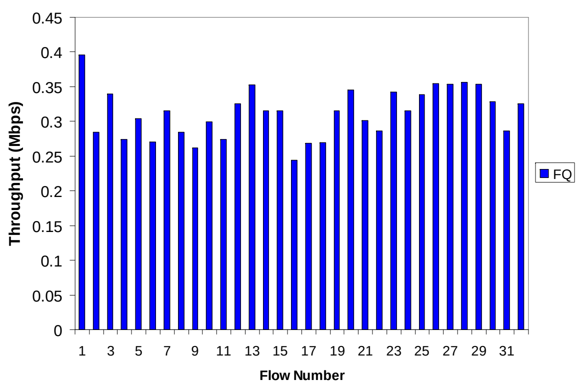
Bit-by-bit RR Example



Fair Queuing Tradeoffs

- Complex computation
  - Classification into flows may be hard
  - Must keep queues sorted by finish times
  - dR/dt changes whenever the flow count changes
- Complex state
  - Must keep queue per flow
    - Hard in routers with many flows (e.g., backbone routers)
    - Flow aggregation is a possibility (e.g. do fairness per domain)
- FQ can control congestion by monitoring flows
  - Non-adaptive flows can still be a problem – why?

Example: Throughput of TCP and UDP Flows With Fair Queueing Router



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- Fair-queuing
- Why QOS?
- Integrated services

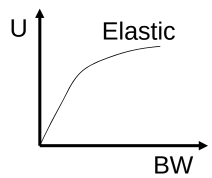
Motivation

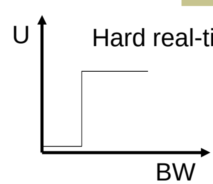
- Internet currently provides one single class of **“best-effort” service**
  - No assurances about delivery
  - All users (packets?) are equal
- At internet design most applications are ***elastic***
  - Tolerate delays and losses
  - Can adapt to congestion
- Today, many “real-time” applications are ***inelastic***

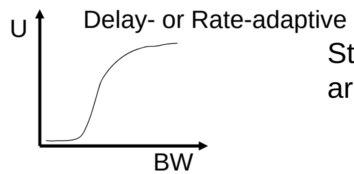
Why a New Service Model?

- What is the **basic objective** of network design?
  - Maximize total bandwidth? Minimize latency?
  - **Maximize user satisfaction** – the total **utility** given to users
- What does utility vs. bandwidth look like?
  - Shape depends on application
  - Must be non-decreasing function

Utility Curve Shapes

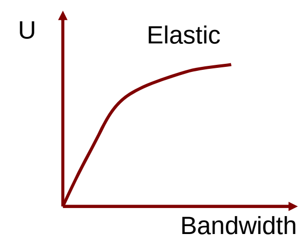






Stay to the right and you are fine for all curves

Utility curve – Elastic traffic

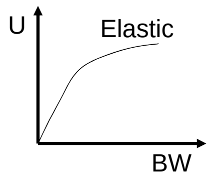


Does equal allocation of bandwidth maximize total utility?

Admission Control

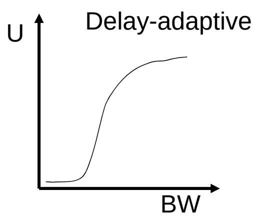
- If U(bandwidth) is concave
  → elastic applications
  - Incremental utility is decreasing with increasing bandwidth
  - Is always advantageous to have more flows with lower bandwidth
    - No need of admission control;

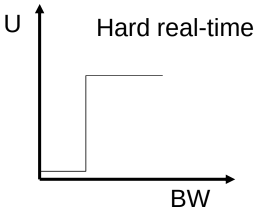
  This is why the Internet works!



41

Utility Curves – Inelastic traffic





Does equal allocation of bandwidth maximize total utility?

