
Online Deduplication for Databases

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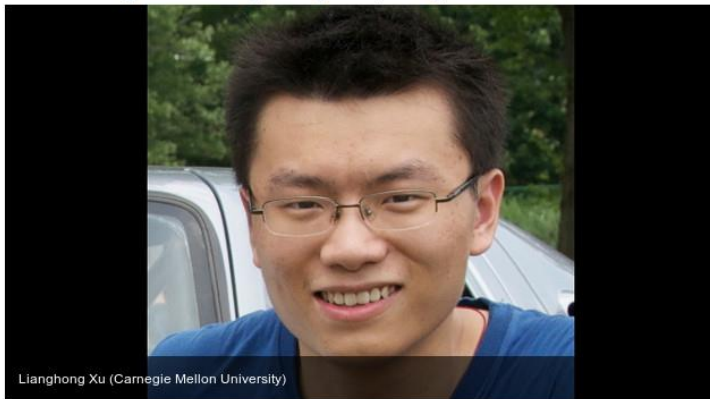
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CMU Database Student Stabbed By Rival DBAs

UPDATED | May 2, 2017 2:00 PM

May 2, 2017 1:54 PM

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Lianghong Xu (Carnegie Mellon University)

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PITTSBURGH (CBSDFW.COM) – Pittsburgh Police say an 26-year-old database PhD student at Carnegie Mellon University was stabbed during a fight with a gang of database administrators (DBAs). This gang has had a long-standing feud with the CMU Database Research group.

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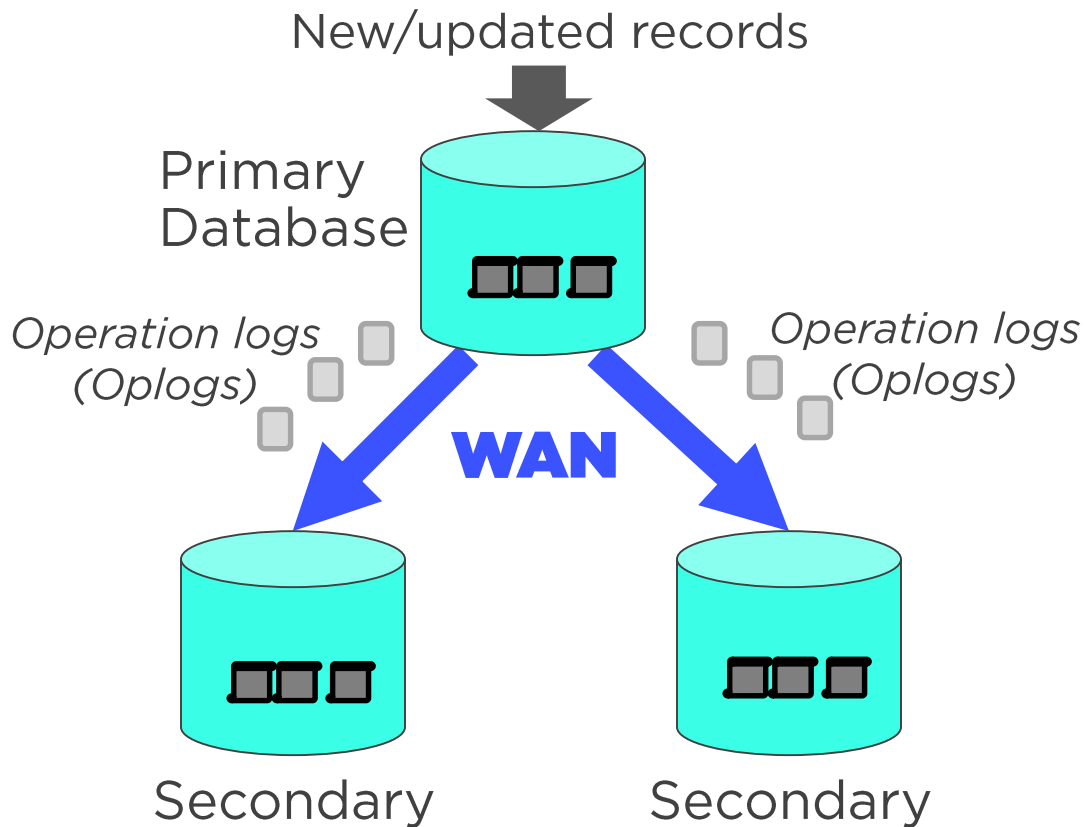
Dallas Police Officer Arrested

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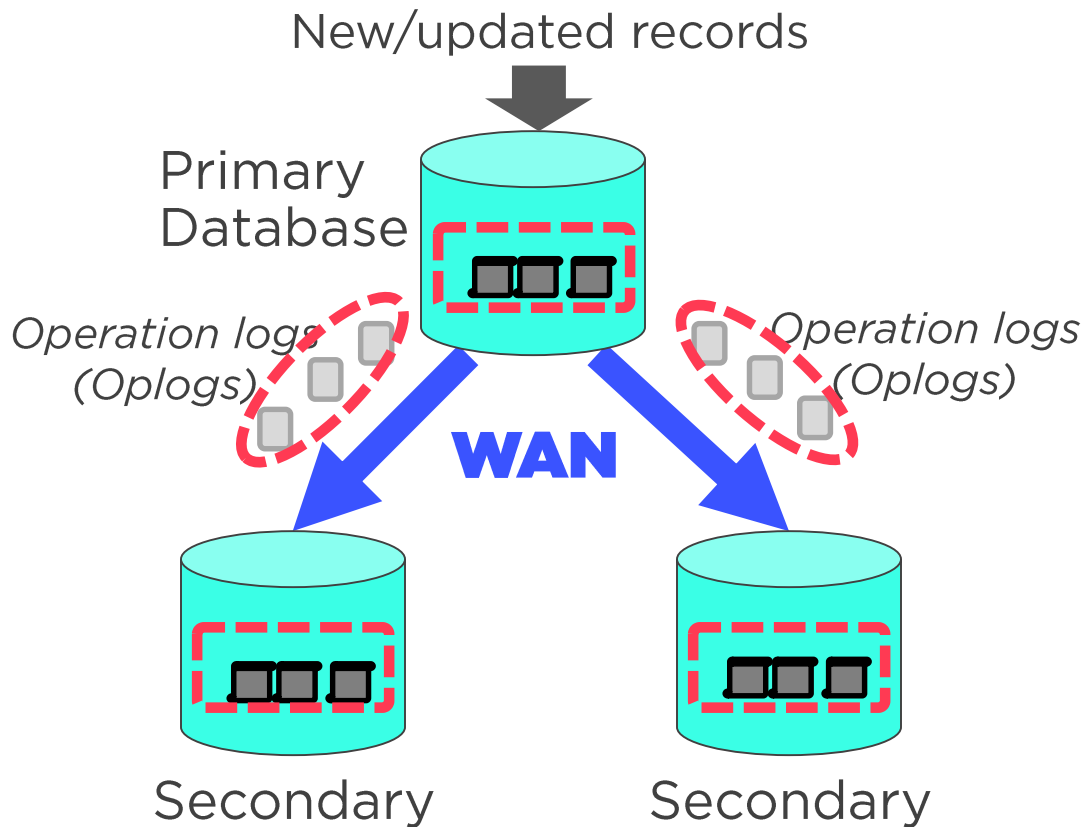


Police: Child Sexually

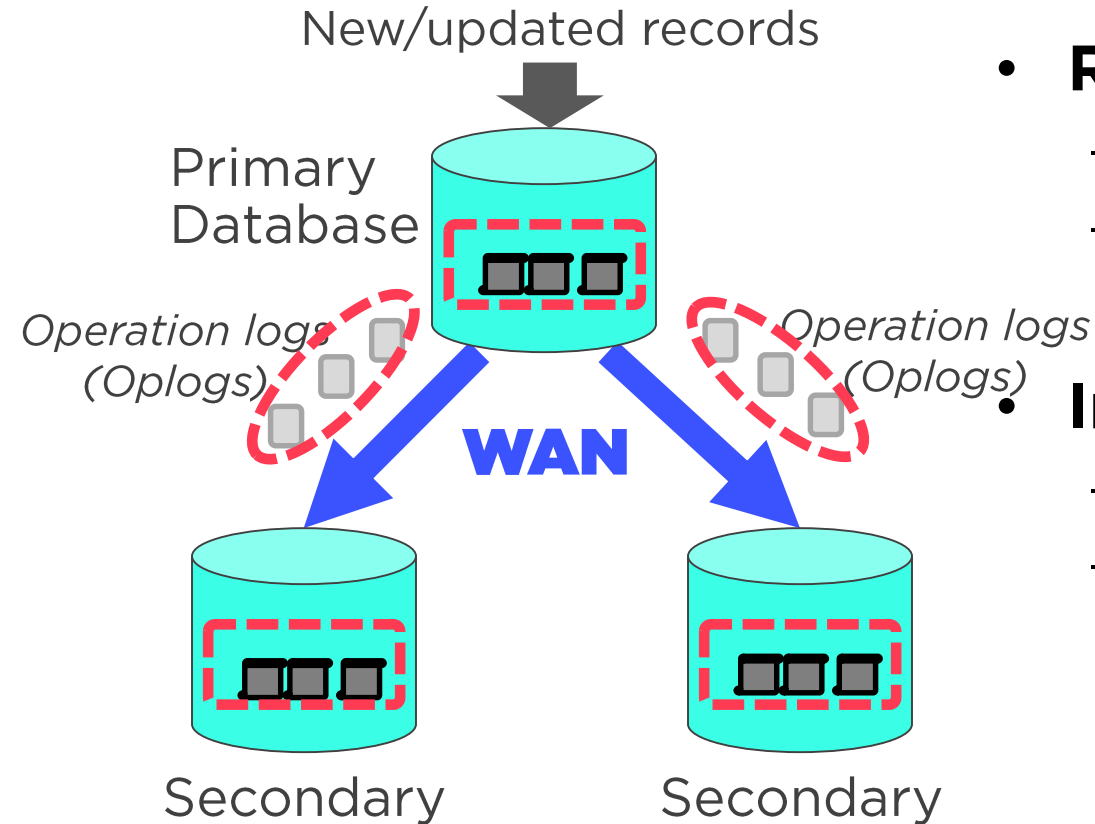
Goal: Data Reduction in DBMSs



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Goal: Data Reduction in DBMSs



- **Reduce cost**

- *Storage & network hardware*
- *Maintenance*

- **Improve Performance**

- *Reduce disk I/Os*
- *Lower replication bandwidth*

Current Reduction Approaches

- Block-level compression
 - *Simple, effective*
 - *Fails to address redundancy across blocks*
- Chunk-based Deduplication
 - *Identify global redundancy across data corpus*
 - *Not suitable for database workloads*
 - Records are relatively small
 - Modifications are small and dispersed

Similarity-based Dedup

- **sDedup** [Xu et al., SoCC'15]
 - *Use byte-level delta compression against similar records to achieve high compression ratio*
 - *Focused on reduction in network bandwidth for replication*
- **dbDedup** (this work)
 - *Achieve superior compression on both network and storage*
 - *Address challenges with accessing delta-encoded storage*
 - *Allow use of dedup in online DBMSs with low perf. overhead*

Outline

- Goal and Motivation
- **Need For Similarity-based Dedup**
- Our Approach: dbDedup
- Evaluation

Duplication Beyond A Block

- Application-level versioning
 - *Few apps take advantage of sys-level versioning*
 - *Revisions of data items seen as unrelated records*
 - *E.g., WordPress, Wikipedia*
- Partial record copying
 - *Inclusion between records not exposed to the DBMS*
 - *E.g., email reply/forward, quotes in message boards*

