



Towards Practical Runtime Type Instantiation

Dispatch on Generics



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Scientific Programming

```
forecast() = do
```

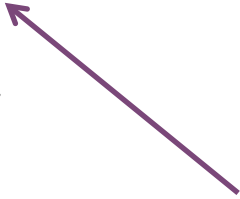
```
... (*) number crunching
```

```
temp = bigMatrix * reallyBigMatrix
```

```
... (*) more number crunching
```

```
end
```

When can this multiplication be optimized?



+ Algorithm choice

`temp = bigMatrix * reallyBigMatrix`

- Depends on the type of the entries

`bigMatrix : Matrix[[Number]]`

`reallyBigMatrix : Matrix[[Number]]`

- General numbers – use standard algorithm
 - Many individual multiplications

`bigMatrix : Matrix[[ $\mathbb{Z}_2$ ]]`

`reallyBigMatrix : Matrix[[ $\mathbb{Z}_2$ ]]`

- Binary Digits – create bit vectors and use bit operators for multiplication
 - Many multiplications at the same time
 - More efficient multiplication



How to Choose Best Algorithm?

```
temp = bigMatrix * reallyBigMatrix
```

- If know statically operands are binary matrixes
 - Can define a separate function for the programmer to call


```
binaryMatrixMult(x : Matrix[ℤ2], y : Matrix[ℤ2]) = ...
```

```
temp = binaryMatrixMult(bigMatrix, reallyBigMatrix)
```
- Otherwise, could use multiple dispatch
 - Define two cases of * operator


```
*(x : Matrix[Number], y : Matrix[Number]) = do
  ... (*) standard algorithm
end
```

```
*(x : Matrix[ℤ2], y : Matrix[ℤ2]) = do
  ... (*) bit vector algorithm
end
```
 - Binary algorithm chosen dynamically if operands are binary matrixes at runtime
 - **Runtime type** must include **generic** information

+ Generic Dispatch

- Instantiated type of objects matters at runtime
 - Used to choose what code is executed – **No type erasure!**
- Instantiated type of generic functions matters
 - Instantiated generic used to create objects – must be tagged at runtime

```
append[X <: Any](l : List[X], n : X) : List[X] = do
  ... (*) duplicate l, add n to duplicate, return duplicate
end
```

- Type of resulting List depends on the instantiation of X
- Want it to be as specific as possible

```
removeNegatives(l : List[Number]) : List[Number] = ... (*) iterate
removeNegatives(l : List[N]) : List[N] = ... (*) no op
```

- Don't always know the best answer statically
 - **Need runtime type instantiation!**

+ Problem – Runtime Instantiation

- Given
 - a generic function signature with type variables X_i
 - the runtime types of the inputs
 - Static return type of the function call
- Find an instantiation of the X_i such that
 - Type variable bounds met (Valid instantiation)
 - Runtime inputs are subtypes of parameter types
(Function applicable to arguments)
 - Return type is a subtype of static return type (Type safe)
 - Minimize return type (As specific as possible)



heuristic

+Fortress Overview

- Symmetric Multiple Dispatch
 - Statically guaranteed a most specific function definition exists
 - Chosen at runtime, static information for optimization of search
- Objects
 - Nominal (declared) subtyping
 - Variant Generics
 - Covariant: If Baz covariant, then $\text{Baz}[[X]] <: \text{Baz}[[Y]]$ if and only if $X <: Y$
 - Invariant: If Baz invariant, then $\text{Baz}[[X]] <: \text{Baz}[[Y]]$ if and only if $X = Y$
 - Types form a lattice
 - Meet (greatest lower bound) and join (least upper bound) defined
- Generic Functions $f[\overline{[X <: \tau_X]}](\overline{[y : \tau_y]}) : \tau_r$
 - Type variables X bound within
 - Upper bounds of type variables τ_X
 - Parameter types τ_y
 - Return type τ_r

+ Problem Example

- List Definition `List[[covariant X <: Any]]`
- Number hierarchy `$\mathbb{Z}_2$  <:  $\mathbb{N}$  <: Number <: Any`
- Function
`append[[X <: Any]](l : List[[X]], n : X) : List[[X]] = do`
 ... () duplicate l, add n to duplicate, return duplicate*
`end`
- Code
`theList : List[[Number]]`
`toAdd : Number`
`...`
`temp = append(theList, toAdd)` Static return type: `List[[Number]]`
`removeNegatives(temp)`



Instantiation Calculation for X

```
append[X <: Any](l : List[X], n : X) : List[X] = do
  ... (*) duplicate l, add n to duplicate, return duplicate
end
```

- Given - $l : \text{List}[\mathbb{N}]$, $n : \mathbb{Z}_2$, $T_r = \text{List}[\text{Number}]$
- Parameters:
 - $\text{List}[\mathbb{N}] <: \text{List}[X]$ implies $\mathbb{N} <: X$ (covariance)
 - $\mathbb{Z}_2 <: X$
- Type Variable Bounds
 - $X <: \text{Any}$
- Return Type Bounds
 - $\text{List}[X] <: \text{List}[\text{Number}]$ implies $X <: \text{Number}$ (covariance)
- Upper Bound
 - $\text{meet}(\text{Any}, \text{Number}) = \text{Number}$
- Lower Bound
 - $\text{join}(\mathbb{N}, \mathbb{Z}_2) = \mathbb{N}$
- Best (most specific) choice = $\mathbb{N}$

+ General Algorithm

Given $f[\overline{X} <: \tau_X](\overline{y} : \tau_y) : \tau_r$, runtime input types T_y ,
and static return type T_r

1. Generate upper and lower bounds for each X from
 - Parameter constraints
 - Return type constraints
 - Type variable bound constraints
2. Let $X_L = \text{join}(\text{lowerBounds}(X))$ and $X_U = \text{meet}(\text{upperBounds}(X))$
3. Check $X_L <: X_U$
4. If true, return X_L as the instantiation

- Single pass algorithm
 - If type variables scoped left to right

$\text{badlyScoped}[X <: Y, Y <: \text{List}[X]](x : X, y : Y) : \text{Any}$

- Otherwise could result in infinite iteration on bound generation

+ Remaining Challenges

1. Multiple potentially applicable function definitions
 - Run this algorithm on each to find most specific applicable
2. Contravariance
 - Function types in Fortress
 - May be undecidable
3. Efficient computation of meets, joins, and comparisons in subtype lattice
4. Enough performance gains to offset overhead?

```
result = forecast() do
```

```
... (*) number crunching
```

```
temp = bigMatrix * reallyBigMatrix
```

```
... (*) more number crunching
```

```
end
```