15-780: Graduate AI Lecture 2. Spatial Search

Geoff Gordon (this lecture) Ziv Bar-Joseph TAs Michael Benisch, Yang Gu

Admin

- WEH 5409, Sep 18, 4:30-5:30pm: matlab tutorial
- Please send your email address to TA
 Michael Benisch (mbenisch at cs), who is
 compiling a class email list
- Please check the website regularly for readings (for Lec. 1–2, Ch. 1–4 of RN)

Last episode, on Grad AI

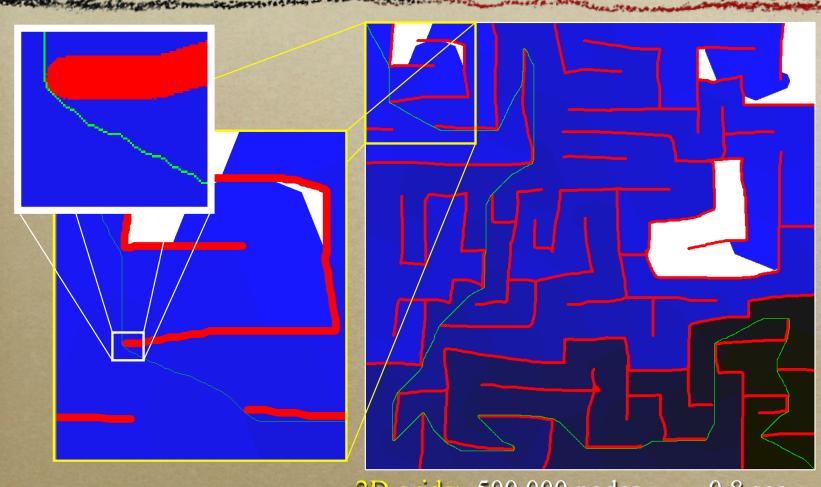
Topics covered

- What is AI? (Be able to discuss an example or two)
- Types of uncertainty & corresponding approaches
- How to set up state space graph for problems like the robotic grad student or path planning

Topics covered

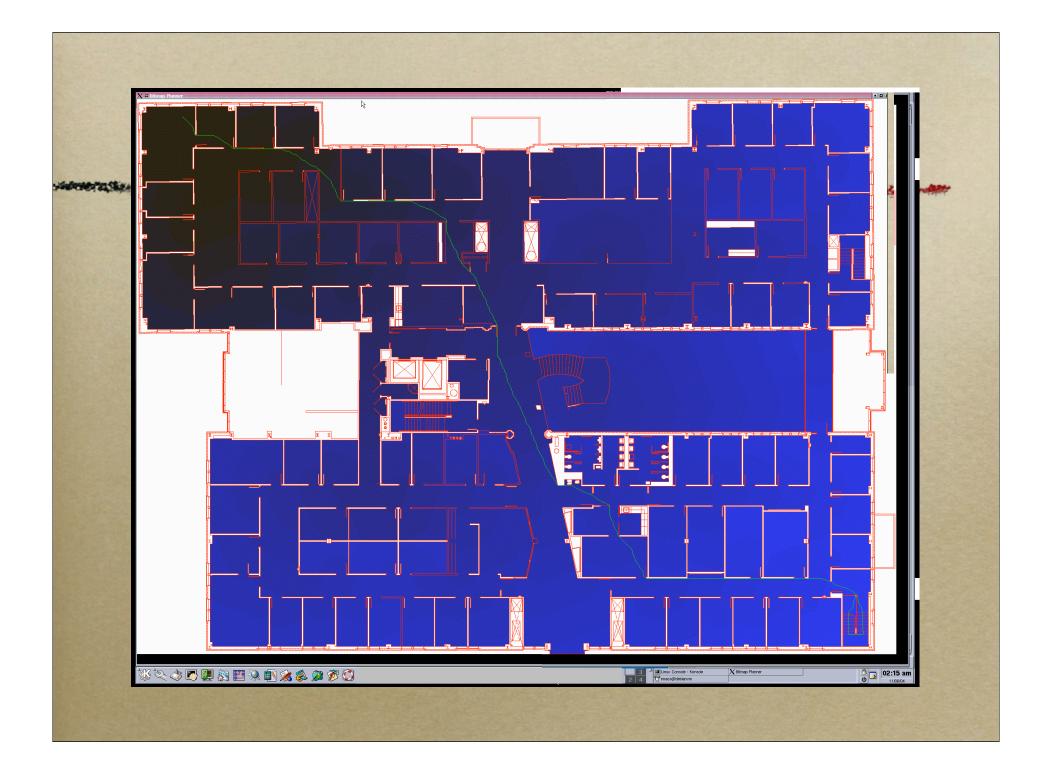
- Generic search algorithm & data structures
- Search methods: be able to simulate
 - BFS, DFS, DFID
 - Heuristic search
 - A*: define admissibility; show optimality, efficiency
- What are advantages of each?

