

Language and Statistics II

Lecture 14: Practical Dynamic
Programming

Noah Smith

Where we left off ...

- Shieber, Schabes, and Pereira: logic programming (deduction) as a way to think about and implement parsers.
- Goodman: add weights!
- No implementation.

Meanwhile, in the “real” world of parsing ...

- People were actually building weighted parsers!
- Crucial: good search strategy.

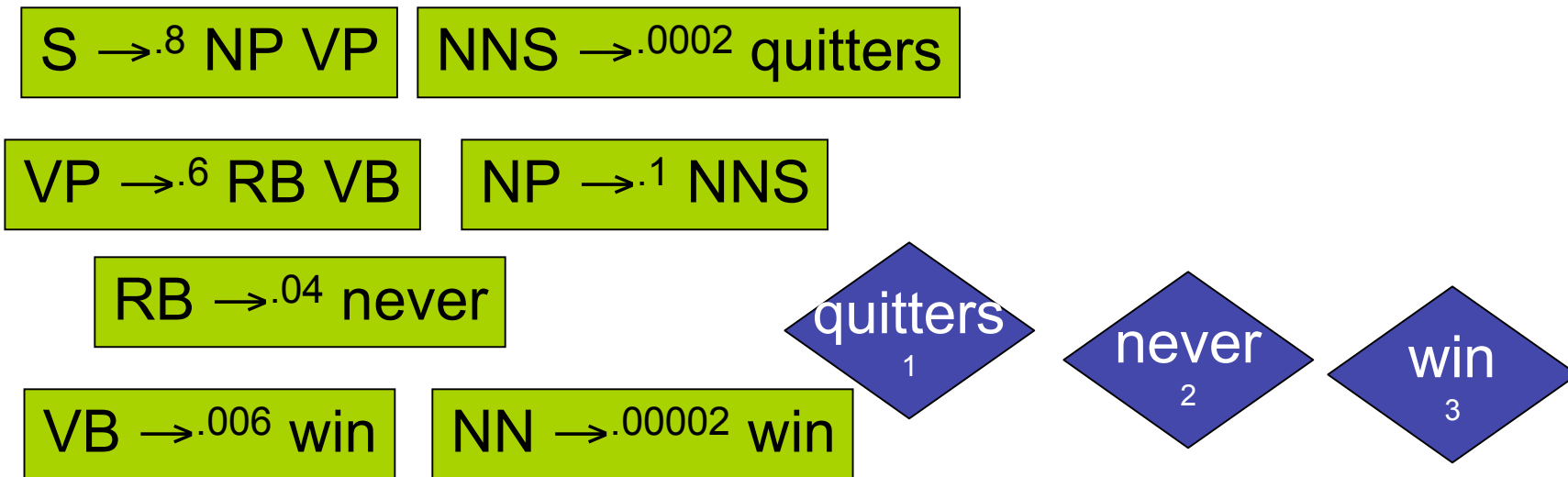
Beyond Goodman (1999)

- Goodman's Algorithm: carry out deduction to build the chart (i.e., fill in items with nonzero value); then compute their values.
 - Tough part: **efficient** ordering of items.
 - For Forward/Viterbi: order by position
 - For CKY: order by width
 - In general?
- Would like efficient **execution strategy** for **arbitrary** programs.
 - Key idea: avoid unnecessary work and repropagation.

Indeed!

- The logic programs *don't tell us how to implement the parser!*
- Is there a **generic** way to go about “compiling” a weighted logic program into a dynamic programming algorithm?
 - Yes: agenda-based DP (for Viterbi).
 - Does this generalize to arbitrary semirings?

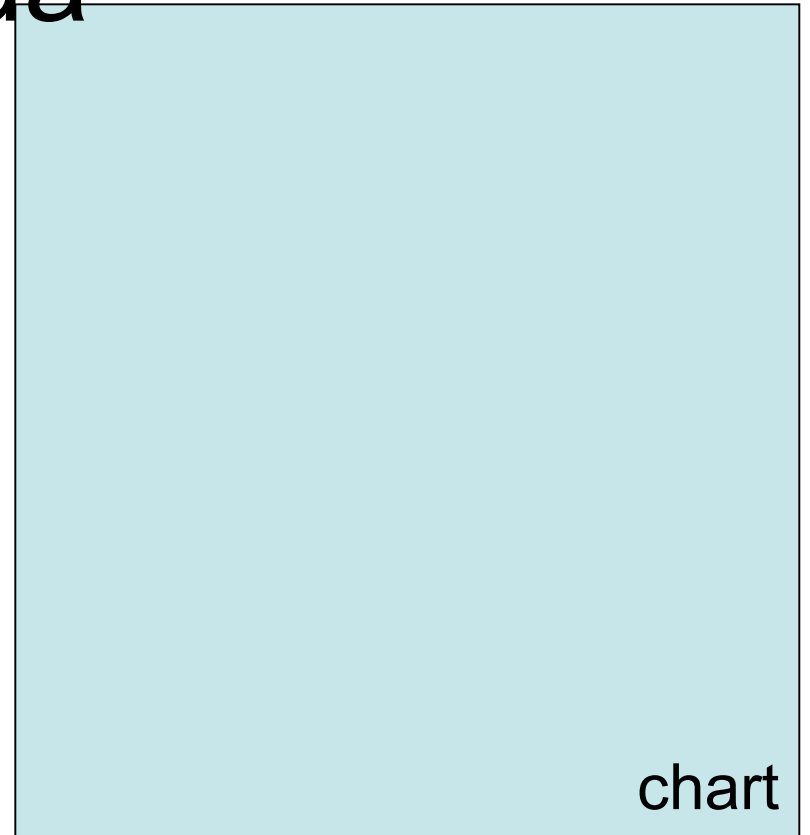
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Agenda

quitters
1



S \rightarrow .8 NP VP NNS \rightarrow .0002 quitters

VP \rightarrow .6 RB VB NP \rightarrow .1 NNS

RB \rightarrow .04 never

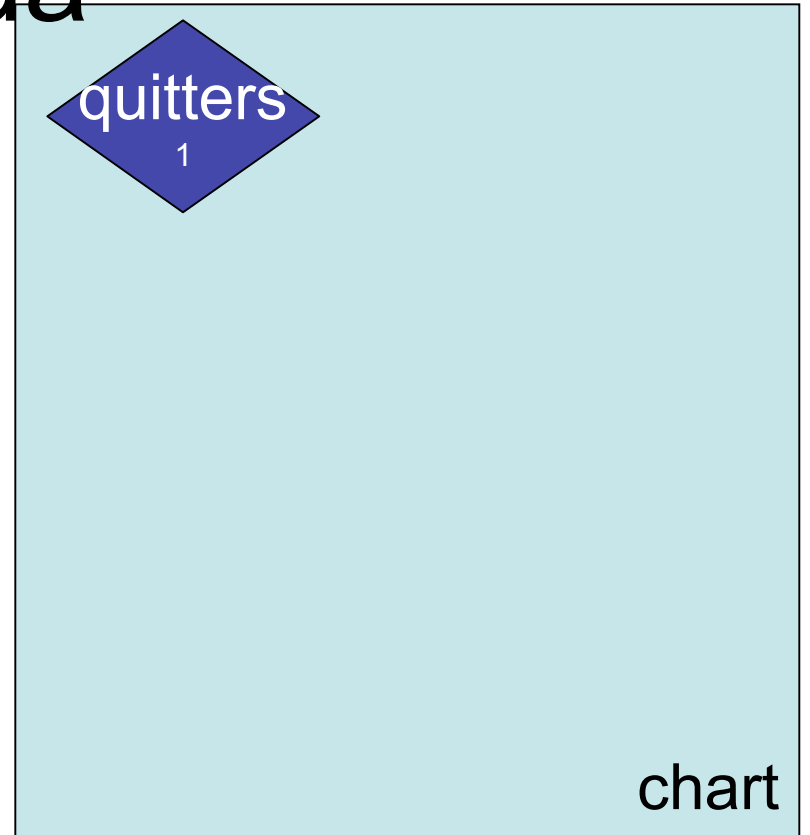
VB \rightarrow .006 win NN \rightarrow .00002 win

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never
2

win
3

Agenda



S \rightarrow .8 NP VP

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VP \rightarrow .6 RB VB

NP \rightarrow .1 NNS

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VB \rightarrow .006 win

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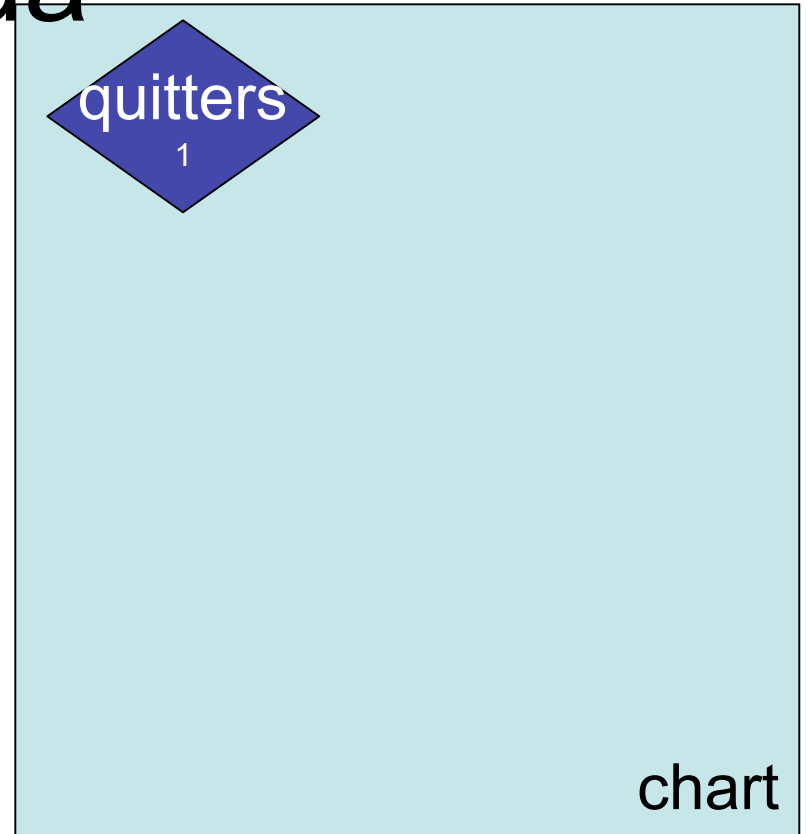
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Agenda

never
2

quitters
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S \rightarrow .8 NP VP

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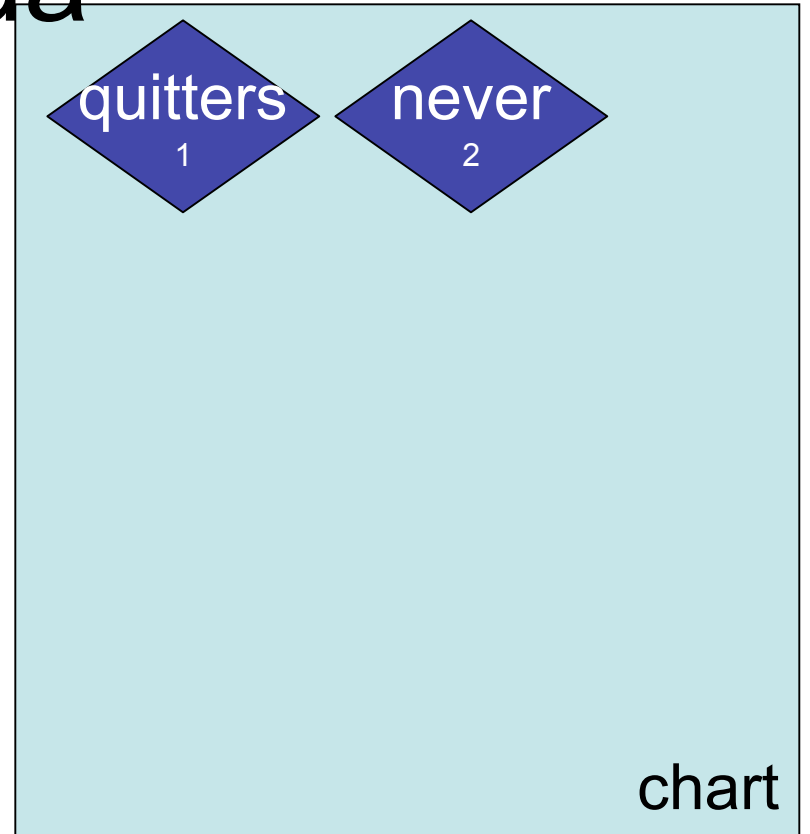
VB \rightarrow .006 win

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win
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Agenda