Traditional Dedup: Ideal Case

| *Chunk Boundary* ■ *Modified Region* ▨ *Duplicate Chunk*

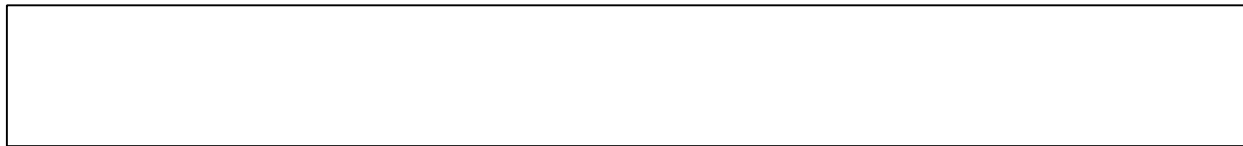
Incoming
Data

{BYTE STREAM}

Traditional Dedup: Ideal Case

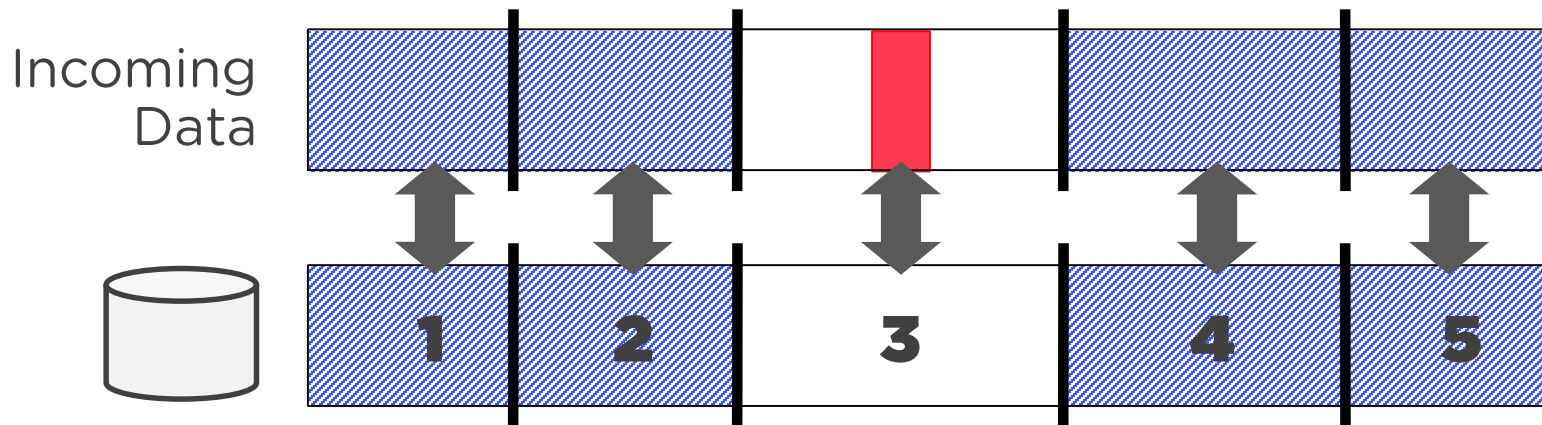
| *Chunk Boundary* ■ *Modified Region* ▨ *Duplicate Chunk*

Incoming
Data



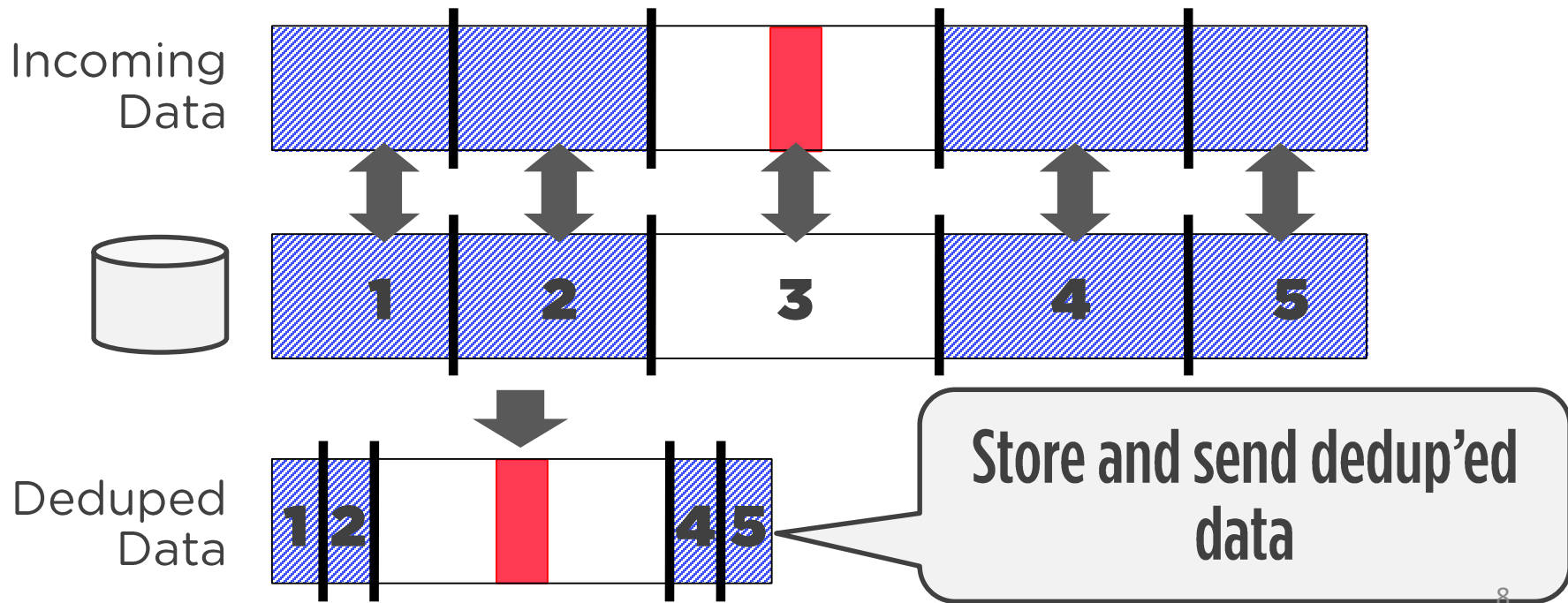
Traditional Dedup: Ideal Case

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Traditional Dedup: Ideal Case

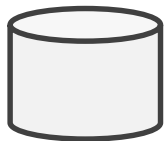
| *Chunk Boundary* █ *Modified Region* *Duplicate Chunk*



Dedup: Common Case in DBs

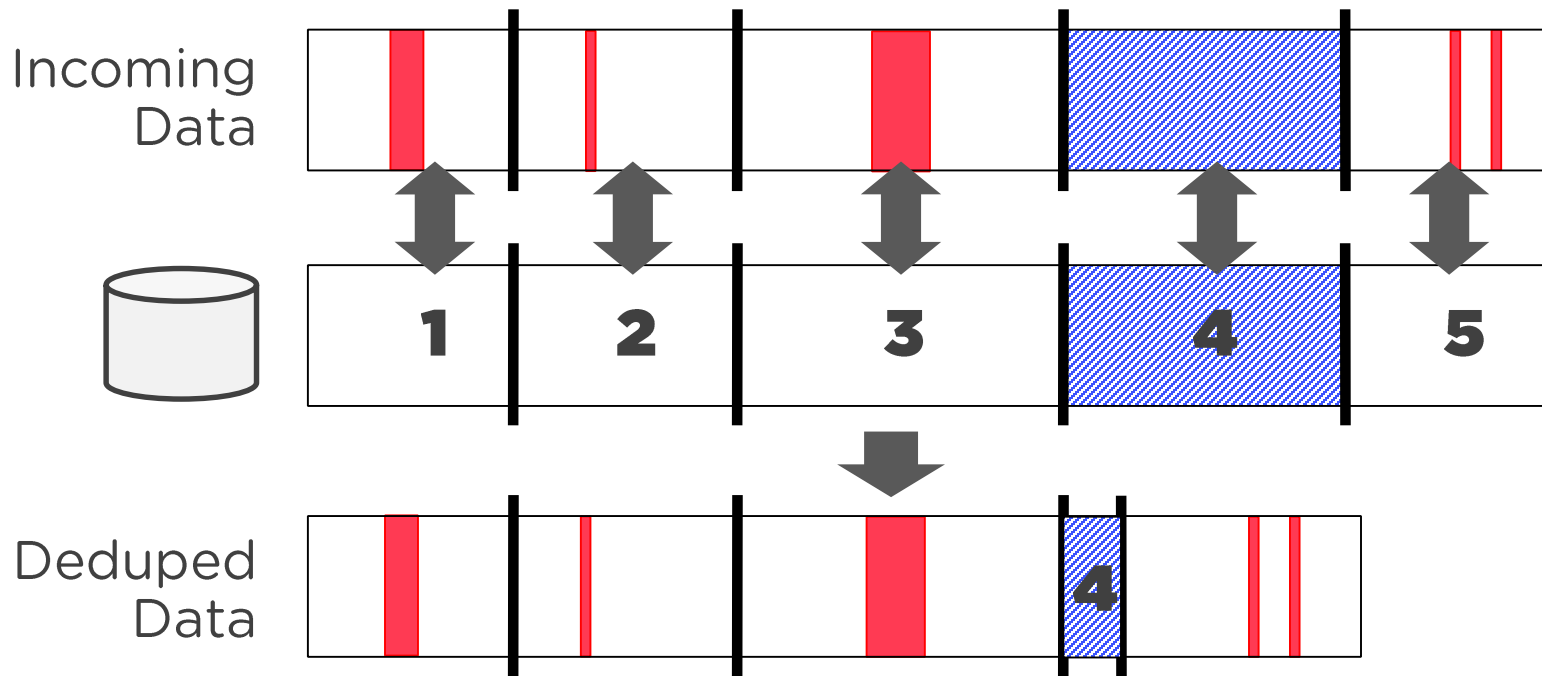
 *Chunk Boundary*  *Modified Region*  *Duplicate Chunk*

