15-780: Grad Al Lecture 15: Planning

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Review

- Planning algorithms
 - reduce to FOL (complications)
 - or use subset of FOL (e.g., STRIPS)
 - linear planner: add op to end of plan
 - partial-order planner (operators, bindings, partial order, guards, open preconditions): resolve open precond
- STRIPS: (world) state, operator =
 { preconditions } + { effects }, variable binding, goals

Reminder

- HWs due today
- Project proposals due Thu

Plan Graphs

Planning & model search

- For a long time, it was thought that SAT-style model search was a non-starter as a planning algorithm
- More recently, people have written fast planners that
 - propositionalize the domain
 - turn it into a CSP or SAT problem
 - search for a model

- Tool for making good CSPs: plan graph
- Encodes a subset of the constraints that plans must satisfy
- Remaining constraints are handled
 - during search (reject solutions that violate them)—needs special-purpose code
 - or by adding extra clauses/constraints

Example

- Start state: have(Cake)
- Goal: have(Cake) ∧ eaten(Cake)
- Operators: bake, eat
- Bake
 - pre: ¬have(Cake)
 - post: have(Cake)

- Eat
 - pre: have(Cake)
 - post: ¬have(Cake), eaten(Cake)

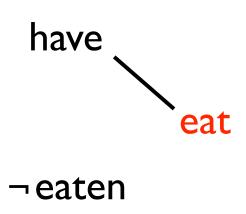
Propositionalizing

- Note: this domain is fully propositional
- If we had a general STRIPS domain, would have to pick a universe and propositionalize
- E.g., eat(x) would become eat(Banana),
 eat(Cake), eat(Fred), ...

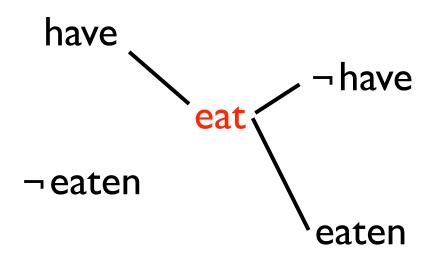
have

¬ eaten

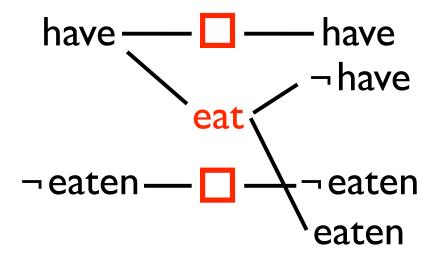
- Alternating levels: states and actions
- First level: initial state



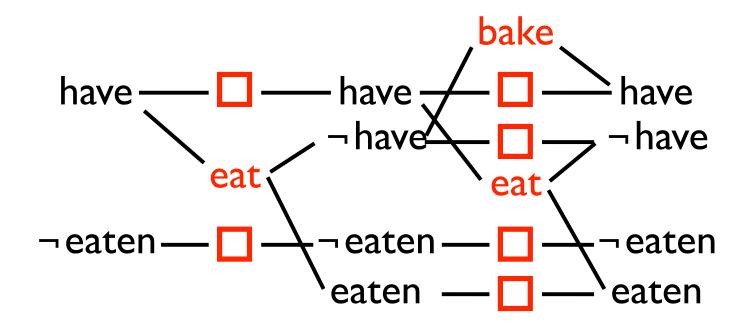
- First action level: all applicable actions
- Linked to their preconditions



 Second state level: add effects of actions to get literals that could hold at step 2

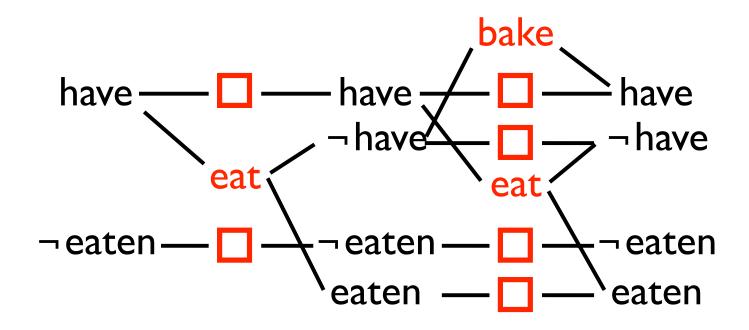


 Also add maintenance actions to represent effect of doing nothing

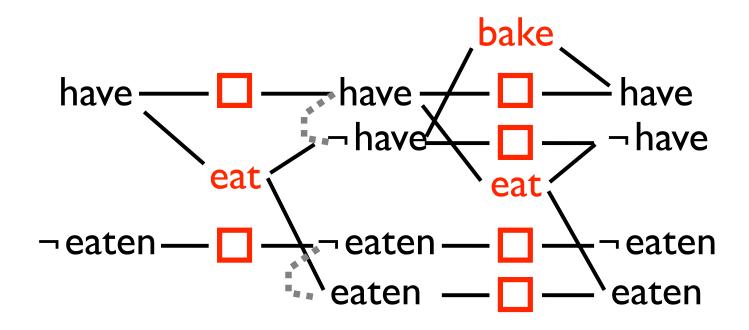


Extend another pair of levels: now bake is a possible action

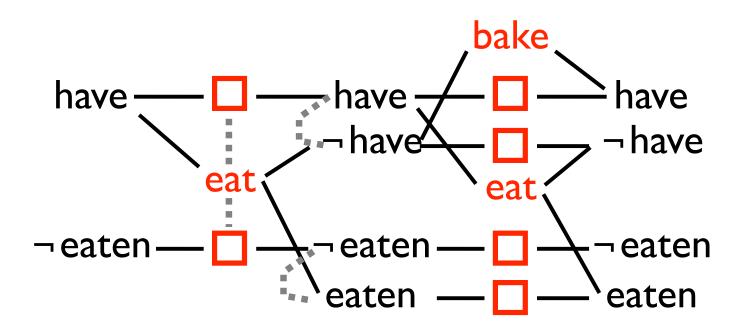
- Can extend as far right as we want
- Plan = subset of the actions at each action level
- Ordering unspecified within a level



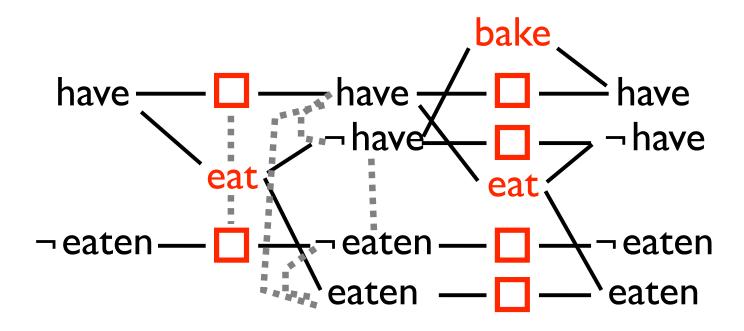
 In addition to the above links, add mutex links to indicate mutually exclusive actions or literals



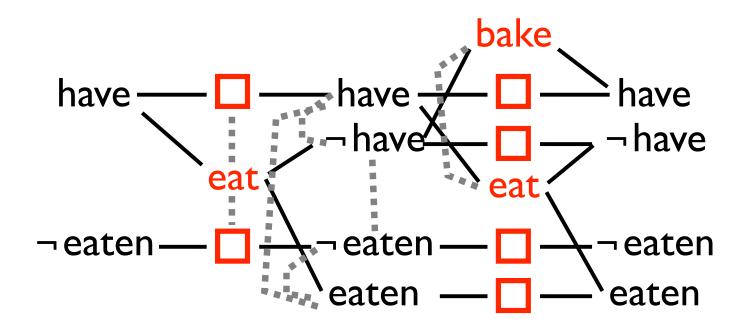
Literals are mutex if they are contradictory



 Actions which assert contradictory literals are mutex (inconsistent effects)



 Literals are also mutex if there is no action or non-mutex pair of actions that could achieve both (inconsistent support)



 Actions are also mutex if one deletes a precondition of other (interference), or if preconditions are mutex (competition)

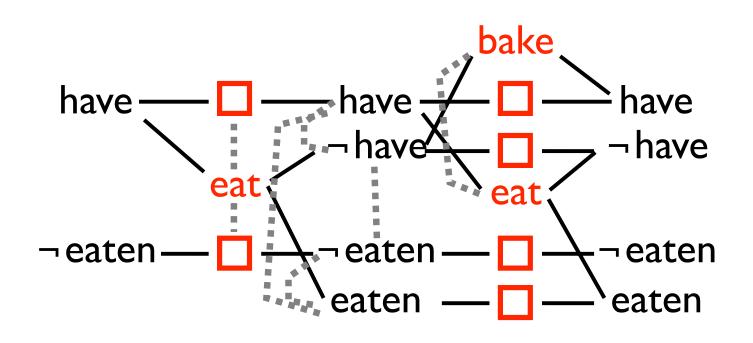
Mutex summary

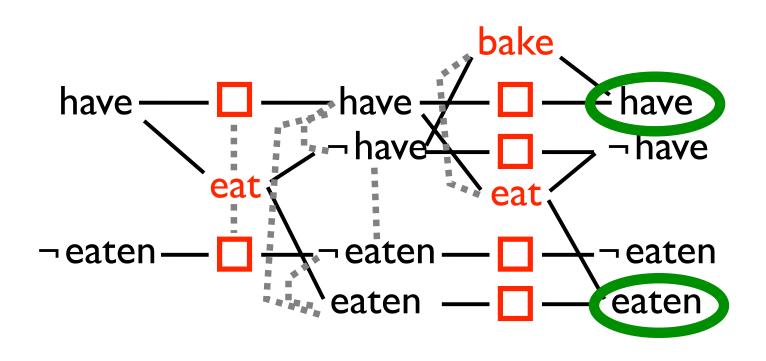
- For each action level, left to right, check pairs of actions A, B (each check linear in rep'n size):
 - inconsistent effects: check each effect of A vs. effects of B
 - interference: effects of A vs. preconds of B
 - competing preconditions: mutex links on preconditions of A, B
- Results at action level L tell us (in)consistent support at proposition level L+I