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Agenda

win
3

quitters
1

never
2

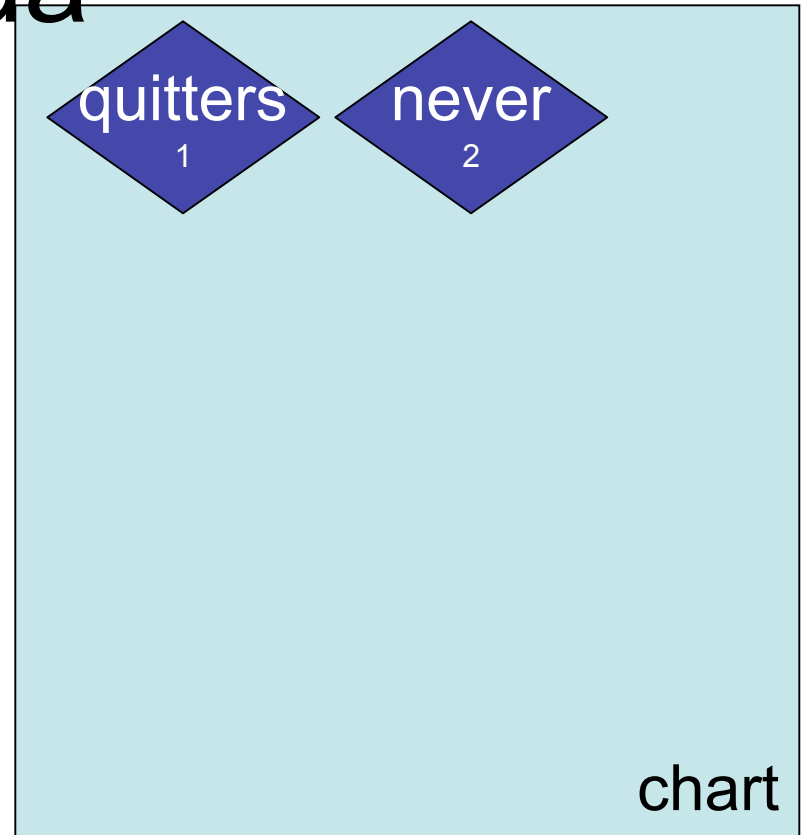
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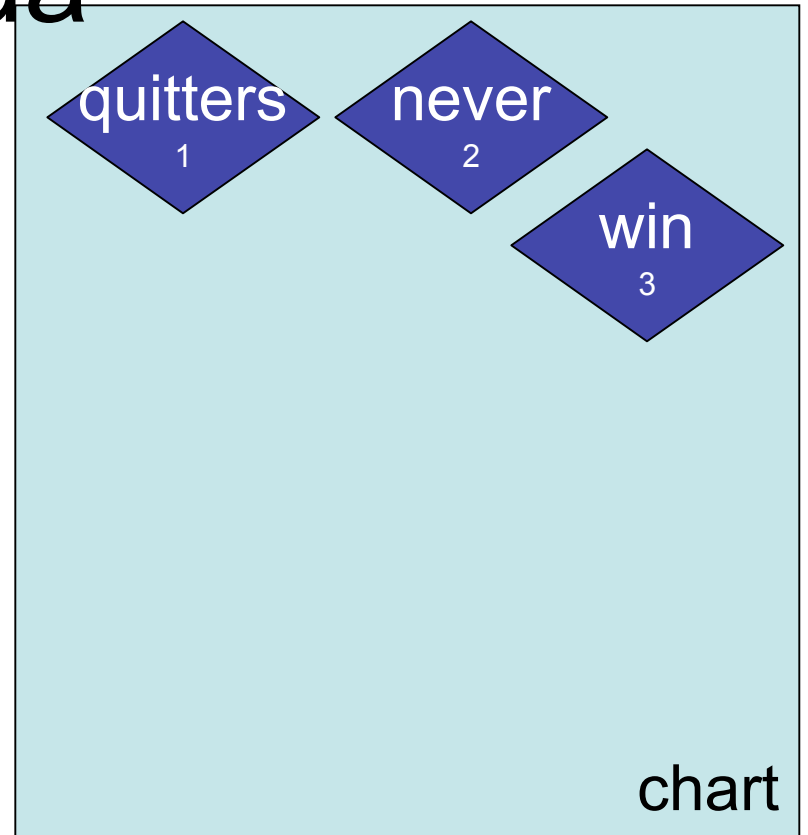
VB \rightarrow .006 win NN \rightarrow .00002 win

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chart

Agenda



S \rightarrow .8 NP VP

NNS \rightarrow .0002 quitters

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S \rightarrow .8 NP VP

NNS \rightarrow .0002 quitters

VP \rightarrow .6 RB VB

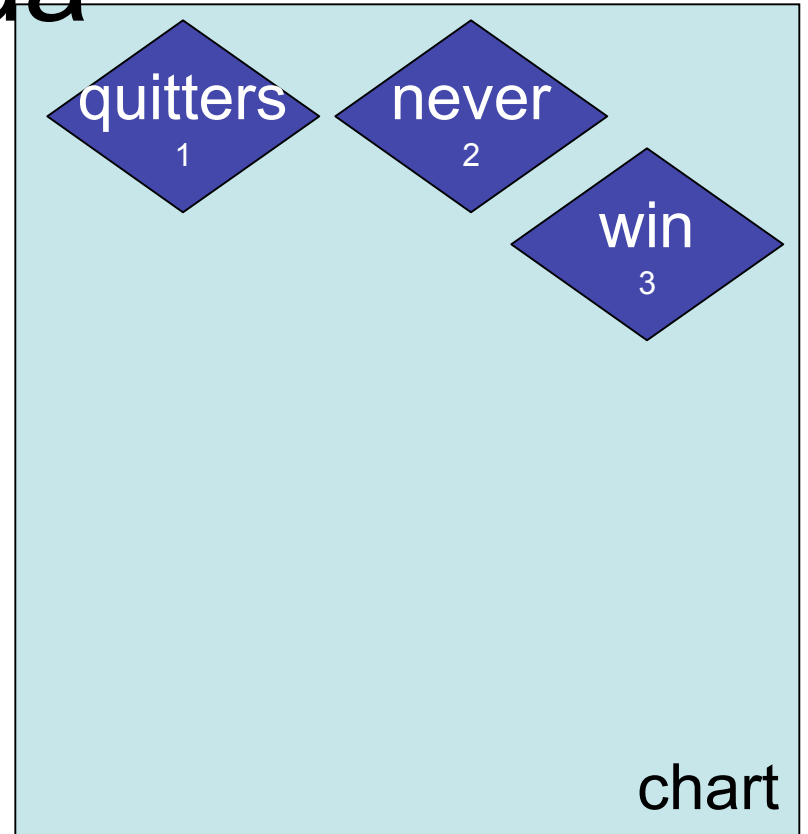
NP \rightarrow .1 NNS

RB \rightarrow .04 never

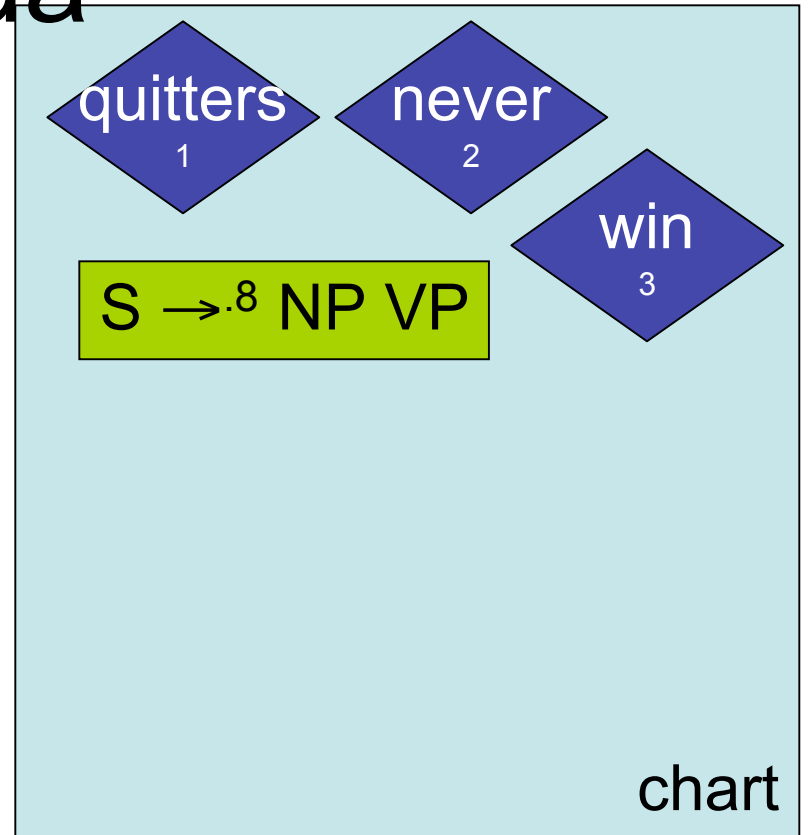
VB \rightarrow .006 win

NN \rightarrow .00002 win

...



Agenda



NNS →.0002 quitters

VP →.6 RB VB

NP →.1 NNS

RB →.04 never

VB →.006 win

NN →.00002 win

...