Incoming
Data



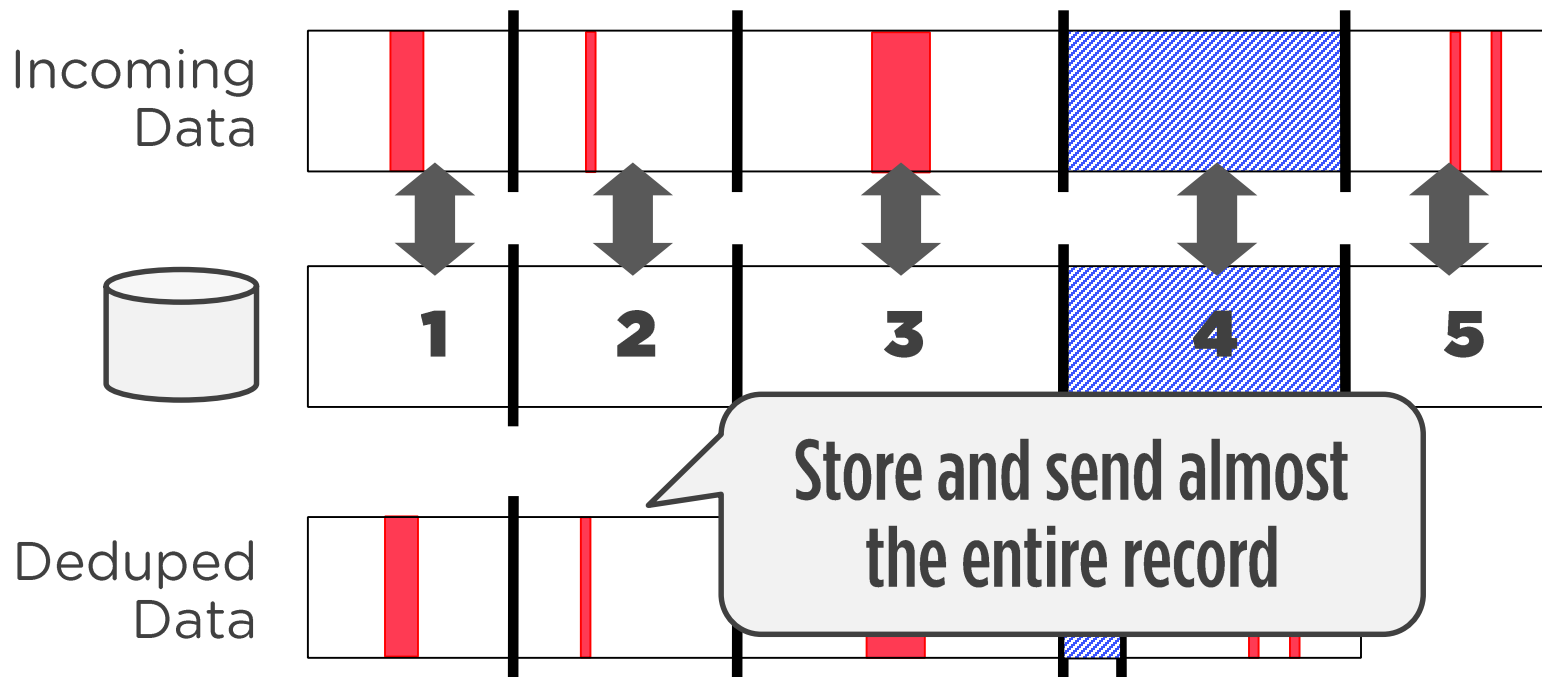
Dedup: Common Case in DBs

| *Chunk Boundary* *Modified Region* *Duplicate Chunk*



Dedup: Common Case in DBs

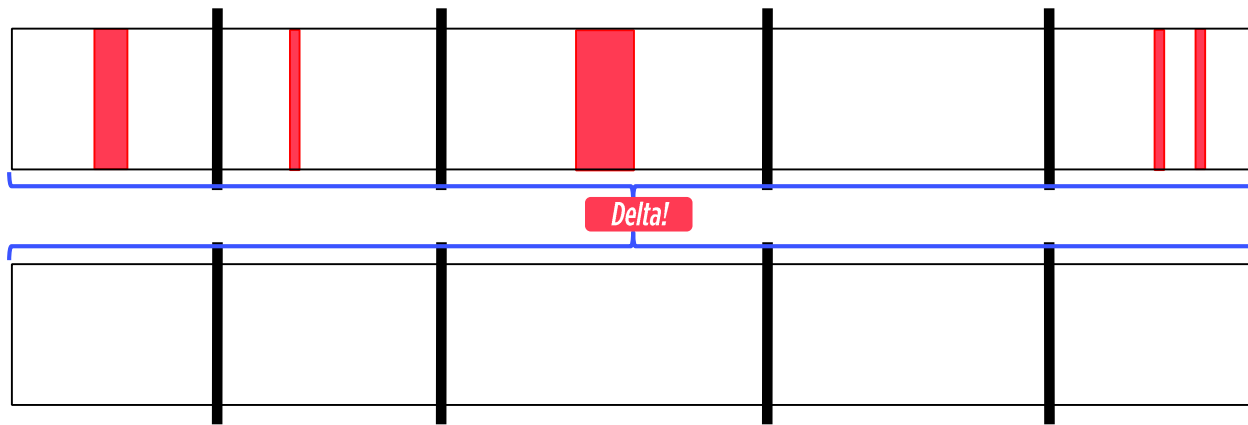
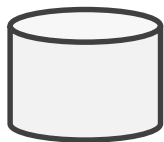
| *Chunk Boundary* █ *Modified Region* ▨ *Duplicate Chunk*



Similarity Dedup (dbDedup)

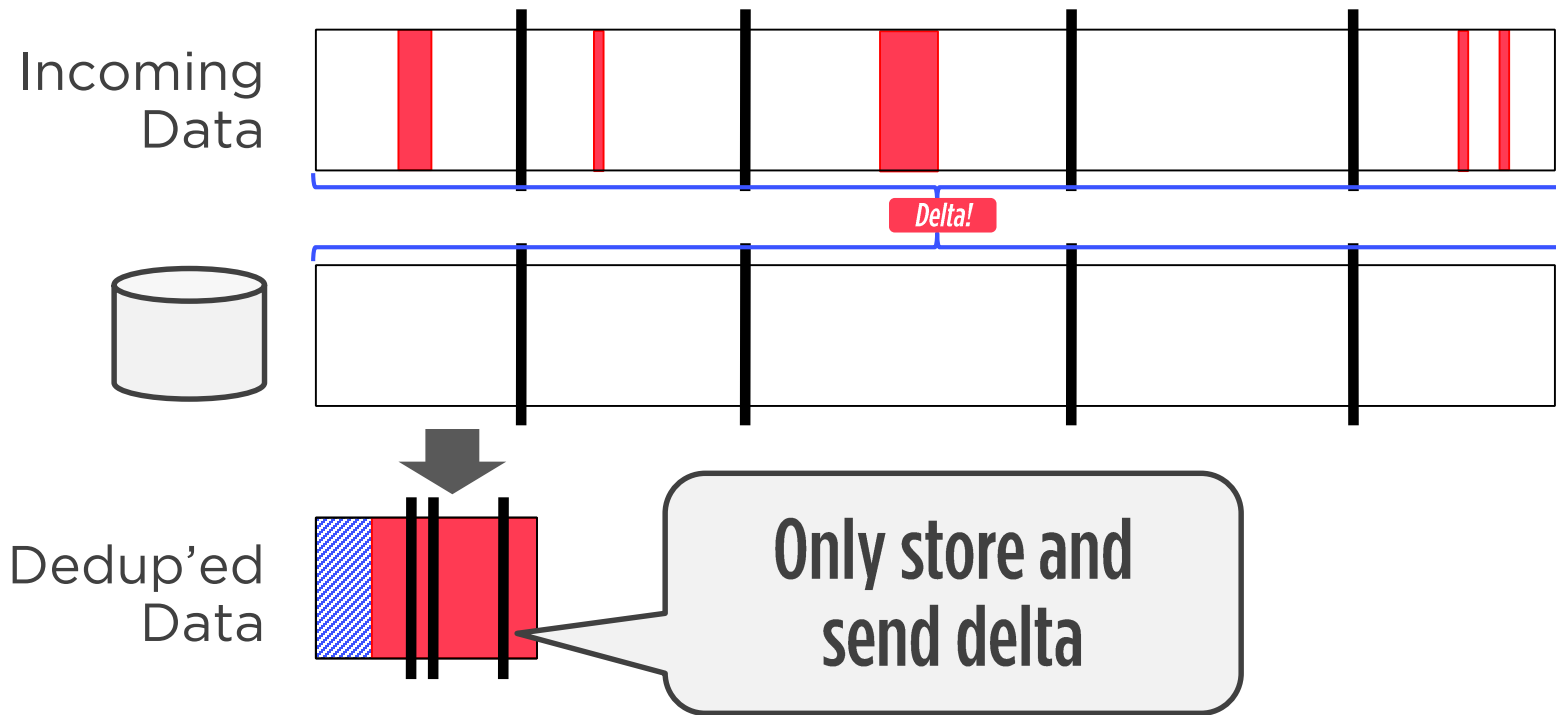
| *Chunk Boundary* ■ *Modified Region* ▨ *Duplicate Chunk*

Incoming
Data

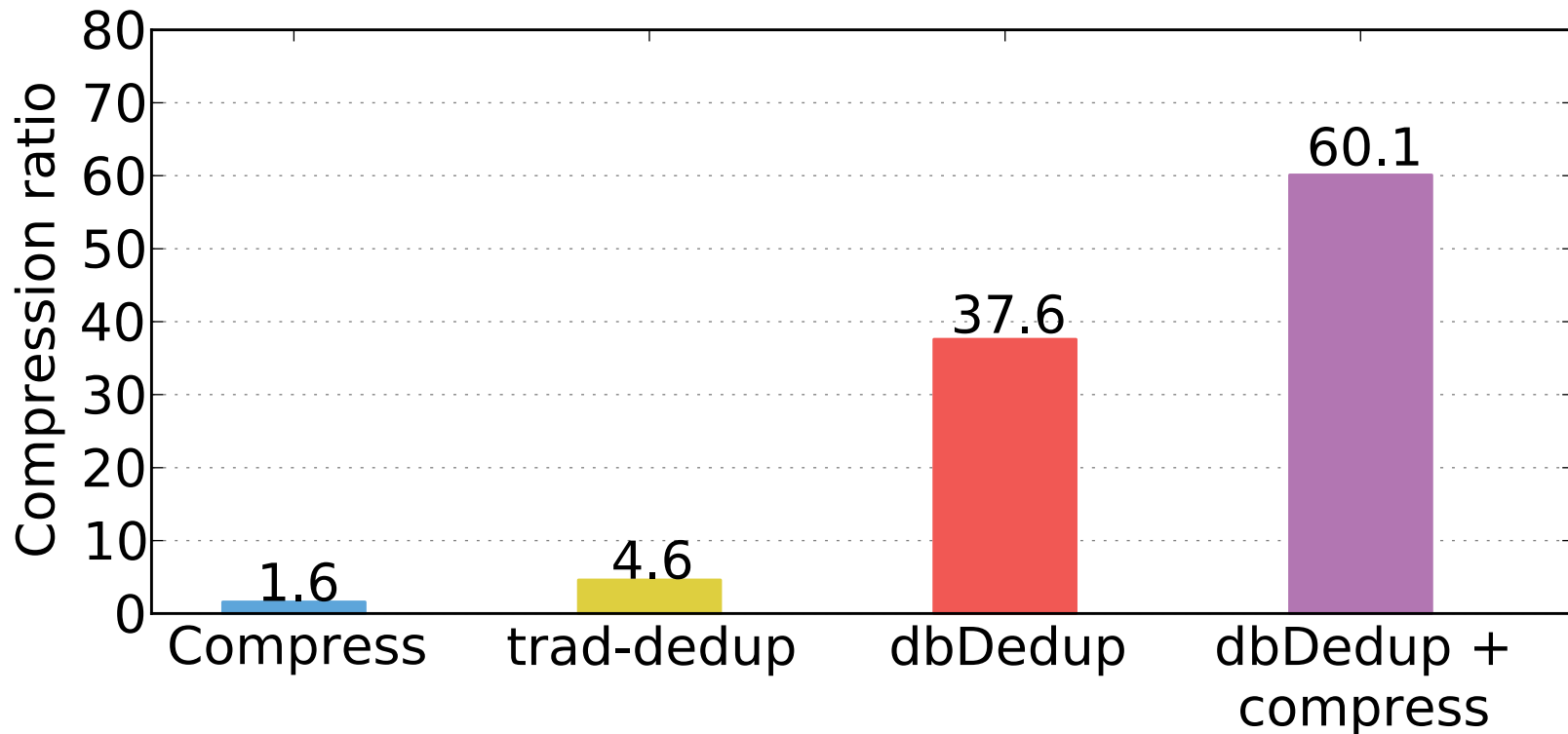


Similarity Dedup (dbDedup)

| *Chunk Boundary* ■ *Modified Region* ▨ *Duplicate Chunk*



Compress vs. Dedup



*Wikipedia Database (20GB sample)
MongoDB v3.1 w/ Snappy Compression*

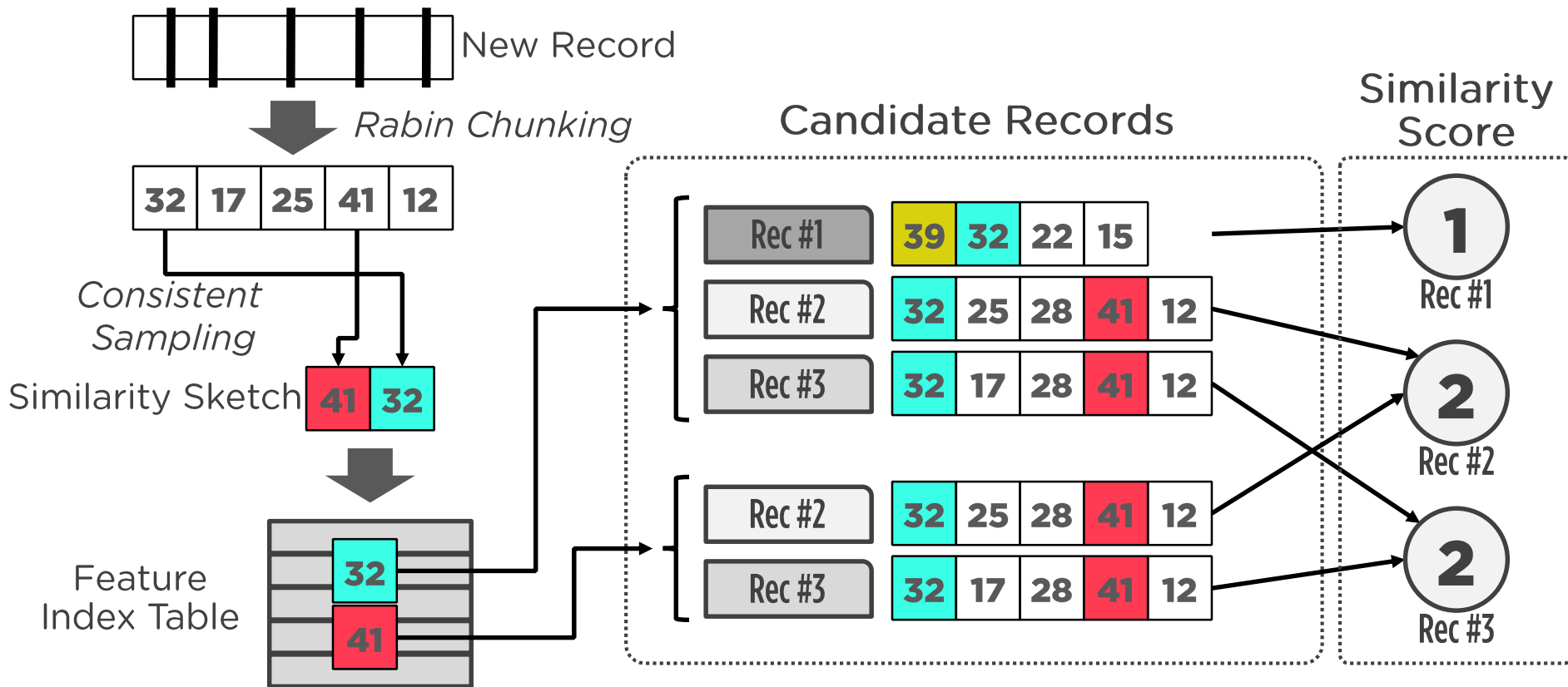
Outline

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- Need for Similarity-based Dedup
- **Our Approach: dbDedup**
- Evaluation

dbDedup Encoding Steps

- For each new record:
 - *Identify Similar Records*
 - *Select the “Best” Match*
 - *Delta Compression*

Identify Similar Records



Select the Best Match

Initial Ranking

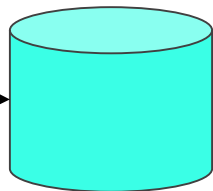
Rank	Candidates	Score
1	Rec #2	2
1	Rec #3	2
2	Rec #1	1

Select the Best Match

Initial Ranking

Rank	Candidates	Score
1	Rec #2	2
1	Rec #3	2
2	Rec #1	1

Is record cached?



