Language and Statistics II

Lecture 7:
$$W = S A S$$

 ϵ T

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Finite-State Technology

- Formally well-understood
 - Regular languages, rational relations
 - Generalizes n-grams, HMMs
- Many applications in NL technologies
 - Speech recognition
 - Lexical, morphological processing
 - Information extraction
 - Translation (!)
 - Parsing (!)
- Several toolkits
- Often determinizable: very fast

Finite-State Automata (Recognizers)

- Automaton that recognizes regular language
- Implementation of a **regular** expression
- Regular languages are closed under numerous operations
 - Concatenation, union, intersection, Kleene *, difference, reverse, complement, …
- Correspond to regular grammars (type 3 in Chomsky hierarchy)
- Pumping lemma: necessary condition for a language to be regular

FSM as a Dictionary

- Example: 850 words in "Basic English"
- Each word is an FSM



Ten-Word Dictionary



Remove *ɛ*-transitions



Determinize



Minimize



Full 850-Word Dictionary

	states	final states	arcs
Union	5303	850	5302
Remove <i>ε</i> -transitions	4454	850	4453
Determinize	2609	848	2608
Minimize	744	42	1535

Algorithms

- Removing ε-transitions
- Determinization
- Minimization

Generalizations

- FS Recognizer is a function from Σ*→{0,1}
 Meaning: fsa(s) = 1 ⇔ s is in the language
- Other rational relations ...
 - FS Transducer: $\Sigma^* \rightarrow \Delta^*$
 - Weighted FSA: $\Sigma^* \rightarrow \mathbb{R}$
 - Weighted FST: $\Sigma^* \rightarrow \Delta^* \times \mathbb{R}$
- WFSAs and WFSTs can be considered probabilistic (but don't have to be)



Finite-State Transducers

- Input alphabet Σ
- Output alphabet Δ
- Set of states Q
- Initial state q₀
- Final states $F \subseteq Q$
- Arcs, $Q \times \Sigma \times Q \times \Delta^*$

Sequential: arcs are functions from Q to $\Sigma \times Q \times \Delta^*$.

p-subsequential: deterministic, **except** each final state has at most *p* output strings after each final state.

Finite-State Transducers

Biggest application: morphology
 – Xerox tools: 20+ languages



Ambiguity and Optionality

leaves →

{leaf +N +pl, leave +V +pres +3p +sing}

- advice +maker → {advisor, adviser}
- inter+ nation +al +ize +ation → {internationalization, internationalisation}

Also, Phonology

- Mapping between pronunciation (phonemes or phonetic symbols) and lexical entries (morpheme sequences or orthography).
- Optionality even more necessary here!

FST Composition

{vouloir \rightarrow veux, vouloir \rightarrow veut, vouloir \rightarrow voulons, vouloir \rightarrow voulez, vouloir \rightarrow veulent, ...}



g

{veux → vœ, veut → vœ, voulons → vulõ, voulez → vule, veulent → vœl, ...}

{vouloir → vœ, vouloir → vœ, vouloir → vulõ, vouloir → vule, vouloir → vœl, ...}

FST Composition

- Formally, (x, z) ∈ f ° g iff there exists y such that (x, y) ∈ f and (y, z) ∈ g.
- Set and relation:

 $(x, z) \in f \circ g \text{ iff } x \in f \text{ and } (x, z) \in g$

• Relation and set:

 $(x, z) \in f^{\circ} g \text{ iff } (x, z) \in f \text{ and } z \in g$

• Set and set (intersection):

 $x \in f^{\circ} g$ iff $x \in f$ and $x \in g$

Basically, treat sets as identity relations.

Why?

- String into transducer (to compute f(s)): string is a set of size one
- Feed set of strings to transducer in parallel!
- Filter a relation by the outputs (compose with a filter set)
- Building a morphological lexicon: define lexicon by a FSA (set), rules by a bunch of transducers (relation)

FST Projection

• Can strip off input or output symbols ... get a FSA. (Might want to determinize.)

Weighted FSAs

Instead of

Is it grammatical (possible)? we might ask,

How grammatical (likely) is it?

- Examples:
 - N-gram models
 - HMMs
 - Acoustic lattices
 - Perfect hash

Weighted Finite-State Acceptors

- Alphabet Σ
- Set of states Q
- Initial weight function, $\pi : Q \rightarrow \mathbb{R}$
- Final weight function, $\xi : Q \rightarrow \mathbb{R}$
- Arcs in $Q \times \Sigma \times Q \times \mathbb{R}$

Unigram model as a WFSA



One state. One arc for every word.

Bigram model as a WFSA



 $|\Sigma|$ + 2 states. One arc for every bigram.

Bigram HMM as a WFSA

- Alphabet Σ (HMM's alphabet)
- Set of states Q (HMM's states)
- Initial weight function, $\pi : \mathbb{Q} \to \mathbb{R}$

(0 for start state, $-\infty$ for others)

• Final weight function, $\xi : Q \rightarrow \mathbb{R}$

log γ(● | q)

• Arcs in $Q \times \Sigma \times Q \times \mathbb{R}$

 $\log \gamma(q' \mid q) + \log \eta(s \mid q)$

Can you tell how to build a WFSA from a trigram HMM?

WFSA as a Log-Linear Model

- WFSAs assign weights to **paths** through the network.
- Can exponentiate, normalize, and interpret this as a joint p(path, output) or conditional p(path | output).
- Feature schemata:
 - Initial state is q
 - Stop state is q
 - Number of times arc (q, q', s) was crossed

Weighted FSTs

- Weighted relation on $\Sigma^* \times \Delta^*$
- Like FSTs, closed under composition
- Examples:
 - Spelling correction
 - Morphological disambiguation
 - Edit distance
 - Machine translation
 - Speech recognition



Toolkits (links on course page)

- FSM libraries (AT&T)
 - Free binaries
 - Implements pretty much everything you need to build weighted and unweighted FS recognizers and transducers ... except training!
- Xerox FS toolkit
 - Web demo; software can be purchased
 - No weights
- RWTH FSA toolkit
 - Newer, open-source
 - Not sure what's implemented

Interesting analogy

- (weighted) Regexps
- Regexp compiler
- (W)FSTs
- Determinization, minimization, ...
- Composition
- Inversion

- Source code
- Compiler
- Object code
- Optimization
- Composition
- Inversion

Summary

- FSAs, FSTs, WFSAs, WFSTs as formal systems, with reference to some key operations, algorithms, and toolkits.
- Next time(s): examples of WFSTs, and learning WFSTs from data
 - Speech and MT
 - Semirings
 - Parameter estimation
 - Grammatical inference