A* Planning on Big Grids



Credit: Kuffner

2D grids: $500,000 \text{ nodes} = \sim 0.8 \text{ sec}$ $10 \text{ million nodes} = \sim 12 \text{ sec}$



Projects

Project ideas



 Plan a path for this robot so that it gets a good view of an object as fast as possible

Project ideas



 Implement a distributed market-based planner and test the contribution of learning to overall performance

Project ideas

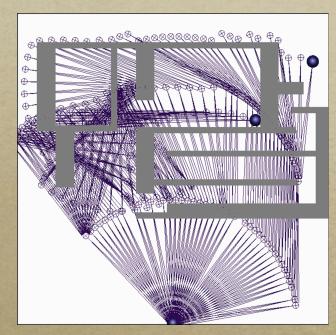




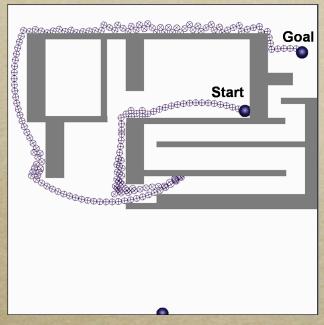
- Give me an excuse to buy the new Lego Mindstorms set
 - plan footstep placements
 - plan how to grip objects

Spatial Planning

Plans in Space...



Optimal Solution



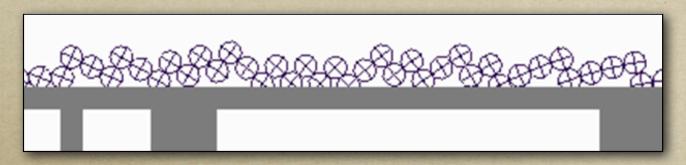
End-effector Trajectory

• Last time, we saw A* for spatial planning

What's wrong w/ A* guarantees?

- (optimality) A* finds a solution of depth g*
- (efficiency) A^* expands no nodes that have $f(node) > g^*$

What's wrong with A*?

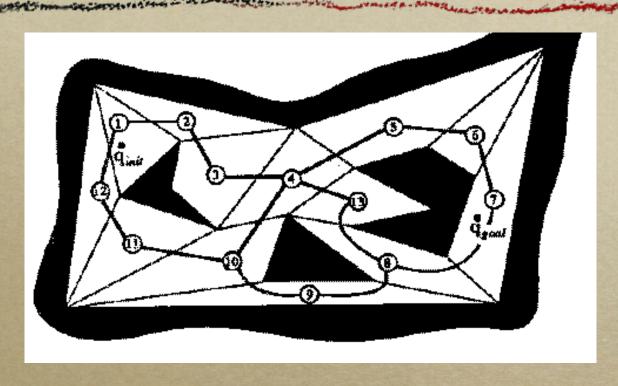


- Discretized space into tiny little chunks
 - o a few degrees rotation of a joint
 - \circ Lots of states \Rightarrow slow
- Discretized actions too
 - o only allowed to move one joint at a time
- Results in jagged paths

What's wrong with A*?



Wouldn't it be nice...

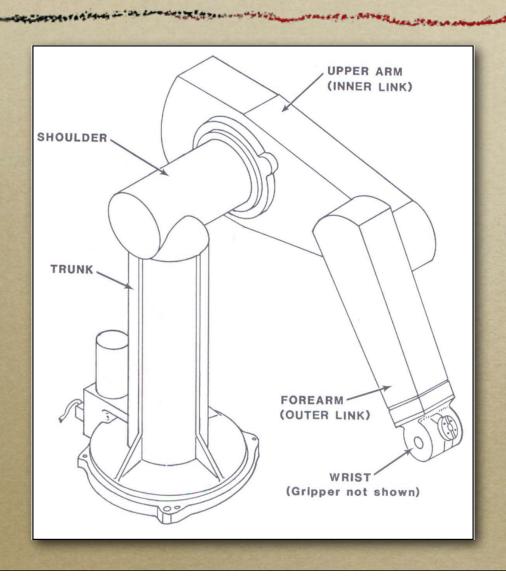


- ... if we could break things up based more on the real geometry of the world?
- Robot Motion Planning by Jean-Claude Latombe

