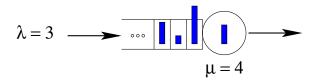
M/G/1 Model with 4 Categories of Scheduling Policies



	Non-Preemptive	Preemptive
Non-Size-Based		
Size/Class-Based		

Table 1: 4 quadrants

Question: Assume that S is continuous. Define these important quantities:

- (a) $\mathbf{E}[S]$
- (b) $\mathbf{E}\left[S^2\right]$
- (c) ρ

Non-Preemptive, Non-Size-Based scheduling policies
Question: Name and define some non-preemptive, non-size-based policies
Question: What did HW 5 teach us on how these compare wrt $\mathbf{E}[T_Q]$? $\mathbf{E}[T]$?
Question: What is some intuition behind this result?

Question: What's the problem with non-preemptive, non-size-based policies?

Non-preemptive Priority Scheduling

Question: When does non-preemptive priority come up?

MODEL:

- \bullet *n* priority classes.
- \bullet Class 1 has _____ priority. Class n has _____ priority. Within a class serve jobs FCFS.
- Arrival rates: $\lambda_1, \lambda_2, \dots, \lambda_n$
- Job size distributions: S_1, S_2, \dots, S_n

Question: What is ρ_i , the load made up of jobs of class i?

Question: How do we denote the total load made up by classes 1 through k?

Non-preemptive Priority Scheduling, cont.

Let $T_Q(k)^{\text{NP-Prio}}$ denote the queueing time for a job of class k. Let $T(k)^{\text{NP-Prio}}$ denote the response time for a job of class k.

$$\mathbf{E}\left[T_{Q}(k)\right]^{\text{\tiny NP-Prio}} = \frac{\rho}{(1 - \sum_{i=1}^{k} \rho_{i})(1 - \sum_{i=1}^{k-1} \rho_{i})} \cdot \frac{\mathbf{E}\left[S^{2}\right]}{2\mathbf{E}\left[S\right]}$$

Question: What is $\mathbf{E}[T_Q(1)]$? How does this compare with $\mathbf{E}[T_Q]^{\text{\tiny FCFS}}$?

Question: What is $\mathbf{E}[T_Q(n)]$? How does this compare with $\mathbf{E}[T_Q]^{\text{\tiny FCFS}}$?

Question: For smaller k, how does $\mathbf{E}[T_Q(k)]^{\text{NP-Prio}}$ compare with $\mathbf{E}[T_Q]^{\text{FCFS}}$?

Question: Why does $\mathbf{E}\left[T_Q(k)\right]$ still see full $\mathbf{E}\left[S^2\right]$?

Non-preemptive Shortest-Job-First (SJF)

Question: Define SJF

Question: Let ρ_x denote the load made up of jobs of size $\leq x$. What is ρ_x ?

<u>Defn</u>: $T_Q(x)^{\text{\tiny SJF}}$ is the queueing time for a job of size x

$$\mathbf{E}\left[T_{Q}(x)\right]^{\text{\tiny SJF}} = \frac{\rho}{(1-\rho_{x})^{2}} \cdot \frac{\mathbf{E}\left[S^{2}\right]}{2\mathbf{E}\left[S\right]}$$

Non-preemptive Shortest-Job-First (SJF), cont.

$$\mathbf{E}\left[T_{Q}(x)\right]^{\text{\tiny SJF}} = \frac{\rho}{(1-\rho_{x})^{2}} \cdot \frac{\mathbf{E}\left[S^{2}\right]}{2\mathbf{E}\left[S\right]}$$

Question: How does $\mathbf{E}[T_Q(x)]^{\text{SJF}}$ compare with $\mathbf{E}[T_Q(x)]^{\text{FCFS}}$ for small x? How about large x?

Question: What is $\mathbf{E}[T_Q]^{\text{SJF}}$? How do you expect it compares to $\mathbf{E}[T_Q]^{\text{FCFS}}$?

Question: Why is $\mathbf{E}[S^2]$ still there?

Preemptive policies

Question: Why do preemptive policies have lower response times than non-preemptive policies? Explain for PS and P-LCFS.

Question: In HW 5 we investigated preemptive policies, PS and P-LCFS. How did these compare? What was their response time?

$$\mathbf{E}\left[T\right]^{\scriptscriptstyle{\mathrm{PS}}} =$$

$$\mathbf{E}\left[T\right]^{ ext{\tiny P-LCFS}} =$$

Question: Let's look under the hood ...

Let T(x) denote the response time for a job of size x.

$$\mathbf{E}\left[T(x)\right]^{\scriptscriptstyle{\mathrm{PS}}} \ = \$$

$$\mathbf{E}\left[T(x)\right]^{\text{\tiny P-LCFS}} =$$

Preemptive, Size-Based Scheduling

r reemptive, bize-based beneduming
Fill in Picture: Response time as a function of load
Question: Can we possibly improve upon PS? What is there left to do?
Three preemptive size/class-based policies:
1. Preemptive Priorities
2. PSJF
3. SRPT

Preemptive Priorities

MODEL:

- \bullet *n* priority classes.
- ullet Class 1 has _____ priority. Class n has _____ priority.
- Arrival rates: $\lambda_1, \lambda_2, \dots, \lambda_n$
- Job size distributions: S_1, S_2, \ldots, S_n
- \bullet $\rho_i = \underline{\hspace{1cm}}$.
- Preempt-resume.

