

Convex optimization

e.g., min s.t.

Linear inequalities:

Positivity:

Terminology

- Feasible
 - Infeasible
 - Unbounded
 - Two convex programs are equivalent if:

Transformations of CPs

- Optimizing some variables
- $\min f(x,y) = y^2 + (x - 2y)^2 / 2$

Transformations of CPs

- Implicit constraints
- $\min x^2$ s.t. $x \geq 1$
- $\min x^2 + I_{x \geq 1}(x)$
 $I_{x \geq 1}(x) =$
- For any set S , $I_S(x) =$

Transformations of CPs

- Epigraph form:
- $\min x^2$ s.t. $x \geq 2$
- $\min t$ s.t.

Optimality

- Optimal value:
 $v^* = \inf \{ f(x) \mid g_i(x) \leq 0 \ (\forall i), Ax = b \}$
- Definition of $v^* = \inf S$:
- For example: $\inf \{ e^x \mid x \text{ real} \} =$

Optimal point

- x^* is optimal iff:
- Is there always an optimal point?
 - ex. 1:
 - ex. 2:

Local optima

- $\min f(x)$ s.t. $g(x) \leq 0$
- x_0 is a local optimum iff
 - In a convex program,

An optimality criterion

- Suppose $f(x)$ is differentiable, convex
- Then $x^* = \arg \min f(x)$ iff
- What about $x^* = \arg \min f(x)$ s.t. $x \in C$?

Proof of criterion

Example: logistic regression

- Data $x_i \in \mathbb{R}$, $y_i \in \{0, 1\}$
- $P(y_i = 1 | x_i, w, b) = s(w^T x_i + b)$
 $s(z) = 1/(1+\exp(-z))$
- $\max_{w,b} P(w, b) P(y | x, w, b) =$
 $\max_{w,b} P(w, b) \prod_i P(y_i | x_i, w, b) =$

Ex: minimize top eigenvalue

- Suppose $A_i \in S^{n \times n}$ for $i = \{1, 2, \dots\}$
- $\min_{\alpha, A} \lambda_1(A)$ s.t.
 $A = \alpha_1 A_1 + \alpha_2 A_2 + \dots$

Ex: minimum volume ellipsoid

- Given points x_1, x_2, \dots, x_k
- $\min_{A, x_C} \text{vol}(A)$ s.t.
$$(x_i - x_C)^T A (x_i - x_C) \leq 1 \quad i = 1, \dots, k$$
$$A \in S^{n \times n}$$
$$A \geq 0$$

Schur complement

- Symmetric block matrix $M =$
- Schur complement is $S =$
- $M \geq 0$ iff

Back to min-volume ellipsoid

- $\max_{A, x_C} \log |A|$ s.t.
 $(x_i - x_C)^T A (x_i - x_C) \leq 1 \quad i = 1, \dots, k$
 $A = A^T, A \geq 0$
- $\max_{A, B, x_C, z} \log |A|$ s.t.