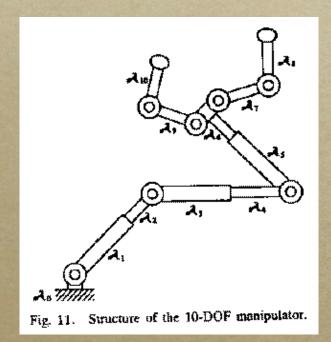
Physical system

- A moderate number of real-valued coordinates
- o Deterministic, continuous dynamics
- Continuous goal set (or a few pieces)
- Cost = time, work, torque, ...

Typical physical system



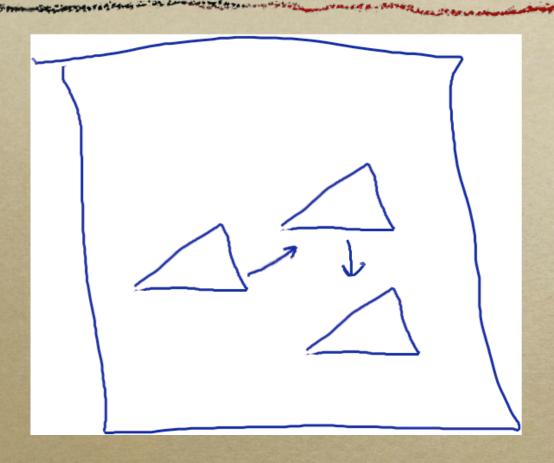
A kinematic chain



- Rigid links connected by joints
 - revolute or prismatic(1 dof each)
- Configuration

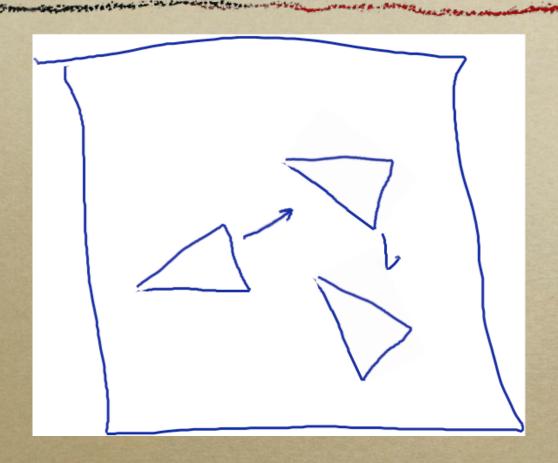
$$\mathbf{q} = (q_1, q_2, ...)$$

Mobile robots



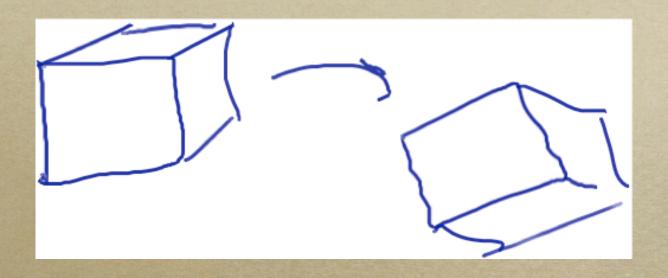
• Translating in space = 2 dof

More mobility

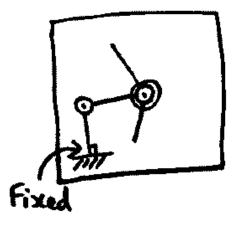


 \circ Translation + rotation = 3 dof

Q: How many dofs?

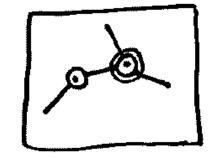


• 3d translation & rotation



How many dofs?

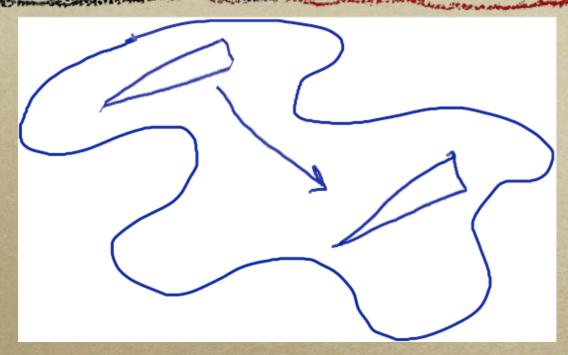
Free flying How many dofs?