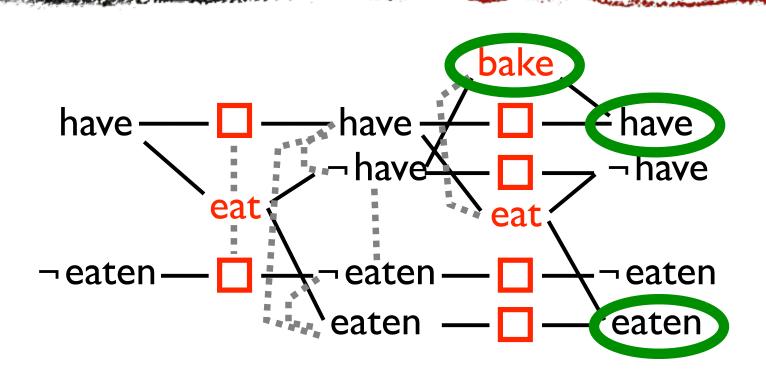
Getting a plan

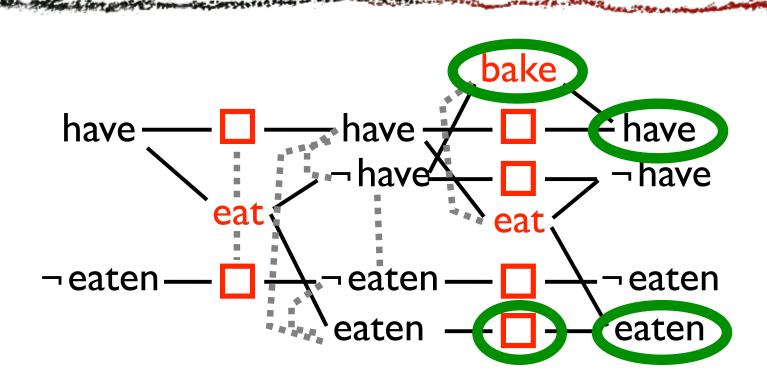
- Build the plan graph out to some length k
- Search:
 - directly on the graph
 - or by translating to SAT or CSP
- If search succeeds, read off the plan
- If not, increment k and try again
- There is a test to see if k is "big enough"

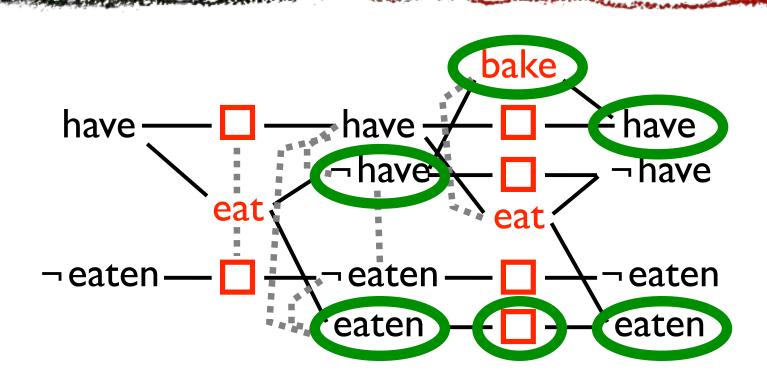
- DFS w/ variable ordering based on plan graph
- Start from last level, fill in last action set, compute necessary preconditions, fill in 2ndto-last action set, etc.
- If at some level there is no way to do any actions, or no way to fill in consistent preconditions, backtrack

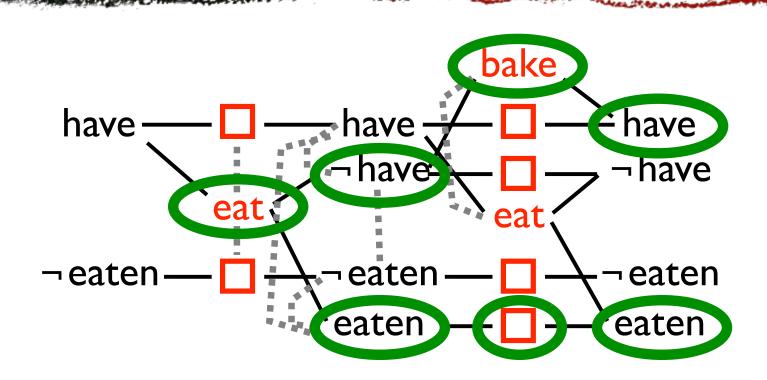


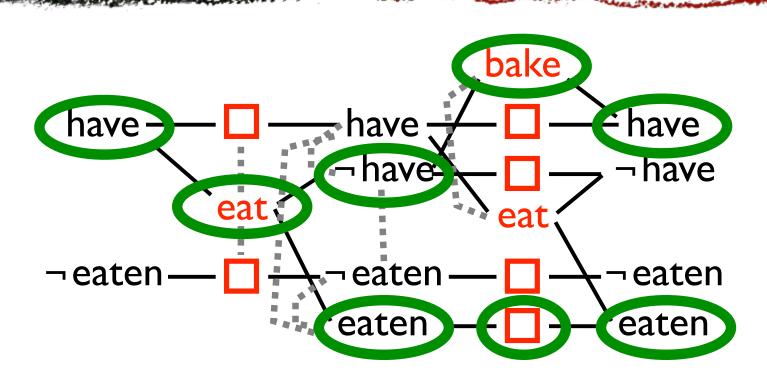












Translation to SAT

- One variable per pair of literals in state levels
- One variable per action in action levels
- Constraints implement STRIPS semantics plus "hints"
- Solution tells us which actions are performed at each action level, which literals are true at each state level

Action constraints

 Each action can only be executed if all of its preconditions are present:

$$act_{t+1} \Rightarrow prel_t \land prel_t \land ...$$

If executed, action asserts its postconditions:

$$act_{t+1} \Rightarrow postl_{t+2} \land post2_{t+2} \land \dots$$

Literal constraints

- In order to achieve a literal, we must execute an action that achieves it
 - ▶ post_{t+2} \Rightarrow act $I_{t+1} \lor$ act $2_{t+1} \lor ...$
- Might be a maintenance action

Initial & goal constraints

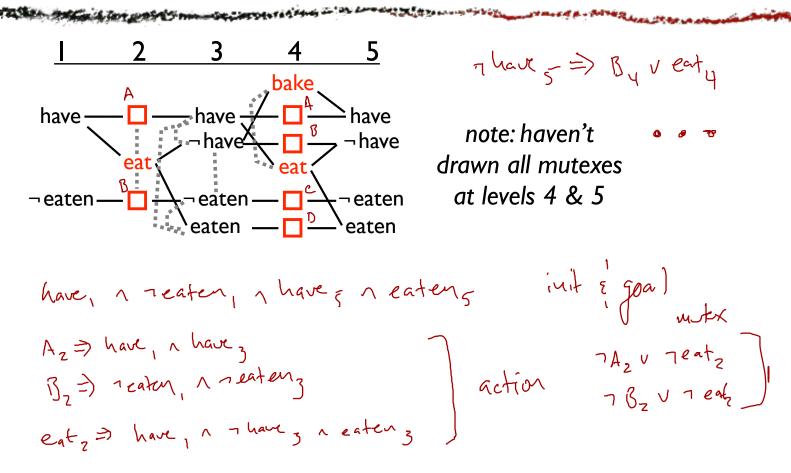
Goals must be satisfied at end:
 goal I T A goal 2T A ...

And initial state holds at beginning:
 init I | \(\lambda \) init 2 | \(\lambda \) ...

Mutex constraints

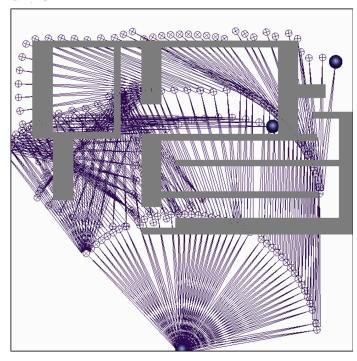
- Mutex constraints between actions or literals: add clause (¬x v ¬y)
- Mutexes are redundant, but help anyway

Translation to SAT: example



Spatial Planning

Plans in Space...



Optimal Solution



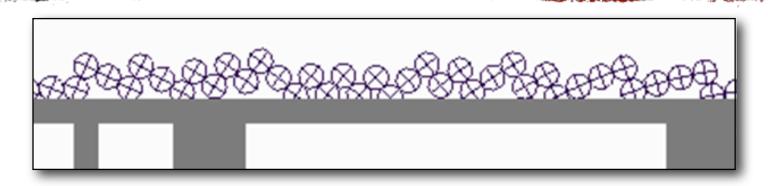
End-effector Trajectory

- A* can be used for many things
- Here, A* for spatial planning (in contrast to, e.g., jobshop scheduling)

What's wrong w/ A*?

- A* guarantees:
 - (optimality) A* finds a solution of cost g*
 - (efficiency) A* expands no nodes that have f(node) > g*

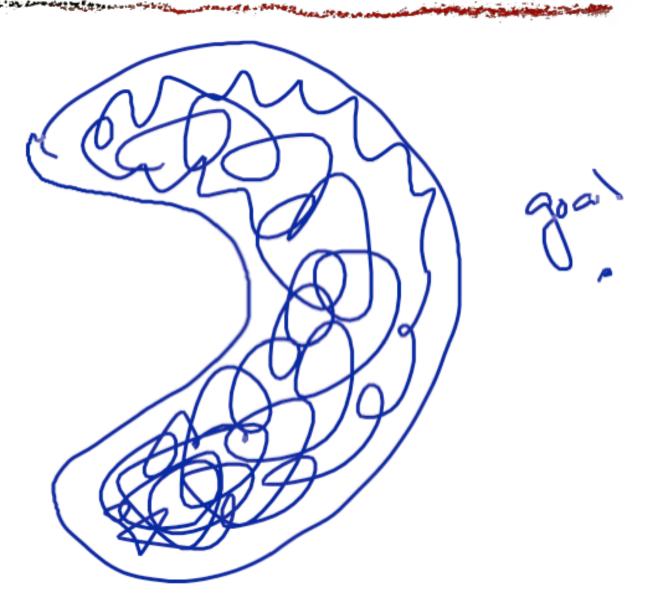
What's wrong with A*?