Agenda

VP \rightarrow .6 RB VB

NNS \rightarrow .0002 quitters

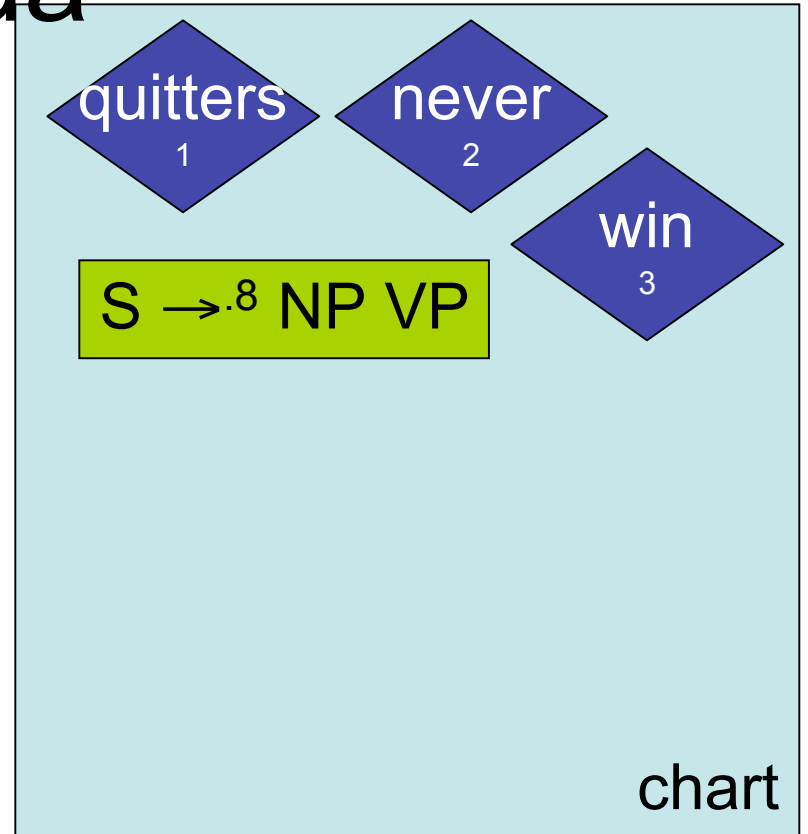
NP \rightarrow .1 NNS

RB \rightarrow .04 never

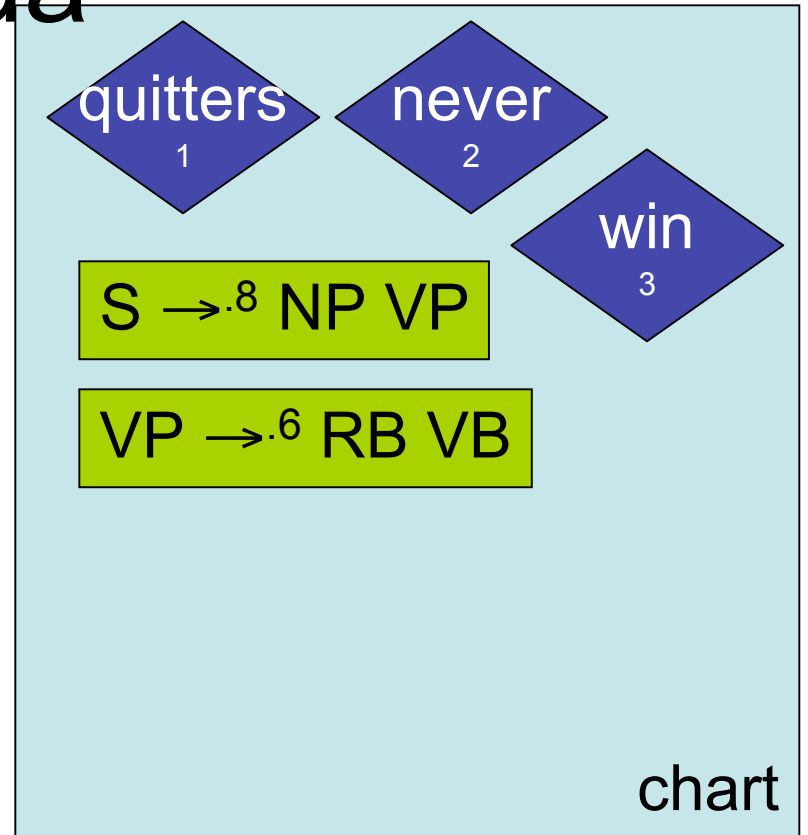
VB \rightarrow .006 win

NN \rightarrow .00002 win

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NNS → .0002 quitters

NP → .1 NNS

RB → .04 never

VB → .006 win

NN → .00002 win

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NP \rightarrow .1 NNS

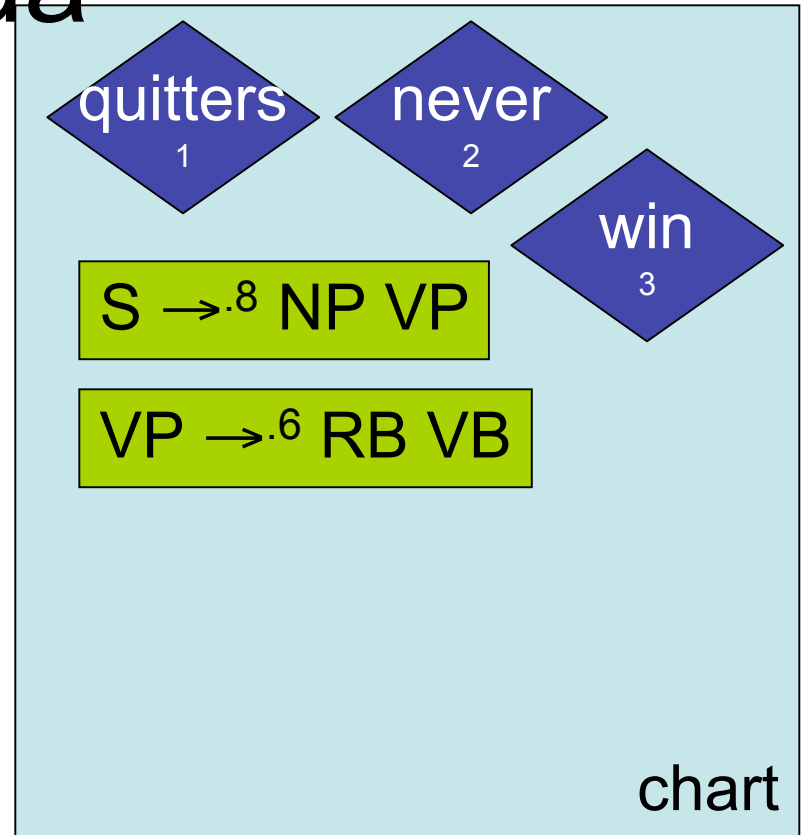
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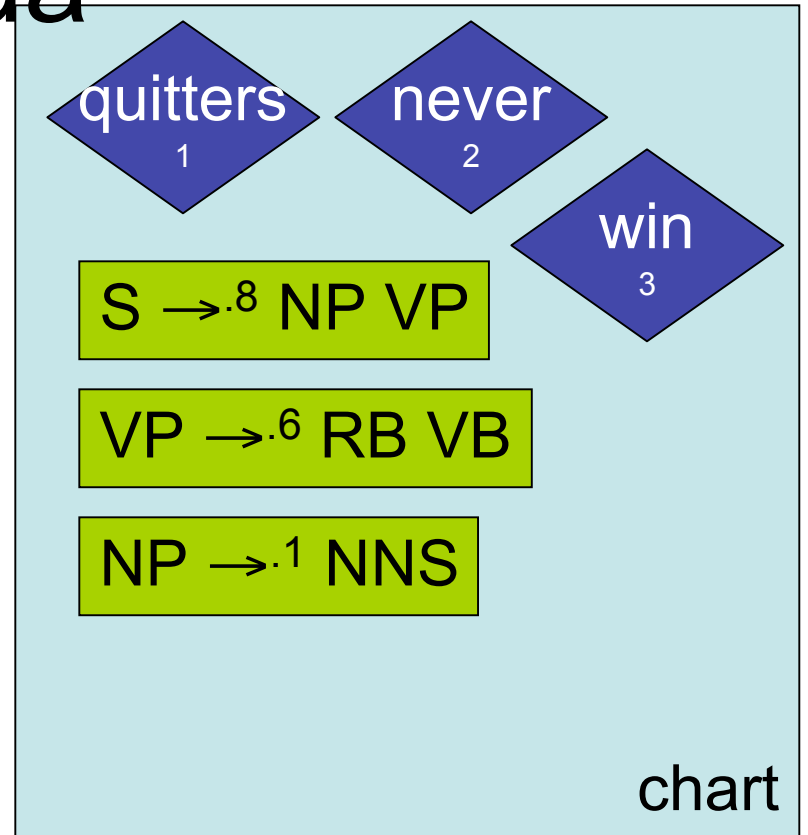
VB \rightarrow .006 win

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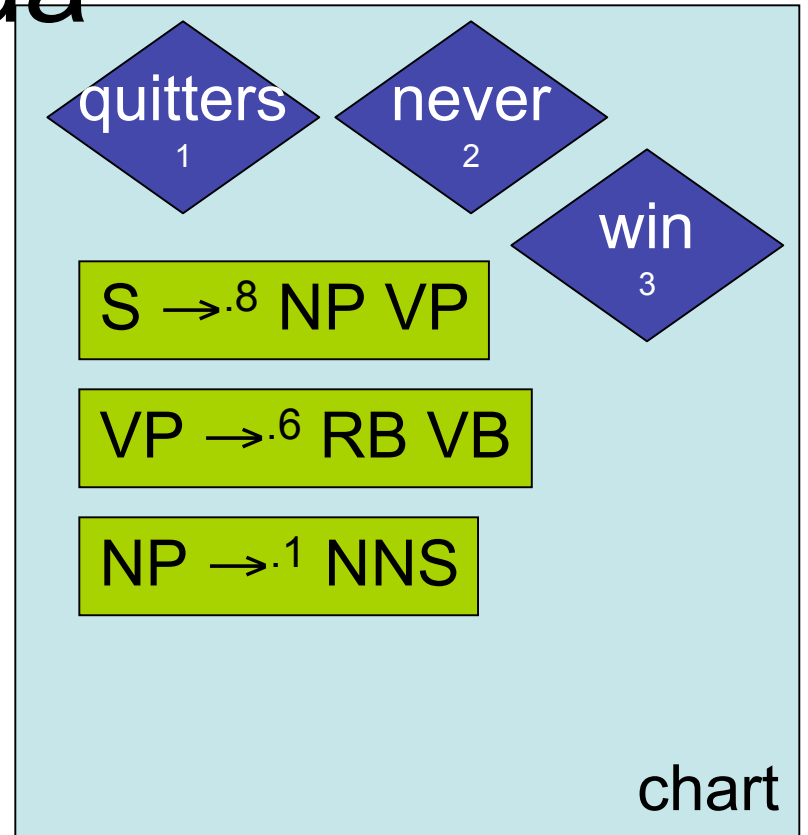
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RB → .04 never

NNS → .0002 quitters

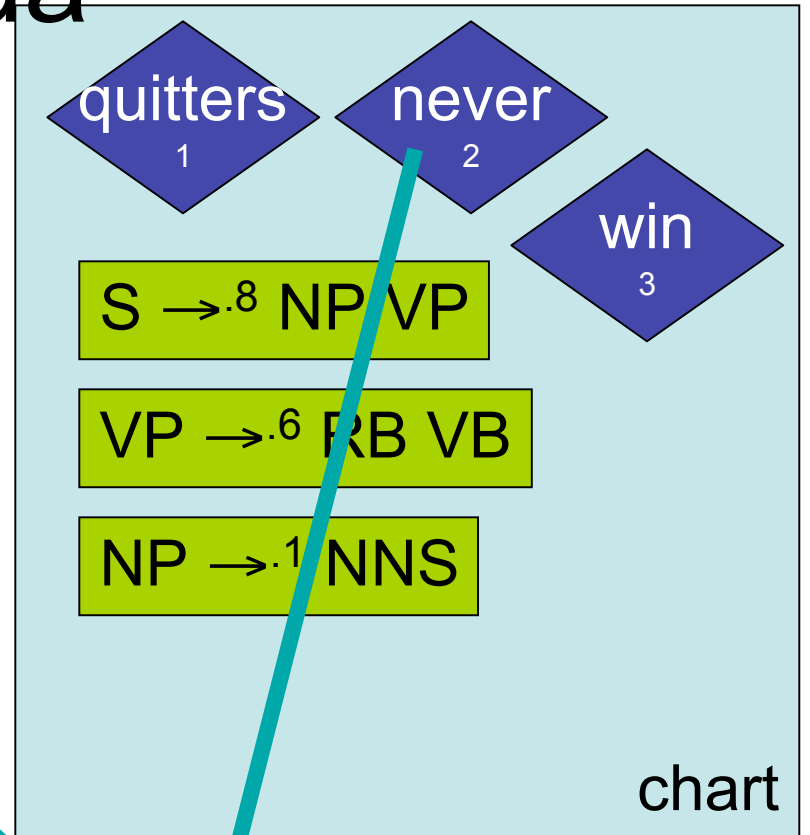


VB → .006 win

NN → .00002 win

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RB → .04 never

NNS → .0002 quitters

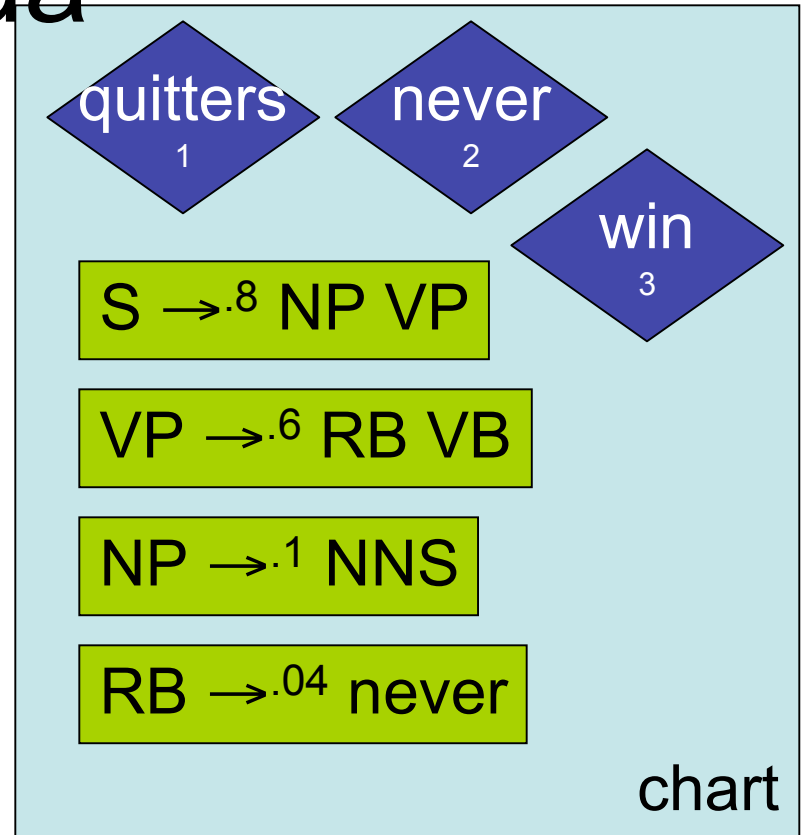
VB → .006 win

NN → .00002 win

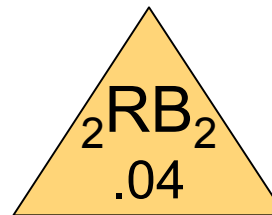
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${}_2\text{RB}_2$
.04

