Homework is due at the start of class on 11/04. Please turn in a hardcopy. It can be handwritten. Please show all computations.

Reading from your textbook: Chapters 28-33 (you can skip the math and just concentrate on the text and intuitions).

Exercises:

1. Analysis Problem: Comparing NP-Prio and P-Prio and FCFS

Assume that you have a file server that downloads files. In this problem you will analytically compare 3 policies for scheduling file downloads.

- (i) Non-preemptive Priority (NP-Prio)
- (ii) Preemptive Priority (P-Prio)
- (iii) First-Come-First-Served (FCFS).

A "job" consists of a file download. The "size" (service requirement) of a job is best approximated by the size of the file being downloaded. Files come in 3 classes:

- Class 1 files can have any size between 0 and ∞ , but tend to be small on average: $S_1 \sim \text{Exp}(1)$.
- Class 2 files range from moderate to larger size: $S_2 \sim \text{Uniform}(10, 60)$.
- Class 3 files are all large, ranging in size from 50 to 1000: $S_3 \sim \text{BoundedPareto}(k = 50, p = 1000, \alpha = 1.2).$

The file size distribution is

$$S = \begin{cases} S_1 & \text{w.p. } \frac{1}{3} \\ S_2 & \text{w.p. } \frac{1}{3} \\ S_3 & \text{w.p. } \frac{1}{3} \end{cases}$$

The system load is $\rho = 0.8$ and the arrival process is a Poisson process.

- (a) What is ρ_1 , the load consisting of class 1 jobs? What is ρ_2 ? What is ρ_3 ?
- (b) For each of the three policies, what is $\mathbf{E}[T(1)]$, $\mathbf{E}[T(2)]$, $\mathbf{E}[T(3)]$, and overall $\mathbf{E}[T]$? Here $\mathbf{E}[T(k)]$ denotes the mean response time of class k jobs.
- (c) How are the three policies ranked with respect to overall mean response time, $\mathbf{E}[T]$? Why intuitively does this make sense?
- (d) For which policy is $\mathbf{E}[T(1)]$ lowest? Why is this?
- (e) For which policy is $\mathbf{E}[T(3)]$ highest? Why is this?

2. Simulation Problem: Scheduling to reduce variability

Shashank decides to build a web server to serve file requests. Arrivals into Shashank's web server come from all over the world and are well modeled by a Poisson process. Shashank has measured his file size distribution and found that it's well-modeled by

$$S \sim \text{BoundedPareto}(k=1333.333, p=10^{10}, \alpha=1.8)$$
 .

- (a) In Shashank's initial attempt, he schedules requests in FCFS order. Simulate Shashank's FCFS queue for loads $\rho = 0.1, 0.2, 0.3, \dots, 0.9$ and plot $\mathbf{E}[T]^{\text{FCFS}}$ as a function of ρ .
- (b) Shashank reasons that mean response time is too high because C_S^2 is high. What is C_S^2 ? To remedy the problem, Shashank decides to instead schedule requests in Processor-Sharing (PS) order. Simulate Shashank's PS queue for loads $\rho = 0.1, 0.2, 0.3, \dots, 0.9$ and plot $\mathbf{E}[T]^{PS}$ as a function of ρ . Put this on the same plot as your FCFS results.
- (c) Shashank is still not pleased with his response times at high load. He decides to try scheduling requests in Shortest-Remaining-Processing-Time (SRPT) order. Simulate Shashank's SRPT queue for loads $\rho = 0.1, 0.2, 0.3, \dots, 0.9$ and plot $\mathbf{E}[T]^{\text{SRPT}}$ as a function of ρ . Put this on the same plot as your FCFS and PS results.
- (d) Unfortunately, Shashank's system only allows 3 priority levels, so he can't implement true SRPT. Instead, as an approximation to SRPT he labels:
 - Files with Remaining size < 4000 are Priority 1;
 - Files with 4000 < Remaining size < 20000 are Priority 2;
 - Files with Remaining size > 20000 are Priority 3.

Note that a file of size 30000 will start out as priority 3; but after the first 10000 bytes are sent it will transition to priority 2; and when the file has only 4000 bytes remaining to be sent, it will transition to priority 1. We call this scheduling policy Bucketed-SRPT (BSRPT). Simulate Shashank's BSRPT queue for loads $\rho = 0.1, 0.2, 0.3, \dots, 0.9$ and plot $\mathbf{E}[T]^{\text{BSRPT}}$ as a function of ρ . Put this on the same plot as your FCFS, PS, and SRPT results.

(e) How did BSRPT do? Is this what you expected? Explain.

Here's a paper that actually implements these policies for a web server and also studies issues of unfairness:

Harchol-Balter, Schroeder, Bansal, Agrawal "Size-based scheduling to improve web performance" ACM Transactions on Computer Systems, Vol. 21, No. 2. pp. 207-233, 2003.