42

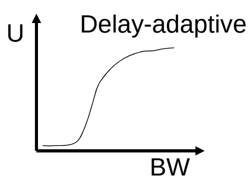
Inelastic Applications

- Continuous media applications
  - **Lower and upper limit** on acceptable performance.
  - BW below which video and audio are not intelligible
  - Internet telephones, teleconferencing with high delay (200 - 300ms) impair human interaction
  - Sometimes called "tolerant real-time" since they can adapt to the performance of the network
- Hard real-time applications
  - Require **hard limits on performance**
  - E.g. control applications

43

Admission Control

- If U is convex → inelastic applications
  - U(number of flows) is no longer monotonically increasing
  - Need admission control to maximize total utility
- **Admission control** → deciding when adding more people would reduce overall utility
  - Basically avoids overload



44

Overview

- Queue management & RED
- Fair-queuing
- Why QOS?
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Components of Integrated Services

1. Type of commitment
   What does the network promise?
2. Packet scheduling
   How does the network meet promises?
3. Service interface
   How does the application describe what it wants?
4. Establishing the guarantee
   How is the promise communicated to/from the network
   How is admission of new applications controlled?

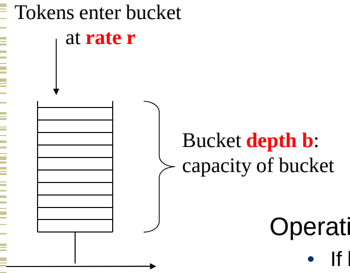
Type of Commitments

- **Guaranteed** service
  - For **hard real-time** applications
  - Fixed guarantee, network meets commitment if clients send at agreed-upon rate
- **Predicted** service
  - For **delay-adaptive** applications
  - Two components
    - If conditions do not change, commit to current service
    - If conditions change, take steps to deliver consistent performance (help apps minimize playback delay)
    - Implicit assumption – network does not change much over time
- **Datagram/best effort service**

Scheduling for Guaranteed Traffic

- Use **token bucket filter** to characterize traffic
  - Described by rate *r* and bucket depth *b*
- Use **Weighted Fair-Queueing** at the routers
- Parekh's bound for worst case queuing delay = b/r

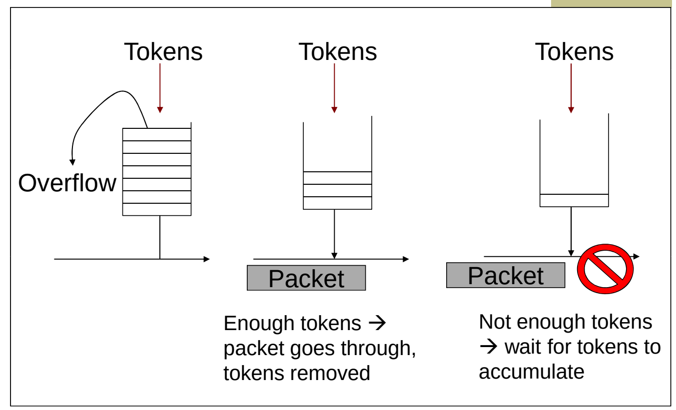
Token Bucket Filter



Operation:

- If bucket fills, tokens are discarded
- Sending a packet of size P uses P tokens
- If bucket has P tokens, packet sent at max rate, else must wait for tokens to accumulate

Token Bucket Operation

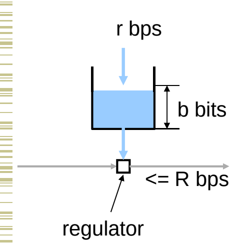


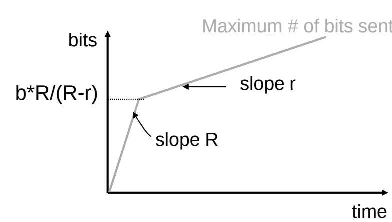
Token Bucket Characteristics

- On the long run, rate is limited to r
- On the short run, a burst of size b can be sent
- Amount of traffic entering at interval T is bounded by:
  - Traffic = b + r*T
- Information useful to admission algorithm