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NNS →.0002 quitters

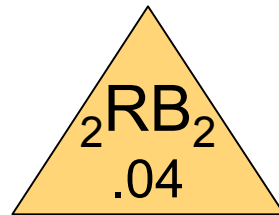


VB →.006 win

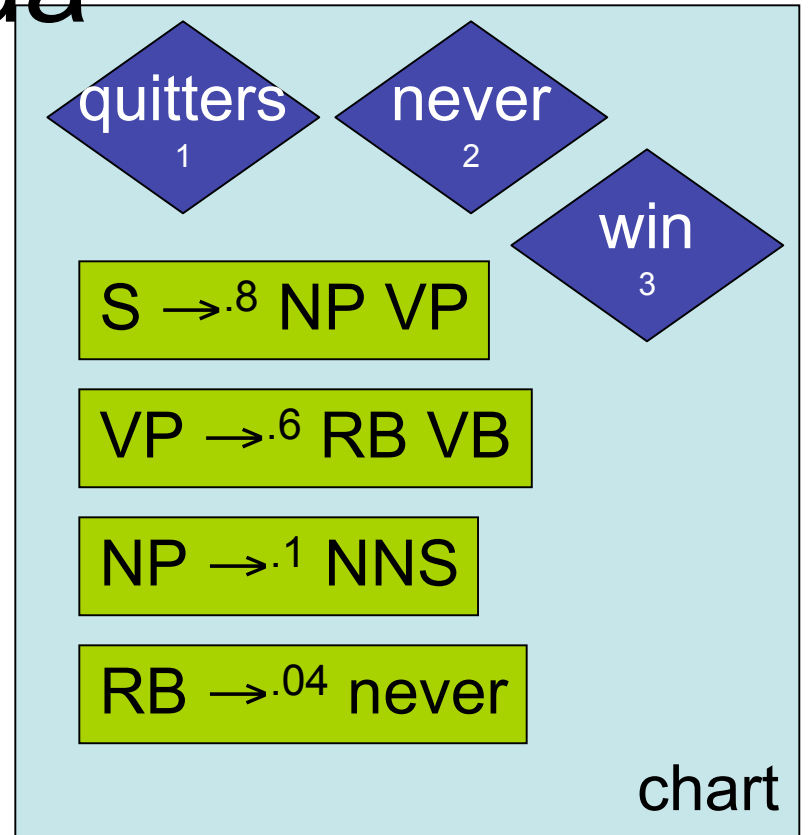
NN →.00002 win

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$NNS \rightarrow .0002$ quitters

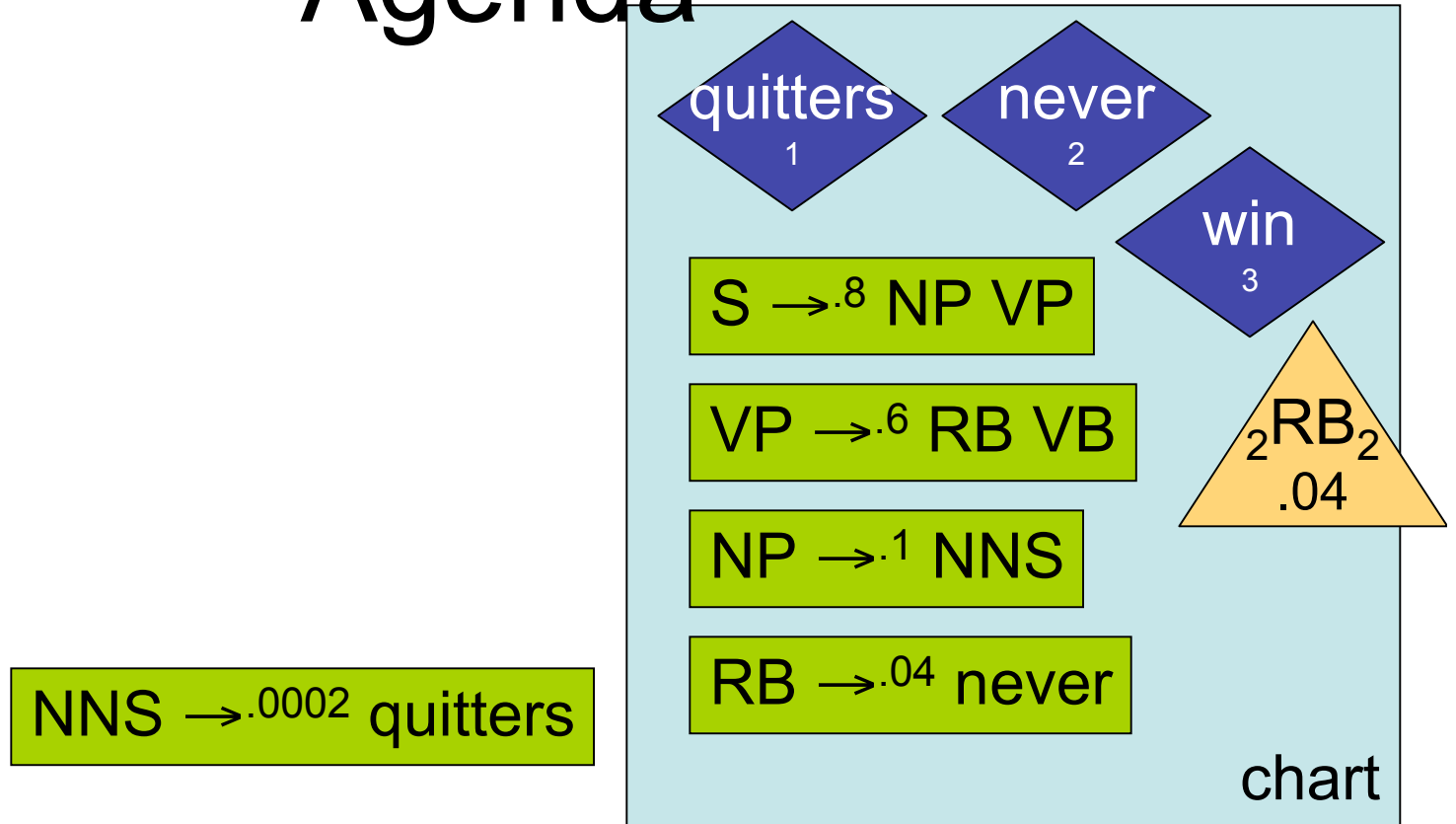


$VB \rightarrow .006$ win

$NN \rightarrow .00002$ win

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VB → .006 win

NN → .00002 win

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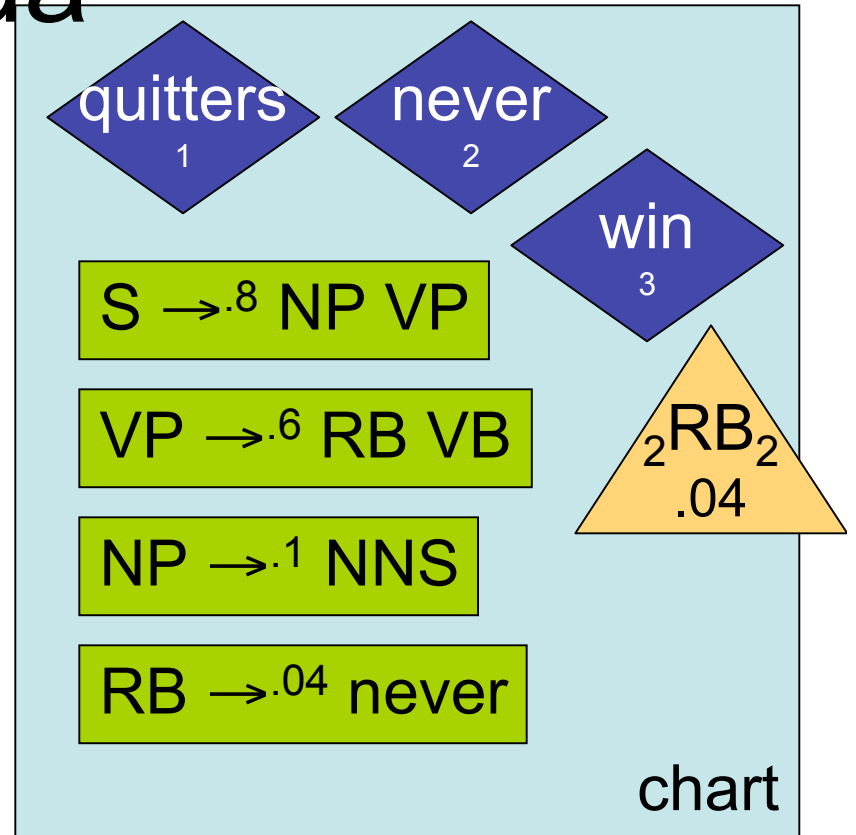
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VB → .006 win

