A Language for Automatically Enforcing Privacy

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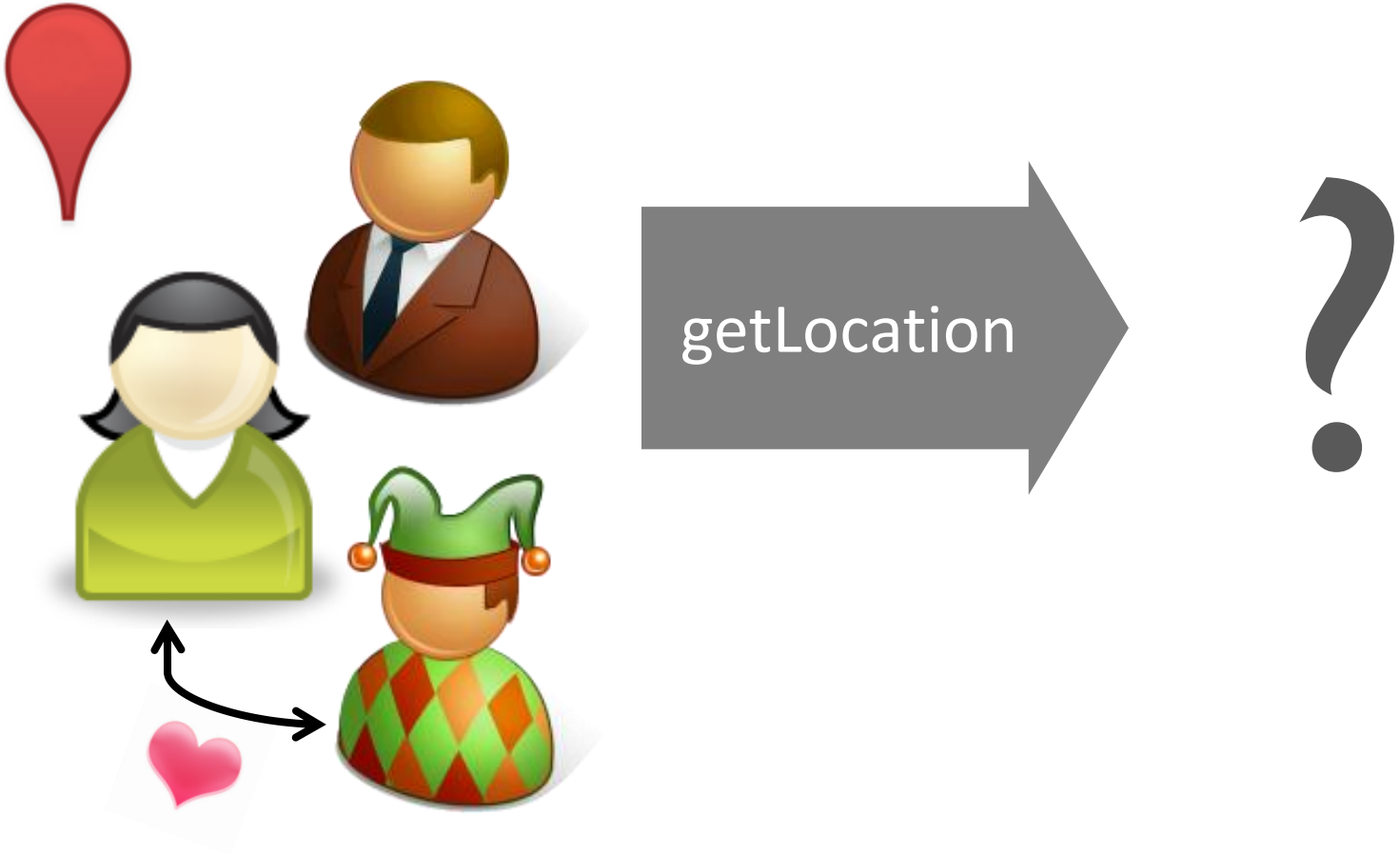
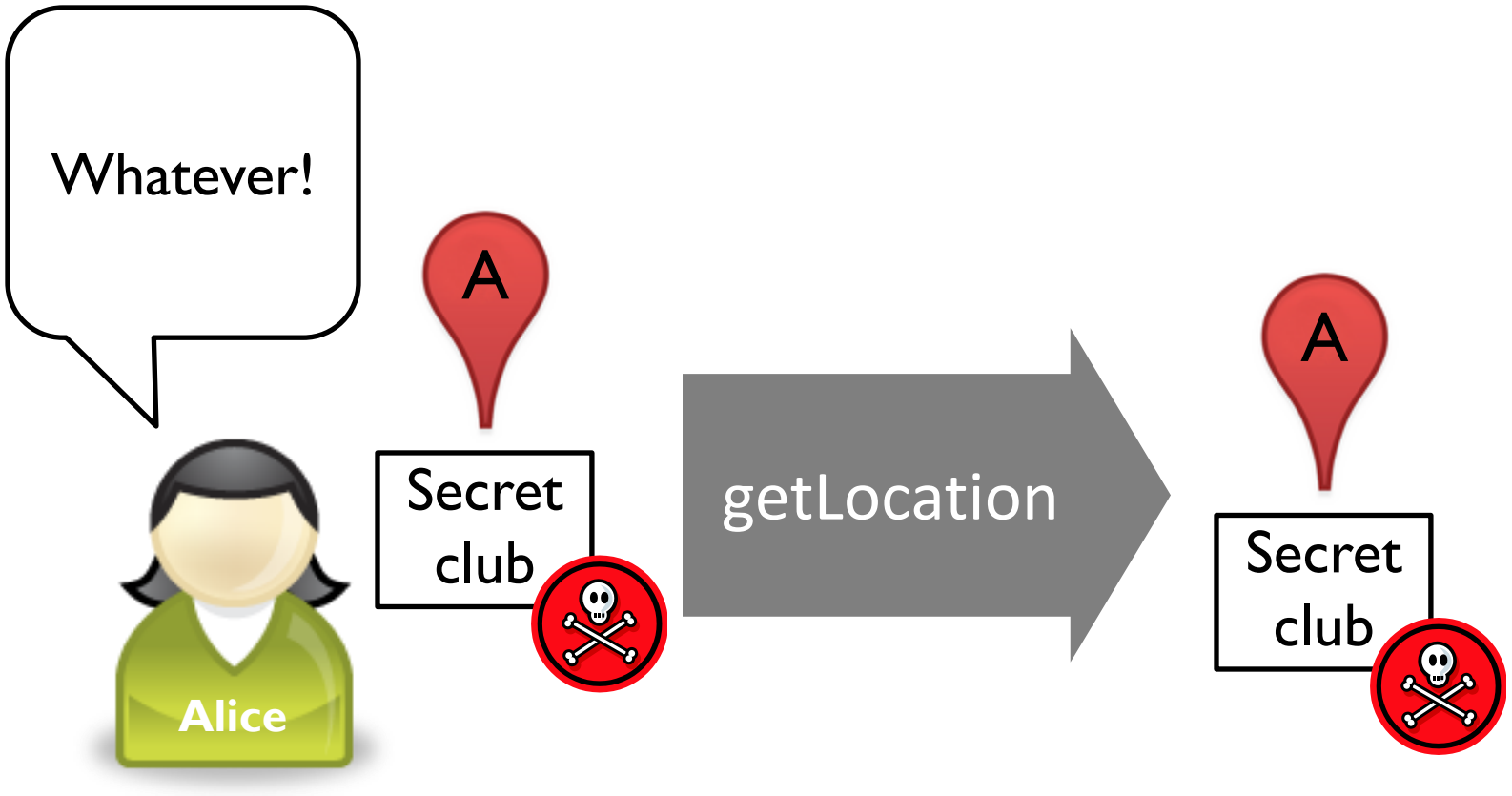
and Armando Solar-Lezama