Source Record
Cache

Select the Best Match

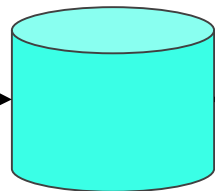
Initial Ranking

Rank	Candidates	Score
1	Rec #2	2
1	Rec #3	2
2	Rec #1	1

Final Ranking

Rank	Candidates	Cached?	Score
1	Rec #3	Yes	4
1	Rec #1	Yes	3
2	Rec #2	No	2

Is record cached?



If yes, reward +2

Cache-aware selection

Source Record
Cache

Delta Compression

- Based on the “xDelta” algorithm
 - *Byte-level diff between source and target records*
- Challenge:
 - *Reading encoded records incurs decoding overhead*
- *Solution: new encoding techniques*
 - *Reduce both freq. and cost of decode*

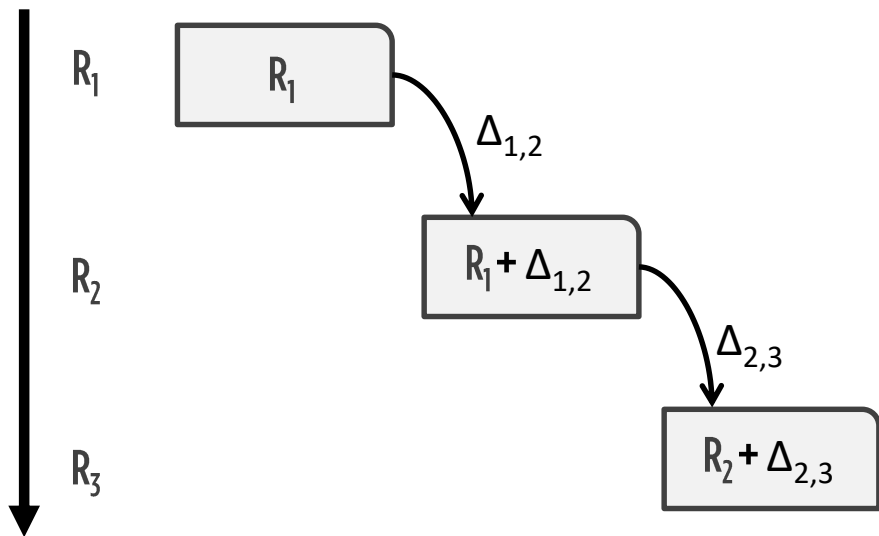
New Encoding Techniques

- Two-way encoding
 - *Forward encoding* for network-level dedup
 - Sending encoding of new records against source
 - Replicas need to use existing record as source
 - *Backward encoding* for storage-level dedup
 - Encode/rewrite source; write unencoded new record
 - Minimize read overhead for the most recent records
- Hop encoding
 - *Reduce worst-case re-construction time*
 - *Maintain compression ratio*

