

# Supporting hyperplane

$$\min 3x + y \text{ s.t.}$$

$$x + y \geq 2$$

$$x \geq 0$$

$$y \geq 0$$

$$\max 2a \text{ s.t.}$$

$$a + b = 3$$

$$a + c = 1$$

$$a, b, c \geq 0$$

# The graduate student nutrition problem

$$\min 3x + y \text{ s.t.}$$

$$x + y \geq 2$$

$$x \geq 0$$

$$y \geq 0$$

$$\max 2a \text{ s.t.}$$

$$a + b = 3$$

$$a + c = 1$$

$$a, b, c \geq 0$$

# The Lagrangian

primal  
↘

$$\begin{array}{ll} \min & 3x + y \quad \text{s.t.} \\ & x + y \geq 2 \\ & x, y \geq 0 \end{array}$$

- $L(a, b, c, x, y) =$

$$[3x + y] - [a(x + y - 2) + bx + cy]$$

- $\min_{x, y} \max_{a, b, c \geq 0} L(a, b, c, x, y)$

dual  
→

$$\begin{array}{ll} \max & 2a \quad \text{st} \\ & a + b = 3 \\ & a + c = 1 \\ & a, b, c \geq 0 \end{array}$$

# Lagrangian cont'd

$$\begin{array}{ll}\min & 3x + y \text{ s.t.} \\ & x + y \geq 2 \\ & x, y \geq 0\end{array}$$

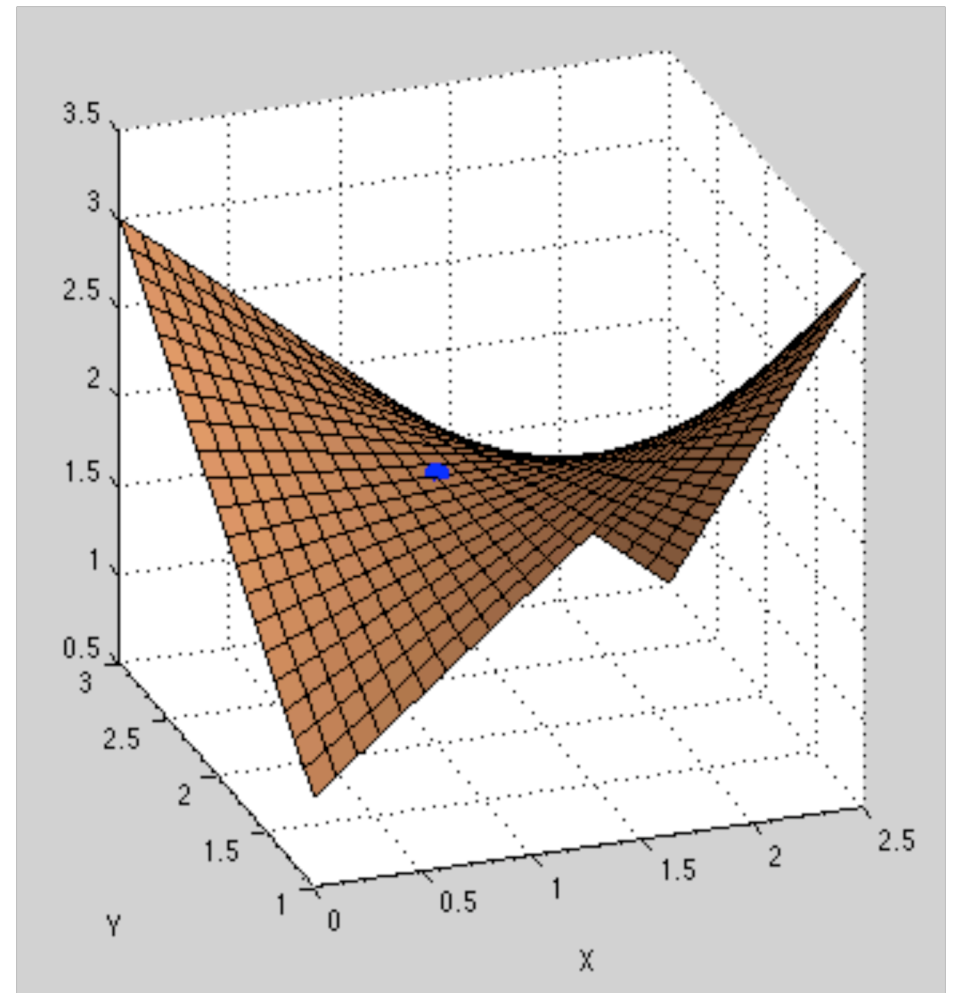
- $L(a, b, c, x, y) =$   
 $[3x + y] - [a(x + y - 2) + bx + cy]$