Displaying User Locations to Other Users

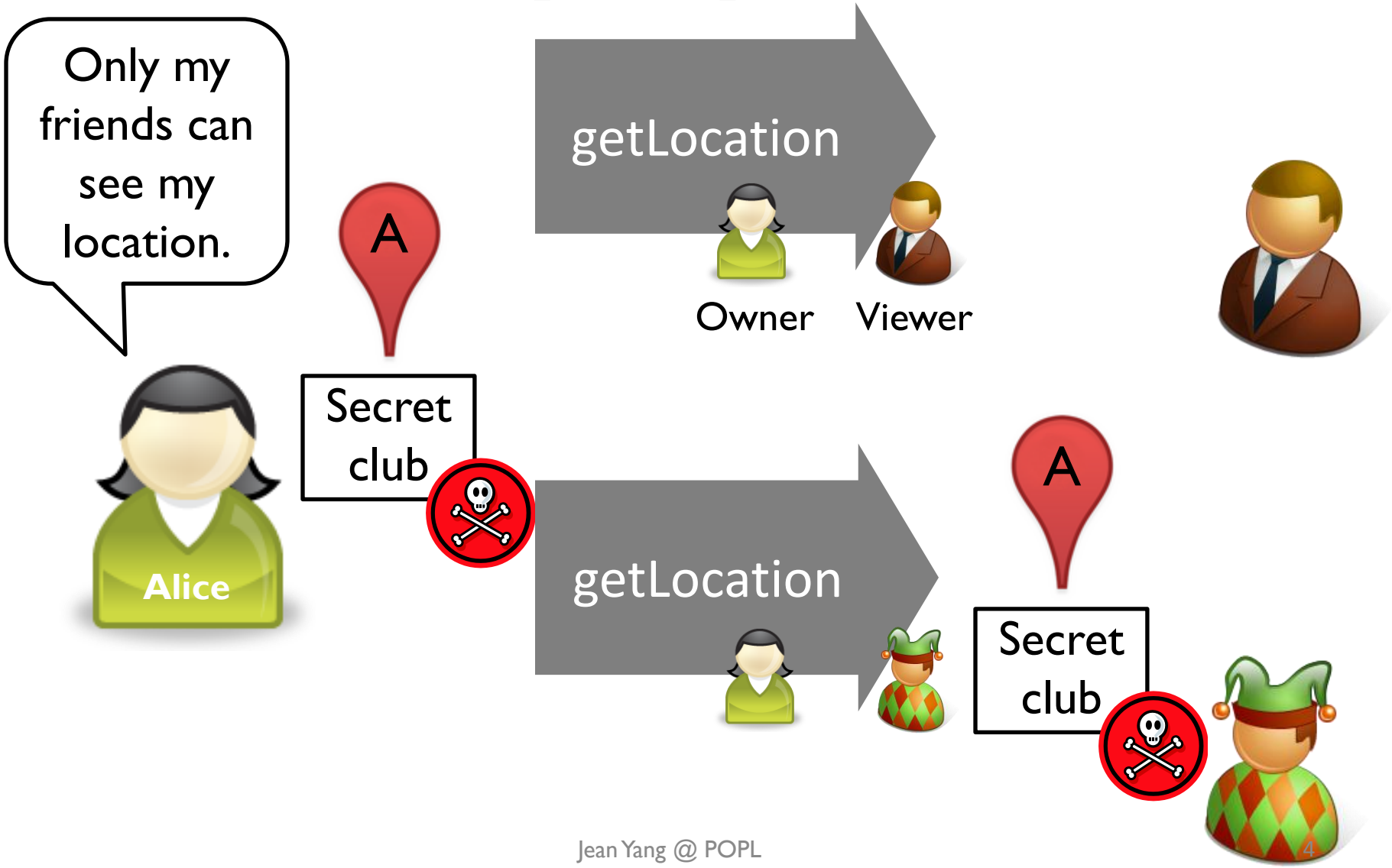


No Privacy Concerns

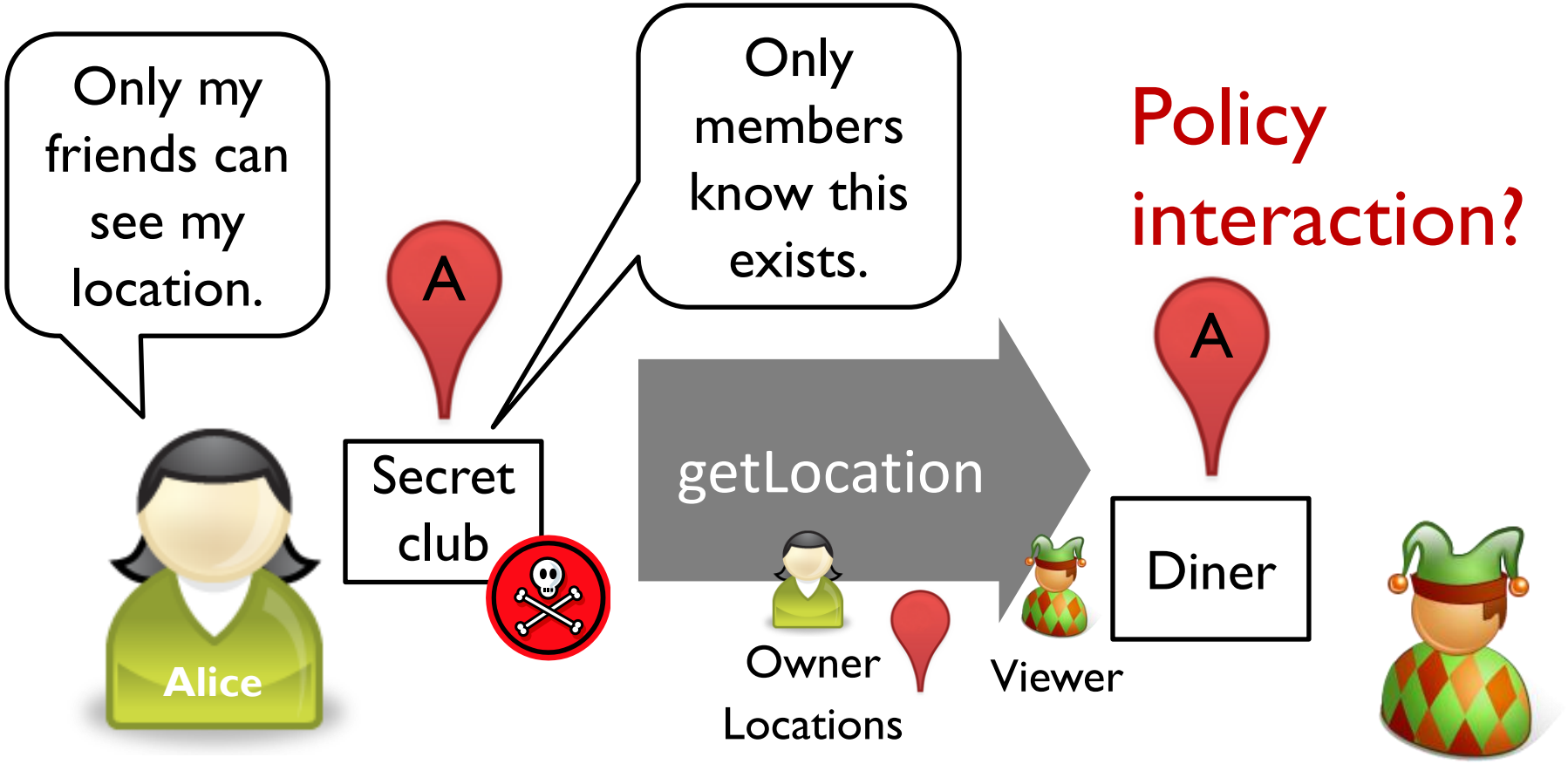



```
def getLocation (user: User): Location = user.location
```