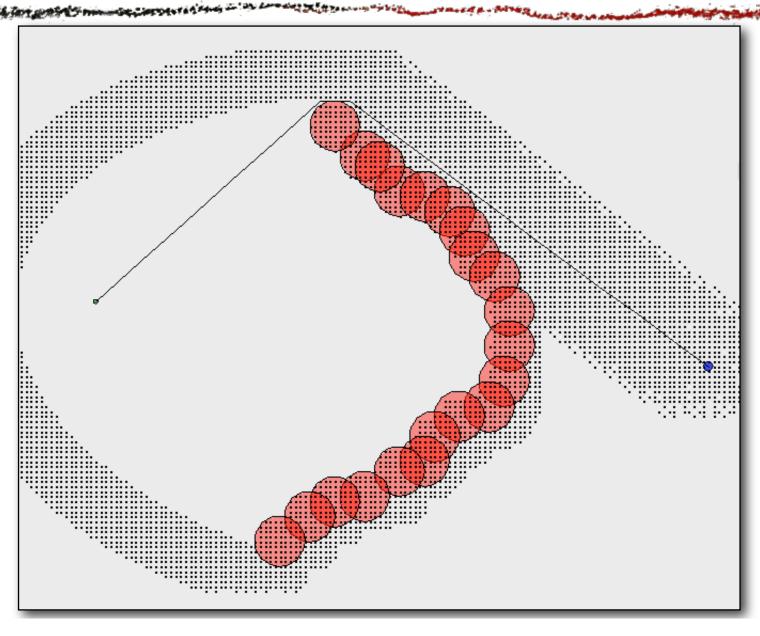
- Discretized space into tiny little chunks
 - a few degrees rotation of a joint
 - ▶ **Lots** of states \Rightarrow lots of states w/ f \leq g*
- Discretized actions too
 - one joint at a time, discrete angles
- Results in jagged paths

What's wrong with A*?

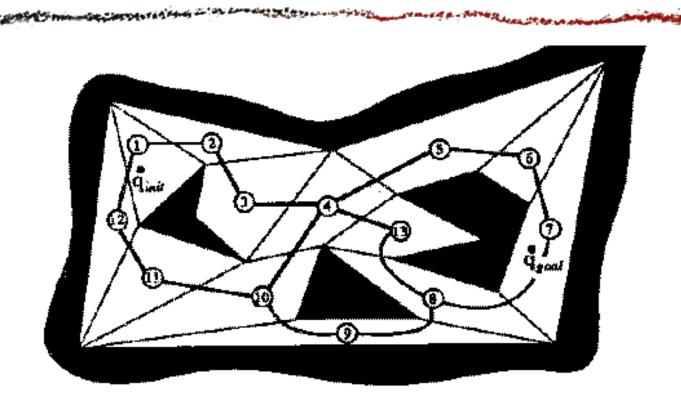
sparx.



Snapshot of A*



Wouldn't it be nice...

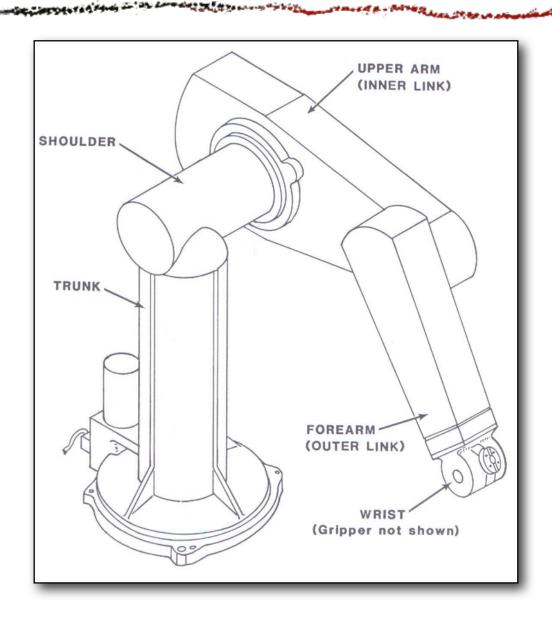


- if we could break things up based more on the real geometry of the world?
- Robot Motion Planning, Jean-Claude Latombe

Physical system

- Moderate number of real-valued coordinates
- Deterministic, continuous dynamics
- Continuous goal set (or a few pieces)
- Cost = time, work, torque, ...

Typical physical system



A kinematic chain

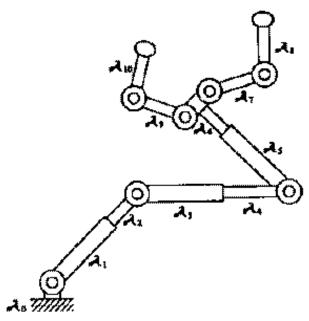


Fig. 11. Structure of the 10-DOF manipulator.

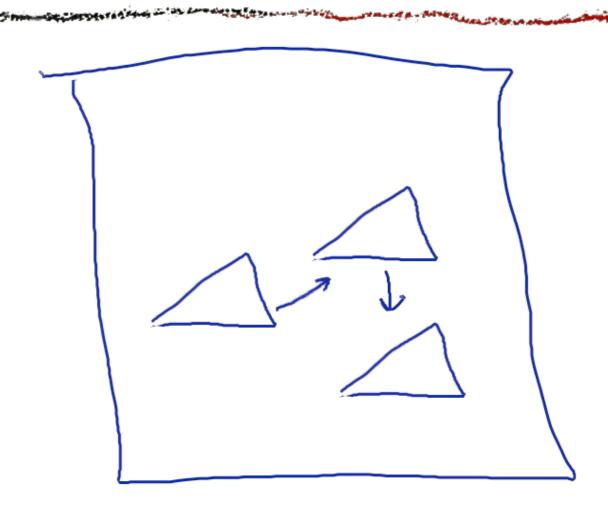
- Rigid links connected by joints
 - revolute or prismatic
- Configuration

$$\mathbf{q} = (q_1, q_2, ...)$$

 q_i = angle or length of joint i

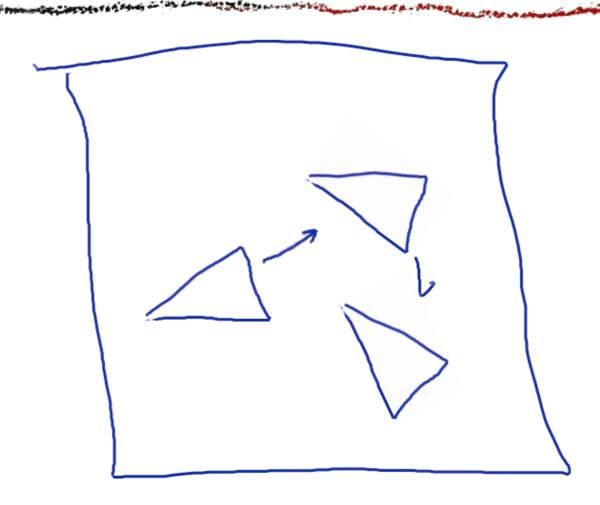
 Dimension of q = "degrees of freedom"

Mobile robots



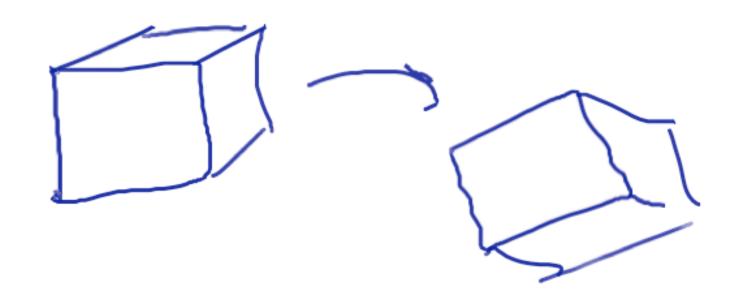
• Translating in space = 2 dof

More mobility

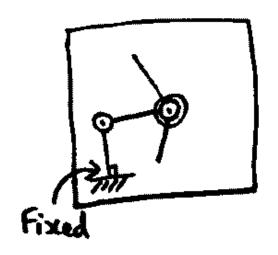


Translation + rotation = 3 dof

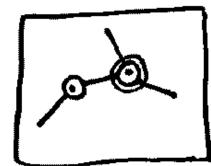
Q: How many dofs?



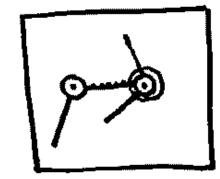
3d translation & rotation



How many dofs?



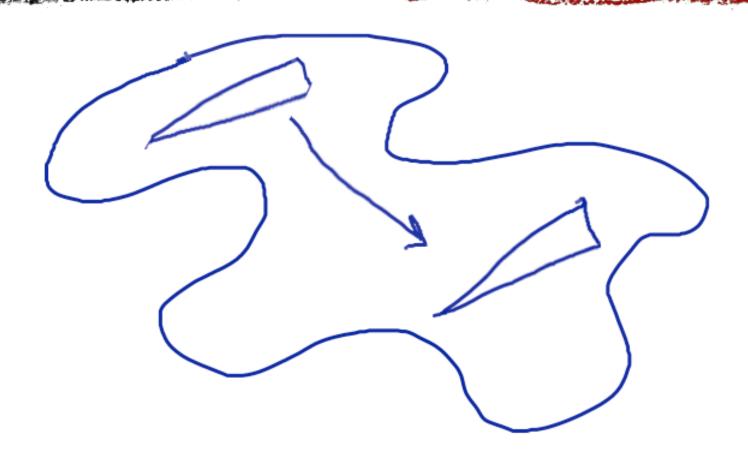
Free flying How many dofs?



Midline Must always be horizontal. How many DOFs?

The configuration of has one real valued entry per DOF.