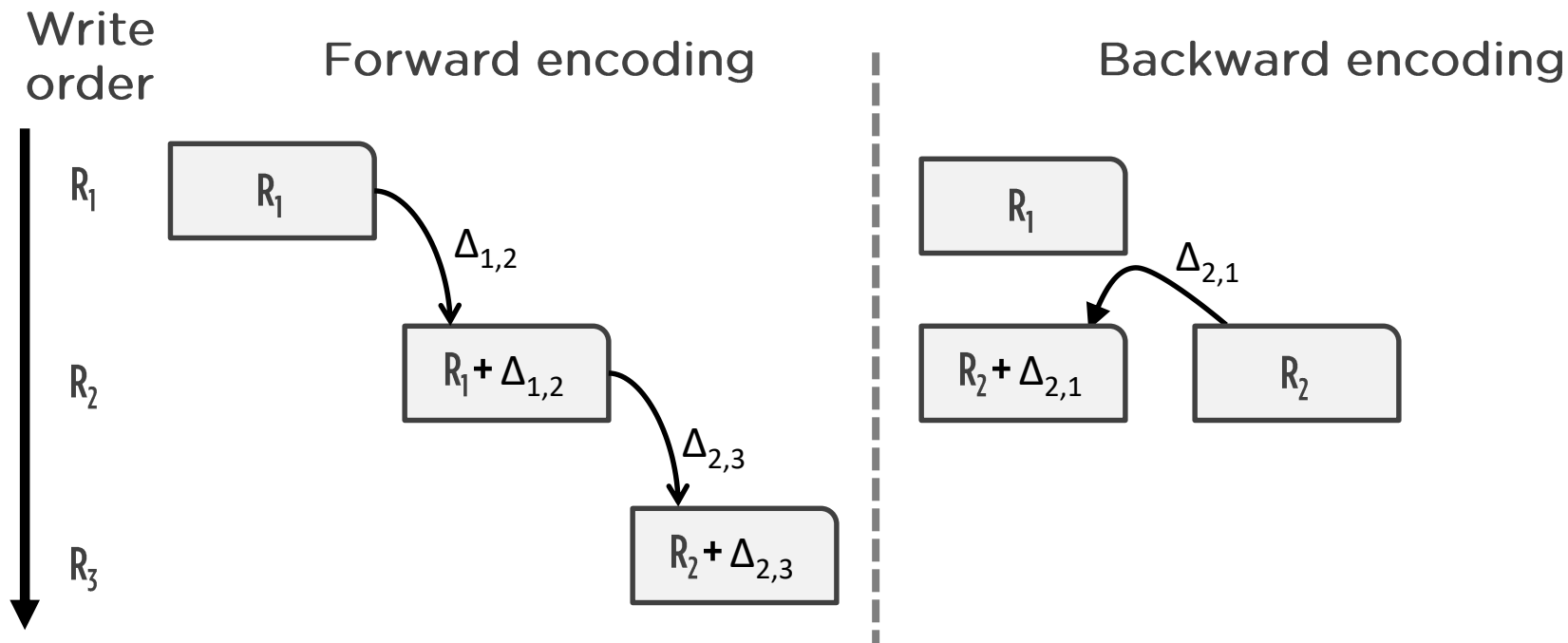
Two-way Encoding

Write
order

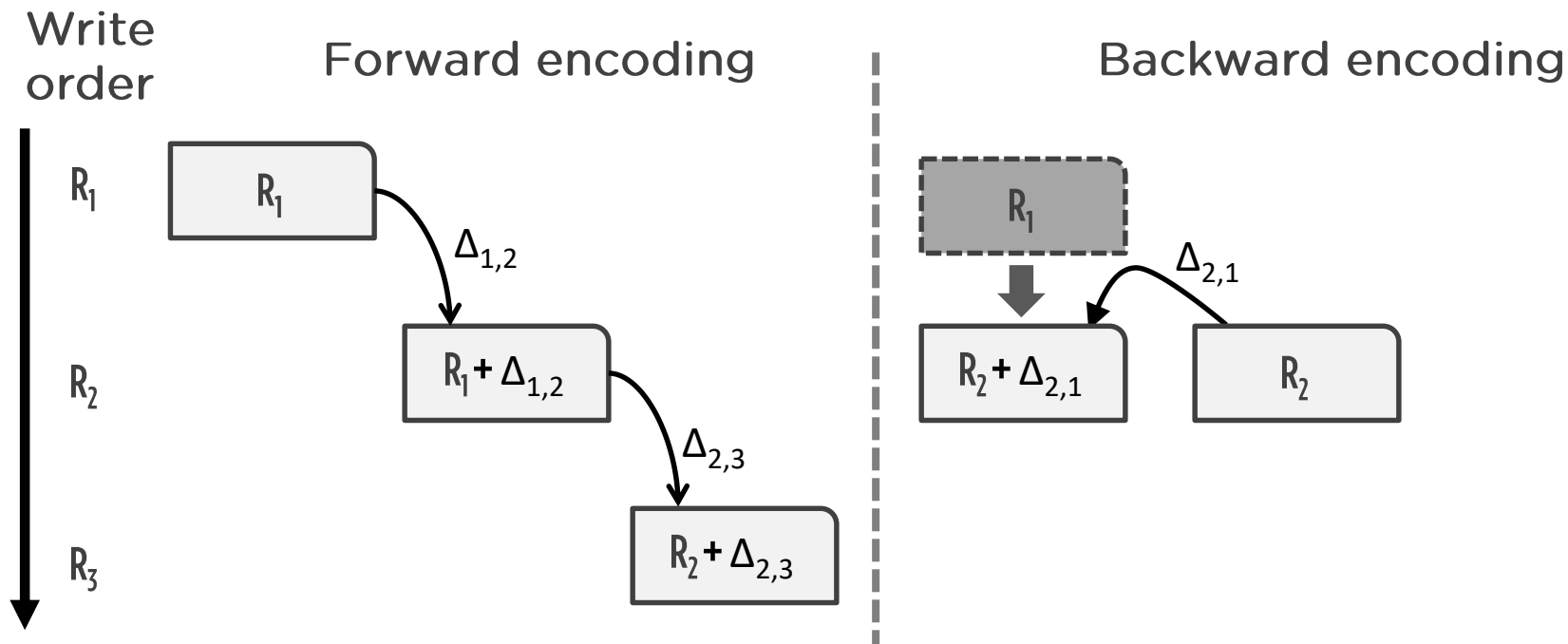
Forward encoding



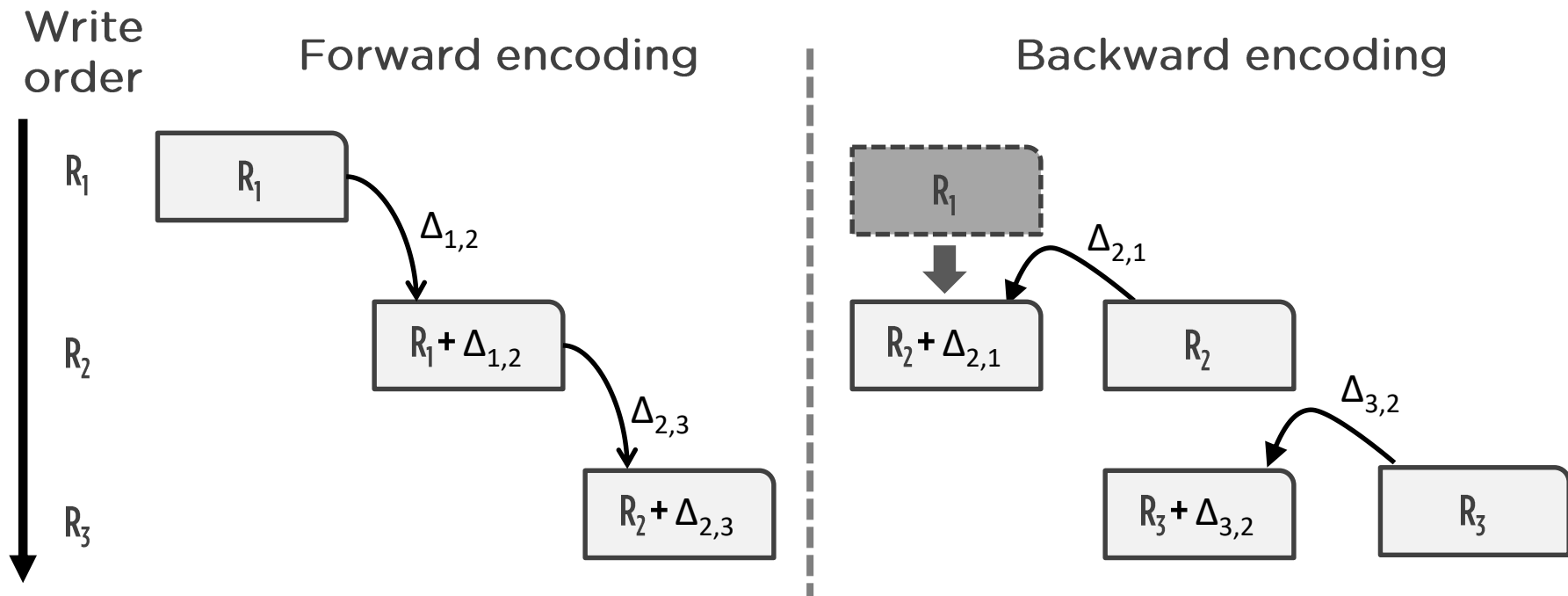
Two-way Encoding



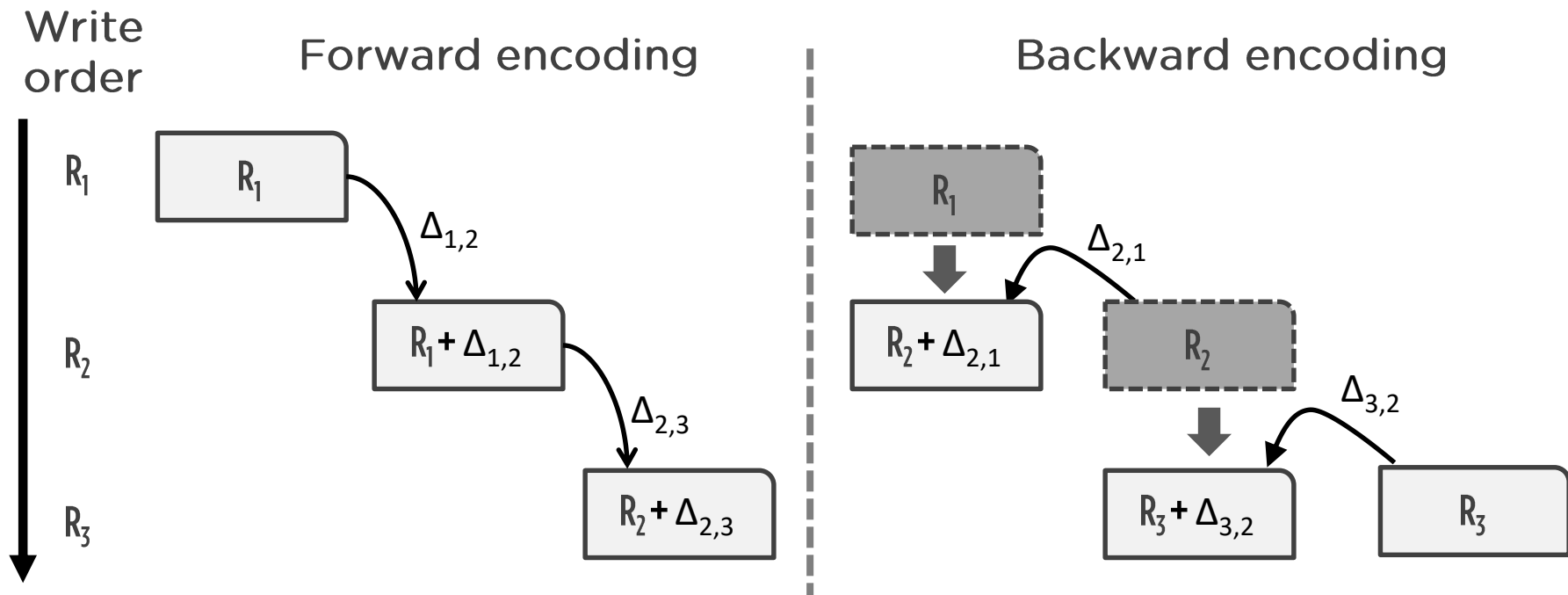
Two-way Encoding



Two-way Encoding

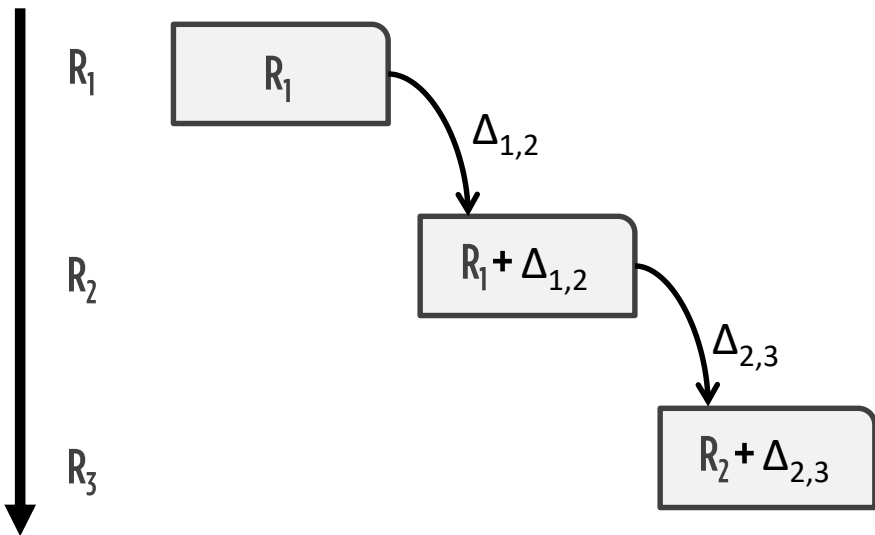


Two-way Encoding



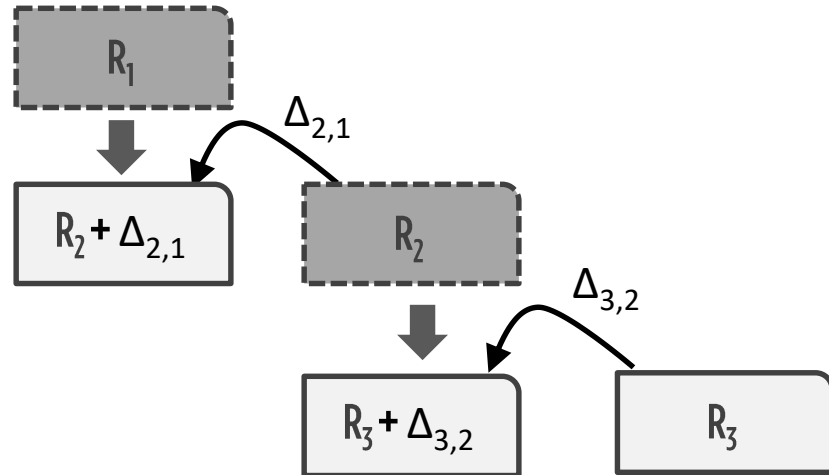
Two-way Encoding

Write
order



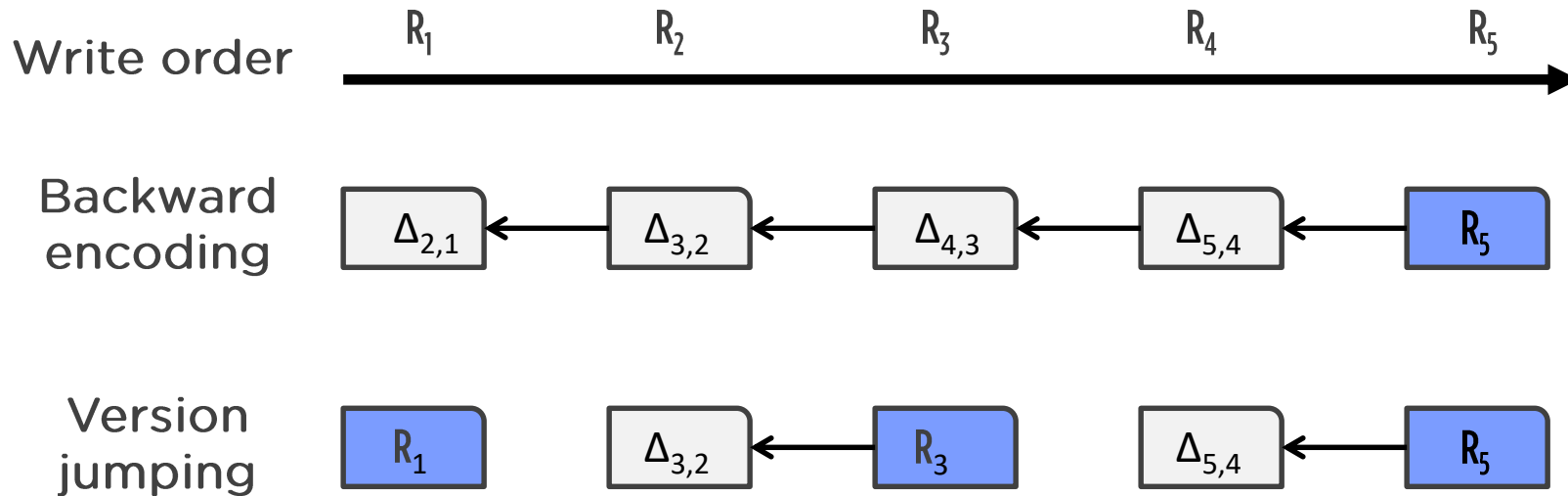
Network-level dedup

Backward encoding

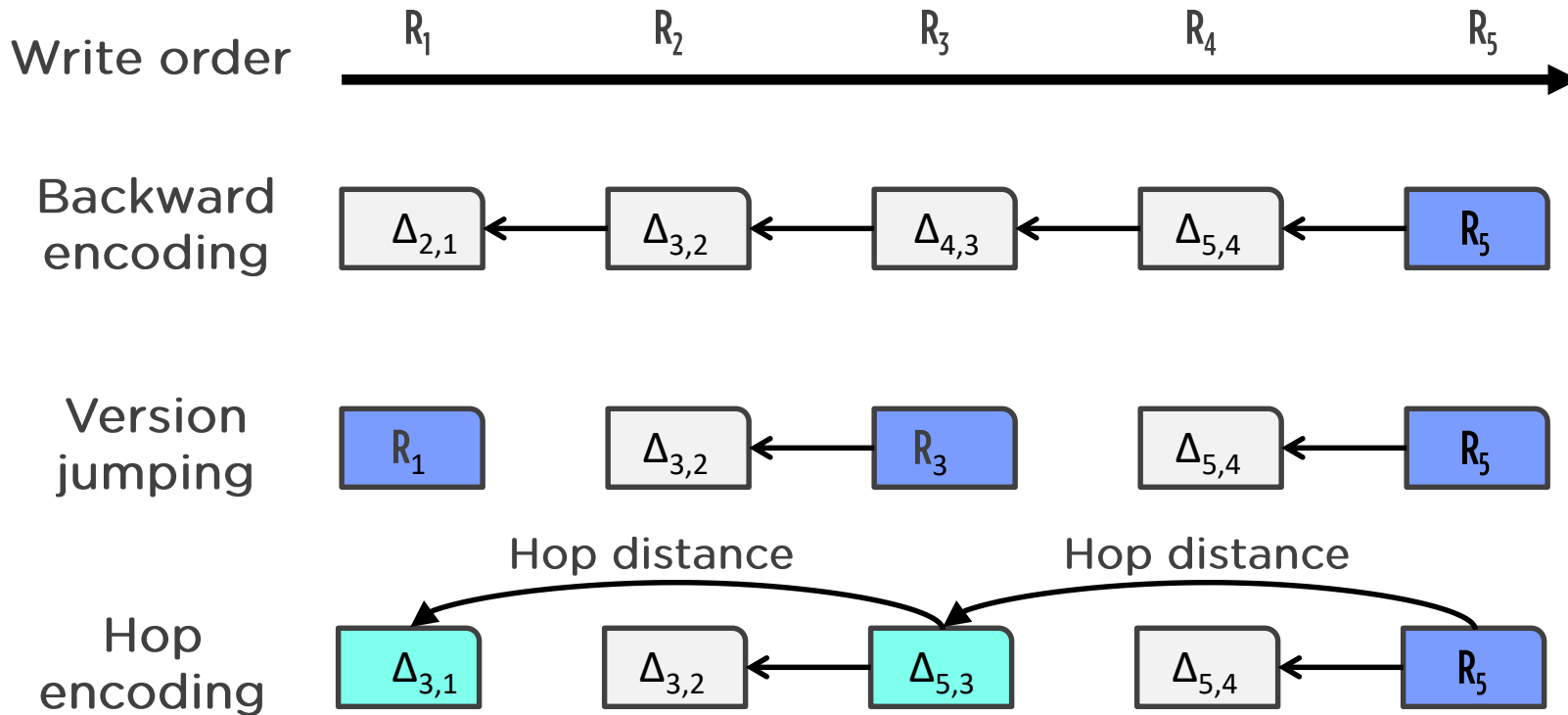


Storage-level dedup

Hop Encoding



Hop Encoding



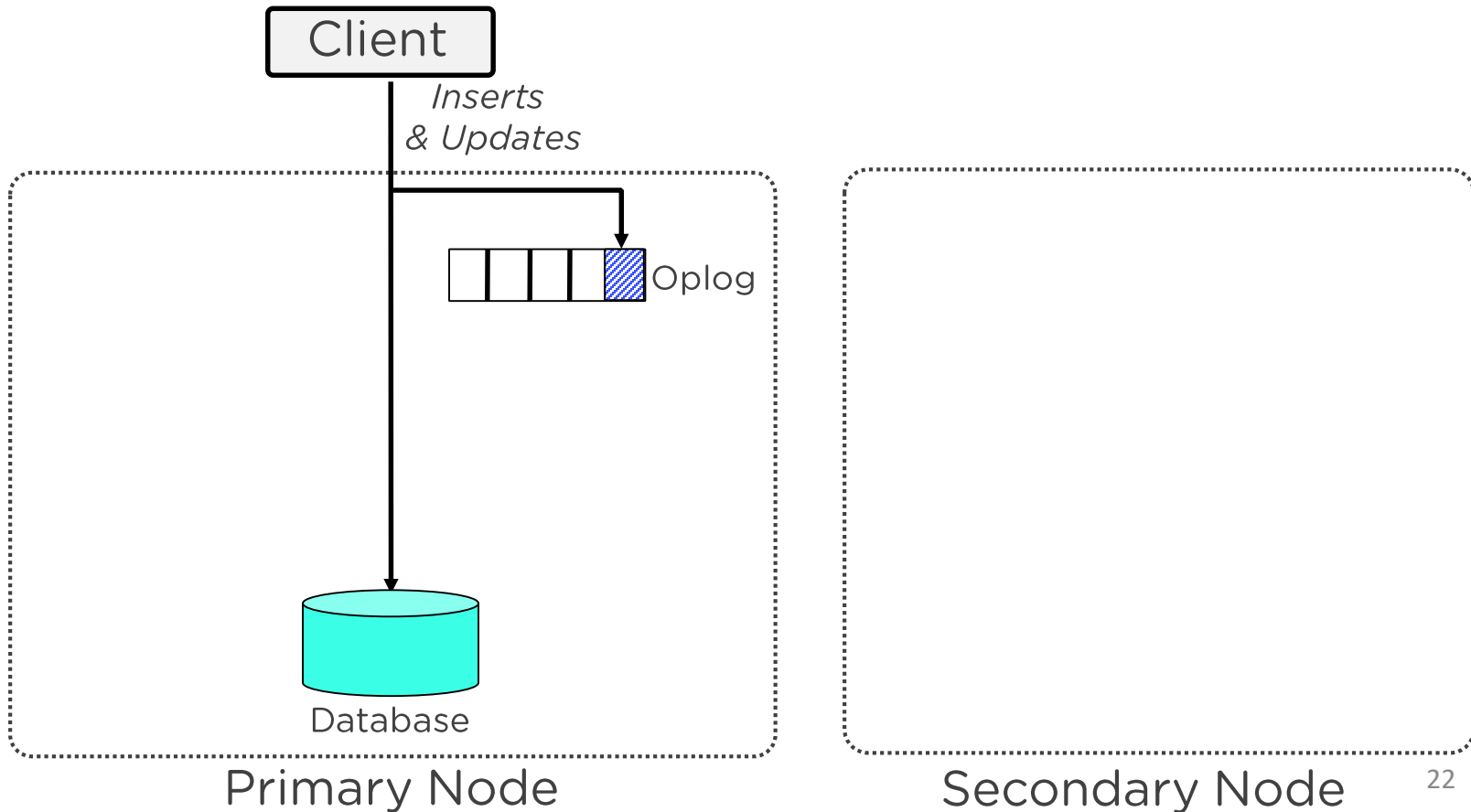