Simple policy



Finer-Grained Policies



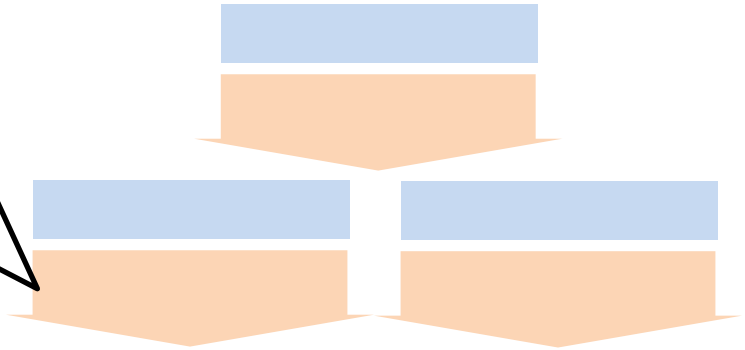
Programmer Burden

```
def getLocation (user: User) (viewer: User)
   : Location = {
   if (isFriends user viewer) {
        if (canSee user.location viewer) {
          user.location;
        } else { scrub(user.location, "Diner"); }
   } else { undisclosedLocation; }
}
```

Output Context

Policies

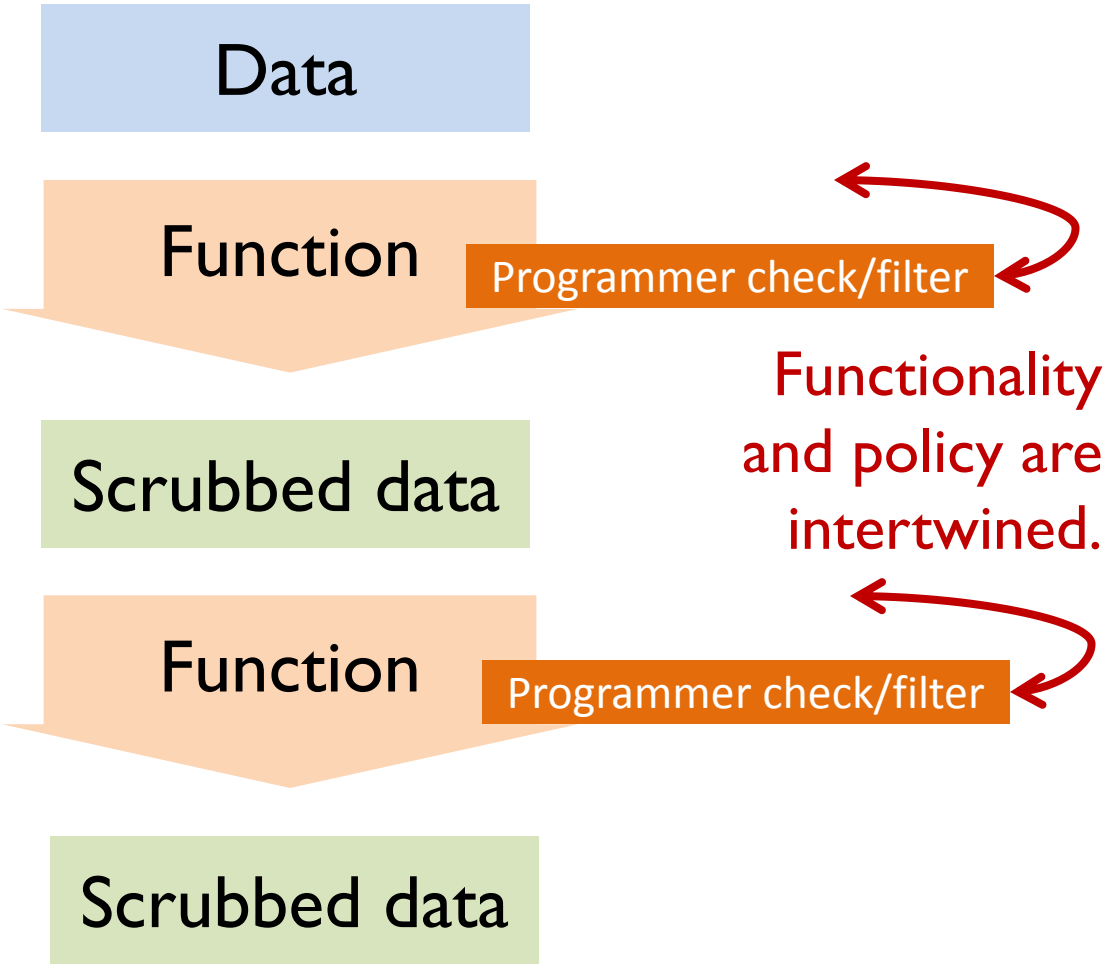
Views of sensitive values



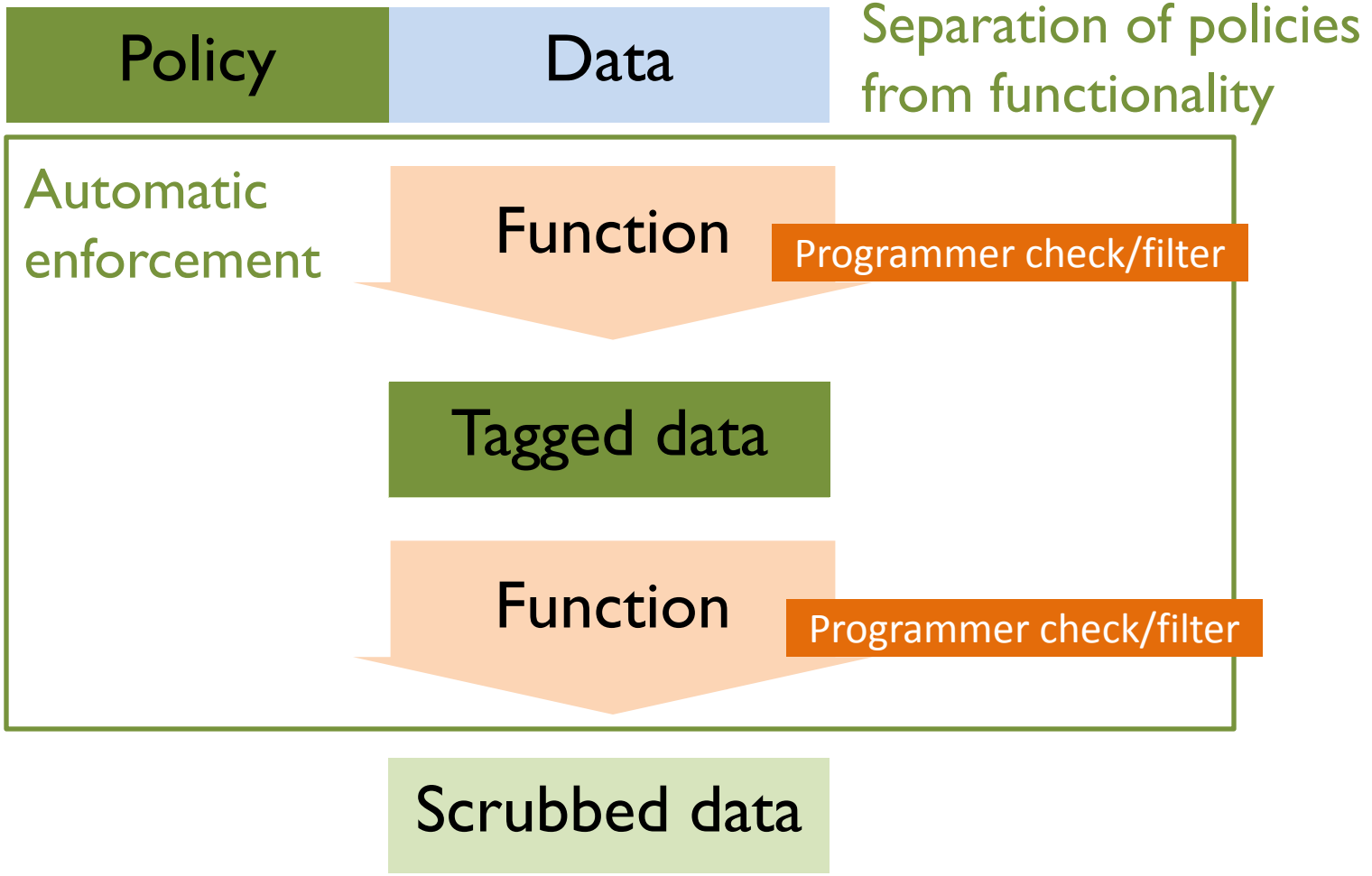
Our Mission

Make it easier for the **programmer** to preserve **confidentiality** of user data.