Token Bucket

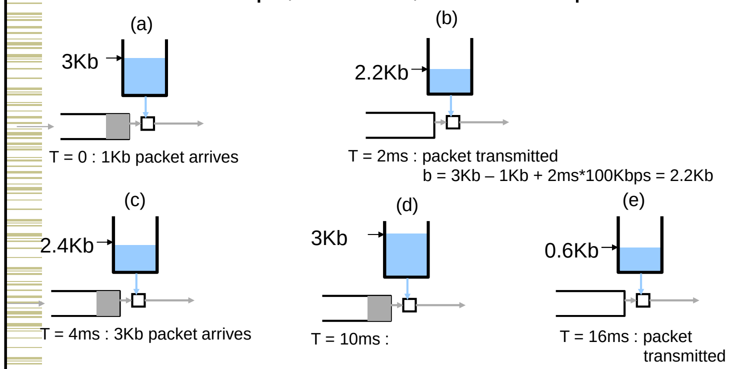
- Parameters
  - r – average rate, i.e., rate at which tokens fill the bucket
  - b – bucket depth
  - R – maximum link capacity or peak rate (optional parameter)
- A bit is transmitted only when there is an available token



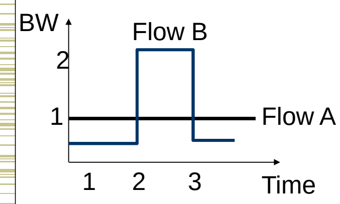


Traffic Enforcement: Example

- r = 100 Kbps; b = 3 Kb; R = 500 Kbps



Token Bucket Specs - Example



Flow A: r = 1 MBps, B=1 byte

Flow B: r = 1 MBps, B=1MB

Guarantee Proven by Parekh

- Given:
  - Flow *i* shaped with token bucket and leaky bucket rate control (depth *b* and rate *r*)
  - Network nodes do WFQ
- Cumulative queuing delay D_i suffered by flow *i* has upper bound
  - **D_i < b/r,** (where r may be much larger than average rate)
  - Assumes that Σr < link speed at any router
  - All sources limiting themselves to *r* will result in no network queuing

Sharing versus Isolation

- Impact of queueing mechanisms:
  - Isolation: Isolates well-behaved from misbehaving sources
  - Sharing: Mixing of different sources in a way beneficial to all
- FIFO: sharing
  - each traffic source impacts other connections directly
    - e.g. malicious user can grab extra bandwidth
  - the simplest and most common queueing discipline
  - averages out the delay across all flows
- Priority queues: one-way sharing
  - high-priority traffic sources have impact on lower priority traffic only
  - has to be combined with admission control and traffic enforcement to avoid starvation of low-priority traffic
- WFQ: two-way isolation
  - provides a guaranteed minimum throughput (and maximum delay)

Putting It All Together

- Assume 3 types of traffic: guaranteed, predictive, best-effort
- Scheduling: use WFQ in routers
- Each guaranteed flow gets its own queue
- All predicted service flows and best effort aggregates in single separate queue
  - Predictive traffic classes
    - Worst case delay for classes separated by order of magnitude
    - When high priority needs extra bandwidth – steals it from lower class
  - Best effort traffic acts as lowest priority class

Service Interfaces

- Guaranteed Traffic
  - Host specifies rate to network
  - Why not bucket size b?
    - If delay not good, ask for higher rate
- Predicted Traffic
  - Specifies (r, b) token bucket parameters
  - Specifies delay D and loss rate L
  - Network assigns priority class
  - Policing at edges to drop or tag packets
    - Needed to provide isolation – why is this not done for guaranteed traffic?
      - WFQ provides this for guaranteed traffic

Lessons

- TCP can use help from routers
  - RED → eliminate lock-out and full-queues problems
  - FQ → heavy-weight but explicitly fair to all
- QoS
  - What type of applications are there? → Elastic, adaptive real-time , and hard real-time.
  - Why do we need admission control → to maximize utility
  - How do token buckets + WFQ provide QoS guarantees?