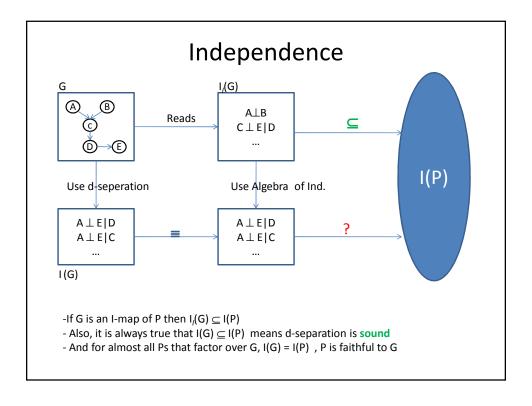
# BN Semantic II d-Separation, PDAGs, etc

Amr Ahmed Sep 25<sup>th</sup>, 2008

#### Outline

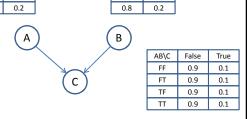
- Independence
- D-separation and Active Trails
- P-dag learning



# **D-Separation**

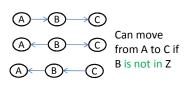
- A graph algorithm for answering independence queries over G
- Sound and complete for almost all Ps that factor according to G
- P is faithful if it doesn't declare extra independence assumption that can't be read from G

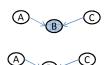
P entails C ⊥ A,B!!
P is not faithful



#### **D-Separation Cont**

- Q: is X⊥ Y | Z?
- · Answer by contradiction
  - Find a way that information flows between X and Y despite the existence of Z
  - Information can flow if there is a path x...y that is not blocked by z (active trail)
- Very simple "local" rules





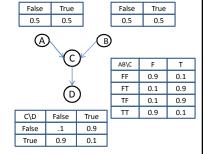
Can move from A to C if B is in Z or one of its decedents

## Understanding the V-structure more

- -C is a noisy X-or of A and B
- If C is not observed then A and B are uniform
- If you observe C, then A and B are dependent
  - -C=1, → A= not B w.h.p
  - -C=0 → A=B w.h.p

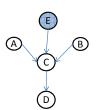
#### Observing a decedent

- -D is a noisy NOT of C
- If you observe D, then w.h.p you know C
- -If you have an idea about C, A and B are dependent



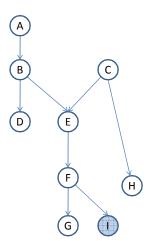
#### Observing an ancestor!!

- -Won't help, E alone shouldn't tell us much about C
- But in some situations it DOES, but later in the semester (context-specific independence)



## **D-separation Example**

- -Given I is  $\mathbf{A} \perp \mathbf{C}$
- Given I is A ⊥ F
- Given I and B is  $A \perp C$



## Why D-separation is useful?

- Intuitively and on an abstract level, when you answer a probabilistic query P(A|B), you would like to consider only those variables that would affect A given B
- Later, when we talk about inference, we will visit this again
- The concept of active trails is really very important in proving and justifying algorithms
  - You should use it in Q2 and Q4

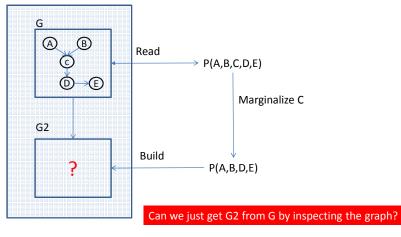
#### **Active Trails**

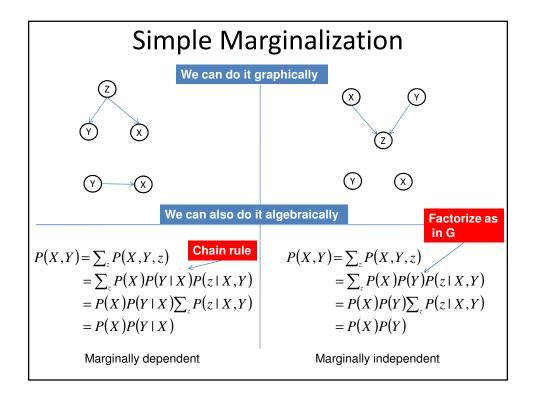
- If it is all about independence, then to show that two graphs, G1 and G2 are equivalent, we need to show that I(G1) = I(G2), or practically:
  - A trail is active in G1 iff it is active in G2\*\*
  - More algorithmically
    - Consider all ways in which some of the variables are observed
    - Show that all active trails in G1 and G2 are the same

\*\* This is only true if G1 and G2 have no triangles, in case of triangles, we require that they agree on the set of **minimal active trails** (see problem 3.16). For this homework, we won't worry too much about this subtlety.

#### Question 4 again

 Marginalization is a key operation, that we will use later in the semester.