What's Hard?


Data
Function
Programmer check/filter
Functionality and policy are intertwined.
Scrubbed data
Function
Programmer check/filter
Scrubbed data

Our Solution



Jeeves Goal

State of the Art

```
def getLocation (user: User) (viewer: User)
  : Location = {
  if (isFriends user viewer) {
      if (canSee user.location viewer) {
        user.location;
      } else obfuscate(user.location, "Work"); }
  } else { undisclosedLocation; }
}
```

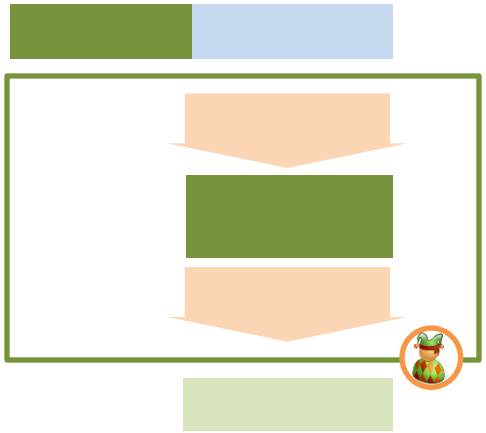
Policy + Functionality

Jeeves

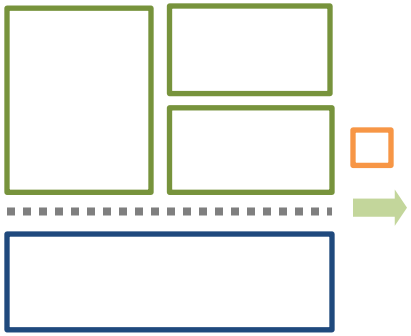
```
def getLocation (user: User): Location =
  user.location
```