The configuration of has one real valued entry per DOF.

Salvesta Salva destruis tori de

Robot kinematic motion planning



- Given a robot (coordinates q)
- ... and a workspace with obstacles
- o ... get from a start to a goal

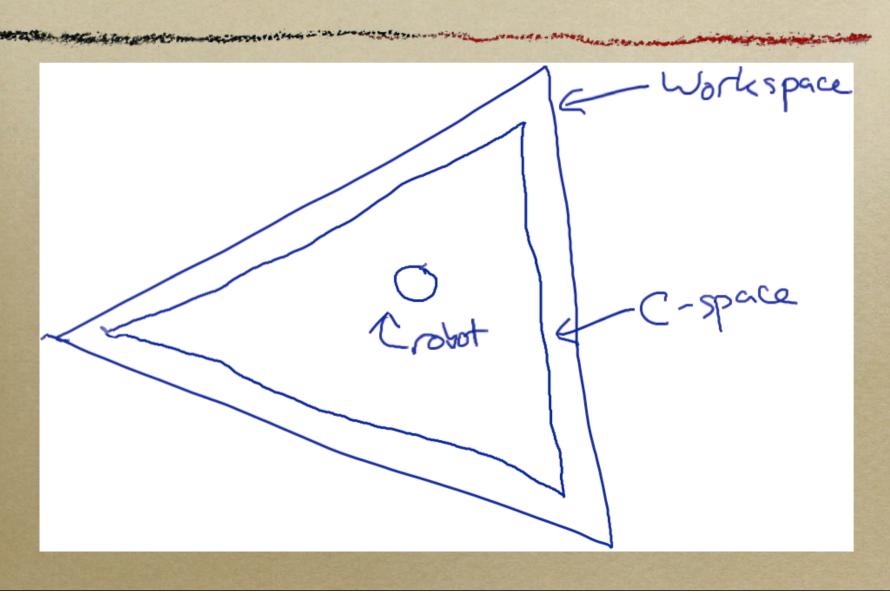
Kinematic planning

- For any configuration q, can test whether it intersects obstacles
- Set of legal configs is "configuration space" C (a subset of \Re^{dofs})
- Path is a continuous function q from [0,1] into C with $q(0) = \mathbf{q_s}$ and $q(1) = \mathbf{q_g}$

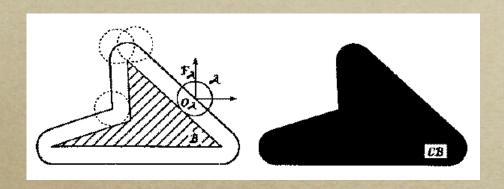
Note: dynamic planning

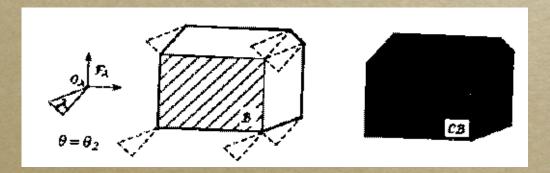
- o Includes inertia as well as configuration
- 。 q, **q**
- Harder, since twice as many dofs
- More later...

C-space example



More C-space examples





Another C-space example

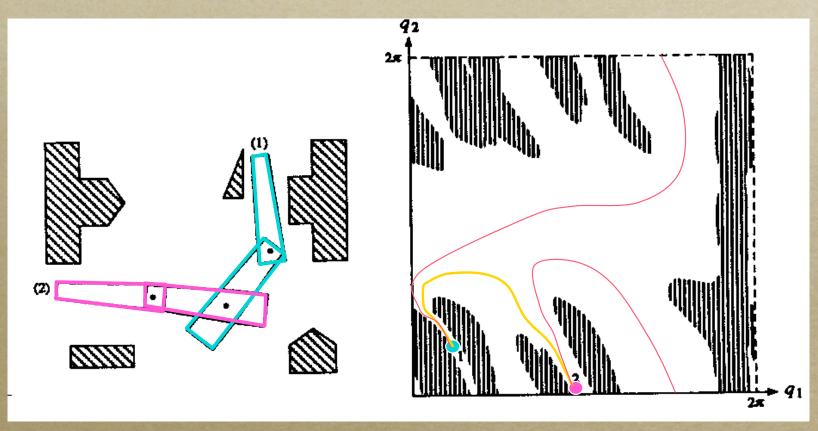
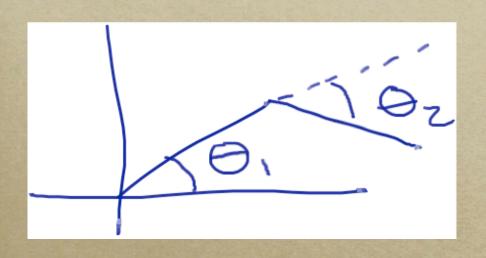


