

Convex optimization

minimize $f(x)$ subject to $Ax = b$

Handwritten notes: $\mathbb{R}^n$ points to $f(x)$, $\mathbb{R}^{m \times n}$ points to A , $\mathbb{R}^m$ points to b . A red arrow labeled "convex" points from $f(x)$ to the word "convex". A red arrow points from $Ax = b$ to the word "subject to". A red parabola is drawn above the text.

e.g., min $x^2 + y^2$ s.t. $g_i(x, y) \leq 0 \quad \forall i \in I$

Handwritten notes: "convex" points to $x^2 + y^2$. $g_i(x, y)$ is written above $1 - x - y \leq 0$.

Linear inequalities: OK
Positivity: OK

Handwritten note: A red bracket groups "Linear inequalities" and "Positivity" with the text "just include in g_i 's".

Terminology

- Feasible $\exists x. g_i(x) \leq 0 \quad \forall i \wedge Ax = b$
- Infeasible *not feasible*
- Unbounded $\forall v \in \mathbb{R} \exists \text{ feasible } x. f(x) \leq v$

- Two convex programs are equivalent if:

Handwritten note: change of vars, slack on linear ineq., max/min or $\leq / \geq$ by negating

Transformations of CPs

- Optimizing some variables
- $\min_{x,y} \underline{f(x,y)} = y^2 + (x - 2y)^2 / 2$
 $\partial f / \partial x = x - 2y = 0 \Rightarrow x = 2y$
 $\min y^2 + \cancel{(2y - 2y)^2 / 2}$

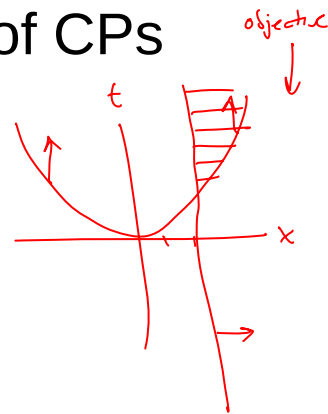
Transformations of CPs

- Implicit constraints
- $\min \underline{x^2} \text{ s.t. } \underline{x \geq 1}$
- $\min \underline{x^2 + I_{x \geq 1}(x)}$ ←

$$I_{x \geq 1}(x) = \begin{cases} 0 & x \geq 1 \\ \infty & \text{o/w} \end{cases}$$
- For any set S , $\underline{I_S(x)} = \begin{cases} 0 & x \in S \\ \infty & x \notin S \end{cases}$

Transformations of CPs

- Epigraph form:
- $\min_x x^2 \text{ s.t. } x \geq 2$
- $\min_{x,t} t \text{ s.t. } x \geq 2 \quad t \geq x^2$

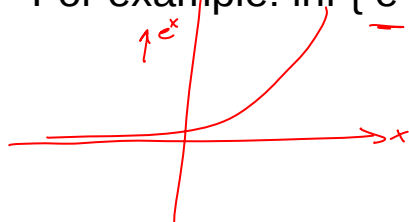


Optimality

- Optimal value:

$$v^* = \inf \{ f(x) \mid g_i(x) \leq 0 \ (\forall i), Ax = b \}$$
- Definition of $v^* = \inf S$: $\rightarrow \subseteq \mathbb{R}$ $\rightarrow$ infimum
 $\Leftrightarrow \forall \epsilon > 0 \exists x_\epsilon \in S. v^* + \epsilon \geq x_\epsilon$
- For example: $\inf \{ e^x \mid x \text{ real} \} =$

Also, $v^* \leq x$
 $\forall x \in S$



Sup (supremum)
 + tightest upper bound

Optimal point

- x^* is optimal iff:
- Is there always an optimal point?
 - ex. 1:
 - ex. 2:

Local optima

- $\min f(x) \text{ s.t. } g(x) \leq 0$
- x_0 is a local optimum iff
- In a convex program,

An optimality criterion

- Suppose $f(x)$ is differentiable, convex
- Then $x^* = \arg \min f(x)$ iff
- What about $x^* = \arg \min f(x)$ s.t. $x \in C$?

Proof of criterion

Example: logistic regression

- Data $x_i \in \mathbb{R}$, $y_i \in \{0,1\}$
- $P(y_i = 1 \mid x_i, w, b) = s(w^T x_i + b)$
 $s(z) = 1/(1+\exp(-z))$
- $\max_{wb} P(w, b) P(y \mid x, w, b) =$
 $\max_{wb} P(w, b) \prod_i P(y_i \mid x_i, w, b) =$

Ex: minimize top eigenvalue

- Suppose $A_i \in S^{n \times n}$ for $i = \{1, 2, \dots\}$
- $\min_{\alpha, A} \lambda_1(A)$ s.t.
 $A = \alpha_1 A_1 + \alpha_2 A_2 + \dots$

Ex: minimum volume ellipsoid

- Given points $x_1, x_2, \dots, x_k$
- $\min_{A, x_c} \text{vol}(A)$ s.t.
 $(x_i - x_c)^T A (x_i - x_c) \leq 1 \quad i = 1, \dots, k$
 $A \in S^{n \times n}$
 $A \succcurlyeq 0$

Schur complement

- Symmetric block matrix $M =$
- Schur complement is $S =$
- $M \succcurlyeq 0$ iff

Back to min-volume ellipsoid

- $\max_{A, x_C} \log |A|$ s.t.
 $(x_i - x_C)^T A (x_i - x_C) \leq 1 \quad i = 1, \dots, k$
 $A = A^T, A \succcurlyeq 0$
- $\max_{A, B, x_C, z} \log |A|$ s.t.

More examples

- SDP: learning a distance matrix (multidimensional scaling)
- learning a kernel matrix
- semidefinite embedding
- maximum variance unfolding