

Unifying Logical Concurrency Control and Physical Data Structure Maintenance in Database Management Systems – CIDR 2021

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Motivation



Arrays Linked Lists Hash Tables Trees

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Motivation

Transaction BEGIN; DELETE FROM table WHERE key = a; COMMIT;

Transaction BEGIN;

SELECT FROM table
WHERE key = a;
COMMIT;



Arrays Linked Lists Hash Tables Trees

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Motivation

Transaction

BEGIN;

DELETE FROM table
WHERE key = a;
COMMIT;

Transaction

BEGIN; SELECT FROM table WHERE key = a; COMMIT;

Must **defer** the removal of physical data until no active txn able to see it.

Arrays Linked Lists Hash Tables Trees

Asynchronous execution framework for internal maintenance actions in a multi-versioned DBMS.

Key Idea: DAF executes actions as transactions with the same visibility mechanisms.

- Single API call: DEFER(action)
- Each action is tagged with txn's commit timestamp.
- Invoke an action when there are no transactions with a start timestamp smaller than tagged timestamp.

Transaction

BEGIN;

DELETE FROM table
WHERE key = a;
COMMIT;

Action Queue

 $\mathsf{OLDEST} \bigcirc = \mathsf{t0}$

Transaction

Action Queue

Transaction

Action Queue

OLDEST \bigcirc = t1

Transaction

DEFER(Unlink Version Chain DEFER(Delete Index Key=a

Generated Actions Action Queue

 $\mathsf{OLDEST} \bigcirc = \mathtt{t1}$

Action Queue

 $\mathsf{OLDEST} \bigcirc = \texttt{t2}$

Action Queue

 $\mathsf{OLDEST} \bigcirc = \texttt{t2}$

Action Queue

OLDEST \bigcirc = t3

Transaction

Transaction

Action Queue

Transaction

Transaction

Action Queue

OLDEST \bigcirc = t4 Unlink Version Chain \bigcirc = t3 Delete Index Key=a \bigcirc = t3

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Transaction BEGIN; $\bigcirc = t1$ DELETE FROM table WHERE key = a; COMMIT; $\bigcirc = t3$

bage 4

Transaction

Action Queue

Transaction		
BEGIN;	💬 = t1	
DELETE	FROM table	
WHERE	key = a;	
COMMIT;	💬 = t3	

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Transaction

Action Queue

Explicit ordering of concurrent actions.

- Example: action A drops a table; action B deletes a tuple in the table.
- Solution: Chained deferral of actions.
 DEFER(DEFER(... DEFER(action)...))

Unwrap and reinsert multi-deferred action to queue with a later timestamp.

Separate its execution with other concurrent actions.

Transaction

Action Queue

 $\mathsf{OLDEST} \bigcirc = \mathsf{t1}$

Transaction

DEFER(DEFER(Delete Table + Index)

Action Queue

 $\mathsf{OLDEST} \bigcirc = \mathsf{t1}$

Transaction BEGIN; $\bigcirc = t1$ DROP TABLE table; COMMIT;

Transaction

Action Queue

 $\mathsf{OLDEST} \bigcirc = \mathsf{t2}$

DEFER(DEFER(Delete Table + Index)

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Action Queue

 $\mathsf{OLDEST} \bigcirc = \mathsf{t2}$

Transaction

Transaction

Action Queue

 $\mathsf{OLDEST} \bigcirc \texttt{=} \texttt{t3}$

DEFER(Unlink Version Chain

DEFER Delete Index Key=a

DEFER (Delete Table + Index)

() = t3

Transaction

Transaction

Action Queue

OLDEST \bigcirc = t4 DEFER (Delete Table + Index) \bigcirc = t3 Unlink Version Chain \bigcirc = t4 Delete Index Key=a \bigcirc = t4

Transaction

Transaction

Action Queue

OLDEST \bigcirc = t5 DEFER (Delete Table + Index) \bigcirc = t3 Unlink Version Chain \bigcirc = t4 Delete Index Key=a \bigcirc = t4

Transaction

Transaction

Action Queue

OLDEST \bigcirc = t5 Unlink Version Chain \bigcirc = t4 Delete Index Key=a \bigcirc = t4 Delete Table + Index \bigcirc = t5

Transaction

Transaction

Action Queue

OLDEST \bigotimes = t5 Delete Index Key=a \bigotimes = t4 Delete Table + Index \bigotimes = t5

At this time the table is not visible to any active txn and is safe to drop.

Action Queue

Applications

Index Cleaning

Cache Invalidation

Data Transformation

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Preliminary Results

Integrated DAF with <u>NoisePage</u> DBMS.

Measure MVCC GC scalability using TPC-C:

- ►One warehouse per worker thread.
- Stored Procedure API.
- ►Write-Ahead Logging Enabled.

Compare three system configurations:
Single GC thread
Dedicated DAF threads
Cooperative DAF threads

MVCC GC Scalability (TPC-C)

→Single-GC →2 DAF →4 DAF →8 DAF →Coop-DAF

So Peloton

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9 10	+ ## UPDATE 2019-03-17 +		
11	11 + The Peloton project is **dead**. We have abandoned this repository and moved on to build a new DBMS. There are a several engineering job at implementing in the second system. 12 +		
13	+ We will not accept pull requests for this repository. We will also not respond to questions or problems that you may bay	0 with	

noise poge

NoisePage Project

In-Memory HTAP DBMS Postgres compatible (wire, SQL, catalog) Apache Arrow compatible columnar storage HyPer-style MVCC (snapshot isolation) Hybrid Vectorization + Pipeline Query Codegen JIT Query Compilation (DSL \rightarrow OpCodes \rightarrow LLVM) Integrated self-driving components

NoisePage Projec

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Summary

DAF shows how to leverage logical data concurrency protocols for physical data structures.

Unifying the notion of visibility makes it easier to integrate new data structures and extend transactions to support maintenance operations.

Acknowledgements

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