Kinematic motion planning



Now let's add obstacles

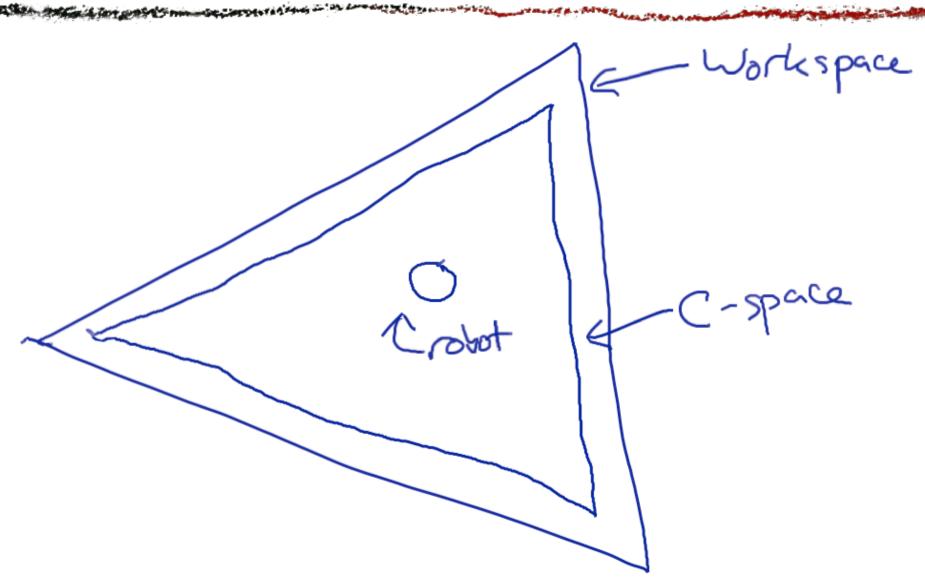
Configuration space

- For any configuration q, can test whether it intersects obstacles
- Set of legal configs is "configuration space"
 C (a subset of a dof-dimensional vector space)
- Path is a continuous function from [0,1] into C with $q(0) = \mathbf{q_s}$ and $q(1) = \mathbf{q_g}$

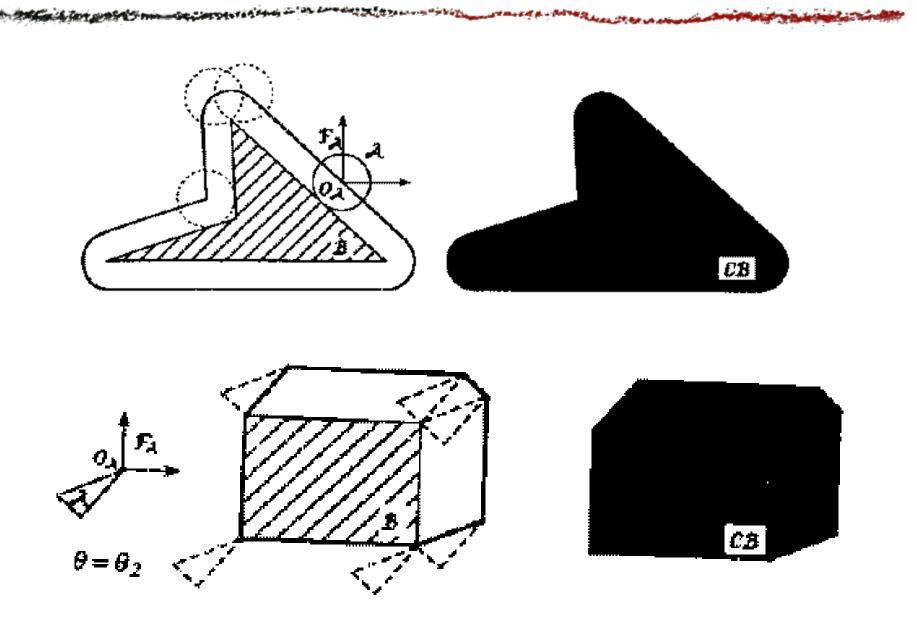
Note: dynamic planning

- Includes inertia as well as configuration
 - q, q
- Harder, since twice as many dofs, and typically stronger constraints
- Won't really cover here...

C-space example



More C-space examples



Another C-space example

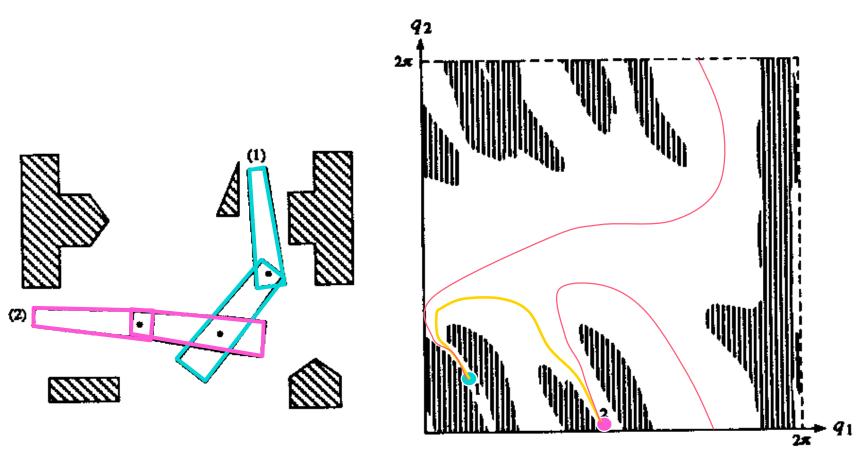
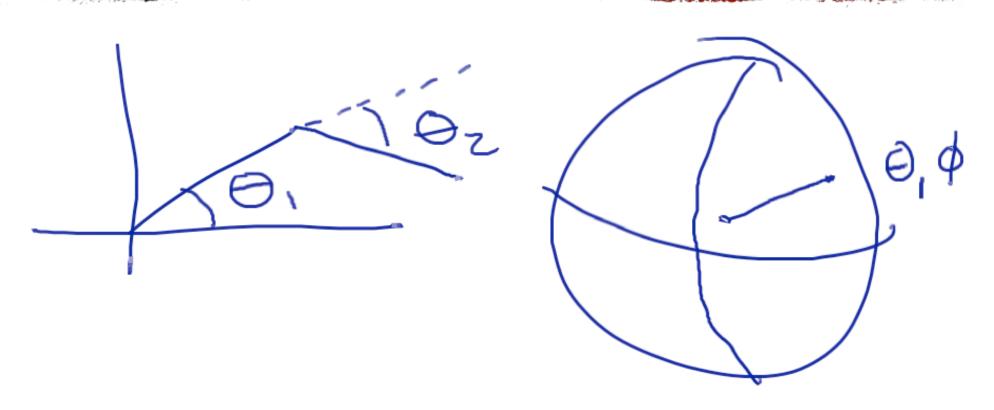


image: J. Kuffner

Topology of C-space

- Topology of C-space can be something other than the familiar Euclidean world
- \circ E.g. set of angles = unit circle = SO(2)
 - not $[0, 2\pi)$!
- Ball & socket joint (3d angle) ⊆ unit sphere
 = SO(3)

Topology example

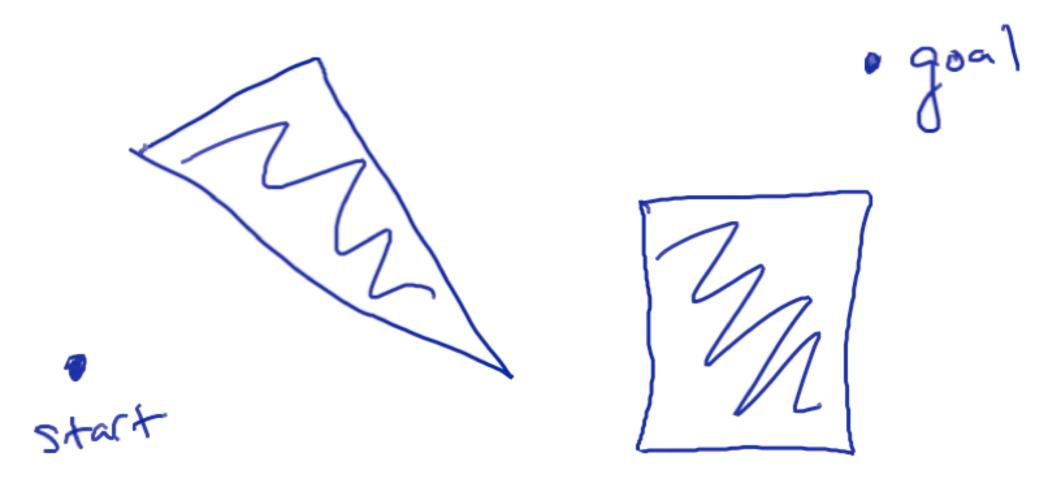


 Compare L to R: 2 planar angles v. one solid angle — both 2 dof (and neither the same as Euclidean 2-space)

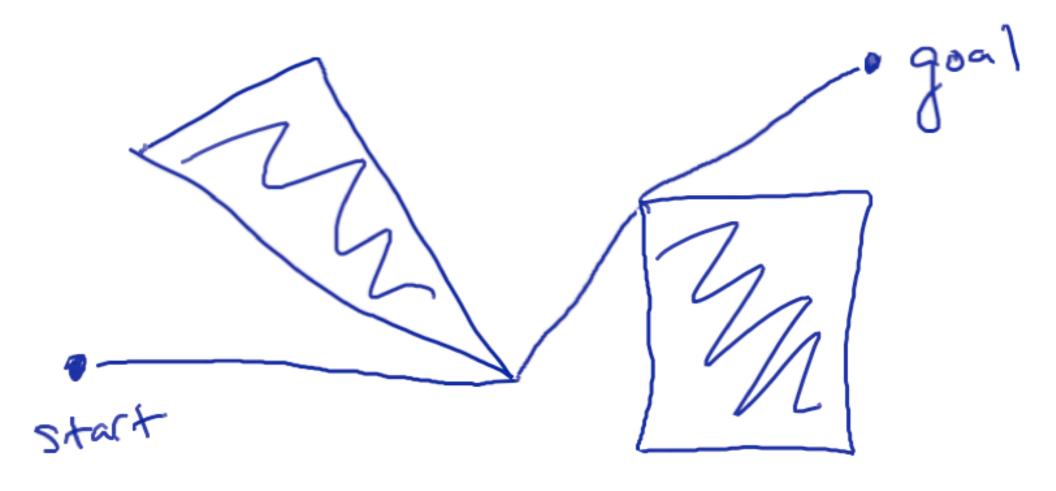
Back to planning

- Complaint with A* was that it didn't break up
 C-space intelligently
- Our How might we do better?
- Lots of roboticists have given lots of answers!

