

## LP duality cheat sheet

$$\min \underline{c}'x + \underline{d}'y \quad \text{s.t.}$$

$$\underline{Ax} + \underline{By} \geq \underline{p} \quad \leftarrow$$

$$\underline{Ex} + \underline{Fy} = \underline{q} \quad \leftarrow$$

$$\underline{x} \text{ free, } \underline{y} \geq 0$$

$$\begin{pmatrix} \underline{A} & \underline{B} \\ \underline{E} & \underline{F} \end{pmatrix}$$

$$\max \underline{p}'v + \underline{q}'w \quad \text{s.t.}$$

$$\underline{A}'v + \underline{E}'w = \underline{c} \quad \leftarrow$$

$$\underline{B}'v + \underline{F}'w \leq \underline{d}$$

$$\underline{v} \geq 0, \underline{w} \text{ free}$$

$$\begin{pmatrix} \underline{A}' & \underline{E}' \\ \underline{B}' & \underline{F}' \end{pmatrix}$$

→ Swap RHS and objective  
Swap max/min ←

Transpose constraint matrix  
+ve vars yield ≤, free vars yield =

## Linear feasibility problem

LP

$$\min \underline{c}'x \quad \text{s.t.}$$

$$\underline{Ax} + \underline{b} \geq 0$$

$$\underline{c}'x^* \in [\underline{L}, \underline{u}]$$

$$\underline{G} = \frac{\underline{L} + \underline{u}}{2}$$

linear feasibility

$$\text{find } x \quad \text{s.t.}$$

$$\underline{Ax} + \underline{b} \geq 0$$

$$\text{Find } x \quad \text{s.t.}$$

$$\underline{Ax} + \underline{b} \geq 0$$

$$\underline{c}'x \leq \underline{G}$$

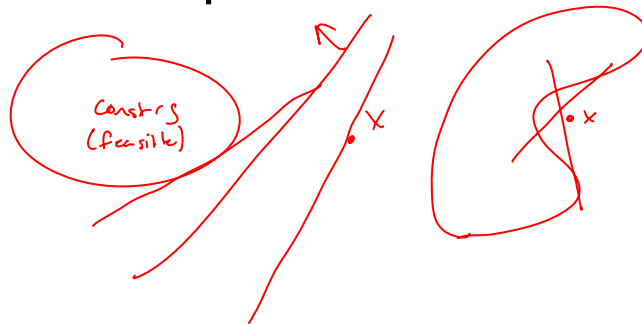
$$\text{feasible} \Rightarrow [\underline{L}, \underline{G}]$$

$$\text{infeasible} \Rightarrow [\underline{G}, \underline{u}]$$

$P(M)$  steps to solve  
↑ polynomial  
↑ bit length of input

$P(M) \log_2(\frac{1}{\epsilon})$  steps  
↑ relative accuracy

## Separation oracle



## Ellipsoid preview



## Difficulties

- How do we get bounding sphere? *later*
- How do we know when to stop?  
 $\forall i \in I \quad a_i^T x + b_i + \eta \geq 0 \quad \eta > 0, \text{ small}$
- Bound region gets complicated—how do we find its center?



## Bounding a partial ellipsoid

- General ellipsoid w/ center  $x_C$ , radius  $R$ :
- Halfspace:  $p^T x + q \leq 0$
- Translate to origin, scale to be spherical

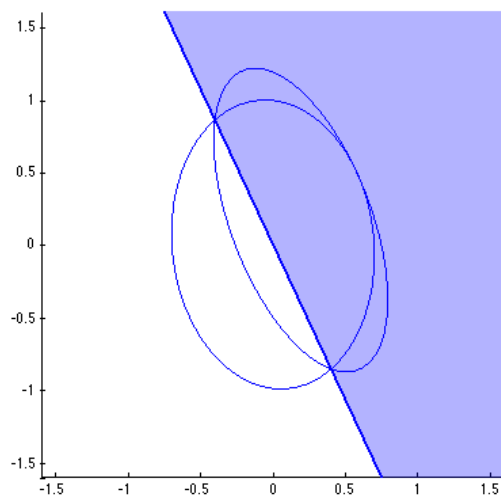
$y =$

$x =$

## Bounding a partial sphere

- Rotate so hyperplane is axis-normal
- New center  $x_c$ :
- New shape A:

For example



## Ellipsoid algorithm

- Want to find  $x$  s.t.  $Ax+b+\eta \geq 0$
- Pick  $R$  s.t.  $\|x^*\| \leq R$
- $E_0 := \{ x \mid \|x\| \leq R \}$
- Repeat:
  - $x_t :=$  center of  $E_t$
  - ask whether  $Ax_t + b \geq 0$ 
    - yes: declare feasible!
    - no: get separating hyperplane
  - $E_{t+1} := \text{bound}(E_t \cap \{ x \mid p_t^T x \leq p_t^T x_t \})$
  - if  $\text{vol}(E_{t+1}) \leq \varepsilon \text{vol}(E_0)$ : declare infeasible!

## Getting bounds

- How big do  $L, U$  need to be?
- How big does  $R$  need to be?
- What should  $\eta$  be?
- How small does  $\varepsilon$  need to be?

## Other algorithms

- Ellipsoid is polynomial, but slow
- Some other algorithms:
  - simplex: exponential in worst case, but often fast in practice
  - randomized simplex: polynomial [Kelner & Spielman, 2006]
  - interior point: polynomial
  - subgradient descent: weakly polynomial, but really simple, and fast for some purposes

What's a subgradient?

## Subgradient descent for SVMs

- $\min_{s,w,b} ||w||^2 + C \sum_i s_i$  s.t.  
 $y_i(x_i^T w - b) \geq 1 - s_i$   
 $s_i \geq 0$
- Equivalently,

## Subgradient in SVM

- $\min_w L(w) = ||w||^2 + C \sum_i h(y_i x_i^T w)$
- Subgradient of  $h(z)$ :
- Subgradient of  $L(w)$  wrt  $w$ :