ANALYSIS:

Waiting Time: Time until job first starts serving

Residence Time: Time from when job first receives service until it departs

$$\mathbf{E} [T(k)]^{\text{\tiny P-Prio}} = \frac{\sum_{i=1}^{k} \rho_{i} \frac{\mathbf{E}[S_{i}^{2}]}{2\mathbf{E}[S_{i}]}}{(1 - \sum_{i=1}^{k-1} \rho_{i})(1 - \sum_{i=1}^{k} \rho_{i})} + \frac{\mathbf{E} [S_{k}]}{1 - \sum_{i=1}^{k-1} \rho_{i}}$$

Question: How do waiting time and residence time compare under P-Prio and NP-Prio?

Preemptive Shortest Job First – PSJF

Question: Define PSJF. How does this differ from SJF?

Question: What is ρ_x ? Why is it relevant? Who is visible to a job of size x?

ANALYSIS:

Waiting Time: Time until job first starts serving

Residence Time: Time from when job first receives service until it departs

$$\mathbf{E} [T(x)]^{\text{\tiny PSJF}} = \frac{\lambda \int_0^x t^2 f_S(t) dt}{2(1 - \rho_x)^2} + \frac{x}{1 - \rho_x}$$

Question: What is $\mathbf{E}[T]^{PSJF}$?

Question: Why should we expect that PSJF is superior to SJF?

SRPT

Question: Define SRPT. How does this differ from PSJF? How do we program SRPT?

ANALYSIS:

Waiting Time: Time until job first starts serving

Residence Time: Time from when job first receives service until it departs

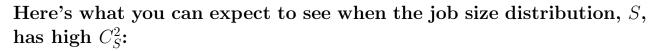
$$\mathbf{E} [T(x)]^{\text{\tiny SRPT}} = \frac{\lambda \int_0^x t^2 f_S(t) dt + \lambda x^2 (\overline{F_S}(x))}{2(1 - \rho_x)^2} + \int_{t=0}^x \frac{dt}{1 - \rho_t}$$

Question: Compare the formulas for $\mathbf{E}[T]^{PSJF}$ and $\mathbf{E}[T]^{SRPT}$.

$$\mathbf{E} [T(x)]^{\text{PSJF}} = \frac{\lambda \int_0^x t^2 f_S(t) dt}{2(1 - \rho_x)^2} + \frac{x}{1 - \rho_x}$$

Theorem [Schrage 1966]: SRPT produces optimal mean response time for any arrival sequence (any arrival times, job sizes).

Comparison of policies



FILL IN GRAPH

Note: This shows $\mathbf{E}[T]$, but actually all moments of T are known (see textbook).

EXTRA: Scheduling and starvation

Question: What is starvation? Does it really exist?

Definition: $Slowdown(x) = \frac{T(x)}{x}$

Question: What is $\mathbf{E}[Slowdown(x)]^{PS}$?

Question: For policies which favor short jobs, what do you think $\mathbf{E}[Slowdown(x)]$ looks like as a function of x?

The All-Can-Win-Theorem: If $\rho \leq \frac{1}{2}$, then for all job sizes x,

$$\mathbf{E}\left[Slowdown(x)\right]^{\text{\tiny SRPT}} \leq \mathbf{E}\left[Slowdown(x)\right]^{\text{\tiny PS}} \ .$$

(Bansal and Harchol-Balter "Analysis of SRPT scheduling: Investigating unfairness" ACM SIGMETRICS 2001.)

EXTRA: Opt scheduling when we don't know job size

1. The FB policy – Optimal when job size distribution has DFR.

- 2. Shortest-Expected-Remaining-Processing-Time (SERPT) Simple and almost always near optimal
 - Scully, Harchol-Balter, Scheller-Wolf "Simple near-optimal scheduling for the $\rm M/G/1$ " SIGMETRICS 2020.

3. Gittins Index Policy – Always optimal, but complex.

First exact analysis of both SERPT and Gittins: Scully, Harchol-Balter, Scheller-Wolf. "SOAP: One clean analysis of all age-based scheduling policies." SIGMETRICS 2018.

EXTRA: Opt scheduling when jobs have value and size
DRAW PICTURE
Question: What scheduling policy makes the most intuitive sense?
The $c\mu$ Rule
Real-world complications
Scully and Harchol-Balter. "The Gittins policy in the $M/G/1$ queue." 19th International Symposium on Modeling and Optimization in Mobile, Ad hoc, and Wireless Networks 2021.

EXTRA: Scheduling when jobs have deadlines

Set up: Each job has a deadline. Job is worthless after its deadline.
Question: What's wrong with Earliest Deadline First?
Better idea: Drop-If-Hopeless (DIH)
Even better idea: Drop-The-Largest-Hopeless (DIH)
(see Discussion in Section 7.2 of this paper: Harchol-Balter. "Open problems in queueing theory inspired by datacenter computing." Queueing Systems, vol. 97, no. 1, pp. 3-37, 2021.