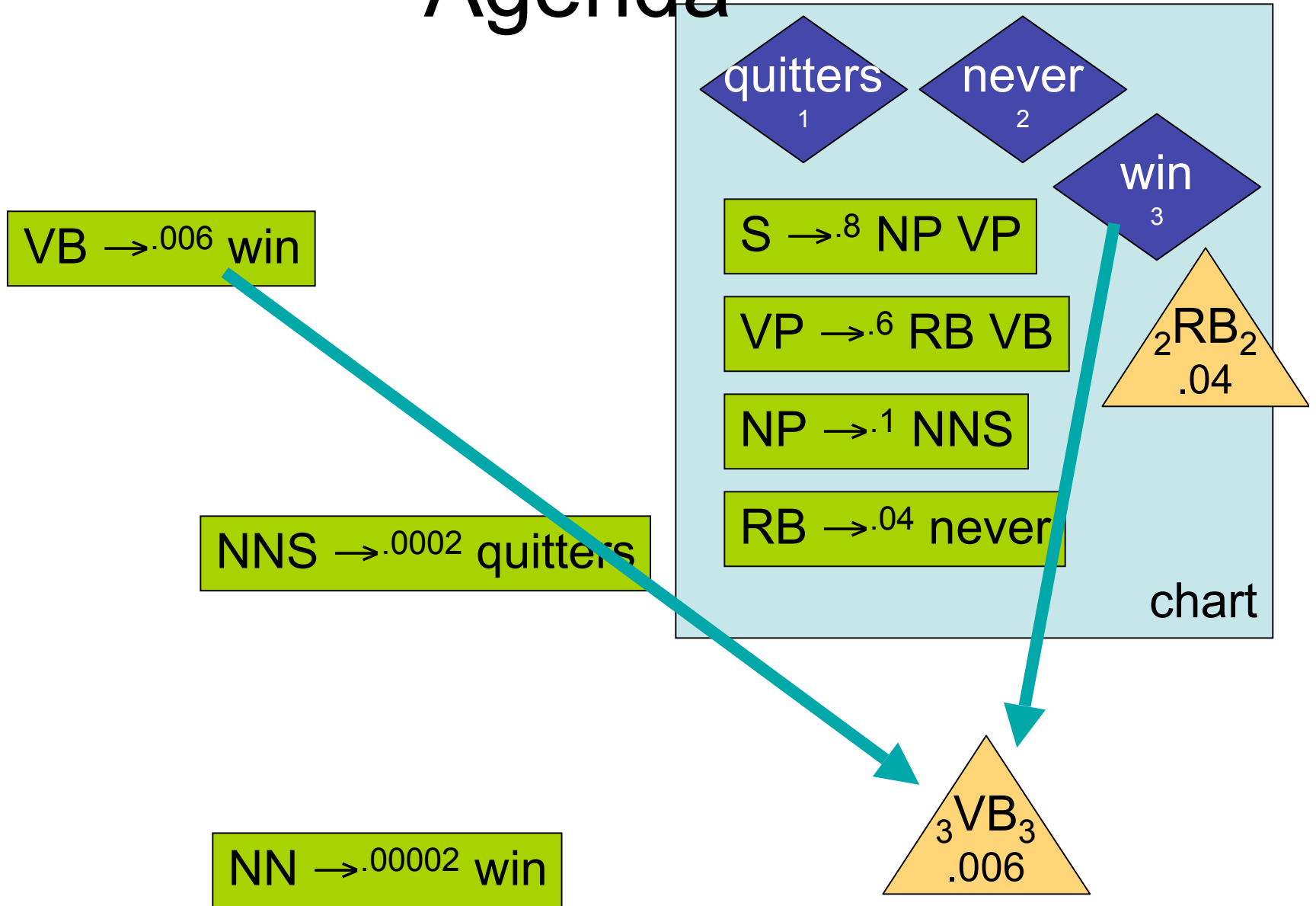
NNS → .0002 quitters

NN → .00002 win

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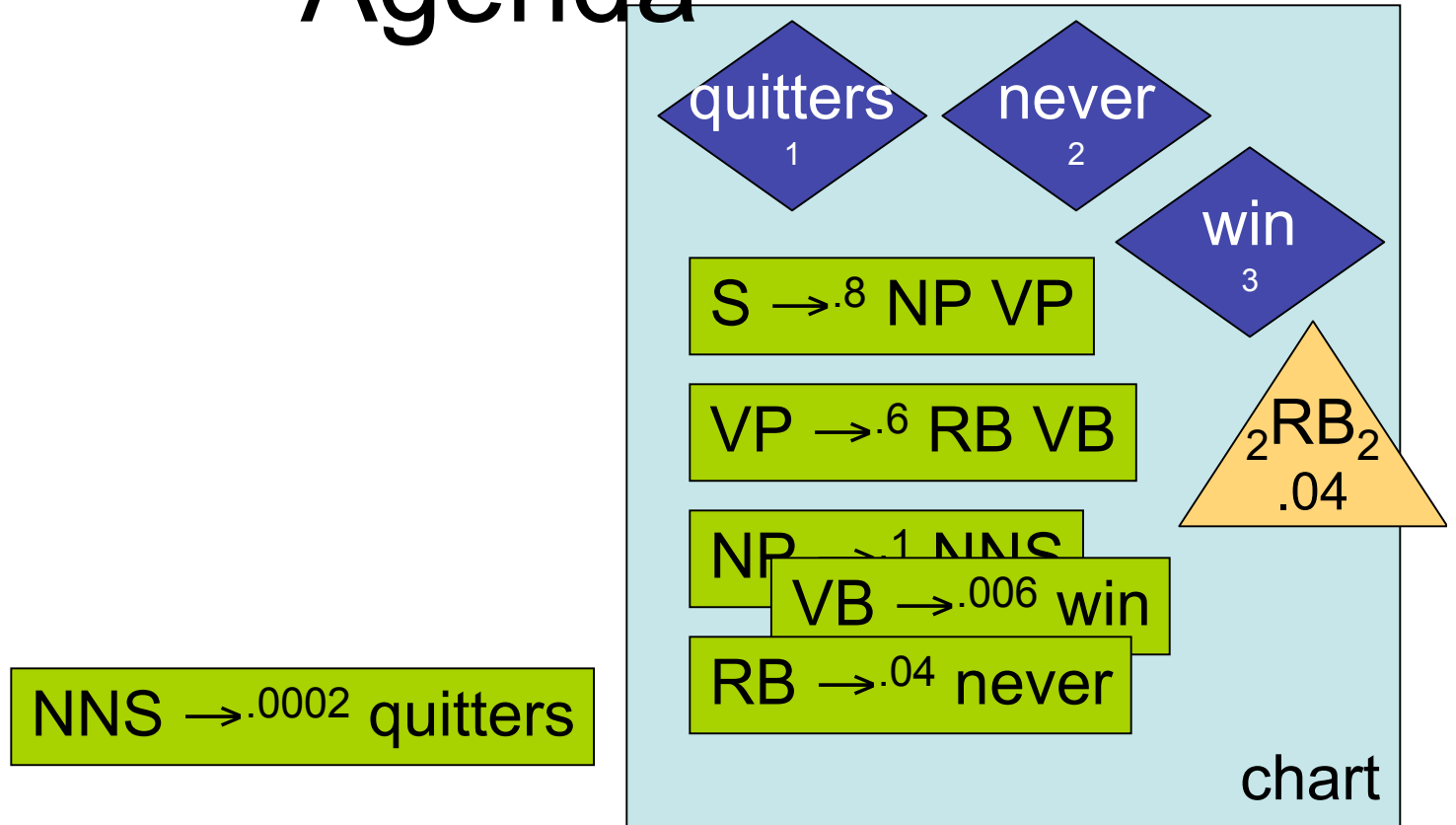


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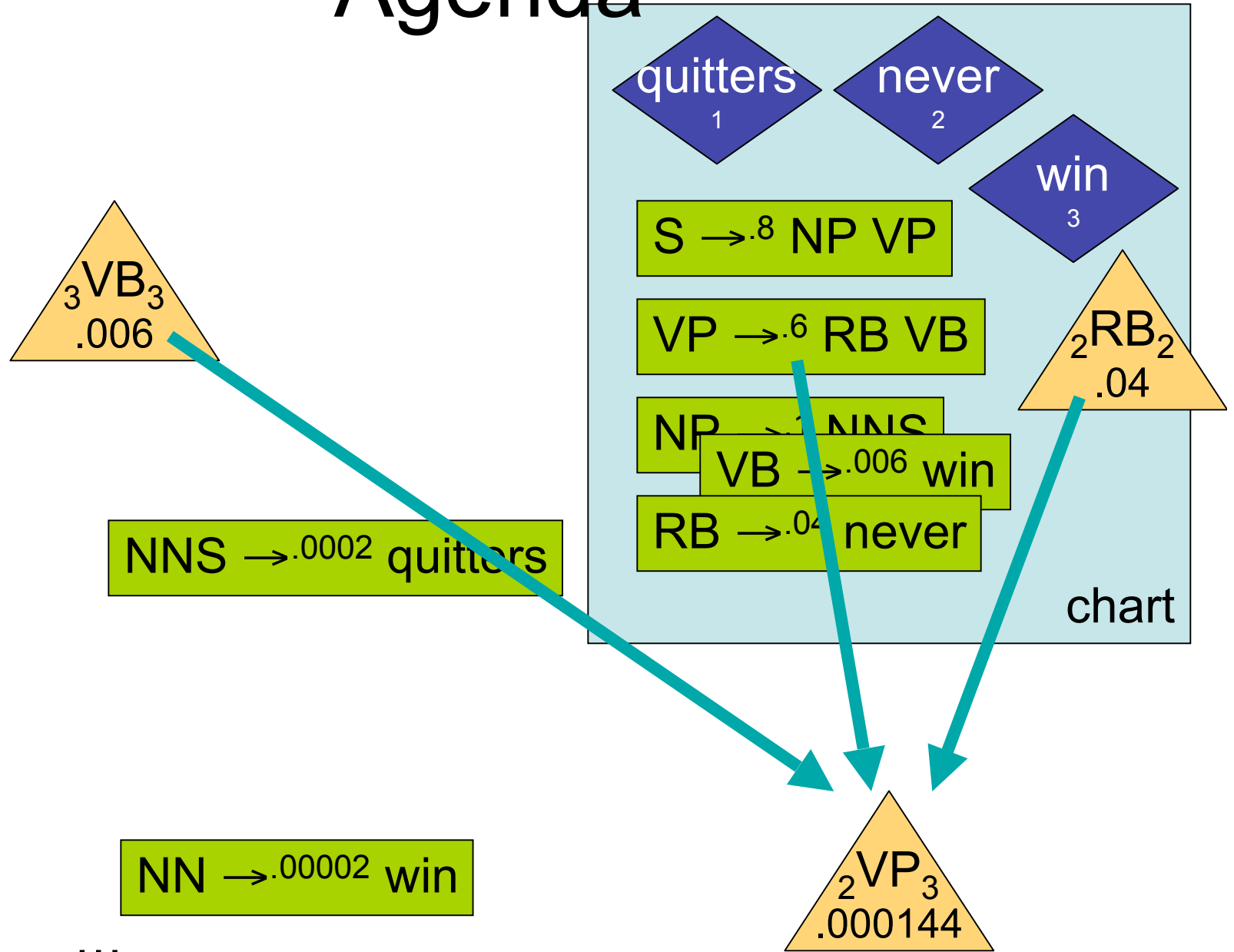


NN → .00002 win

³VB₃ .006

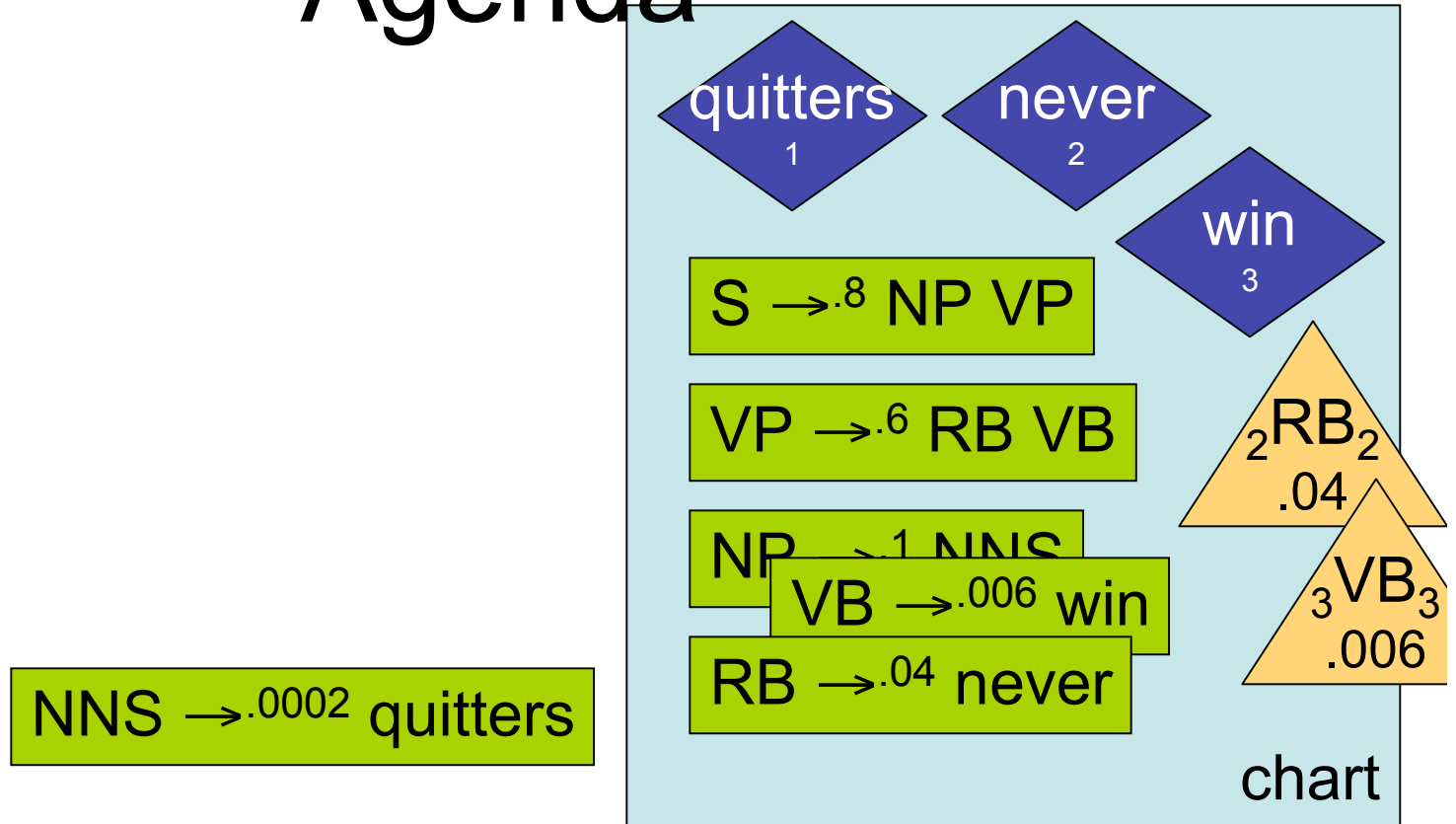
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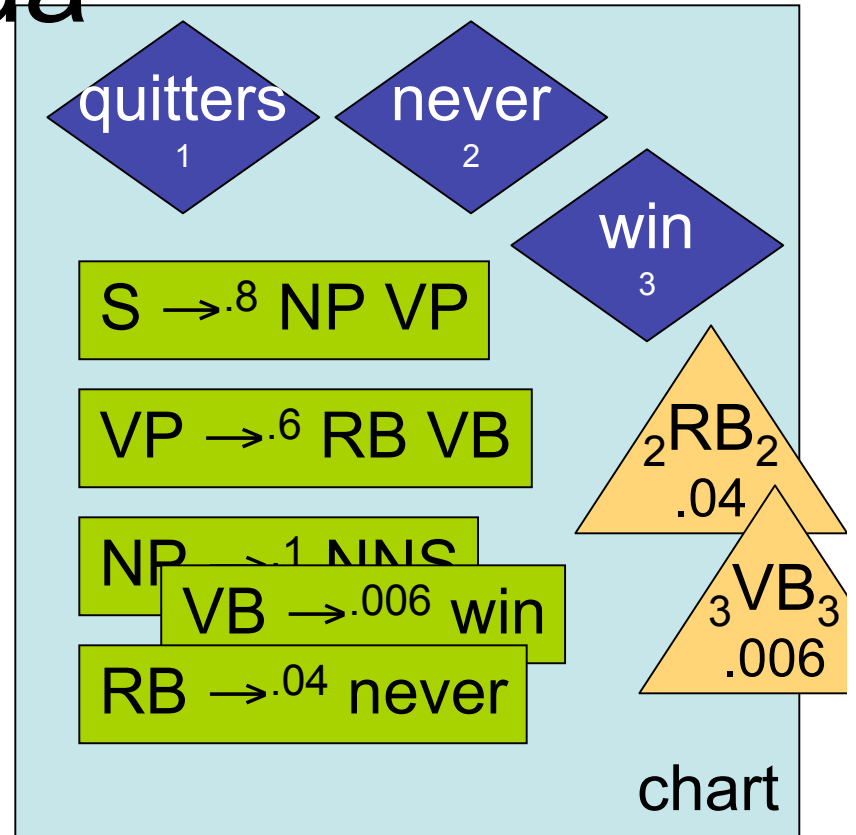
NN → .00002 win