Shortest path in C-space



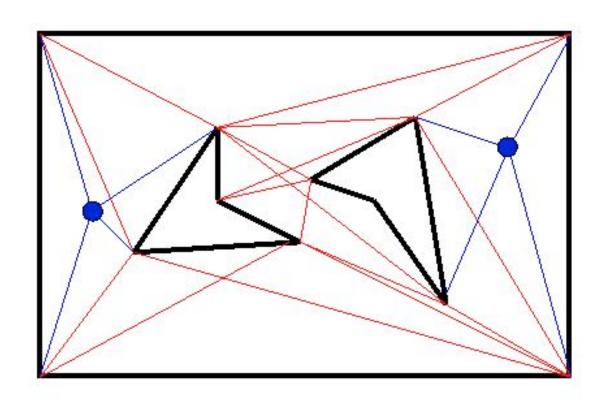
Shortest path in C-space



Shortest path

- Suppose a planar polygonal C-space
- Shortest path in C-space is a sequence of line segments
- Each segment's ends are either start or goal or one of the vertices in C-space
- In 3-d or higher, might lie on edge, face, hyperface, ...

Visibility graph



http://www.cse.psu.edu/~rsharma/robotics/notes/notes2.html

Naive algorithm

```
For i = I ... points
  For j = I ... points
  included = t
  For k = I ... edges
  if segment ij intersects edge k
  included = f
```

Complexity

- Naive algorithm is O(n³) in planar C-space
- For faster algorithms, O(n²) or O(k+n log(n)),
 see [Latombe, p. 157]
 - k = number of edges that wind up in visibility graph
- In dimension d, graph gets much bigger, more complex; speedup tricks stop working
- Once we have graph, search it!

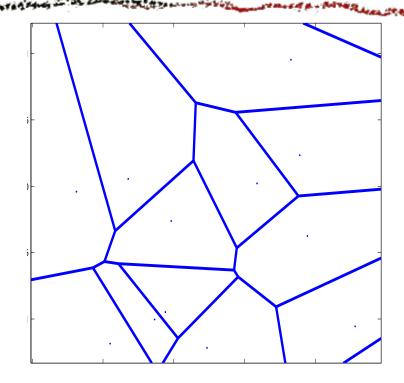
Discussion of visibility graph

- Good: finds shortest path
- Bad: complex C-space yields long runtime, even if problem is easy
 - get my 23-dof manipulator to move Imm when nearest obstacle is Im
- Bad: no margin for error

Getting bigger margins

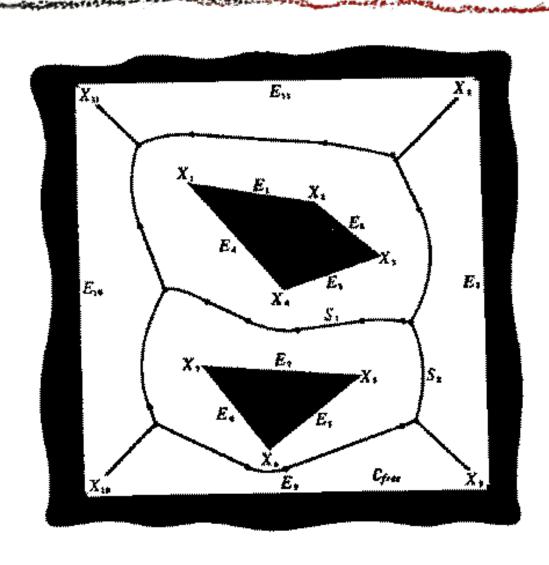
- Could just pad obstacles
 - but how much is enough? might make infeasible...
- What if we try to stay as far away from obstacles as possible?

Voronoi graph



- Set of all places equidistant from two or more obstacles: Voronoi graph
 - point obstacles: network of line segments
 - nonzero extent: graph may include curves

Voronoi w/ polygonal C-space

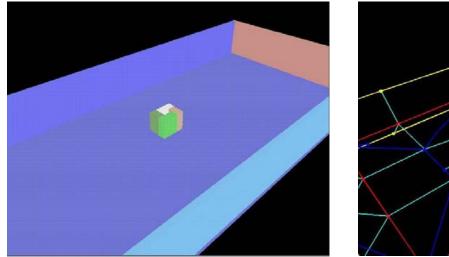


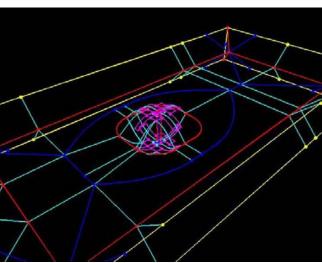
Voronoi method for planning

- Compute Voronoi diagram of C-space
- Go straight from start to nearest point on diagram
- Plan within diagram to get near goal (A*)
- Go straight to goal

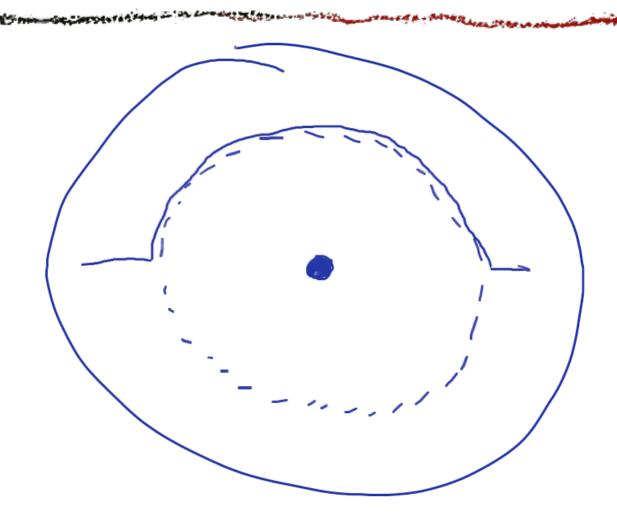
Voronoi discussion

- Good: stays far away from obstacles
- Bad: assumes polygons
- Bad: gets kind of hard in higher dimensions (but see Howie Choset's web page and book)