Caching Mechanisms

- Unique challenge for delta compression
 - *Extra disk I/Os to read/update records*
- Solution: specialized caches
 - *Source Record cache*
 - Reduce disk reads to source records
 - Keep only the latest version of a record
 - *Lossy write-back cache*

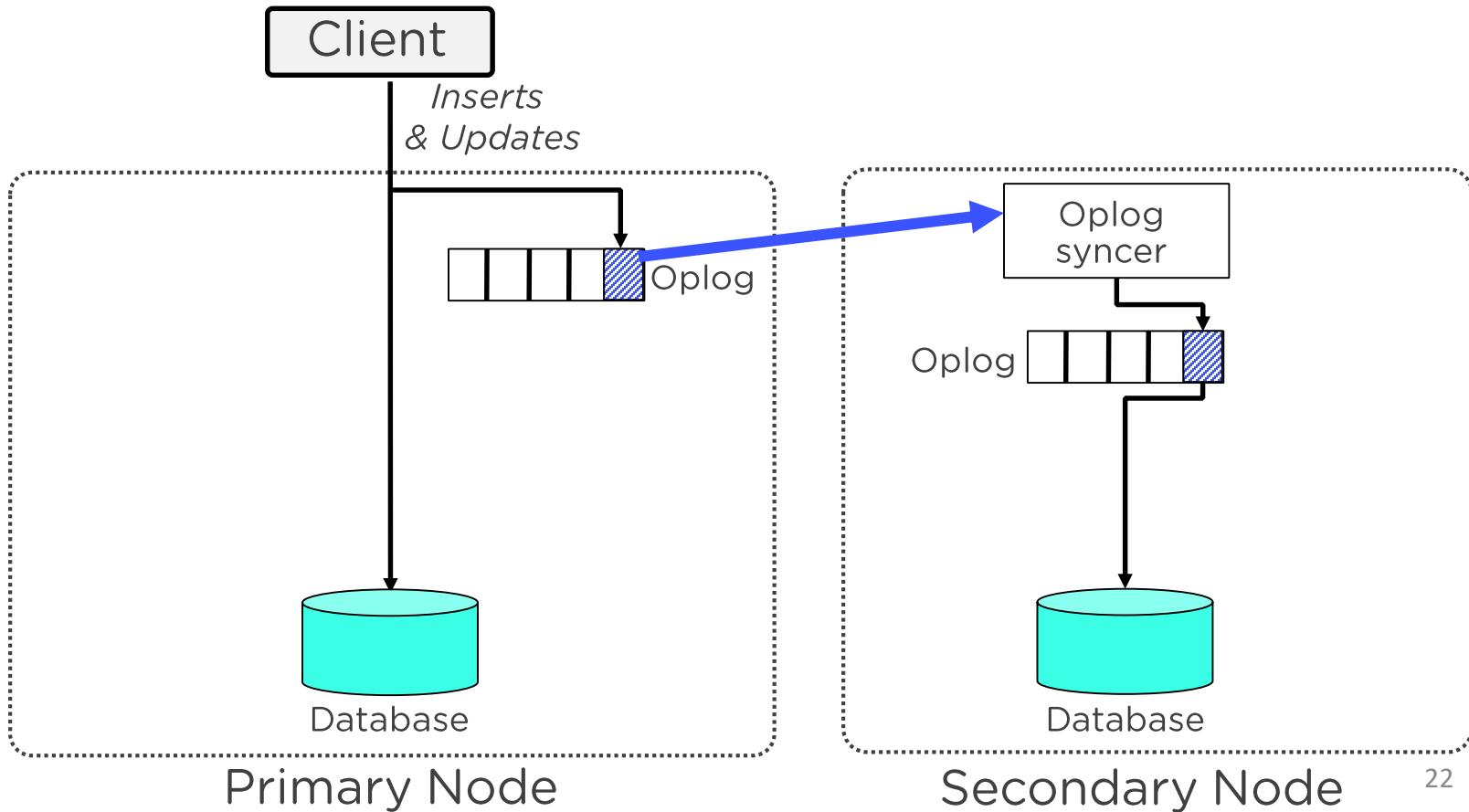
Lossy Write-back Cache

- Reduce overhead of writing backward-encoded sources
 - *Delay updates until system is idle*
 - *Prioritize updates based on space savings*
- “Lossy” property
 - *Unapplied updates can be safely discarded*
 - *Records remain un-encoded in such case*