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2VP₃ .000144

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NNS → .0002 quitters

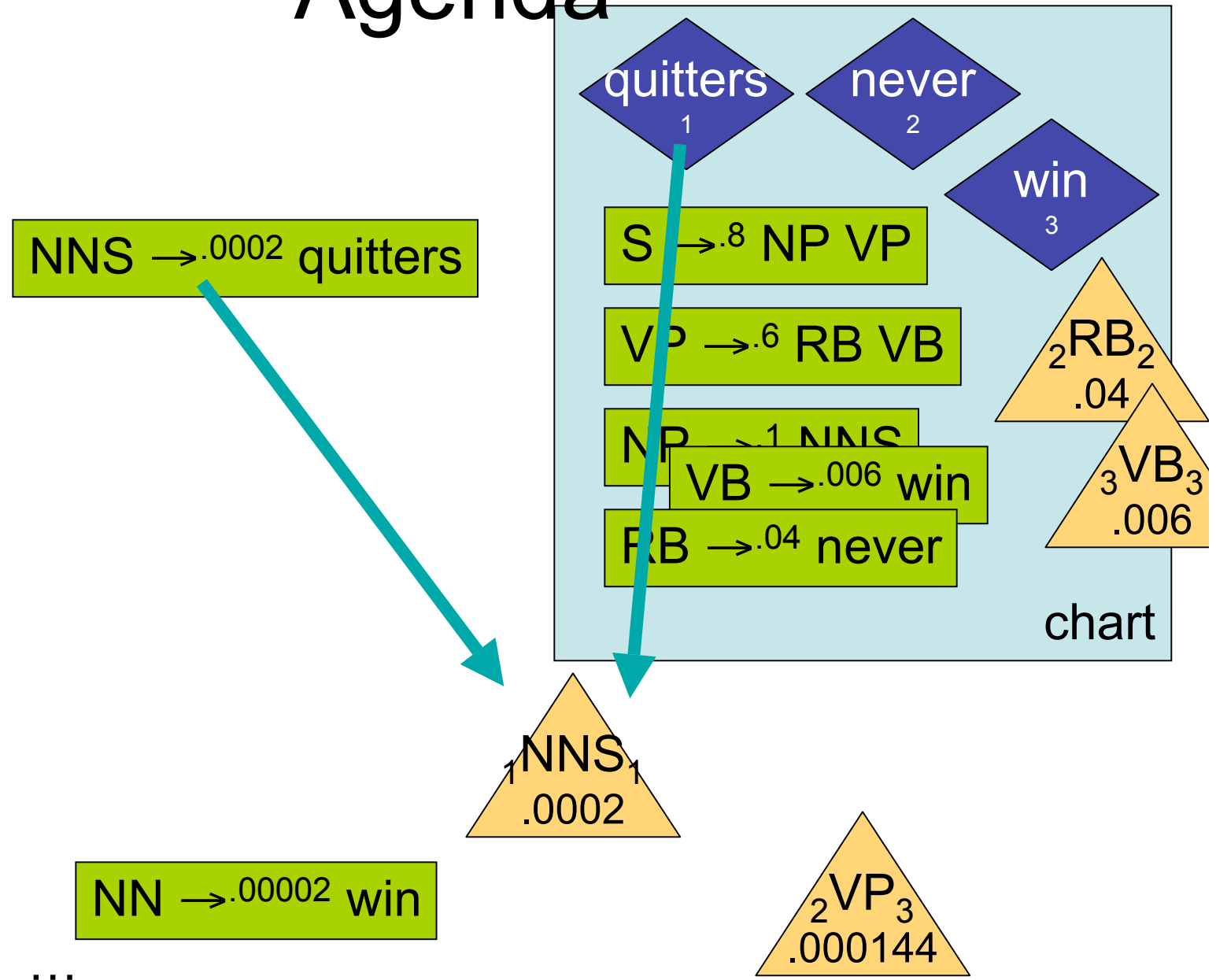


NN → .00002 win

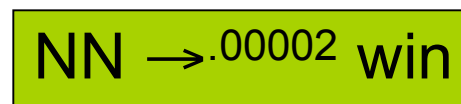
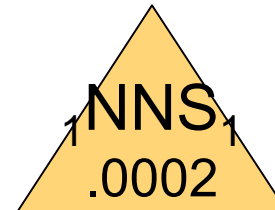
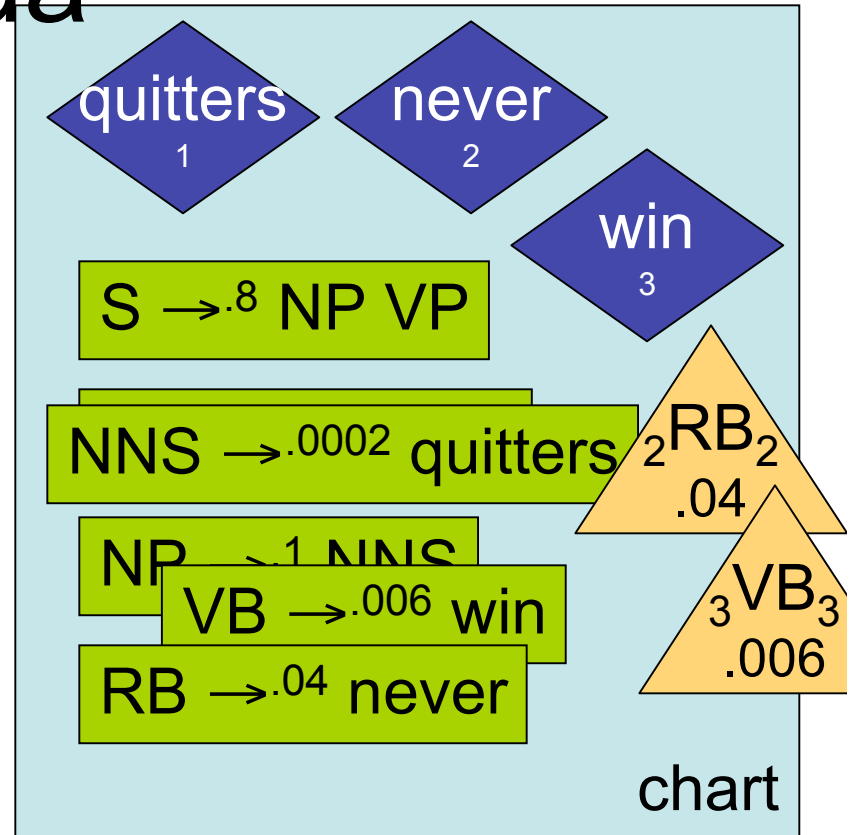
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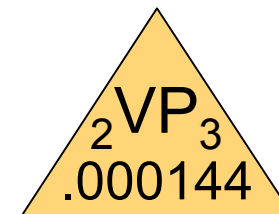
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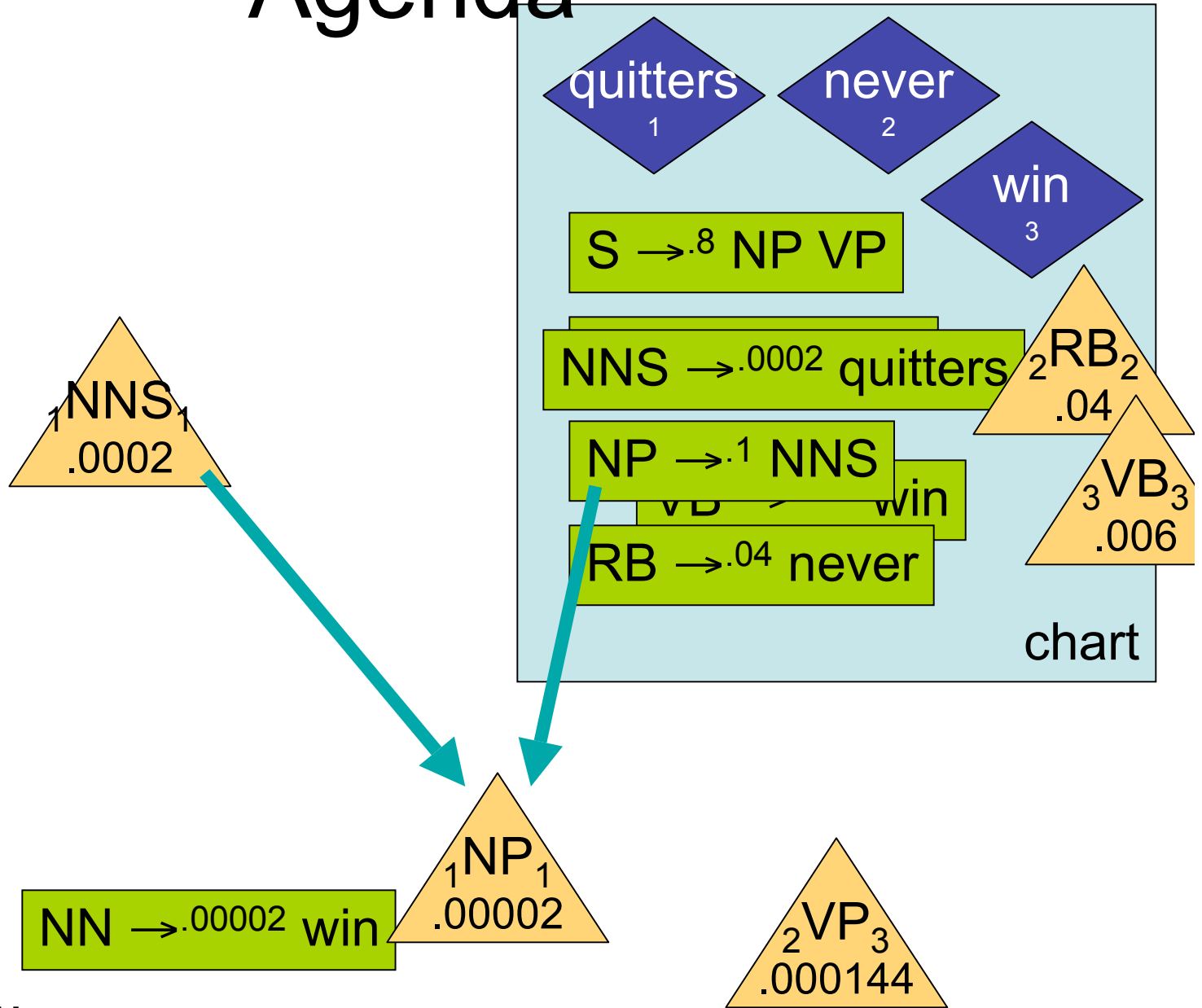
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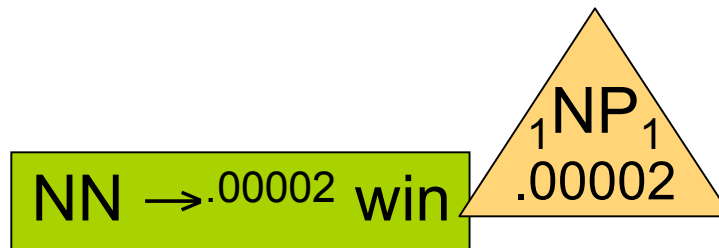
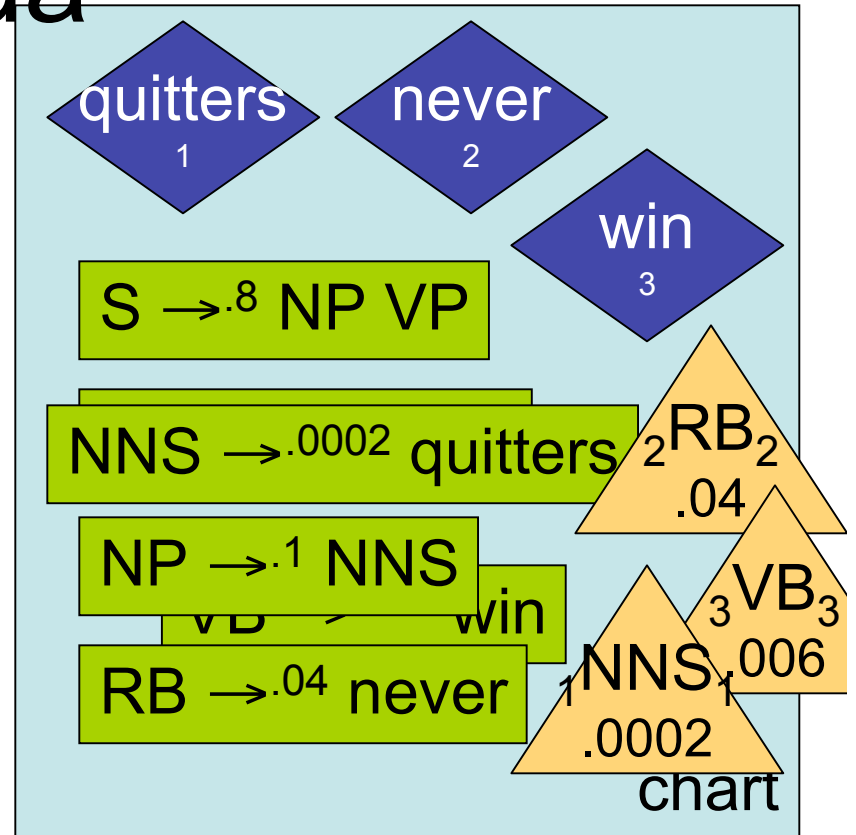
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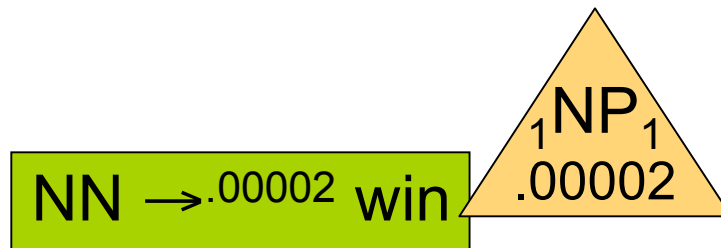
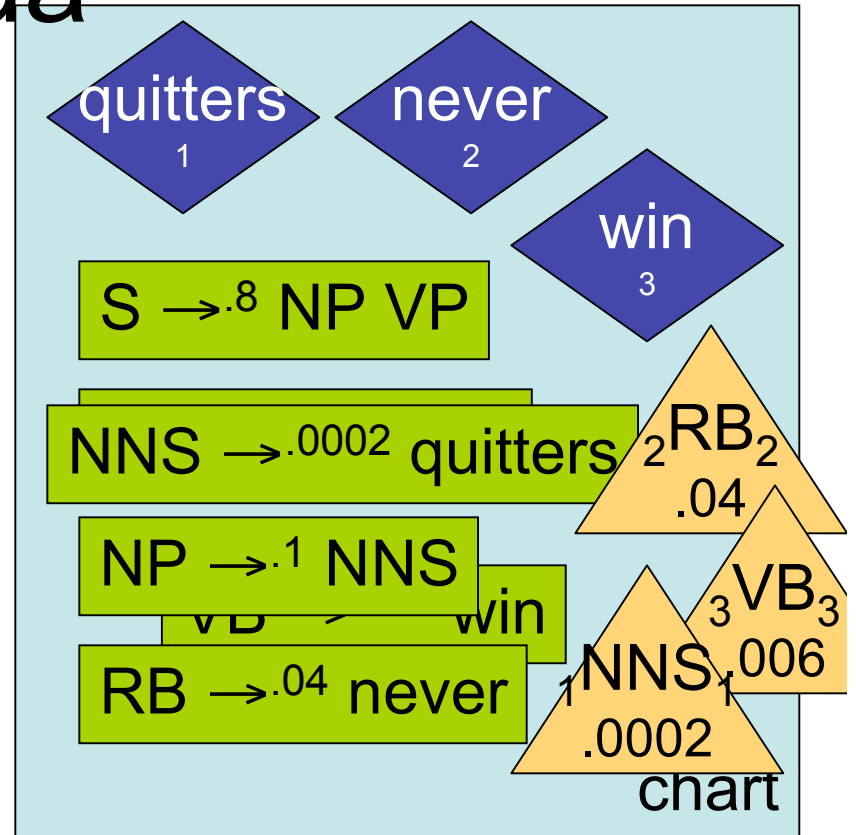