- $\min_{x, y} \max_{a, b, c \geq 0} L(a, b, c, x, y)$

$$\begin{array}{ll}\max & 2a \text{ st} \\ & a + b = 3 \\ & a + c = 1 \\ & a, b, c \geq 0\end{array}$$

# Saddle-point picture

- $\min y$  s.t.  $y \geq 2$



# + vs – in Lagrangian

$$\min y \quad \text{s.t.} \quad 2 \leq y \leq 4$$

$$\max y \quad \text{s.t.} \quad 2 \leq y \leq 4$$

# Duality summary

## Primal

min problem

max problem

constraint

$\leq$  constraint

$\geq$  constraint

= constraint

variable

## Dual

# Duality summary

## Primal

tight constraint

slack constraint

zero/nonzero variable

infeasible problem

unbounded problem

finite optimal value

## Dual



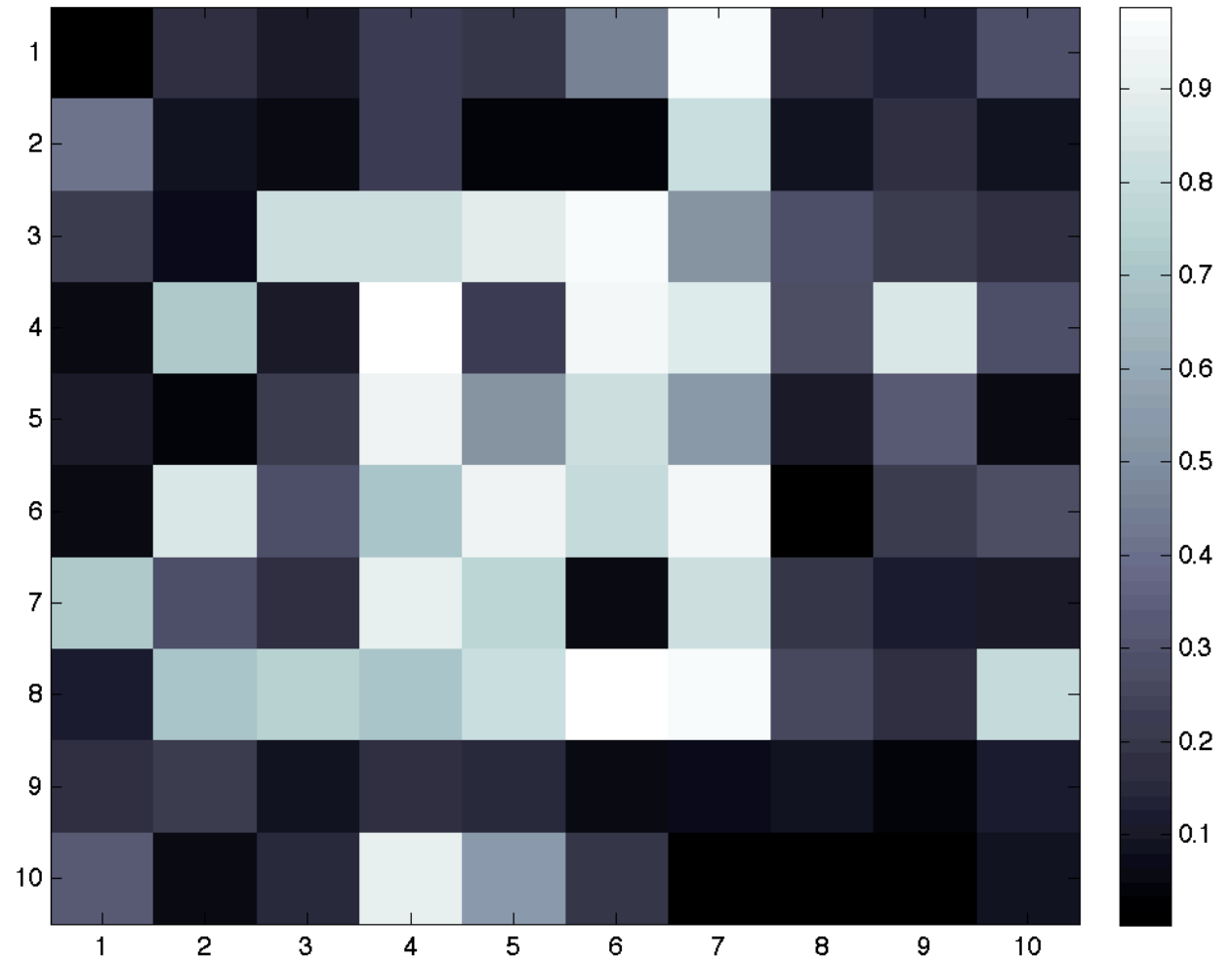
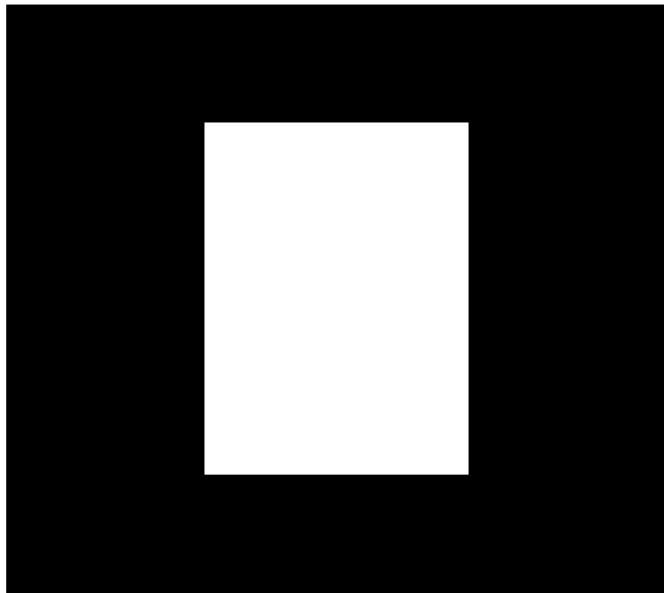
# Example: max flow