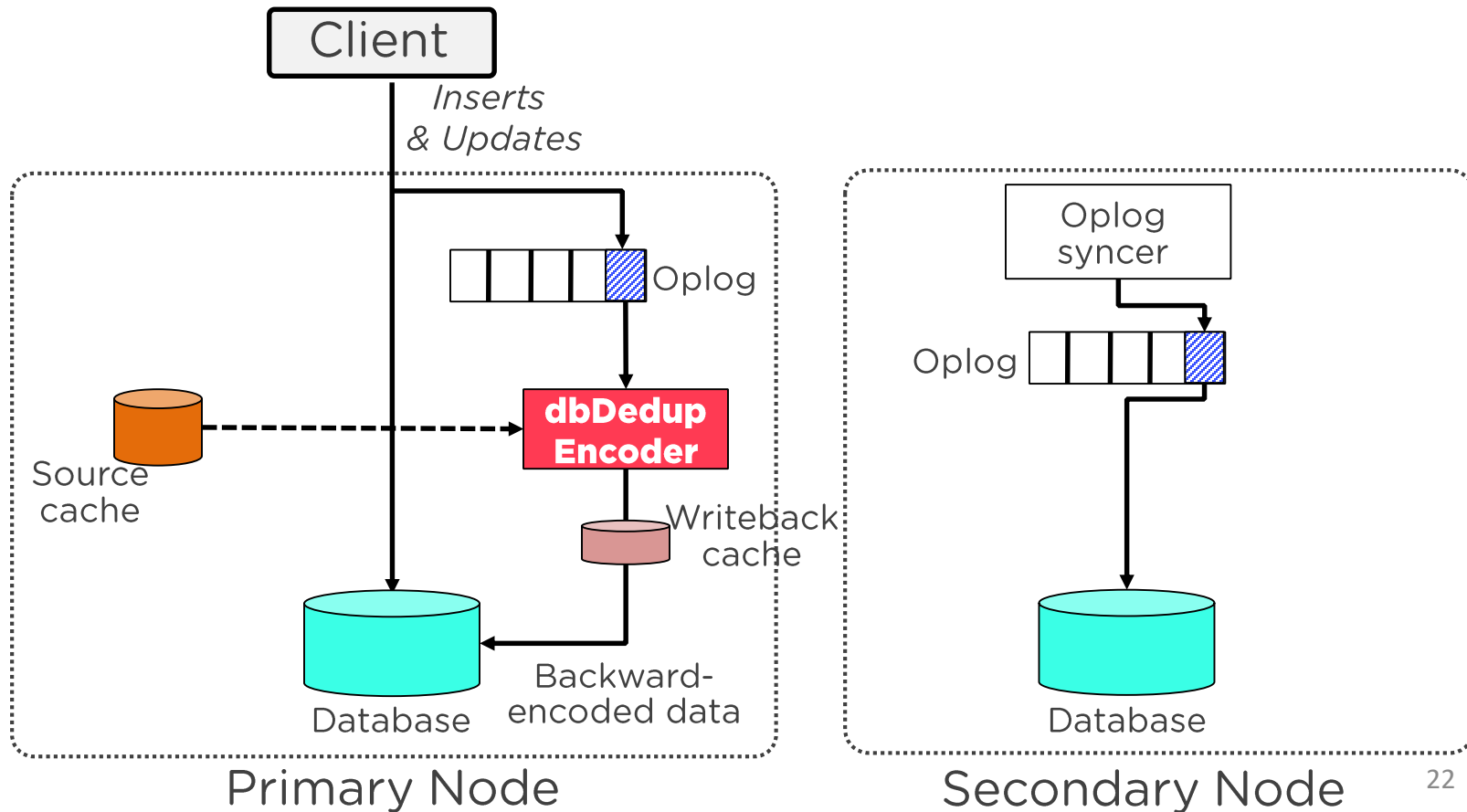
Integration into DBMS



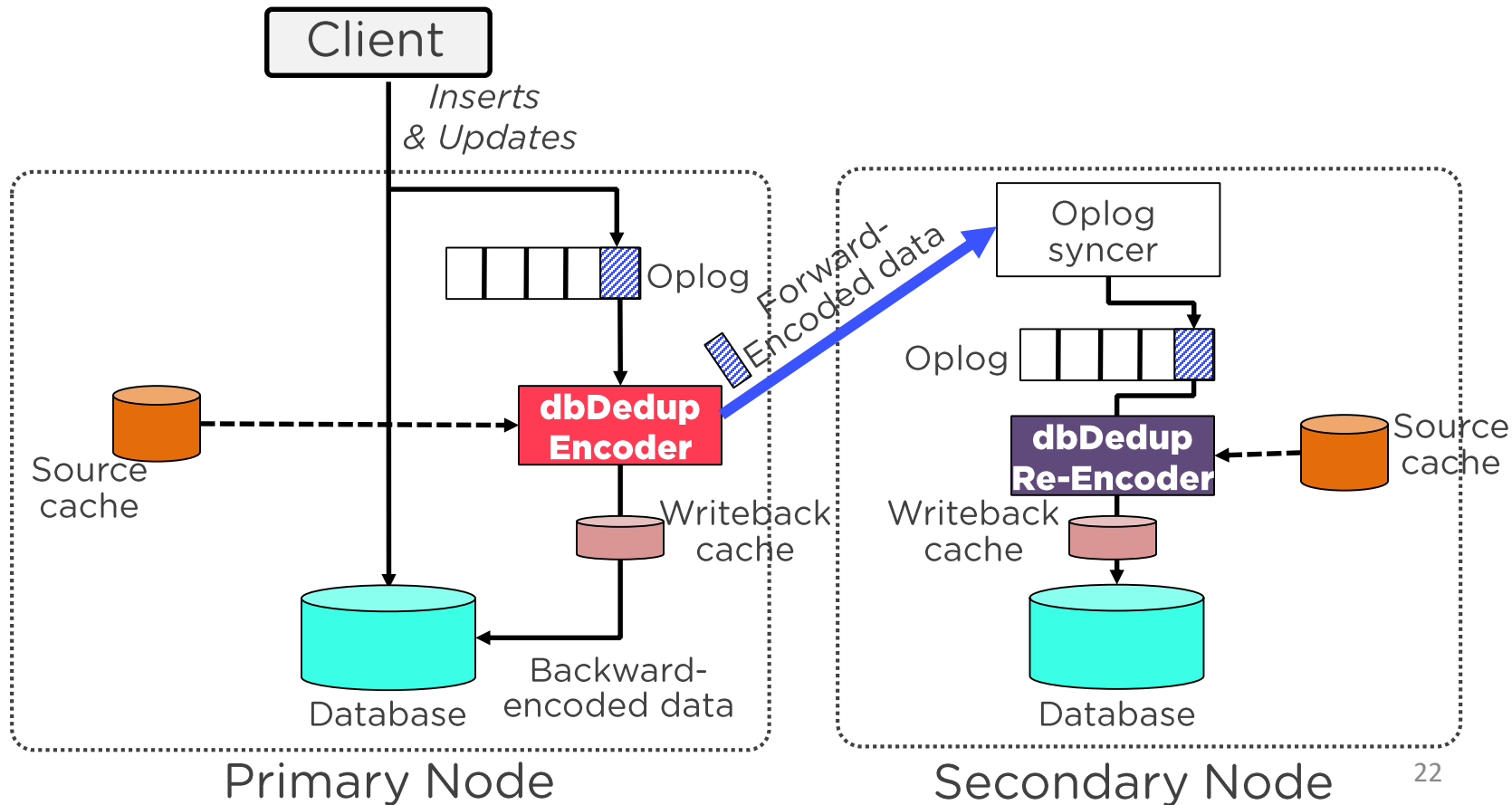
Integration into DBMS



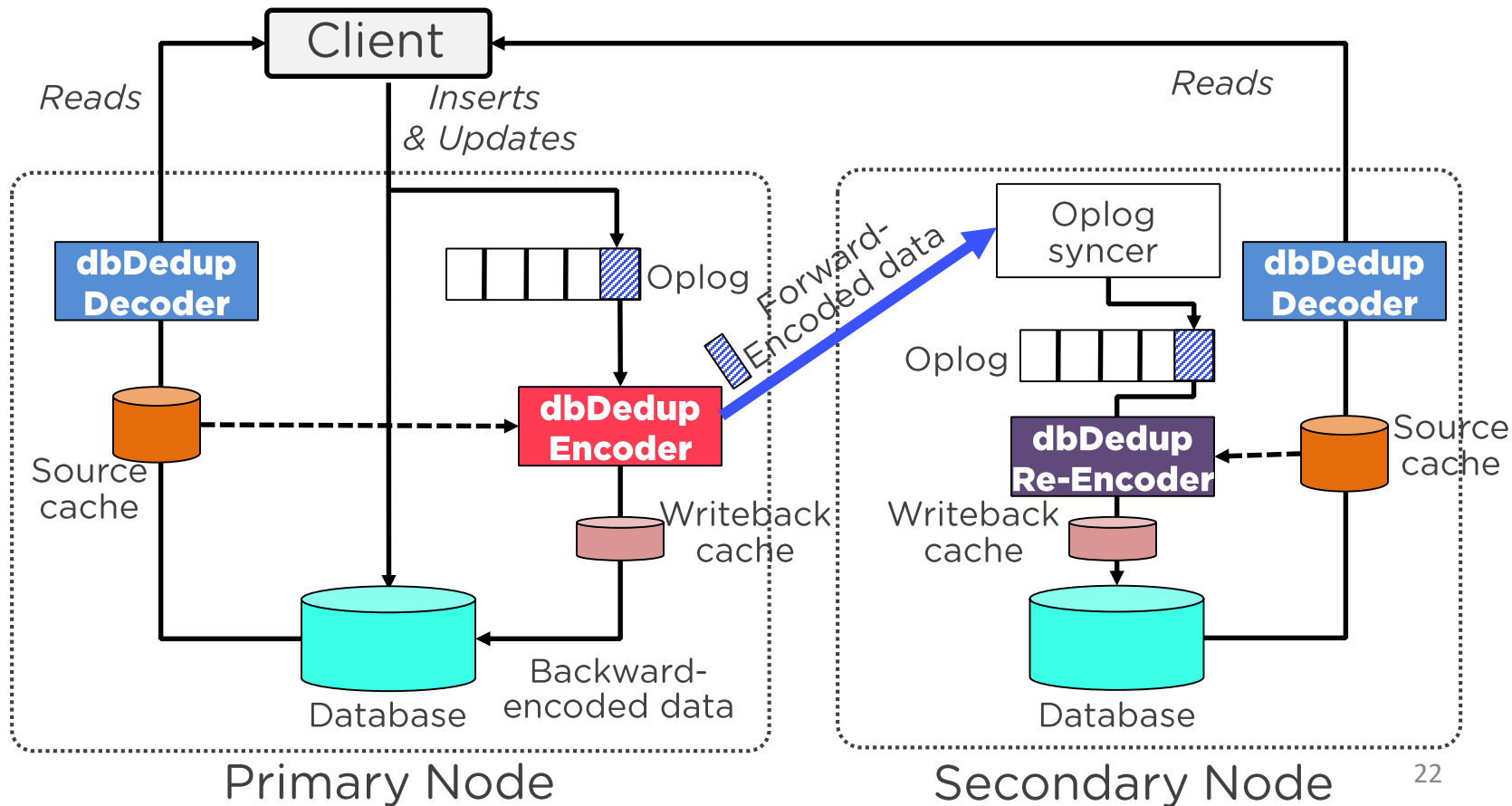
Integration into DBMS



Integration into DBMS



Integration into DBMS



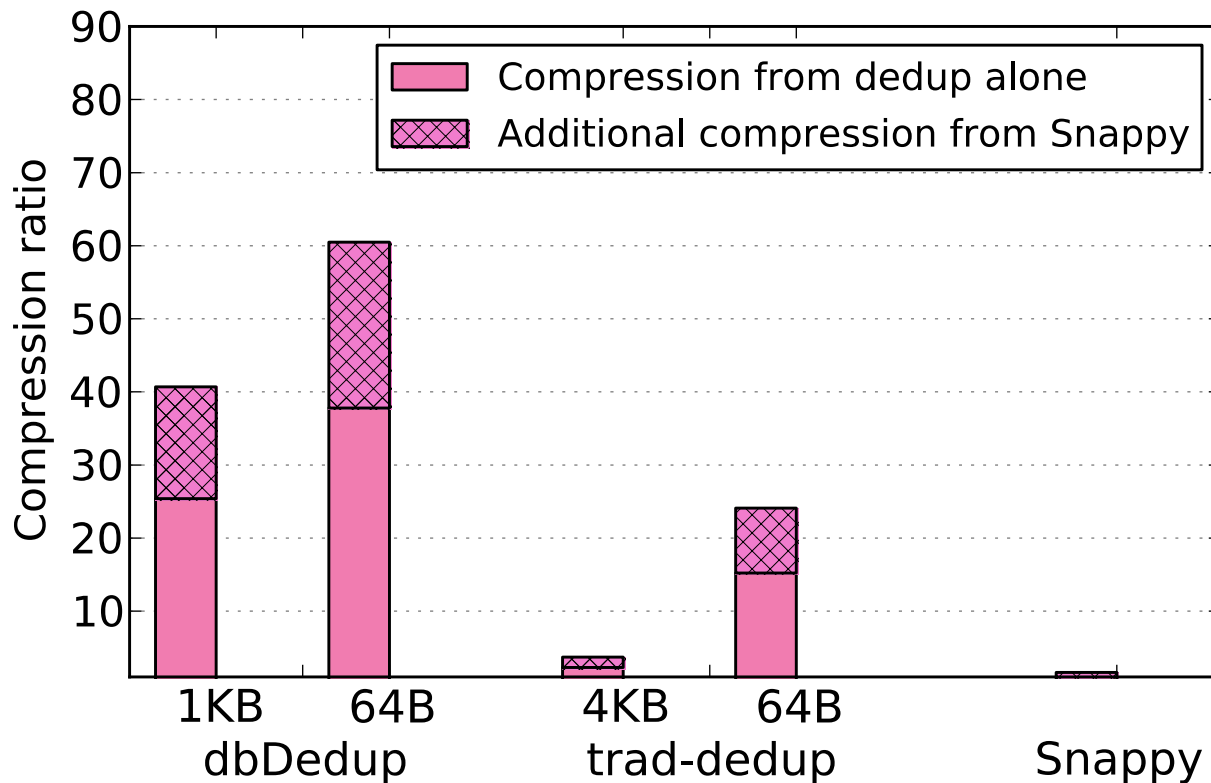
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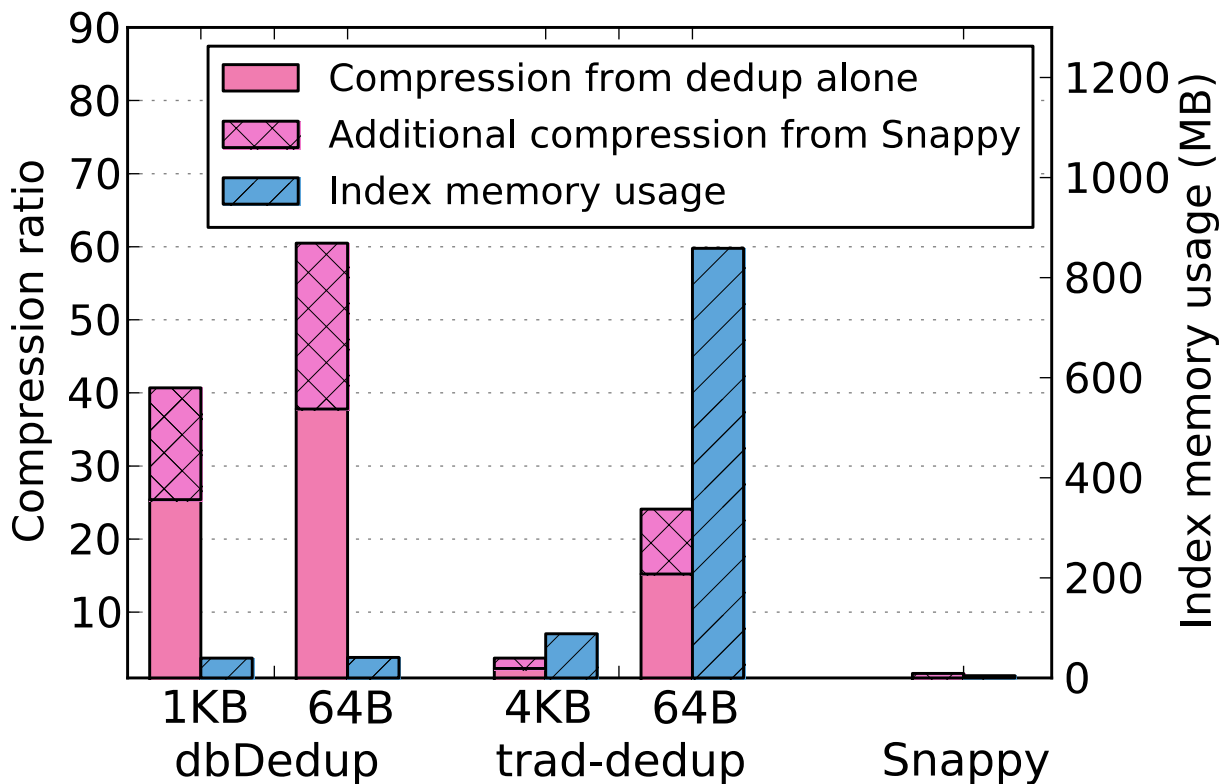
Evaluation