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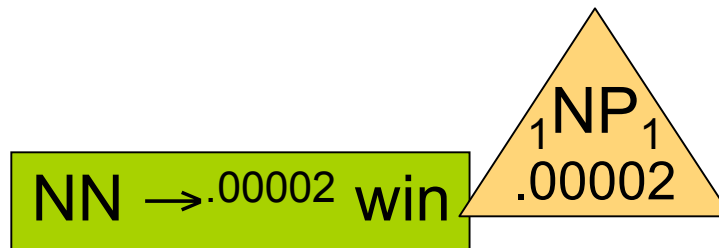
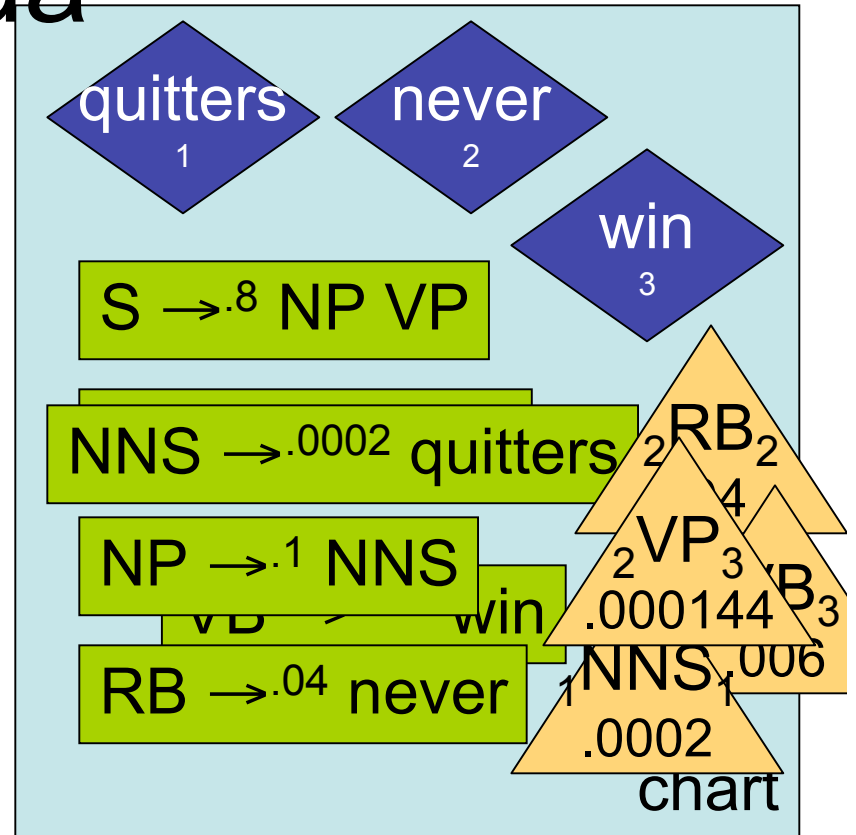
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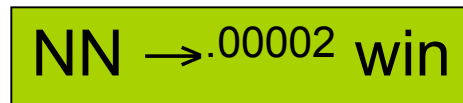
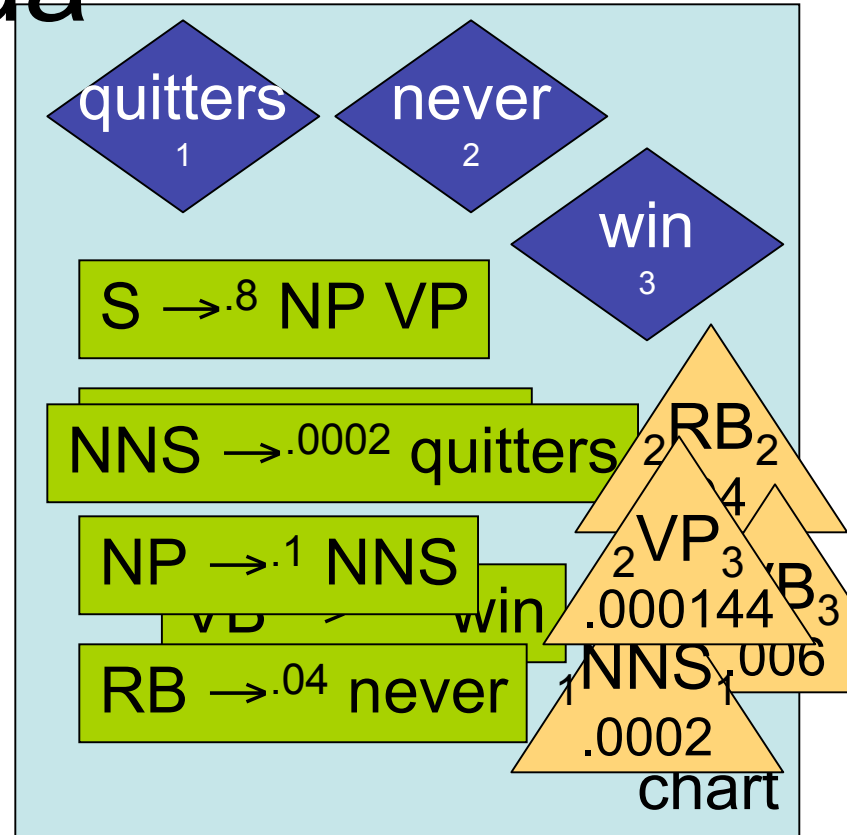
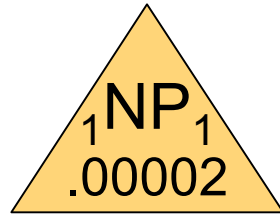
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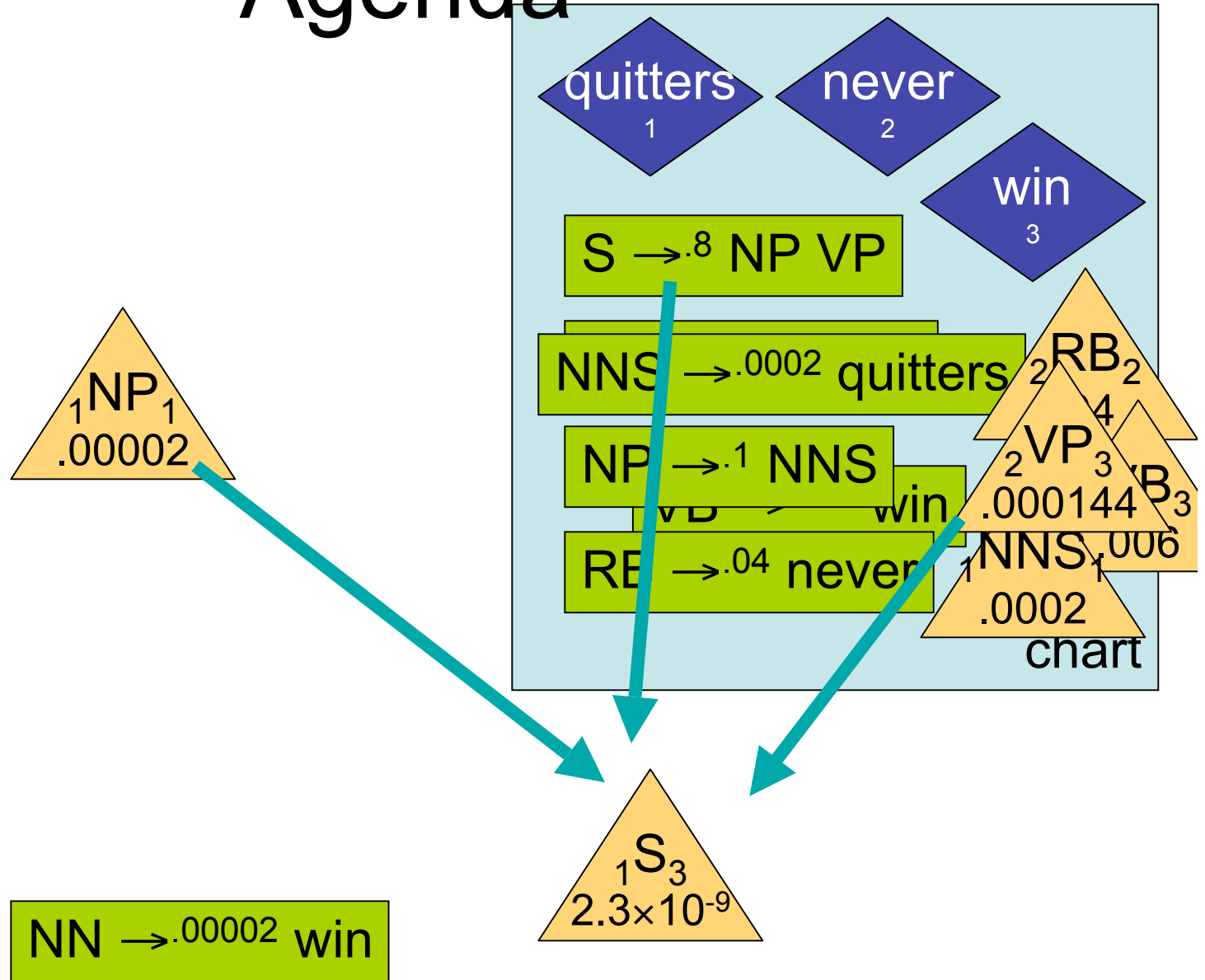
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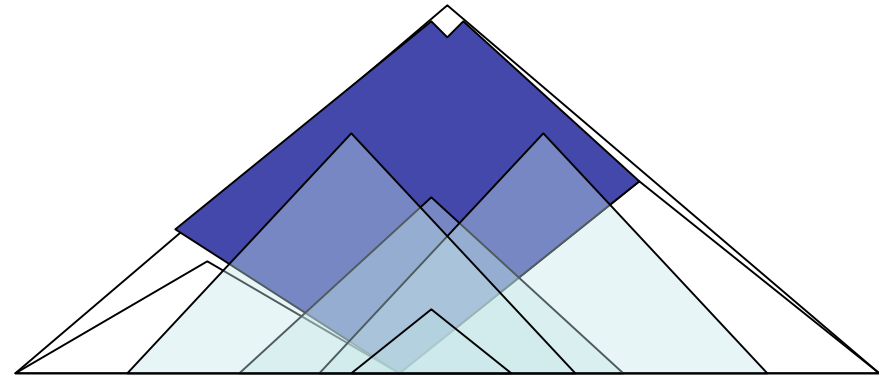
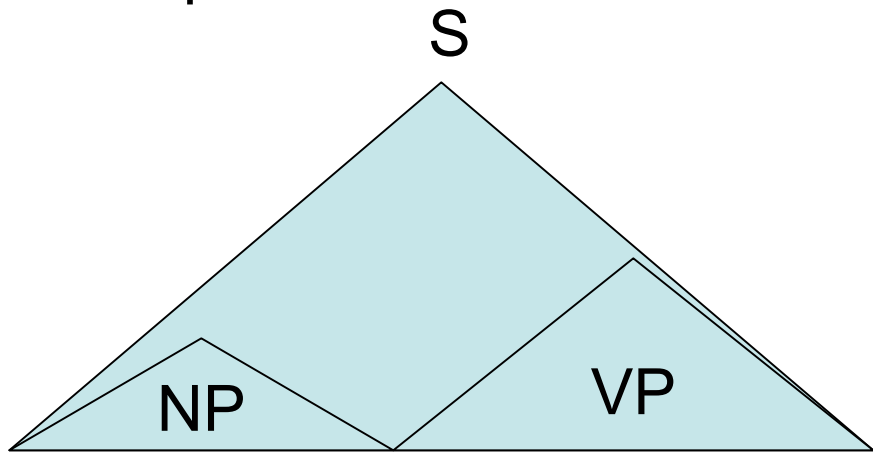