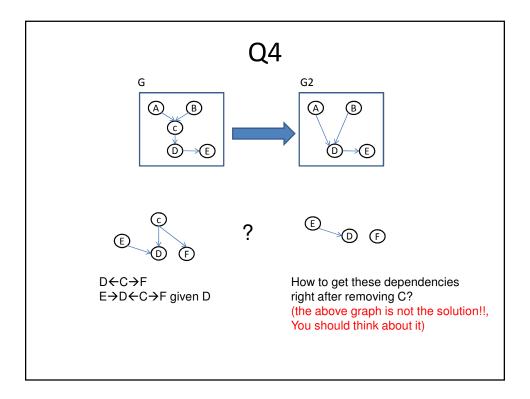


#### Q4 again

- Removing C, introduces new independence assumptions not in G like A ⊥ D
- We need to add more edges to compensate
- You need to consider what active trails are enabled by C
  - $-A \rightarrow C \rightarrow D$
  - But also, A → C ← B given D?
- Think what need to be done to make sure that the end variables are still dependent in G2 under the same conditions but when C is marginalized
- Sometimes fixing a trail will fix the other ones (we need to add the minimum number of edges)
- Hint: think first about trails that don't require observing any other variables as fixing them might fix the others!







#### Outline

- Independence
- D-separation and Active Trails
- P-dag learning

#### **PDAG**

- PDAG is a compact way of representing equivalent graphs
- Orient edges only if they must be this way
- Undirected edges can be either way
  - Remember key is active trails
  - For some active trails (other than v-structure --immoralities--) edge direction is not important
  - $A \rightarrow B \rightarrow C$
  - $A \leftarrow B \rightarrow C$

Can move from A to C if B is not in Z





#### 

# **Learning P-DAGs**

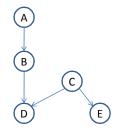
- Learning the skeleton
- Discovering immoralities
- Orienting edges (this is straightforward)

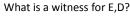
## Learning the skeleton

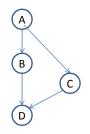
- There is an edge between X,Y if you can not stop information flow between them
- You can stop information flow if you can block all paths between X,Y
- You can block a path, if you observe some variables (possibly empty set) = U
- The test:
  - Can you find U such that  $X \perp Y \mid U$ ?
    - If NO  $\rightarrow$  then x-Y
    - If Yes → then there is no edge

#### Step 1: Learning the skeleton

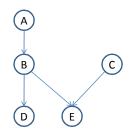
- Test: Can you find U such that X ⊥ Y | U?
  - What is U? subset of all variables- {X,Y}
  - Can go up to size d (max fan in, or degree)? Why?
  - You don't have to go over all possible U
    - A witness is all what you need to answer YES







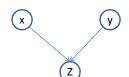
What is a witness for A,D?



What is a witness for B,C?

## Step 2: Discover Immoralities

- For immoralities, we must direct edges in a certain way, so we should discover them
- A v-structure with no married parents
- Simple test:
  - Is X dependent on Y given Z?
    - If yes, ...
    - If no, ...



#### Step 2: Discover Immoralities

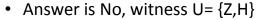
- This simple test will introduce false positives
- If it is a true immorality, we are OK
- But what about:



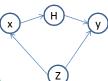
- Given Z, X and Y are dependent
  - But not via Z, unfo. via another path

#### Step 2: Discover Immoralities

- Simple test: Is X dependent on Y given Z? that fails
- Should be: Is X dependent on Y given Z via a path that goes only through Z?
  - Practically we should block all other paths that lead from X to Y
  - In addition to observing Z, we might observe as many other variables as possible
- Test: Is X dependent on Y for all U, z in U?
  - Yes → immorality
  - $-NO \rightarrow$  not immorality



- · We are really asking the same questions
- Skeleton: Can you find U such that X ⊥ Y | U?
  - Yes → no edge,
  - No → an edge
- Immorality: Can you find U such that X ⊥ Y | U?
  - Yes, and z in U → not immorality
  - Yes, and z not in U → immorality
  - The answer here can NOT be NO, why?
- This has been exploited via cashing in the book (but see the extra credit problem)
  - As instructed, you shouldn't cache in your solution.
  - You should consider all U that contains Z until you find a witness (if there is one)



## Some Hints to the programming problem

- A suggestion about representing PDAG
  - G(a,b)= 1, G(b,a) = 1 if a-b
  - G(a,b) = 2 and G(b,a) = 0 if  $a \rightarrow b$
  - -G(a,b) = 0 and G(b,a) = 2 if  $a \leftarrow b$
  - Makes life easier, ex. Check if a→b
    - If (G(a,b) ==2)
    - In old representation, if(G(a,b)==1 && G(b,a)==0)
- Size of U in witness test
  - You need only up to d
  - But it won't hurt to go up to 2\*d, why?
    - After all you are looking for a witness.