Talk Outline

Jeeves language

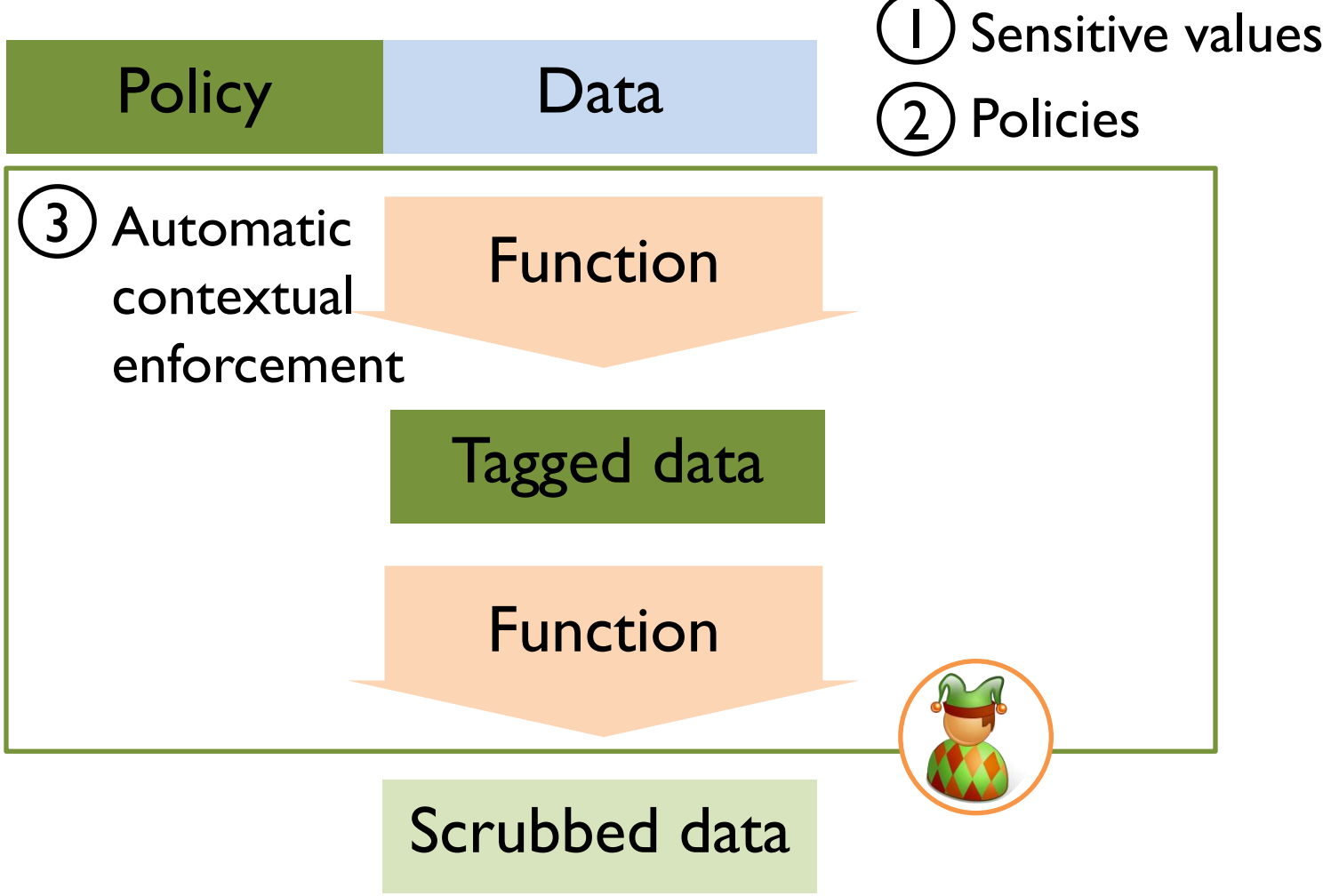


How it works

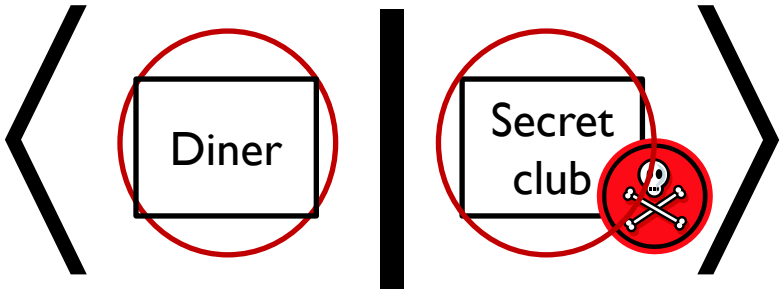


Coding in Jeeves

Jeeves Language



Jeeves for Locations



Low confidentiality High confidentiality

Diner

Secret club

Using Jeeves

Sensitive Values

level a in { **low, high** }

val location: String = <"school" | "MIT"> *a*

Policies

policy a: **context** != alice ➔ **low**

Core Functionality

val msg: String = "Alice is at " + location

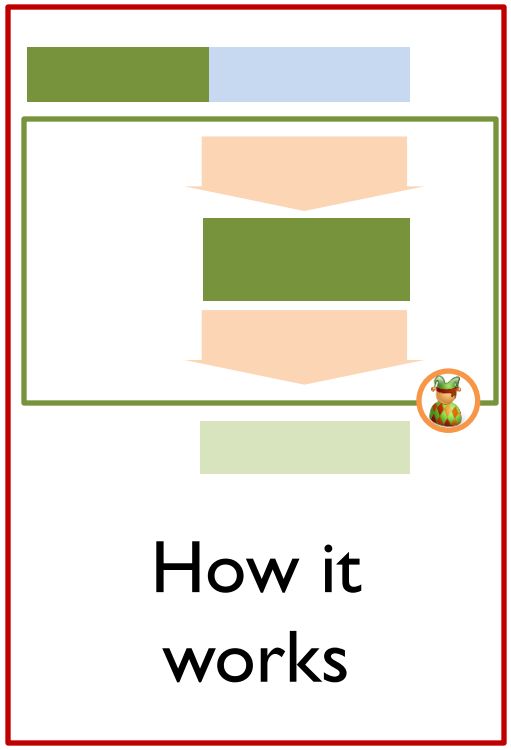
Contextual Enforcement

print {alice} msg /* *"Alice is at MIT"* */

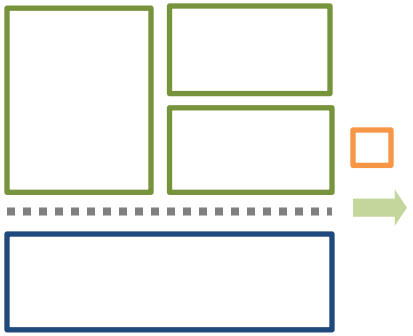
print {bob} msg /* *"Alice is at school"* */

Talk Outline

Jeeves language

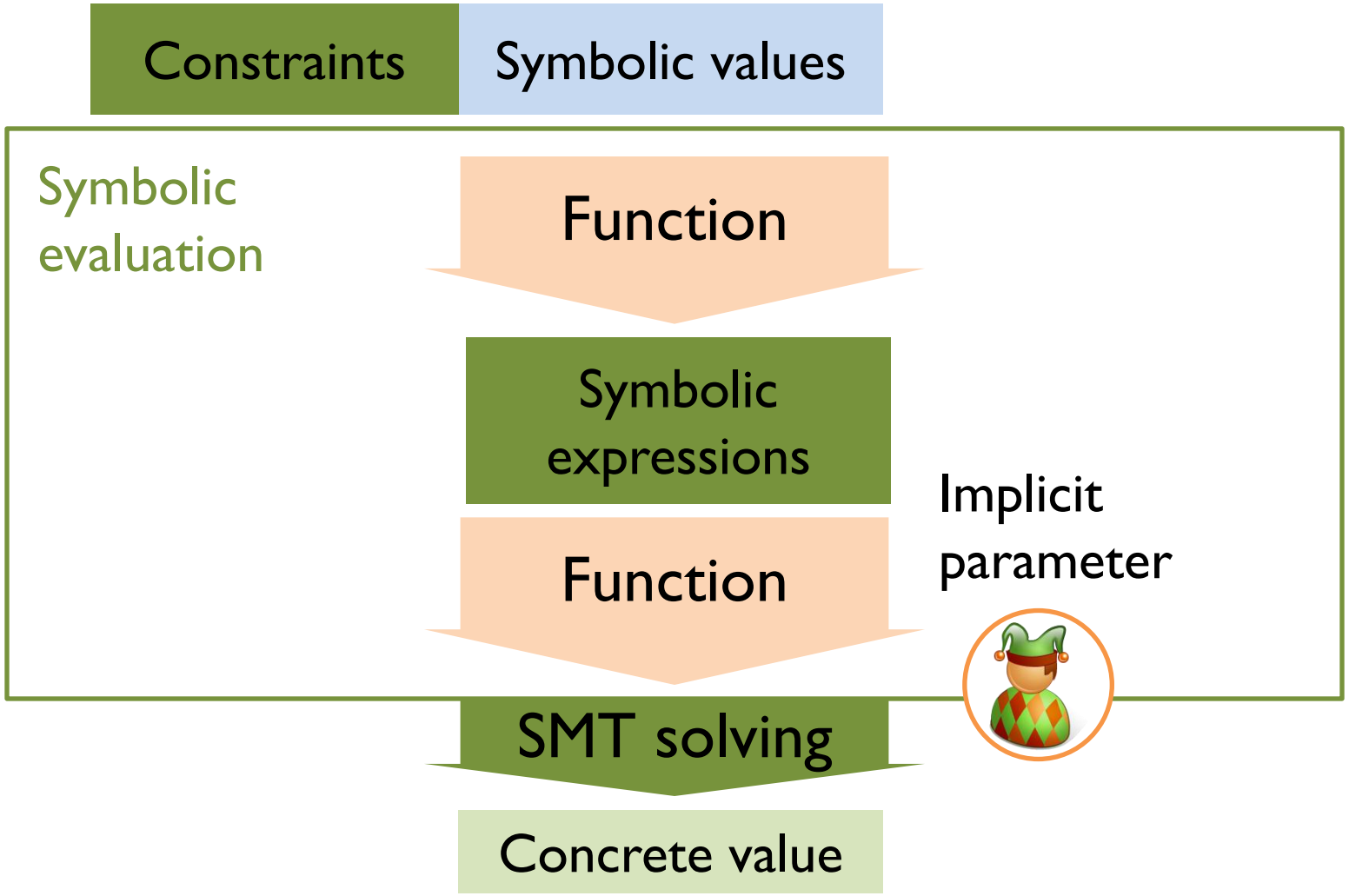


How it works



Coding in Jeeves

How Jeeves Works



Representing Sensitive Values in Jeeves

Without Jeeves

Name	Location
Alice	MIT
Bob	POPL
Claire	POPL

Jeeves

Name	Location	
Alice	$\langle ? \mid \text{MIT} \rangle_a$	Policy
Bob	POPL	
Claire	$\langle ? \mid \text{POPL} \rangle_b$	Policy

Symbolic Evaluation for Information Flow

Name	Location
Alice	$\langle\ \vert\ \rangle_a$
Bob	POPL
Claire	$\langle\ \vert\ \rangle_b$

How many people are at POPL?