- Given a directed graph
  - edges  $(i,j) \in E$
  - flows  $f_{ij}$ , capacities  $c_{ij}$
  - source  $s$ , terminal  $t$  ( $c_{ts} = \infty$ )
- $\max f_{ts}$  s.t.
  - positive flow
  - capacity
  - flow conservation

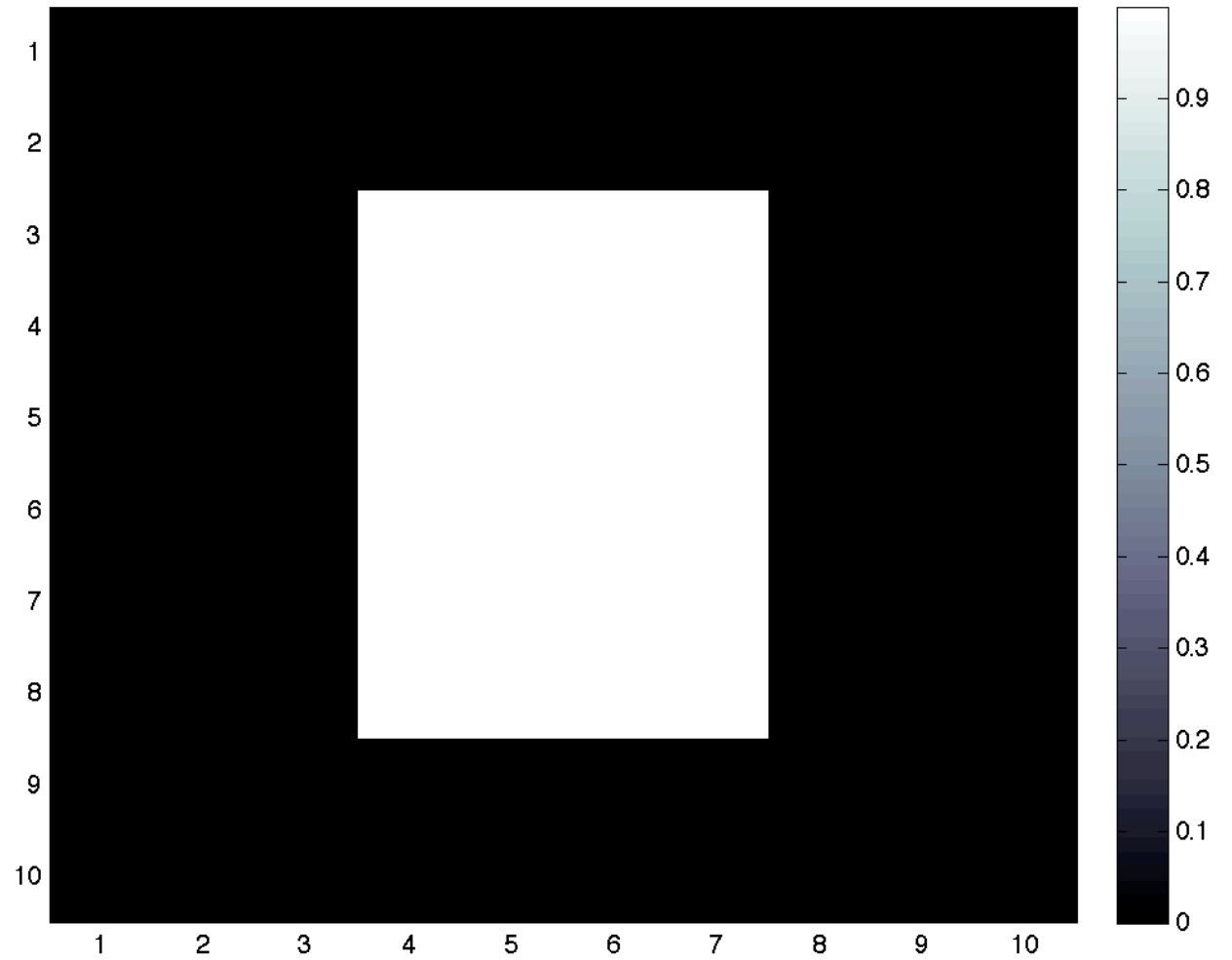
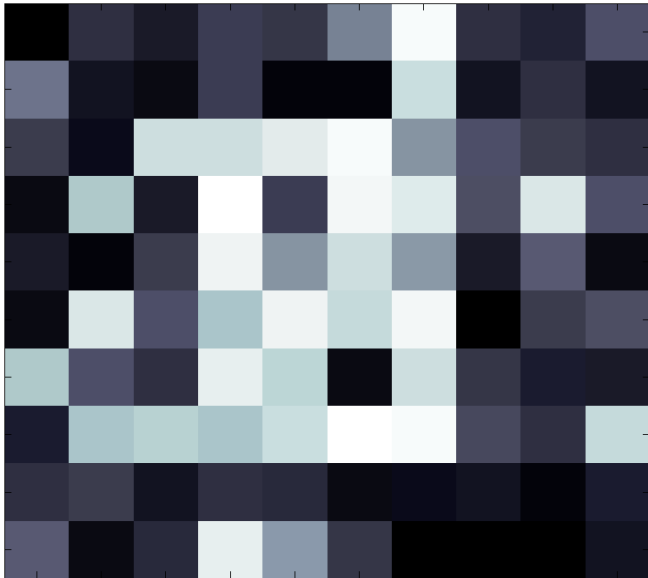
# Dual of max flow

# Interpreting dual

# min cut: image segmentation



# Solution



# What about QP duality?

- $\min x^2 + y^2$  s.t.  
 $x + 2y \geq 2$   
 $x, y \geq 0$
- How can we lower-bound OPT?

# Works at other points too

- $\min x^2 + y^2$  s.t.  
 $x + 2y \geq 2$   
 $x, y \geq 0$
- Try Taylor @  $(x, y) = (v, w)$

# SVM duality

- Recall: min  $\quad$  s.t.
- Taylor bound objective:
- Generic constraint:
  
- To get bound, need:



# SVM dual

- $\max_{\alpha, v} \sum_i \alpha_i - ||v||^2/2 \quad \text{s.t.}$

$$\sum_i \alpha_i y_i = 0$$

$$\sum_i \alpha_i y_i x_{ij} = v_j \quad \forall j$$

$$\alpha_i \geq 0 \quad \forall i$$