

## 8803 Connections between Learning, Game Theory, and Optimization

Homework # 2

Due: October 28th 2010

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This homework is due by the start of class on October 28th 2010. You can either submit the homework via the course page on T-Square or hand it in at the beginning of the class on October 28th 2010. Start early!

### Groundrules:

- Your work will be graded on correctness, clarity, and conciseness.
- You may collaborate with others on this problem set. However, you must *write your own solutions* and *list your collaborators/sources* for each problem.

### Problems:

1. **VC-dimension of linear separators:** In this problem you will prove that the VC-dimension of the class  $H_n$  of halfspaces (another term for linear threshold functions) in  $n$  dimensions is  $n + 1$ . We will use the following definition: The *convex hull* of a set of points  $S$  is the set of all convex combinations of points in  $S$ ; this is the set of all points that can be written as  $\sum_{x_i \in S} \lambda_i x_i$ , where each  $\lambda_i \geq 0$ , and  $\sum_i \lambda_i = 1$ . It is not hard to see that if a halfspace has all points from a set  $S$  on one side, then the entire convex hull of  $S$  must be on that side as well.
  - (a) [**lower bound**] Prove that  $\text{VC-dim}(H_n) \geq n + 1$  by presenting a set of  $n + 1$  points in  $n$ -dimensional space such that one can partition that set with halfspaces in all possible ways. (And, show how one can partition the set in any desired way.)
  - (b) [**upper bound part 1**] The following is “Radon’s Theorem,” from the 1920’s.

**Theorem.** *Let  $S$  be a set of  $n + 2$  points in  $n$  dimensions. Then  $S$  can be partitioned into two (disjoint) subsets  $S_1$  and  $S_2$  whose convex hulls intersect.*

Show that Radon’s Theorem implies that the VC-dimension of halfspaces is *at most*  $n + 1$ . Conclude that  $\text{VC-dim}(H_n) = n + 1$ .
  - (c) [**upper bound part 2**] Now we prove Radon’s Theorem. We will need the following standard fact from linear algebra. If  $x_1, \dots, x_{n+1}$  are  $n + 1$  points in  $n$ -dimensional space, then they are linearly dependent. That is, there exist real values  $\lambda_1, \dots, \lambda_{n+1}$  *not all zero* such that  $\lambda_1 x_1 + \dots + \lambda_{n+1} x_{n+1} = 0$ .

You may now prove Radon’s Theorem however you wish. However, as a suggested first step, prove the following. For any set of  $n + 2$  points  $x_1, \dots, x_{n+2}$  in  $n$ -dimensional space, there exist  $\lambda_1, \dots, \lambda_{n+2}$  *not all zero* such that  $\sum_i \lambda_i x_i = 0$  and  $\sum_i \lambda_i = 0$ . (This is called *affine dependence*.)
2. **Symmetric congestion games:** Consider a network congestion game defined on a directed graph  $G = (V, E)$ . Assume each player  $i$  wants to get from source  $o$  and to destination  $t$  (so all players have the same set of strategies, which is the set of paths from  $o$  to  $t$ ). Assume that the cost of each edge is monotonically increasing with the number of players using that

edge. Show that these games have a pure Nash equilibrium and provide a polynomial time algorithm to compute a pure Nash equilibrium.

**Hints:** Try to compute the solution that minimizes the potential function. Think of min-cost network flow.

**Extra Credit:**

1. **Boosting and Game Theory:** Think about the boosting process in relation to the decision theoretic experts model and derive a boosting algorithm based on this connection.

**Note:** You are welcome to read the paper “Game Theory, On-line Prediction and Boosting” by Rob Schapire and Yoav Freund, but please write the solution in your own words.