- MongoDB (v3.1)
 - *1 primary, 1 secondary node, 1 client*
 - *Node Config: 4 cores, 8GB RAM, 100GB HDD storage*
- Datasets:
 - *Wikipedia dump (20GB out of ~12TB)*
 - *Additional datasets evaluated in the dissertation*

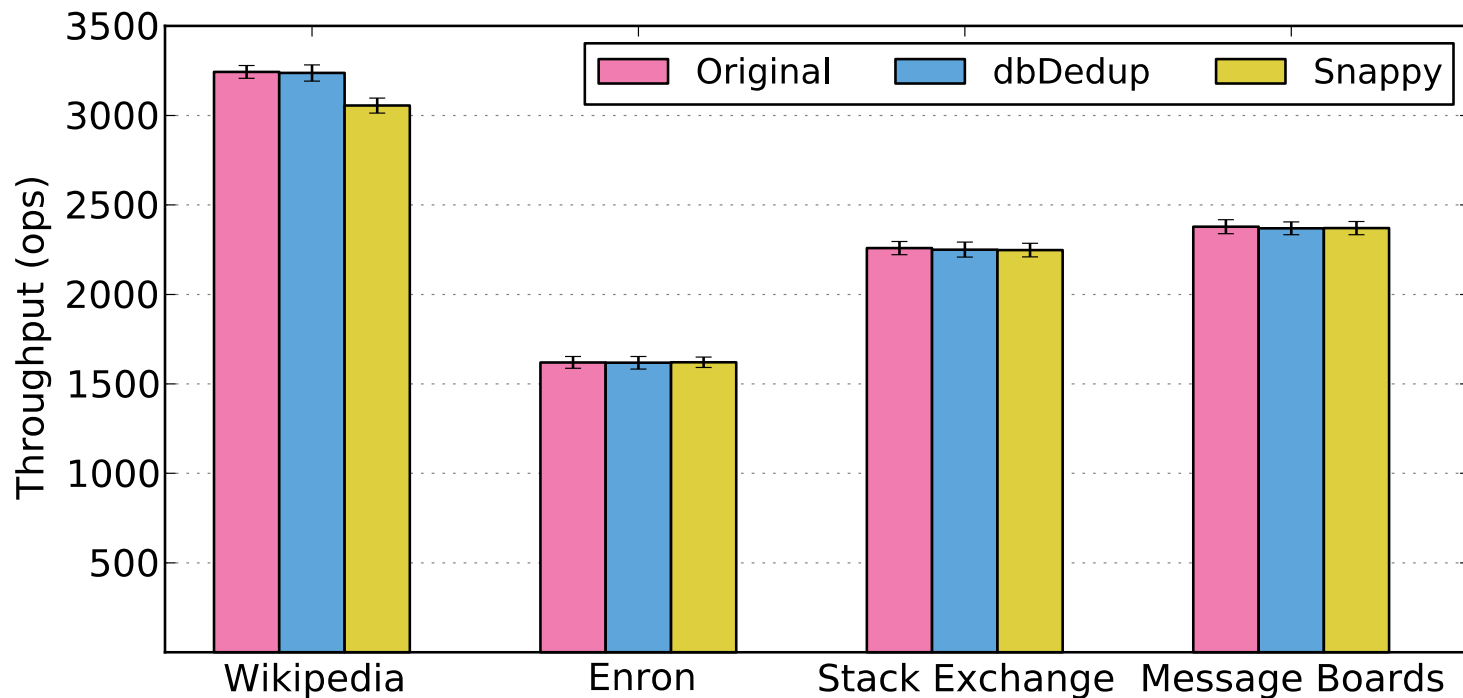
Compression Ratio



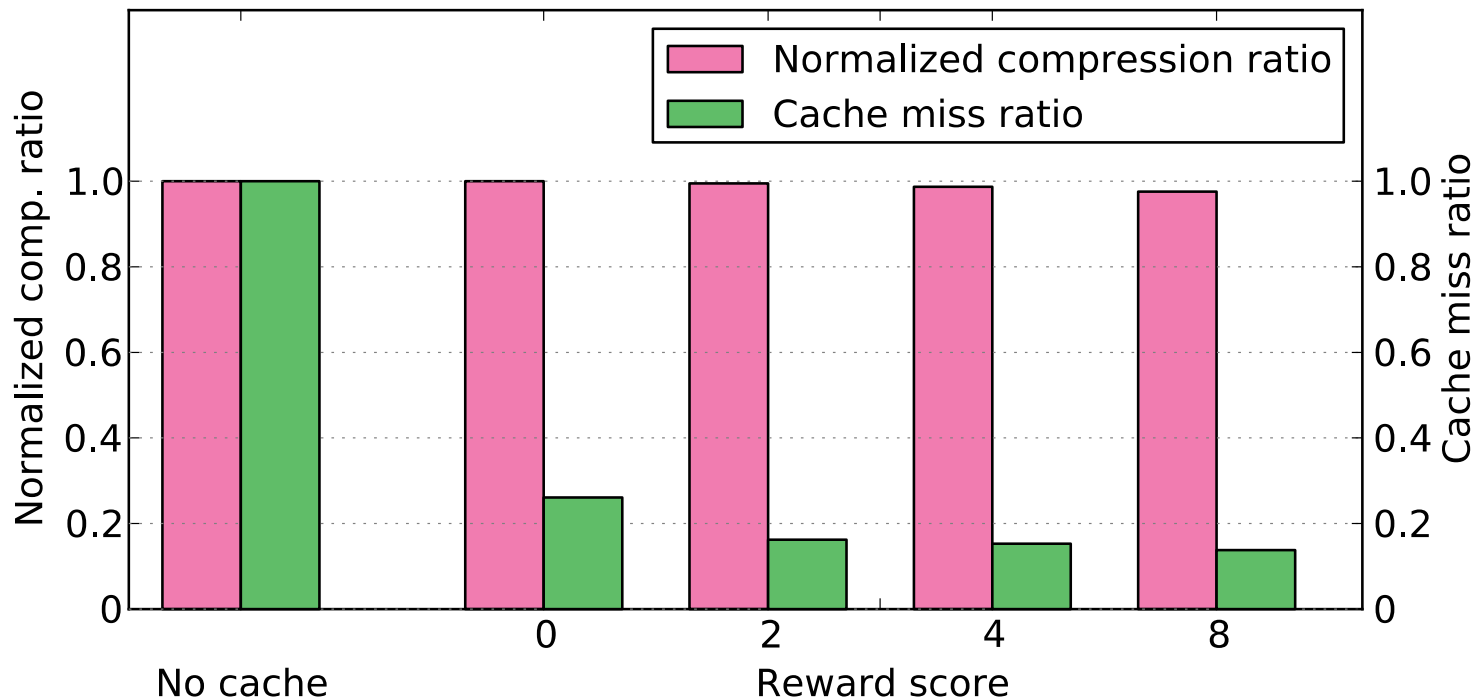
Comp. Ratio & Index Memory



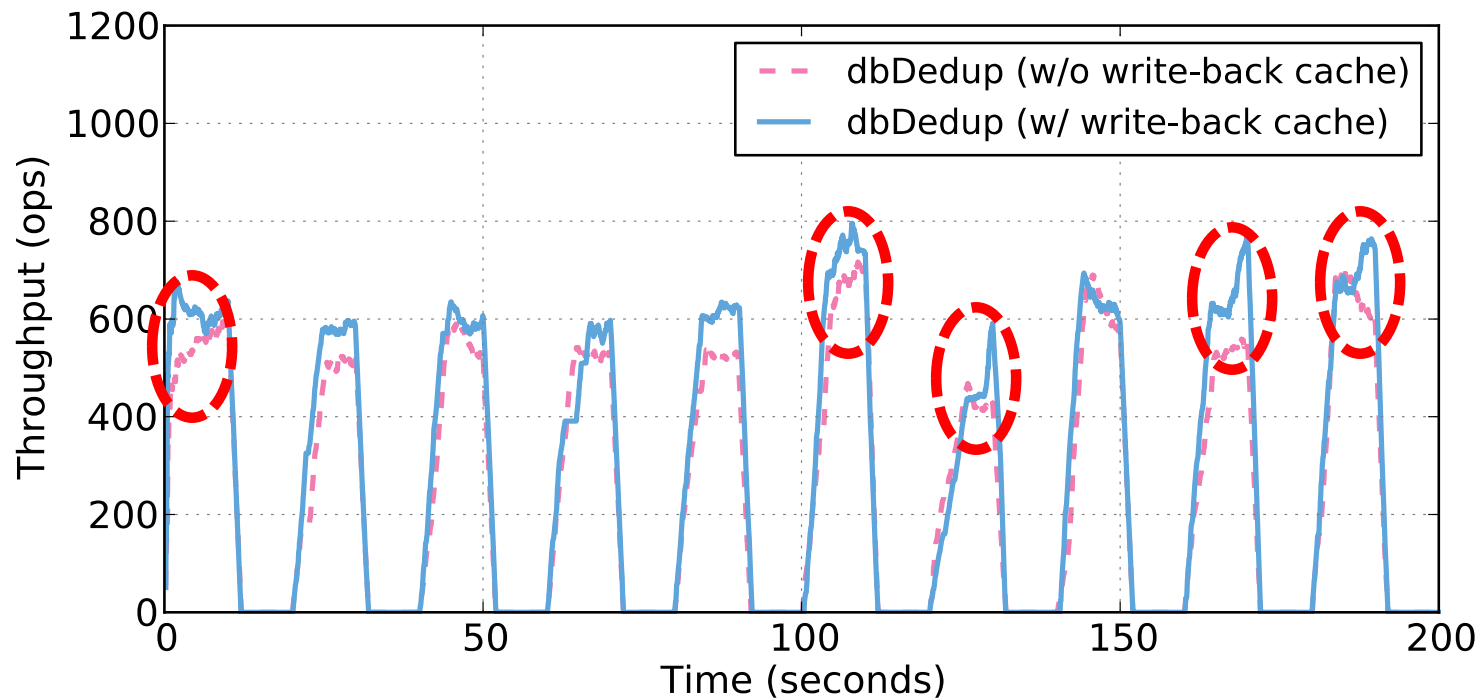
Throughput Overhead



Effect of Source Cache



Effect of Lossy Write-back



Conclusion

- **dbDedup**: Similarity-based dedup for DBMSs
 - *First DB dedup system for both storage and network layers*
 - *Much greater data reduction than traditional dedup*
 - *Up to 38x compression ratio for Wikipedia*
 - *Resource-efficient design with negligible overhead*