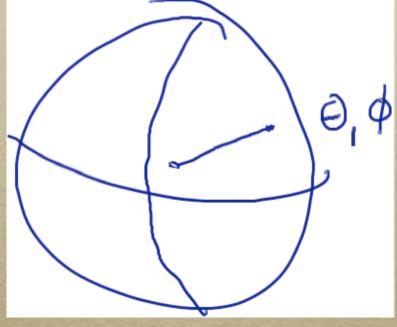
image: J Kuffner

Topology of C-space

- Topology of C-space can be something other than the familiar Euclidean world
- E.g. set of angles = unit circle = SO(2)
 not [0, 2π)!
- Ball & socket joint (3d angle) \subseteq unit sphere = SO(3)

Topology example



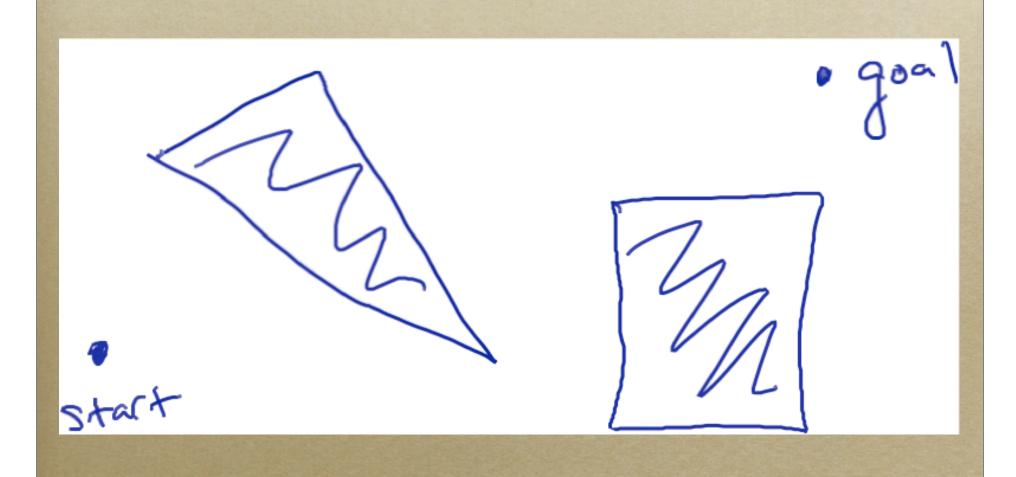


 Compare L to R: 2 planar angles v. one solid angle — both 2 dof (and neither the same as Euclidean 2-space)

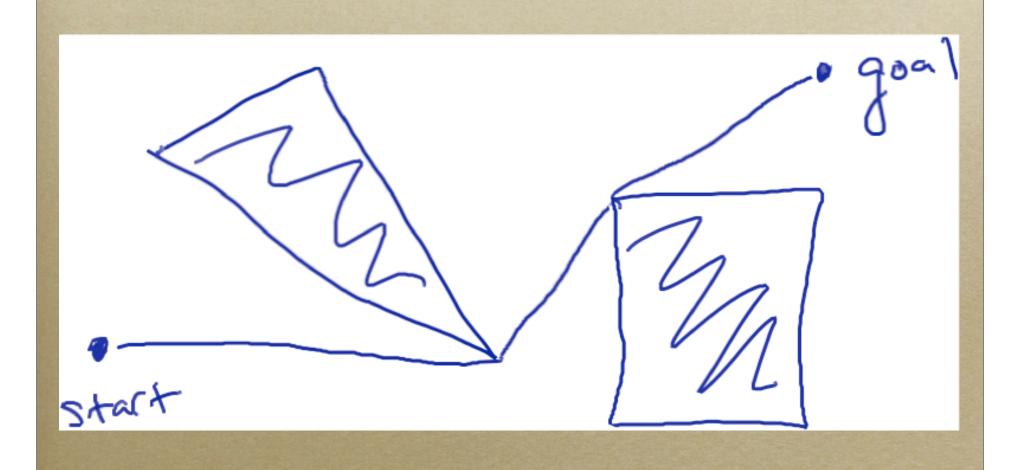
Back to planning

- Complaint with A* was that it didn't break up space intelligently
- How might we do better?
- Lots of roboticists have given lots of answers!

