Loop Notes

```
\mathbf{Input} \colon \mathrm{Minimal} \ \mathrm{annotated} \ \mathrm{consistent} \ \mathrm{partial} \ \mathrm{order} \ \mathcal{P}
Output: Template T representing \mathcal{P}
\mathbf{procedure}\ \mathrm{Convert\_To\_Template}(\mathcal{P}) \colon
     T.conditions \leftarrow Find_Relevant_Current(\mathcal{P}) + Find_Relevant_Goal()
     T.\text{body} \leftarrow \mathcal{P}
     Identify_Loops(T)
procedure Identify_Loops(T):
     change \leftarrow \mathbf{true}
      while (change) do
           \mathrm{change} \leftarrow \mathbf{false}
           ∀ fan outs
                 \forall fans with same sequences:
                       identify varying parameter(s),
                       introduce new variable for them
                       if fans then have same init conds & results then
                            newloop \leftarrow empty while loop
                            newloop.conditions \leftarrow fans.init conds & results
                            newloop.body \leftarrow fan.sequence
                            replace similar fans with newloop
                             reconnect condition and result arcs
                             change \leftarrow \mathbf{true}
           \forall sequences
                 use string-matching algs to find repeated seqs
                 \forall repeated sequences
                       if last repetition has different outcome then
                            identify varying parameter(s) (if any)
                            introduce new variable for them
                            newloop \leftarrow empty while loop
                            newloop.conditions \leftarrow not last outcome and any common conditions
                             newloop.body \leftarrow sequence
                             change \leftarrow true
```

Table 1: Converting a plan into a template with loops

Worrisome issues:

```
while ((in_current_state (at(v?1:obj v?2:loc)) or
             in_current_state (inside(v?1:obj v?4:rocket))) and
         (in_goal_state (at(v?1:obj v?3:loc)) or
             in_goal_state (inside(v?1:obj v?4:rocket)))) do
    if (in_current_state (at(?1:obj ?2:loc)) and
             in_current_state (at(?4:rocket ?5:loc))) then
         move(?4:rocket ?5:loc ?2:loc)
    while (in_current_state (at(v?6:obj ?2:loc)) and
             in_current_state (at(?4:rocket ?2:loc)) and
             (in_goal_state (at(v?6:obj v?7:loc)) or
                  in_goal_state (inside(v?6:obj ?4:rocket)))) do
         load(?6:obj ?4:rocket ?2:loc)
    while (in_current_state (inside(v?6:obj ?4:rocket)) and
             in_current_state (at(?4:rocket ?2:loc)) and
             in_goal_state (at(v?6:obj ?2:loc))) do
         unload(?6:obj ?4:rocket ?2:loc)
while (in_current_state (inside(v?1:obj v?2:rocket)) and
         in_current_state (at(v?2:rocket v?3:loc)) and
         in_goal_state (at(v?1:obj v?4:loc))) do
    move(?2:rocket ?3:loc ?4:loc)
    while (in_current_state (inside(v?5:obj ?2:rocket)) and
             in_goal_state (at(v?5:obj ?4:loc))) do
         unload(?5:obj ?2:rocket ?4:loc)
while (in_current_state (at(v?1:rocket v?2:loc)) and
         in_goal_state (at(v?1:rocket v?3:loc))) do
    move(?1:rocket ?2:loc ?3:loc)
```

Table 2: Rocket domain template

- How to handle whiles with not-quite-matching bodies (whiles with nested ifs)? —there's a line somewhere about how aggressively we should merge things, but not sure where it is.
- If this is going to work as an watch-and-learn system WITHOUT guaranteed super-nice-and-wise teachers, we have to relax requirements on form of observed examples. How can we learn templates from inconsistent examples, examples with messy/difficult orderings, optimal examples, etc?
- Class of problems/domains attacking?
- Will we be able to say ANYTHING about efficiency?

Professed goals:

- Learn dom-spec planner from super-nice-and-wise teacher
- Learn to solve "some" problems faster than g-p planning

- Learn to solve more of "some" problems than g-p planning (horizon!)
- Use for agent modelling (develop example domain(s))
- Any impact/help in multi-agent situation(s)? (not sure how)
- Can we learn dom-spec planner from non-super-nice-and-wise teacher?

Issues to remind folks of:

- Folks have done work on revealing domain operators (mei wang)
- Folks have done work on speeding up subop plans (knoblock)

Schedule (a.k.a. To Do list):