$1 + ((x_1 = \text{POPL})\ ?\ 1 : 0) + ((x_2 = \text{POPL})\ ?\ 1 : 0)$

Runtime Environment

context != alice → a = low

… → b = low

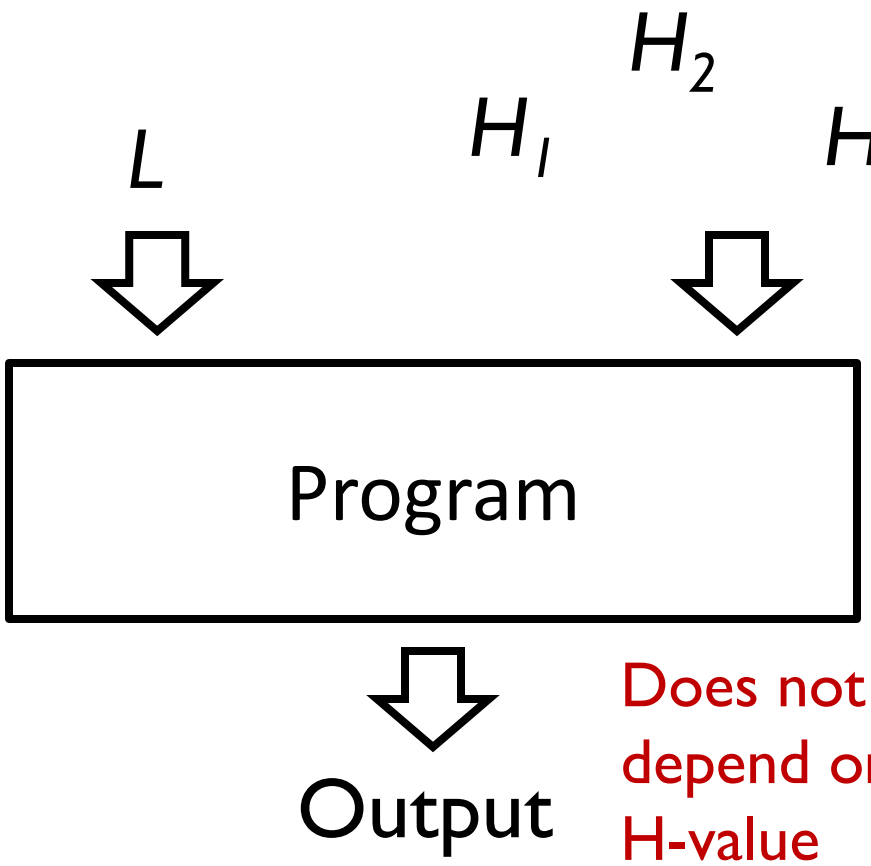
Outputs computed from sensitive values are **symbolic** & concretized under the policy environment.

Jeeves Non-Interference Guarantee

Consider the **sensitive value**

Given a fixed *L*, all executions where *a* must be **low** produce equivalent outputs no matter the value of *H*.

Standard Non-Interference

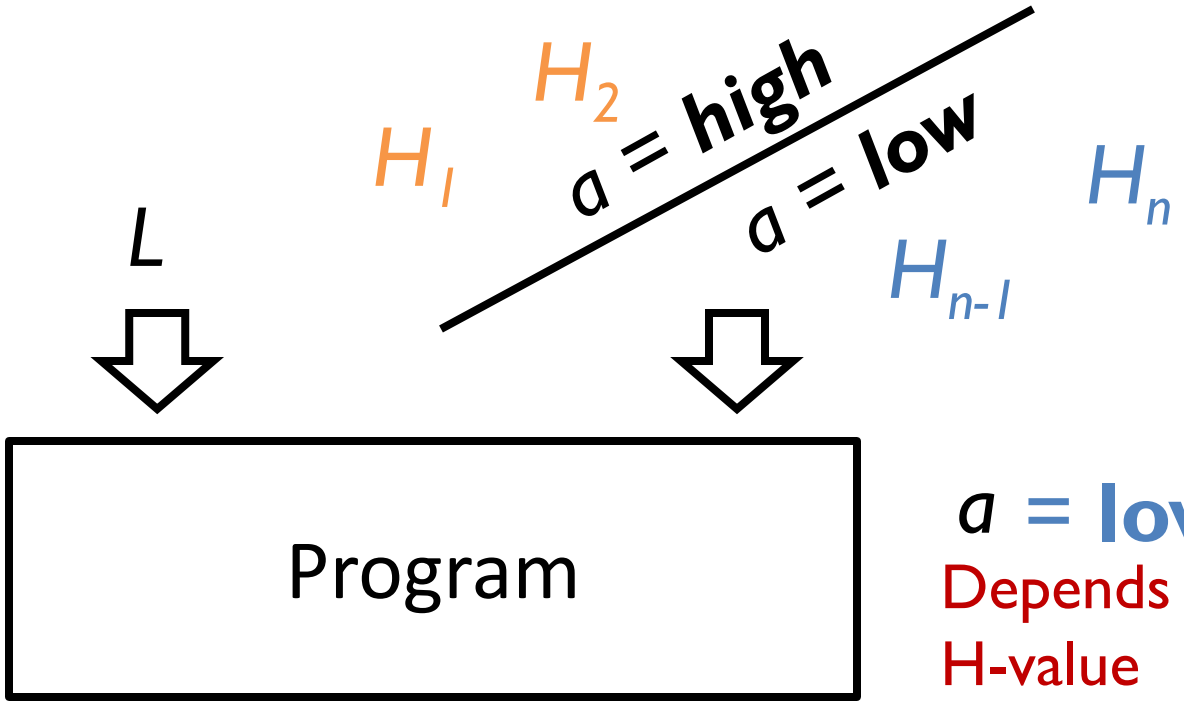


a = **low**

Does not depend on the H-value

Does not depend on the H-value

Jeeves Non-Interference

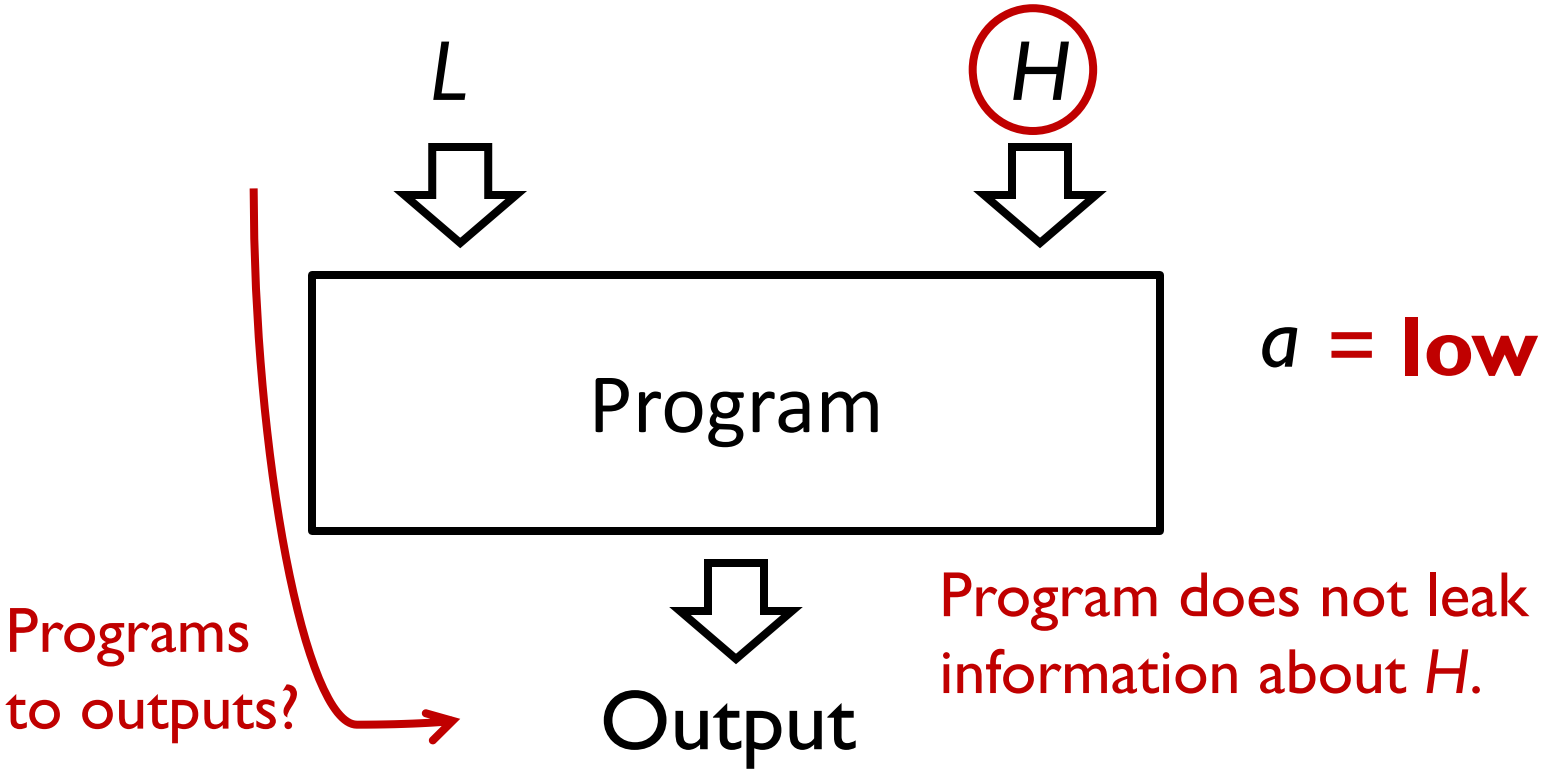


a = **low**

Depends on the H-value

Cannot distinguish between H-values that imply *a* = **low**

Jeeves Non-Interference



Language Restrictions

Constraints | Symbolic values

Primitives and objects.
No functions.