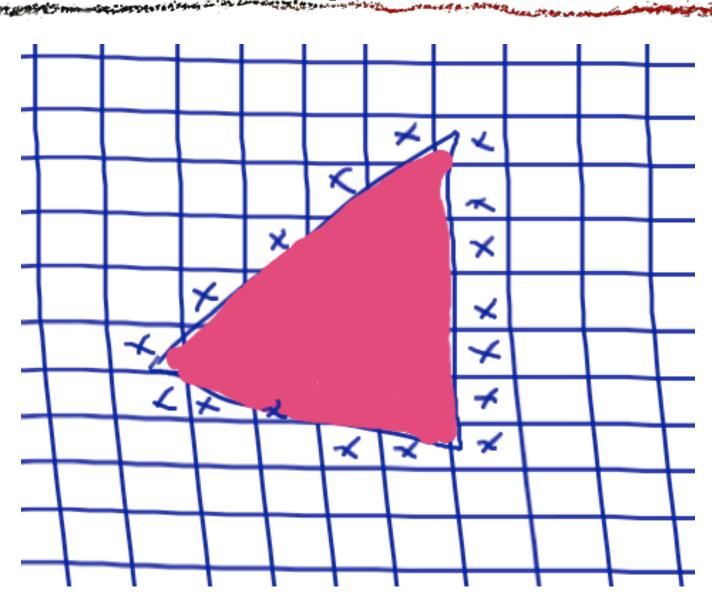


Voronoi discussion



Bad: kind of gun-shy about obstacles

(Approximate) cell decompositions

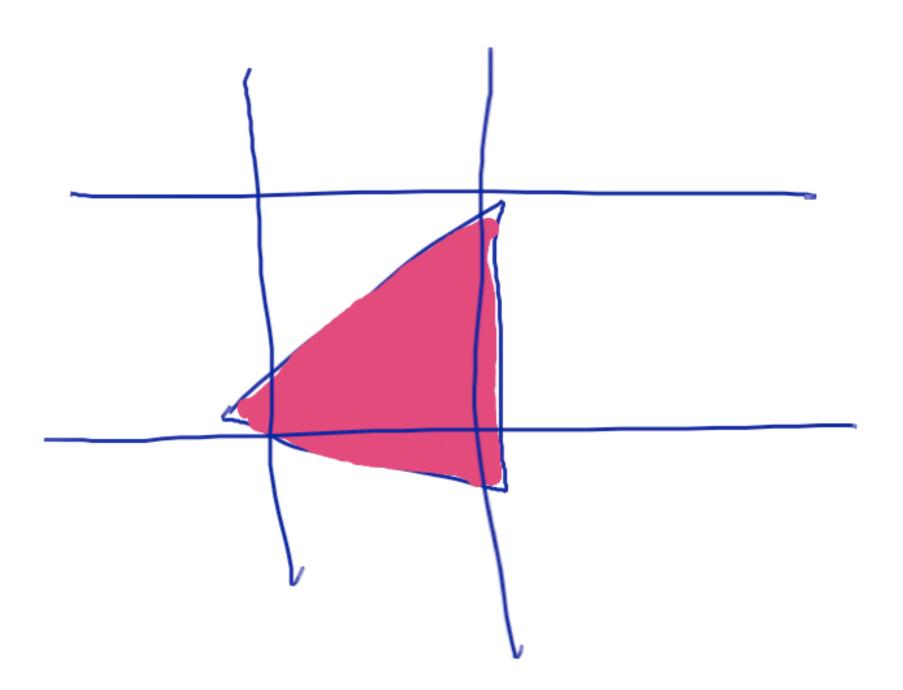


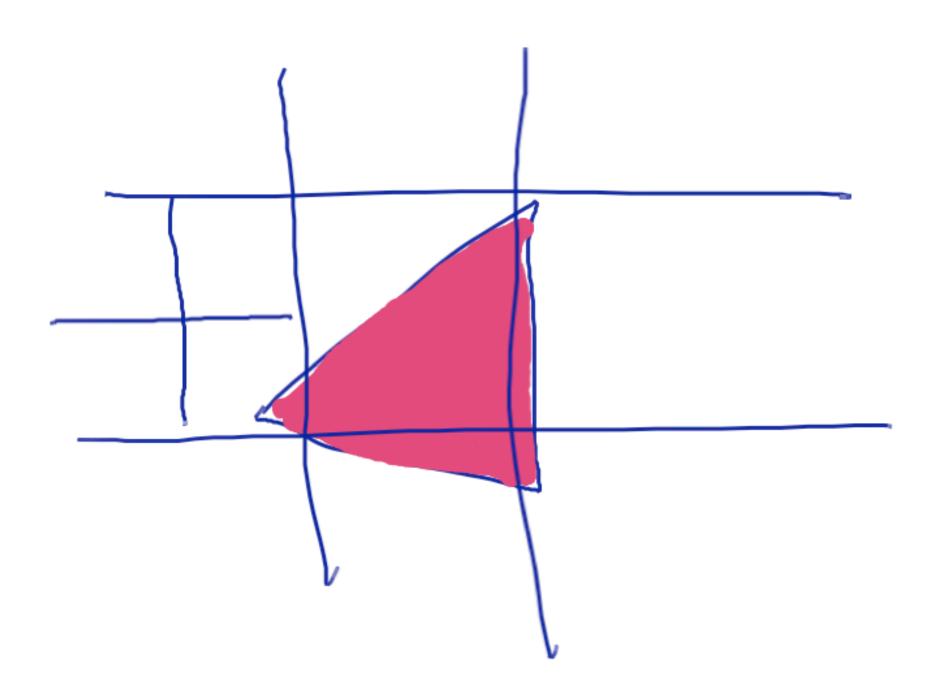
Planning algorithm

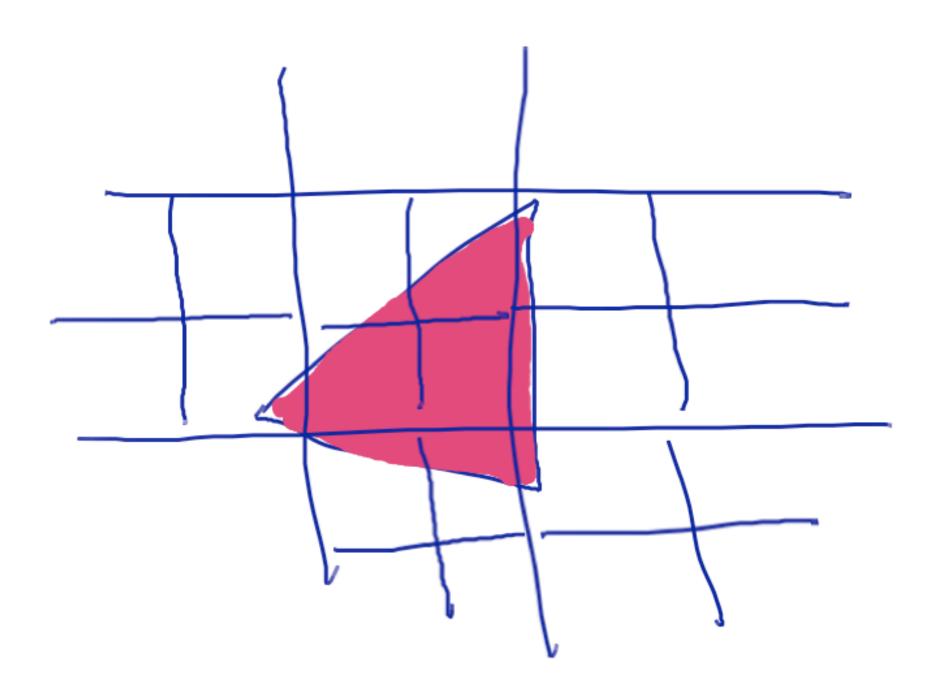
- Lay down a grid in C-space
- Delete cells that intersect obstacles
- Connect neighbors
- **A***
- If no path, double resolution and try again
 - never know when we're done
 - resolution: want high near obstacles, low everywhere else

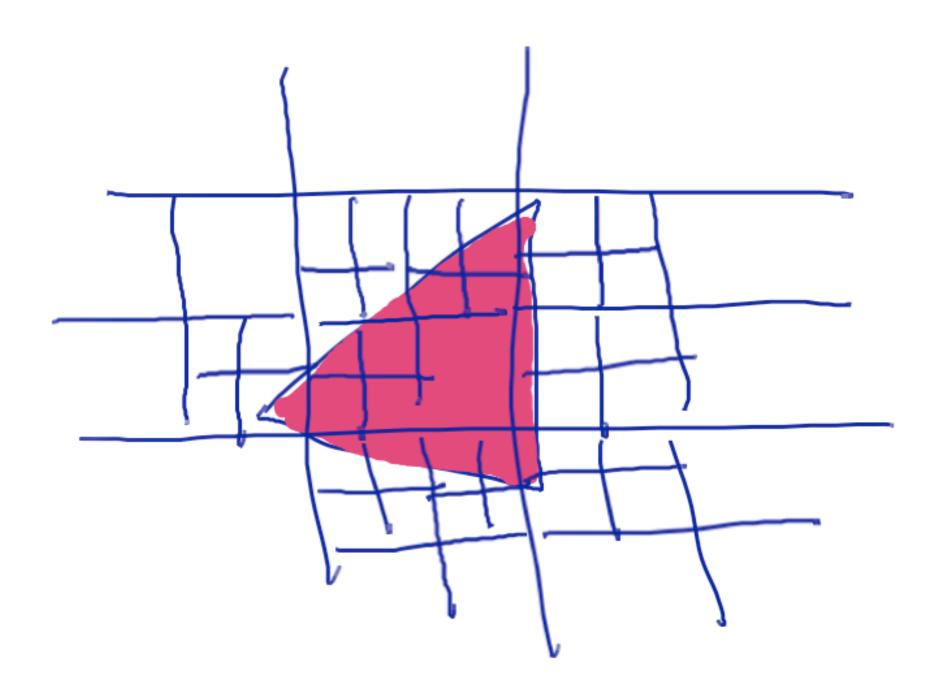
Fix: variable resolution

- Lay down a coarse grid
- Split cells that intersect obstacle borders
 - empty cells good
 - full cells also don't need splitting
- Stop at fine resolution
- Data structure: quadtree









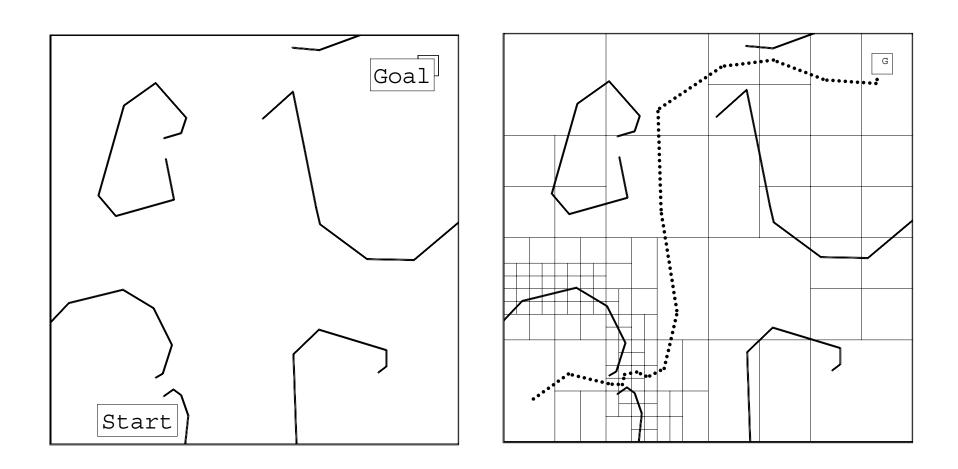
Discussion

- Works pretty well, except:
 - Still don't know when to stop
 - Won't find shortest path
 - Still doesn't really scale to high-d

Better yet

- Adaptive decomposition
- Split only cells that actually make a difference
 - are on path from start
 - make a difference to our policy

An adaptive splitter: parti-game

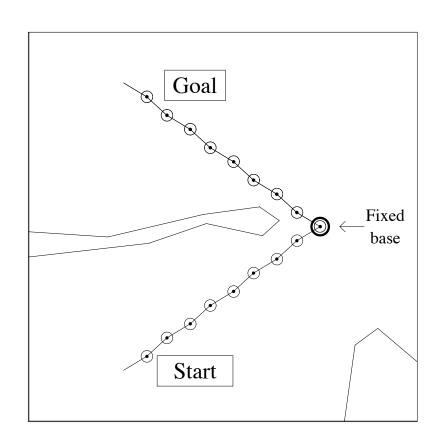


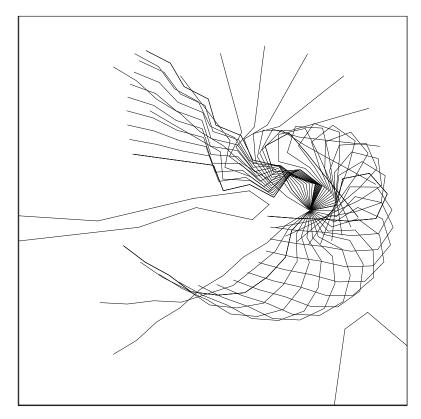
Andrew Moore and Chris Atkeson. The Parti-game Algorithm for Variable Resolution Reinforcement Learning in Multidimensional State-spaces. http://www.autonlab.org/autonweb/14699.html

Parti-game algorithm

- Sample actions from several points per cell
- Try to plan a path from start to goal
- On the way, pretend an opponent gets to choose which outcome happens (out of all that have been observed in this cell)
- If we can get to goal, we win
- Otherwise we can split a cell

9dof planar arm





85 partitions total

Randomness in search

Rapidly-exploring Random Trees

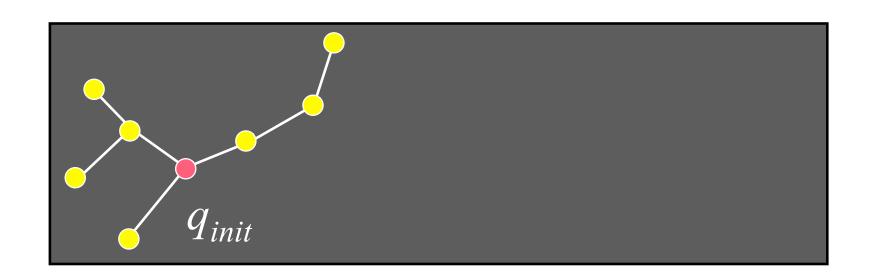
- Break up C-space into Voronoi regions around random landmarks
- Invariant: landmarks always form a tree
 - known path to root
- Subject to this requirement, placed in a way that tends to split large Voronoi regions
 - coarse-to-fine search
- Goal: feasibility not optimality (*)

RRT: required subroutines

- RANDOM_CONFIG
 - samples from C-space
- EXTEND(q, q')
 - local controller, heads toward q' from q
 - stops before hitting obstacle (and perhaps also after bound on time or distance)
- FIND_NEAREST(q, Q)
 - searches current tree Q for point near q