- Rewrite Sprawl paper to include poly solution (almost done!)
- \bullet Hand-write templates for all '00 & 02 domains
- Decide what to do about difficult issue #1 (handling loops of non-identical steps (whiles with nested ifs)) &/or if it needs to be addressed right now
- Lure unsuspecting committee members (Manuela (hah! already trapped!), Reid, Avrim or Steve Smith, Craig) with appealing extended abstract
- Write up proposal
- Read a bunch more papers! (including the ones suggested by Dan Weld & by Avrim)
- Give proposal talk
- Clean up & mass-market Sprawl
 - Detach from FF
 - Implement poly solution
 - Add saving needs tree
- Implement proposed while loop stuff (hah!)
- Clean up template language
- Save template in form of c program
- Hang moon
- DEFEND!
- Parrr-tay

```
while (not (in_current_state (broken(v?1:package))) and
         in_current_state (at(v?1:package v?2:loc)) and
         in_goal_state (broken(v?1:package))) do
    if (not (in_goal_state (blownup(?2:loc))) or
             (in_goal_state (blownup(?2:loc)) and
                 in_current_state (blownup(?2:loc)))) then
         if (in_goal_state (at(?1:package ?3:loc)) and
                 in_goal_state (blownup(?3:loc)) and
                 not (in_current_state (blownup(?3:loc)))) then
             move(?1:package ?2:loc ?3:loc)
         else if (in_goal_state (blownup(?3:loc)) and
                 not (in_current_state (blownup(?3:loc)))) then
             move(?1:package ?2:loc ?3:loc)
         else if (in_goal_stantet((intg(doadwstant(e?2)noot)))bloomdup(?3:loc))))) then
             move(?1:package ?2:loc ?3:loc)
while (in_goal_state (blownup(v?1:loc)) and
         not (in_current_state (blownup(v?1:loc)))) do
    while (in_current_state (at(v?2:package ?1:loc)) and
             in_goal_state (not (broken(?2:package)))) do
         if (in_goal_state (at(?2:package ?3:loc)) and
                 (not (in_goal_state (blownup(?3:loc))) or
                      in_current_state (blownup(?3:loc)))) then
             move(?2:package ?1:loc ?3:loc)
         else if (in_current_state (blownup(?3:loc))) then
             move(?2:package ?1:loc ?3:loc)
         else if (not (in_goal_state (blownup(?3:loc))))
             move(?2:package ?1:loc ?3:loc)
         move(?2:package ?1:loc ?3:loc)
    blow(?1:loc) while (in_current_state (at(v?1:package v?2:loc)) and
         in_goal_state (broken(v?1:package)) and
         not (in_current_state (broken(v?1:package)))) then
    while (in_current_state (at(v?3:package v?2:loc)) and
             in_goal_state (not (broken(v?3:package)))) then
         move(?3:package ?2:loc ?4:loc)
    blow(?2:loc)
while (in_current_state (at(v?1:package v?2:loc)) and
         in_goal_state (at(v?1:package v?3:loc))) do
    move(?1:package ?2:loc ?3:loc)
```

Table 3: Bomb domain template

```
generate intermediate states b<br/>tw steps by propagating init state forward identify CEs that do occur
```

 $\mathbf{for} \ \mathrm{each} \ \mathrm{step} \ \mathbf{do}$

for each precondition of it do
find last provider of precondition
add link between provider and this step
if precondition added by active CE then
add conditions of active CE to that step's precs

Table 4: Poly-time algorithm for finding MAC POs $\,$

```
Input: Minimal annotated consistent partial order \mathcal{P},
          current template T_i.
Output: New template T_{i+1}, updated with \mathcal{P}
procedure DISTILL (\mathcal{P}, T_i):
      \mathcal{A} \leftarrow \text{Find\_Variable\_Assignment}(\mathcal{P}, T_i.variables, \emptyset)
      until match or can't match do
             if A = \emptyset then
                   can't match
             else
                   \mathcal{N} \leftarrow \text{Make\_New\_If\_Statement}(\text{Assign}(\mathcal{P}, \mathcal{A}))
                   \operatorname{match} \leftarrow \operatorname{Is\_A\_Match}(\mathcal{N}, T_i)
             if not can't match and not match then
                   \mathcal{A} \leftarrow \text{Find\_Variable\_Assignment}(\mathcal{P}, T_i.variables, \mathcal{A})
      if can't match then
             \mathcal{A} \leftarrow \text{Find\_Variable\_Assignment}(\mathcal{P}, T_i.variables, \emptyset)
            \mathcal{N} \leftarrow \text{Make\_New\_If\_Statement}(\text{Assign}(\mathcal{P}, \mathcal{A}))
      T_{i+1} \leftarrow \text{Add\_To\_Template}(\mathcal{N}, T_i)
procedure Make_New_If_Statement(\mathcal{P}_{\mathcal{A}}):
      N \leftarrow \text{empty if statement}
      for all terms t_m in initial state of \mathcal{P}_{\mathcal{A}} do
             if exists a step s_n in plan body of \mathcal{P}_{\mathcal{A}} such that
                          s_n needs t_m or goal state of \mathcal{P}_{\mathcal{A}} needs t_m then
                   Add_To_Conditions(N, in_current_state (t_m))
      for all terms t_m in goal state of \mathcal{P}_{\mathcal{A}} do
             if exists a step s_n in plan body of \mathcal{P}_{\mathcal{A}} such that
                          t_m relies on s_n then
                   Add_To_Conditions(N, in_goal_state (t_m))
      for all steps s_n in plan body of \mathcal{P}_{\mathcal{A}} do
             Add_To_Body(N, s_n)
      \mathbf{return}\ N
procedure Is_A_Match(\mathcal{N}, T_i):
      for all if-statements I_n in T_i do
             if \mathcal{N} matches of I_n then
                   return true
procedure Add_To_Template(\mathcal{N}, T_i):
      for all if-statements I_n in T_i do
             if \mathcal{N} matches I_n then
                   I_n \leftarrow \text{Combine}(I_n, \mathcal{N})
                   return
      if \mathcal{N} is unmatched then
             Add_To_End(\mathcal{N}, T_i)
```

Table 5: The current DISTILL algorithm: updating a template with a new observed plan.