Unnecessary Work

- If you only want the best derivation, you don't want to build items that aren't in it!
- But you don't know which items to build until you have the best parse.
- Key idea in the agenda:
 - Intelligently order **updates** to items' weights.
 - Roughly analogous to trading depth and breadth in **search**.
- Note: for exact inside/outside, all of the work **is** necessary!

Unnecessary Work

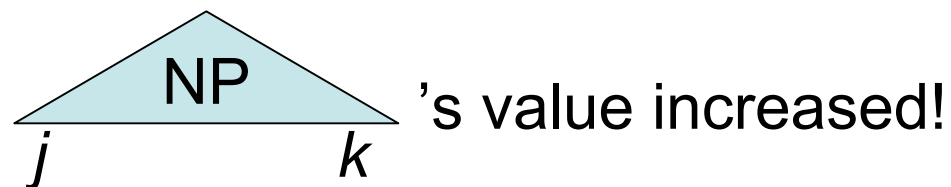
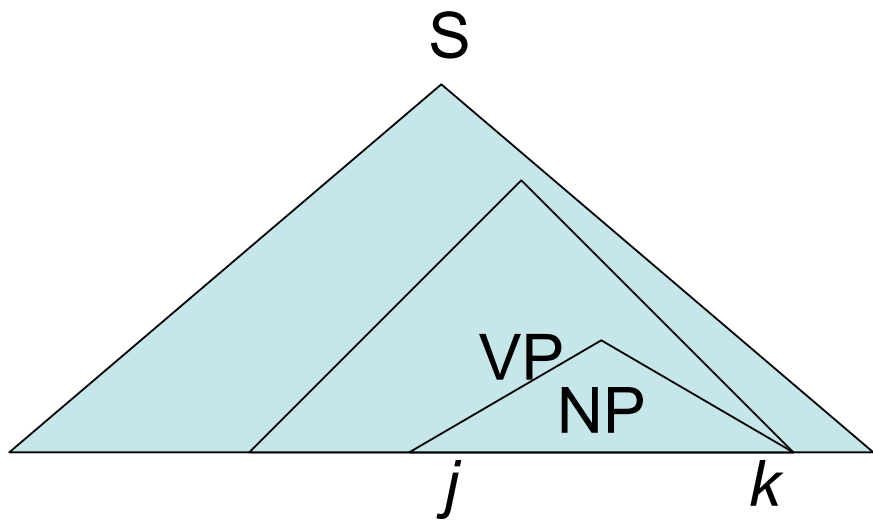
best parse:



Repropagation

- Suppose $(NP, 4, 7)$ currently has a weight of 0.3, constructed by $(DT, 4, 5) \otimes (NP, 5, 7)$.
- Now suppose we find that a **better** way to build $(NP, 4, 7)$: $(DT, 4, 5) \otimes (NNP, 5, 6) \otimes (NNP, 6, 7)$ with value 0.31.
- Maybe now we have a better way to build $(VP, 3, 7)$! (Or anything else that used $(NP, 4, 7)$).
- Have to re-build all of those consequents, and compare again, and recursively repropagate to consequents of any item whose value changes.
- May not be $O(n^3)$ anymore!

Repropagation



So any **consequent** of (NP, j, k) might also increase.

Best-First Parsing

- **Viterbi** semiring (find the best parse): see Nederhof (2003) and Knuth (1977).
 - Cf. Goodman, build the chart and fill in weights at the same time.
- Many parsers in practice: prune, prune, prune.
- Alternative: order items by their weights.
 - “**Uniform cost search**”
 - Guarantee: the first time **goal** is popped from the agenda, you have the optimal parse.
- Charniak et al., 1998: heuristics to speed this up. “Figures of Merit” (big speed payoff).

Priorities

$\text{priority}(I) = \text{weight of the best parse that uses } I$

$$= \text{inside}(I) \otimes \text{outside}(I)$$

$$\leq \text{inside}(I) \otimes \underbrace{\text{estimate}(\text{outside}(I))}_{\text{uniform cost search:1}}$$

- Klein and Manning (2003): Blocked more than 90% of edges!
- Generalization of A* search (for hypergraphs instead of graphs).
- Heuristics? Computed by simpler, cheaper dynamic programs!
- Caveat emptor: only for Viterbi (max) semirings!

Dyna (Eisner et al., 2005)

- **Dyna** is a high-level programming language (like Prolog) for weighted deduction.
- Source code looks like Prolog.
- Compiles into C++.
- Core algorithms:
 - Generalized weighted, **prioritized** agenda.
 - Allows the use of heuristics, including A*
 - Handles repropagation if required
 - Efficient “tape” mechanism for **reverse** computation.
 - Very similar to backpropagation.

Dyna Programs

```
constit(X,I,J) += word(W,I,J) * rewrite(X,W).  
constit(X,I,J) += constit(Y,I,Mid) * constit(Z,Mid,J) * rewrite(X,Y,Z).  
goal           += constit("s",0,N) whenever length(N).
```

```
constit(X,I,J) max= word(W,I,J) * rewrite(X,W).  
constit(X,I,J) max= constit(Y,I,Mid) * constit(Z,Mid,J) * rewrite(X,Y,Z).  
goal           max= constit("s",0,N) whenever length(N).
```

```
constit(X,I,J) max= word(W,I,J) * rewrite(X,W).  
constit(X,I,J) max= constit(Y,I,Mid) * inter(X, Z, Mid, J).  
inter(X, Z, Mid, J) max= constit(Z,Mid,J) * rewrite(X,Y,Z).  
goal           max= constit("s",0,N) whenever length(N).
```

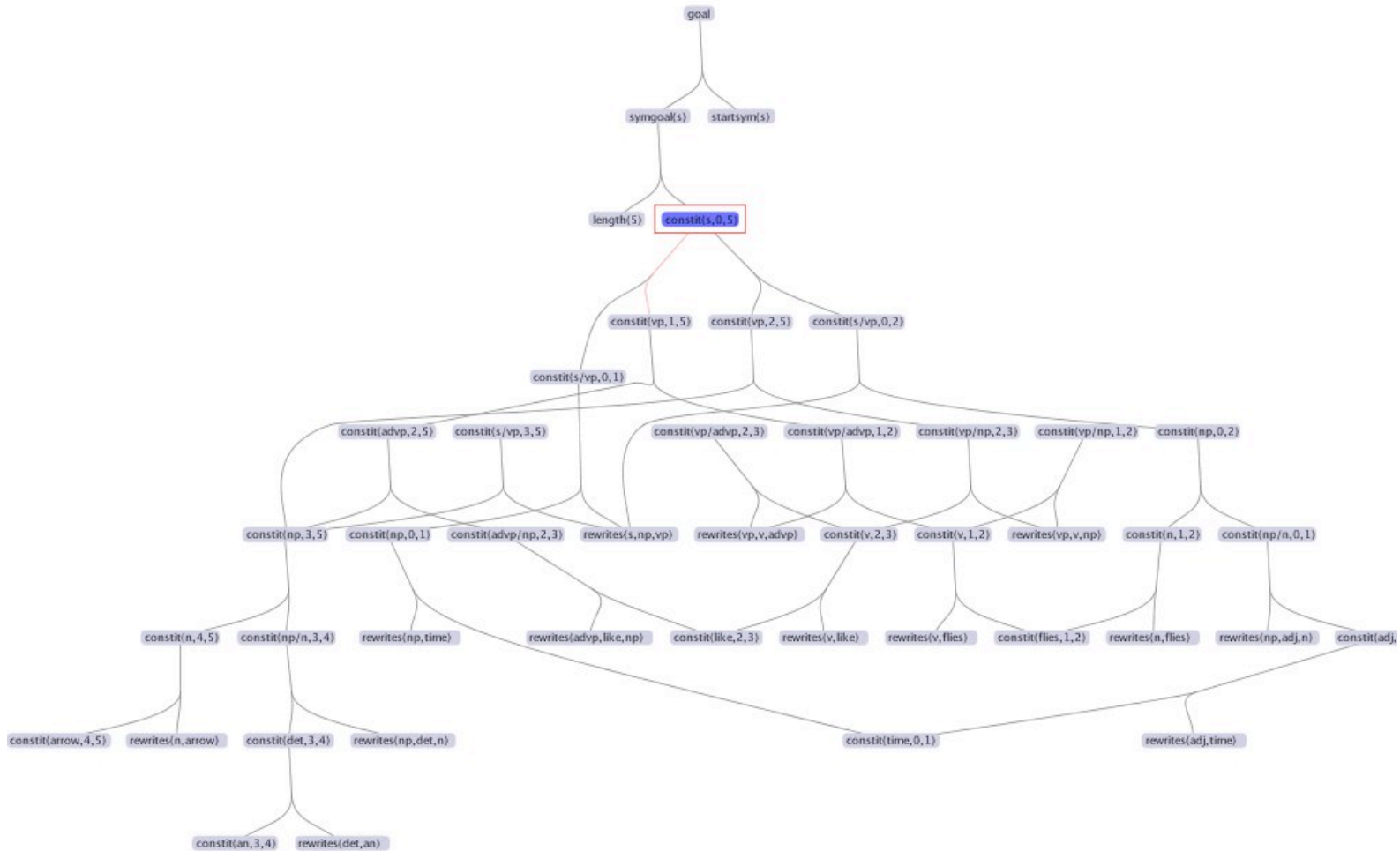
Dyna Programs

```
constit(X,I,J) += word(W,I,J) * rewrite(X,W).  
constit(X,I,J) += constit(Y,I,Mid) * constit(Z,Mid,J) * rewrite(X,Y,Z).  
goal           += constit("s",0,N) whenever length(N).
```

```
reverseconstit(Y,I,Mid) += reverseconstit(X,I,J) * constit(Z,Mid,J) * rewrite(X,Y,Z).  
reverseconstit(Z,Mid,J) += constit(Y,I,Mid) * reverseconstit(X,I,J) * rewrite(X,Y,Z).  
reverseconstit("s",0,N) += 1.
```

Dyna Debugger

File Display Selection Preferences



const(s,0,5)

Shift - select nodes
LeftButton - scroll/select

Alt - stop at edges
RightButton - zoom

Ctrl - move pink trail

Parting Shots

- Weighted deduction as a convenient way to
design,
improve,
understand,
analyze,
unify,
transform,
and implement
otherwise tricky dynamic programming algorithms.