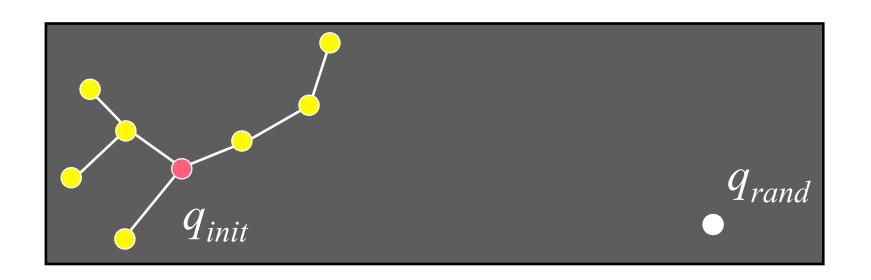
```
BUILT\_RRT(q_{init}) \{ \\ T = q_{init} \\ for k = 1 to K \{ \\ q_{rand} = RANDOM\_CONFIG() \\ EXTEND(T, q_{rand}); \\ \}
```

```
EXTEND(T, q) {
    qnear = FIND_NEAREST(q, T)
    qnew = EXTEND(qnear, q)
    T = T + (qnear, qnew)
}
```



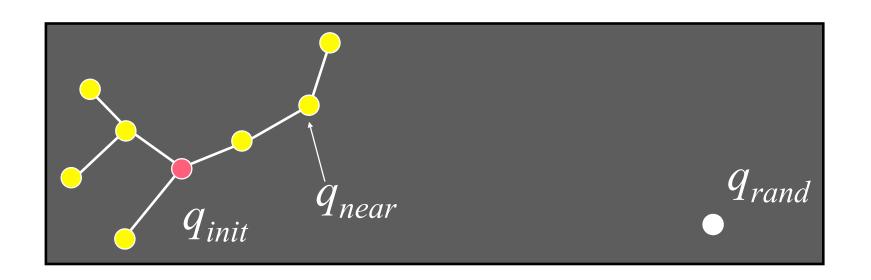
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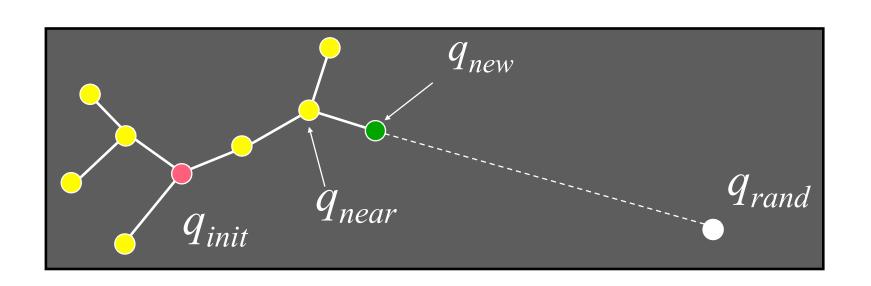
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```
BUILT\_RRT(q_{init}) \{ \\ T = q_{init} \\ for k = 1 to K \{ \\ q_{rand} = RANDOM\_CONFIG() \\ EXTEND(T, q_{rand}); \\ \}
```

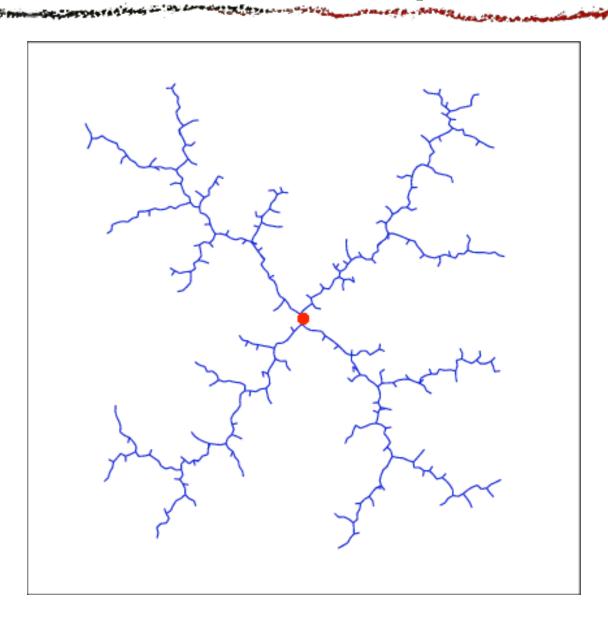
```
EXTEND(T, q) {
    qnear = FIND_NEAREST(q, T)
    qnew = EXTEND(qnear, q)
    T = T + (qnear, qnew)
}
```



```
BUILT\_RRT(q_{init}) \{ \\ T = q_{init} \\ for k = 1 to K \{ \\ q_{rand} = RANDOM\_CONFIG() \\ EXTEND(T, q_{rand}); \\ \}
```