Arithmetic and Boolean constraints with conditionals & implications.

No functions, quantifiers, or theory of lists.

Function

Symbolic expressions

Function

SMT solving

Concrete value

Static Checks

Constraints | Symbolic values

Symbolic evaluation

Function

Symbolic expressions

Function

SMT solving

Concrete value

Symbolic values flow only where expected.

Evaluation does not introduce nontermination.

Contexts are well-formed.

Outputs are concrete.

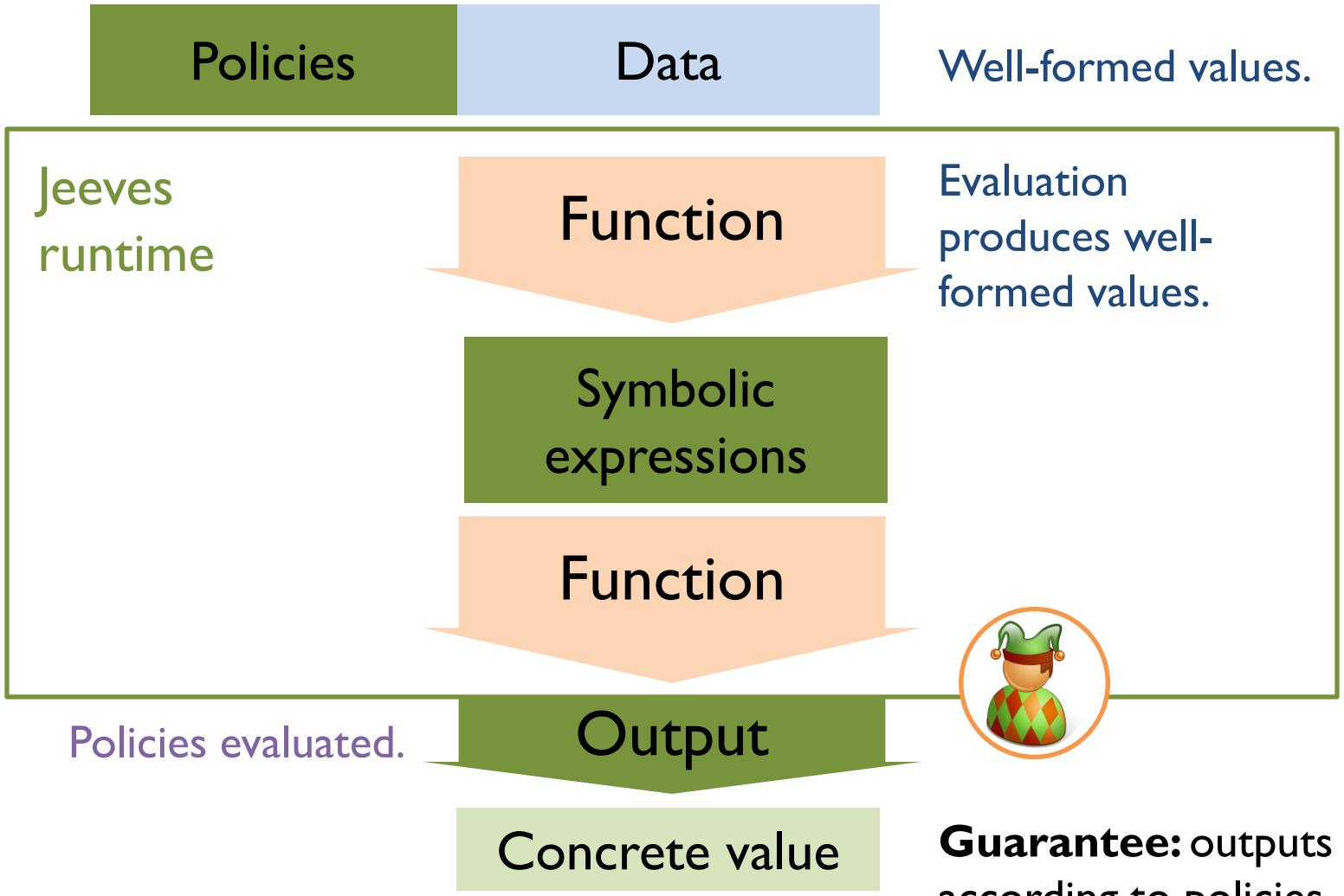
Stateful Policies


```
policy a:
  (distance context alice
    ≥ radius ) → low
```

But Alice's location is changing…

Jeeves: Delay policy evaluation until **output**.

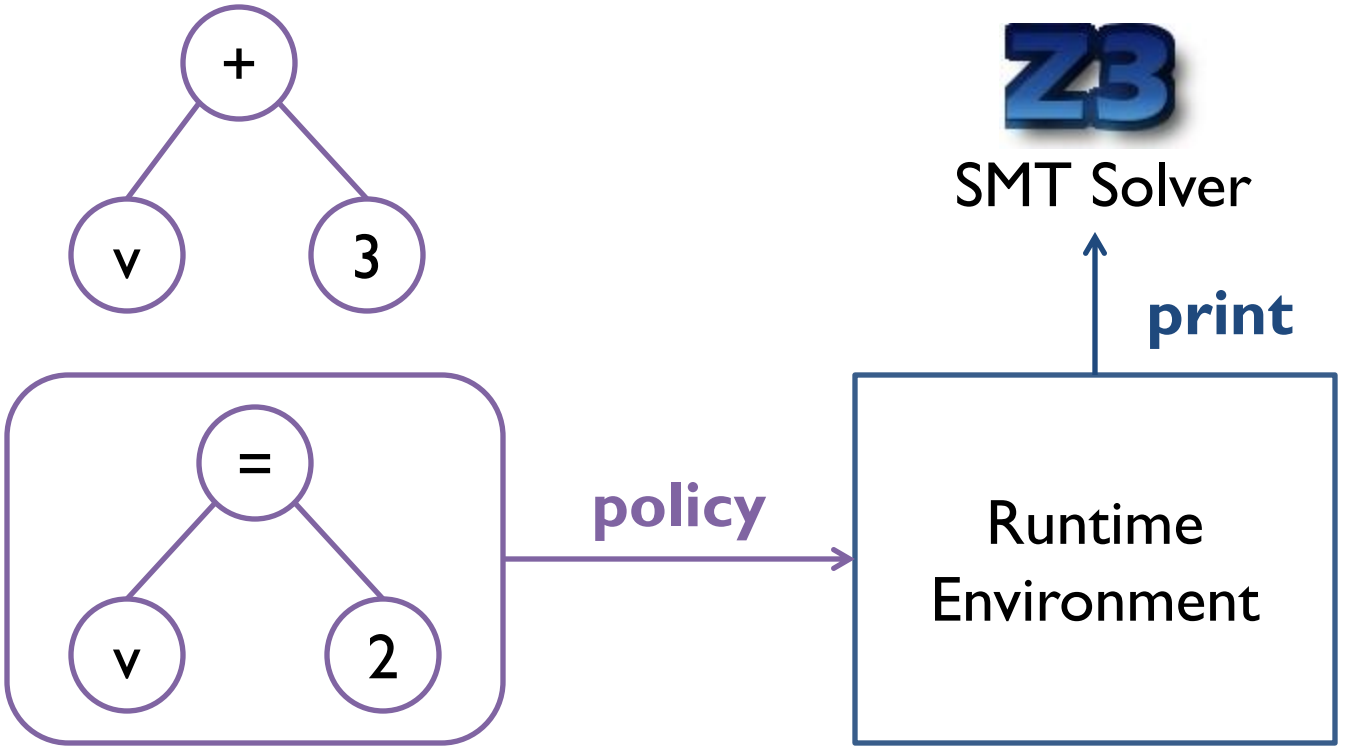
Jeeves System



Guarantee: outputs shown according to policies.