```
EXTEND(T, q) {
    qnear = FIND_NEAREST(q, T)
    qnew = EXTEND(qnear, q)
    T = T + (qnear, qnew)
}
```

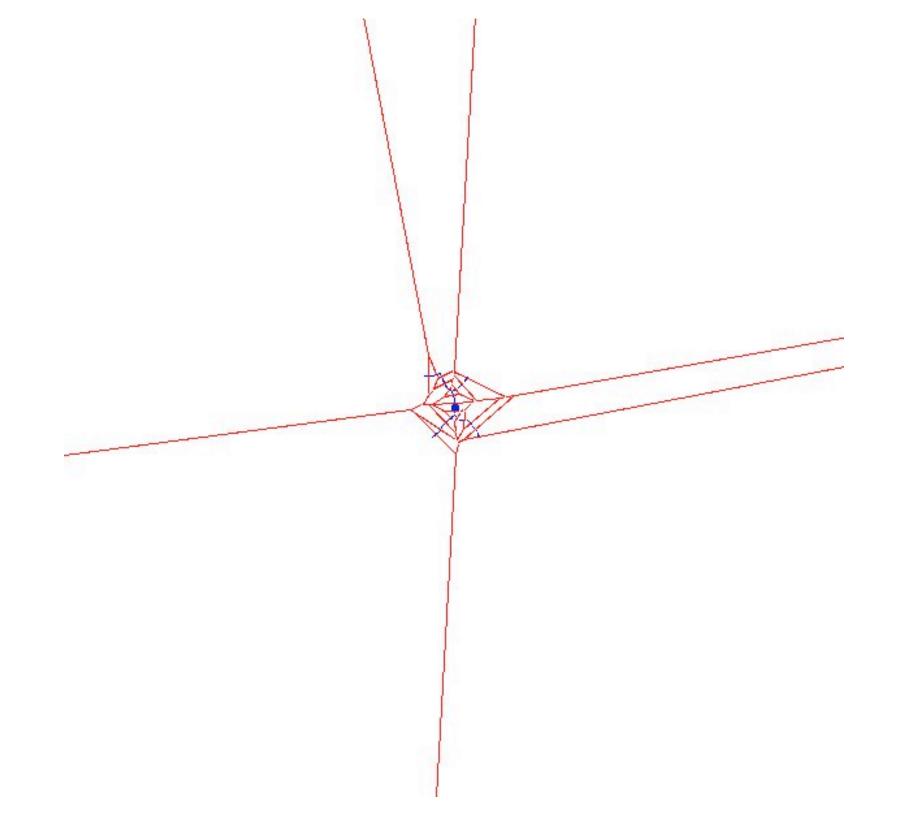
RRT example

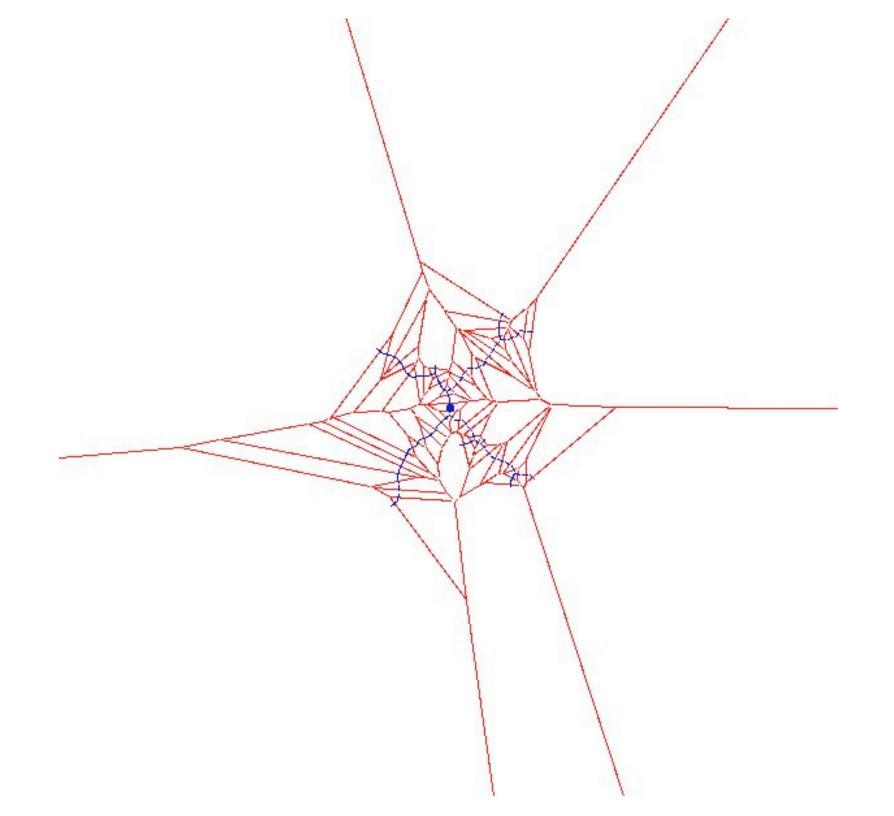


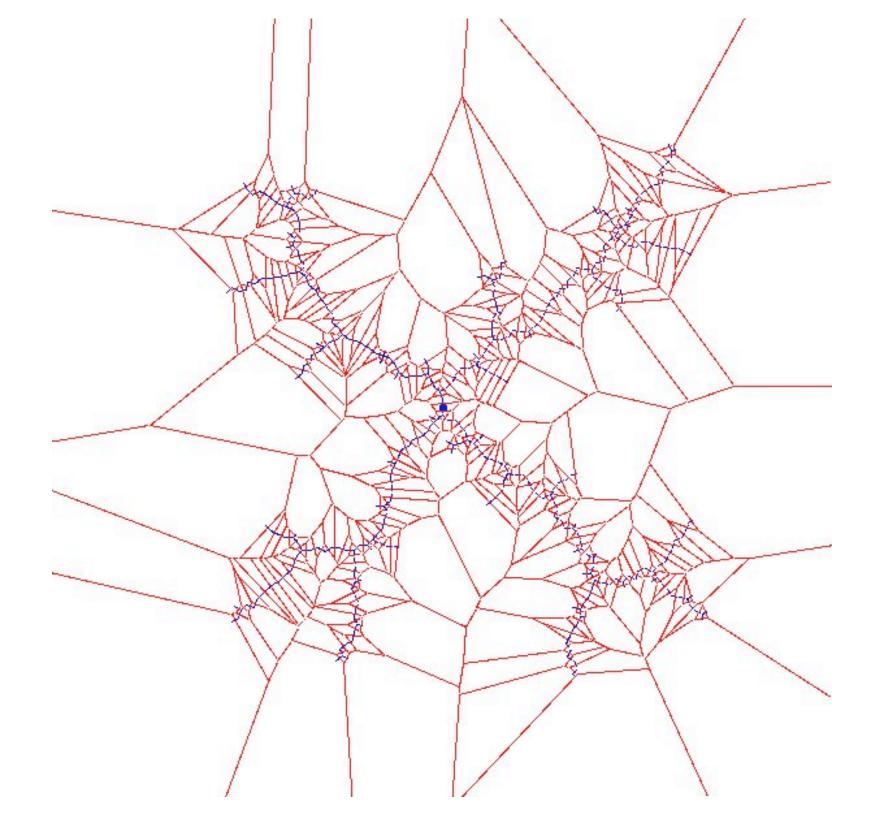
Planar holonomic robot

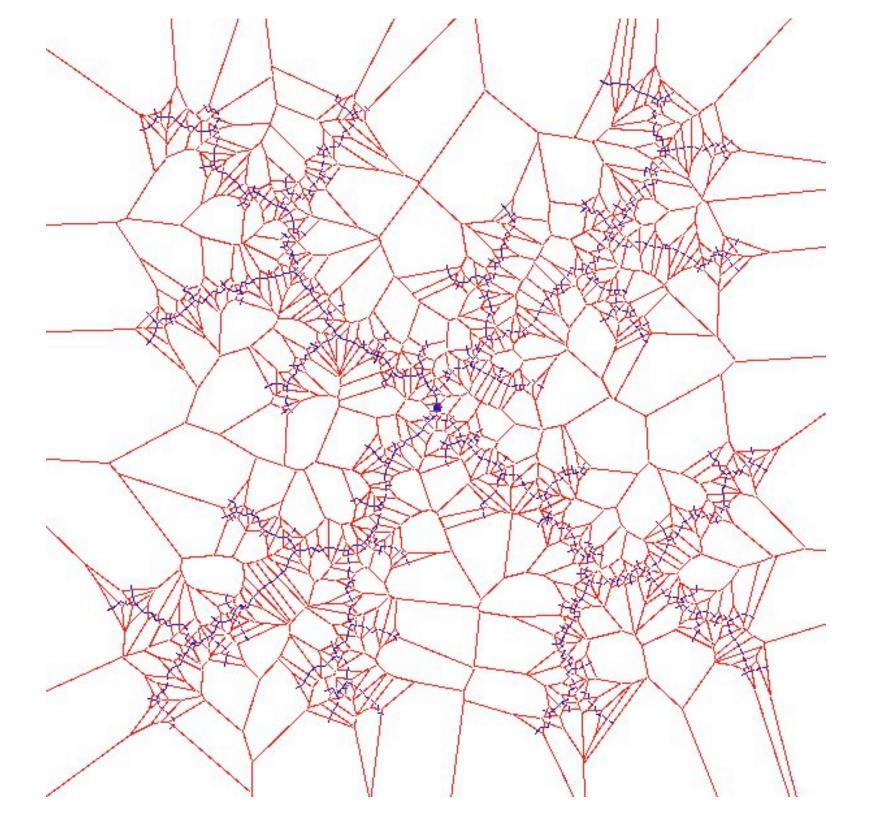
RRTs explore coarse to fine

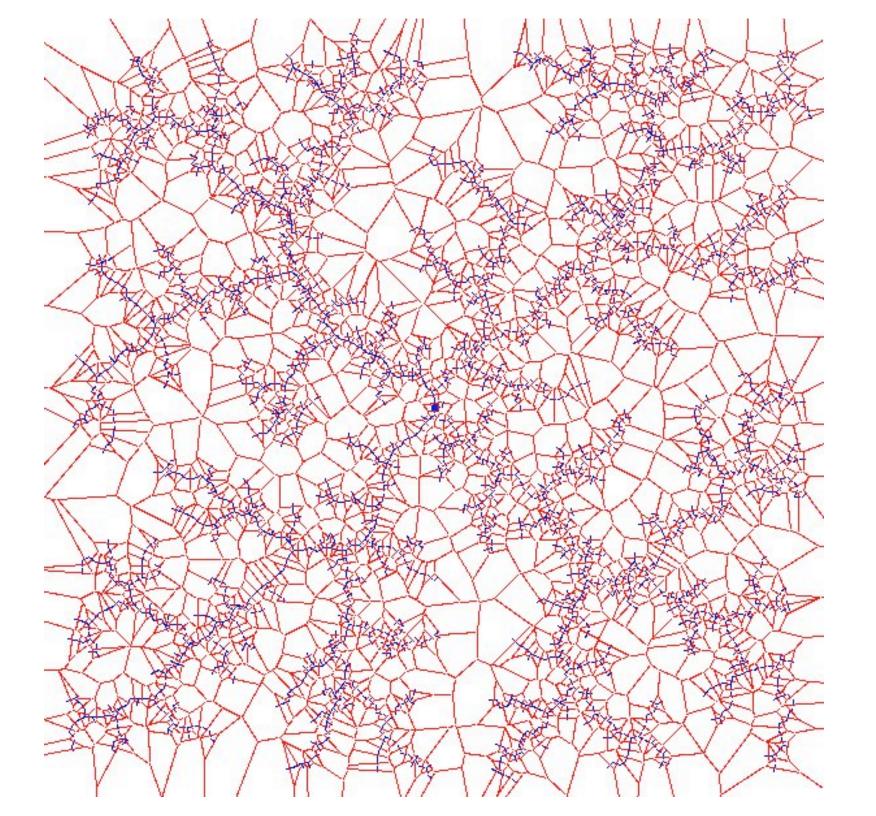
- Tend to break up large Voronoi regions
 - higher probability of q_{rand} being in them
- Limiting distribution of vertices given by RANDOM_CONFIG
 - as RRT grows, probability that q_{rand} is reachable with local controller (and so immediately becomes a new vertex) approaches I



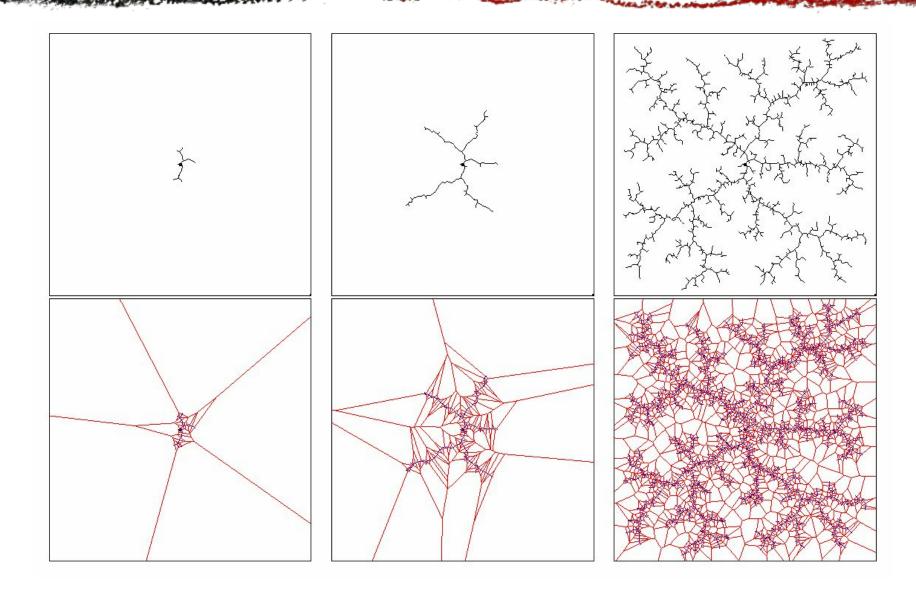




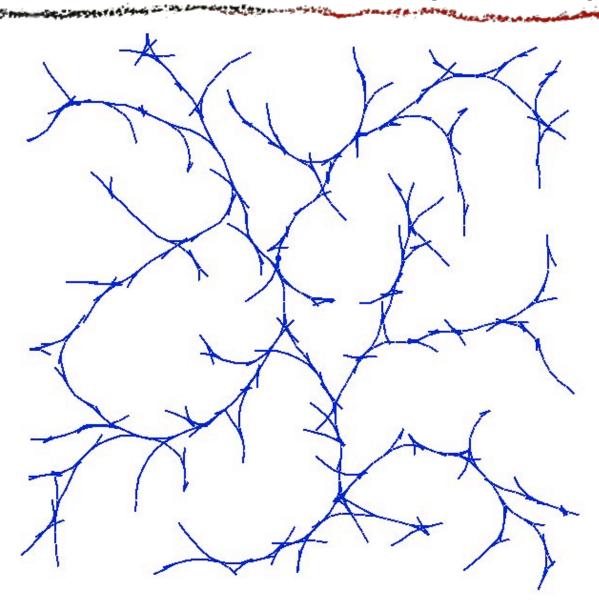




RRT example



RRT for a car (3 dof)



Planning with RRTs

- Build RRT from start until we add a node that can reach goal using local controller
- \circ (Unique) path: root \rightarrow last node \rightarrow goal
- Optional: "rewire" tree during growth by testing connectivity to more than just closest node
- Optional: grow forward and backward