Scala Implementation

Overload operators to create symbolic expressions.



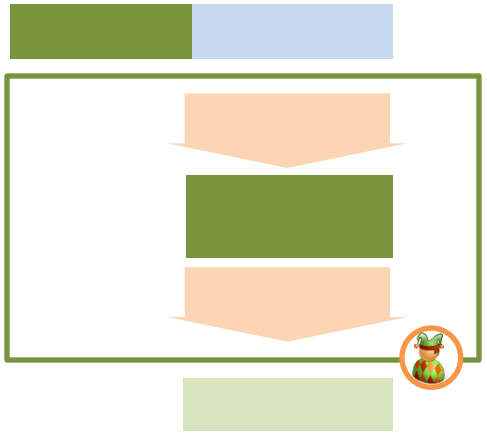
Use an SMT solver as a model finder.

Delay evaluation of policies until output.

Propagate policies.

Talk Outline

Jeeves language



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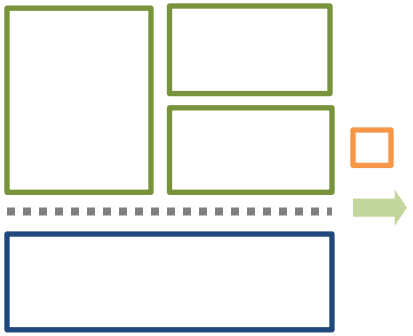
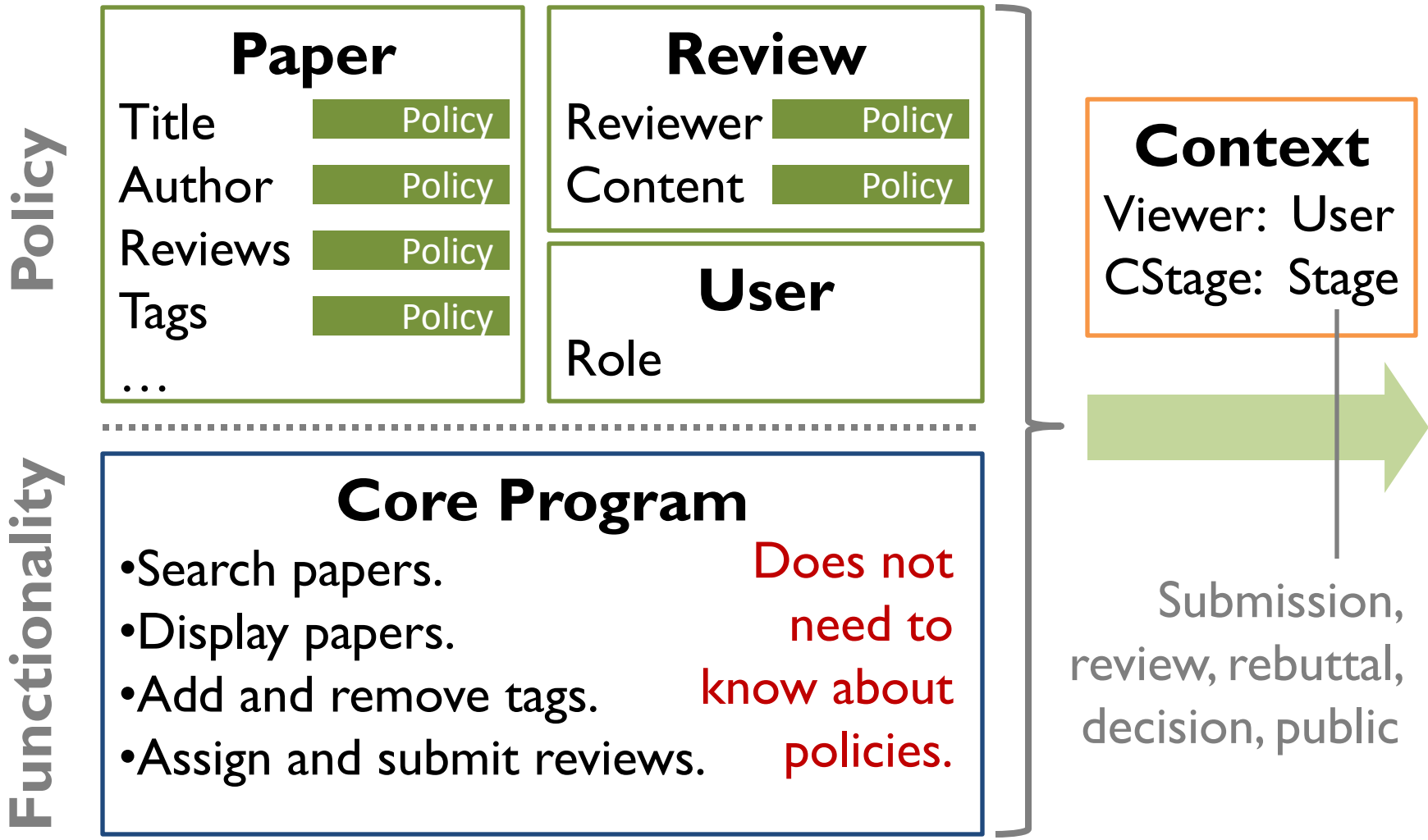


Coding in Jeeves

JConf Architecture

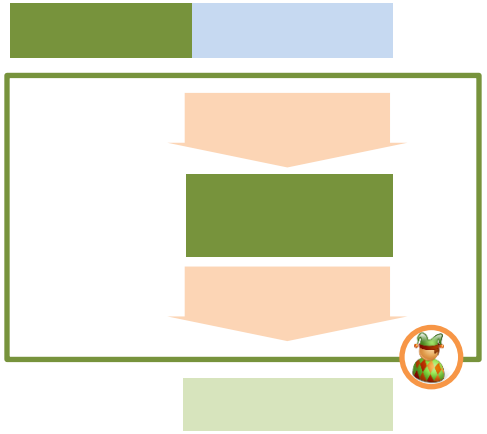


Functionality vs. Policy

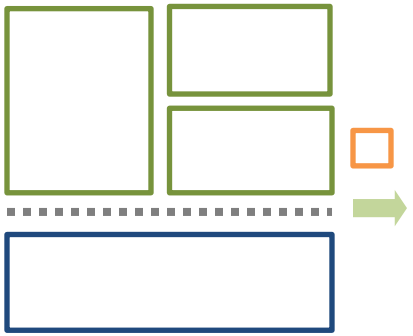
File	Total LOC	Policy LOC
ConfUser.scala	59	17
PaperRecord.scala	103	48
PaperReview.scala	21	11
ConfContext.scala	6	0
Backend	123	0
Frontend (Scalatra)	161	0
Total	**473**	**76**

Conclusions

The Jeeves language: pushing responsibility of privacy to the runtime.



How we designed a language with constraints using symbolic evaluation to provide execution guarantees.



Evaluation of Jeeves in practice: conference management example.

Website:	**sites.google.com/site/jeevesprogramming**
Google Code:	**code.google.com/p/scalasmt**
Contact:	**jeanyang@mit.edu**