Shortest path in C-space



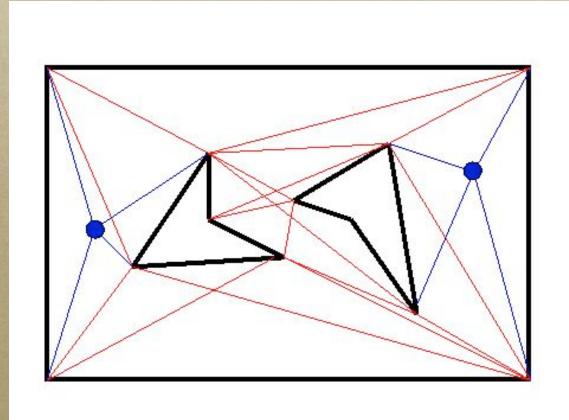
Shortest path in C-space



Shortest path

- Suppose a polygonal C-space
- Shortest path in C-space is a sequence of line segments
- Each segment's ends are either start or goal or one of the vertices in C-space
- In 3-d or higher, might lie on edge, face, hyperface, ...

Visibility graph



http://www.cse.psu.edu/~rsharma/robotics/notes/notes2.html

Naive algorithm

```
For i = 1 \dots points

For j = 1 \dots points

included = t

For k = 1 \dots faces

if segment ij intersects face k

included = f
```

Complexity

- Naive algorithm is O(n³) in planar Cspace (grows fast with d!)
- For algorithms that run faster, $O(n^2)$ and $O(k + n \log n)$, see [Latombe, pg 157]
 - k = number of edges that wind up in visibility graph
- Once we have graph, search it!

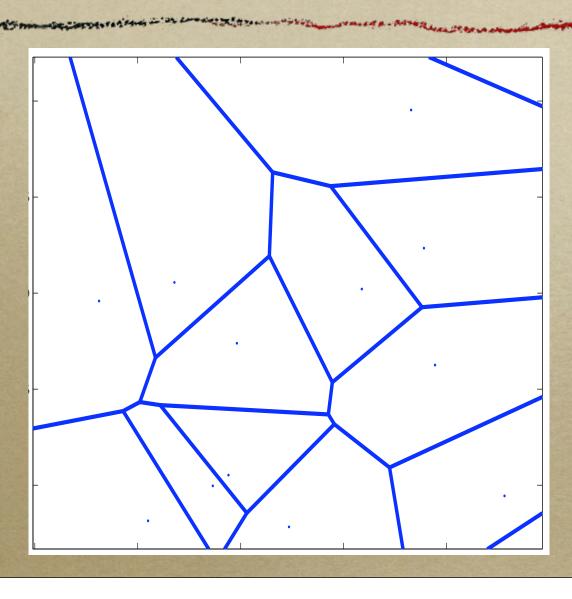
Discussion of visibility graph

- Good: finds shortest path
- Bad: complex C-space yields long runtime, even if problem is easy
 - get my 23-dof manipulator to move
 1mm when nearest obstacle is 1m
- Bad: no margin for error

Getting bigger margins

- Could just pad obstacles
 - but how much is enough? might make infeasible...
- What if we try to stay as far away from obstacles as possible?

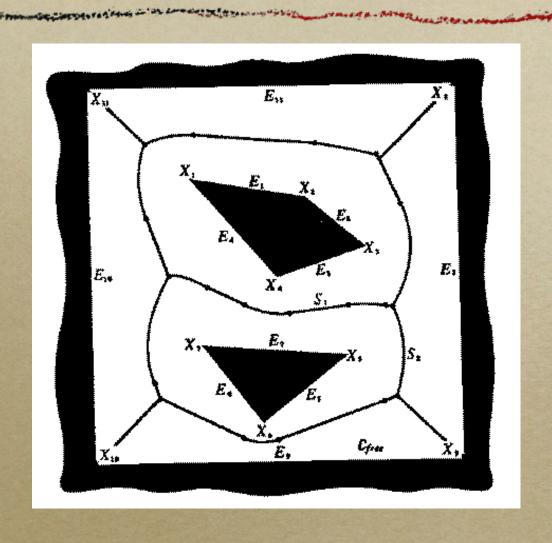
Voronoi



Voronoi

- o Given a set of point obstacles
- Find all places that are equidistant from two or more of them
- Result: network of line segments
- Called Voronoi graph
- Each line stays as far away as possible from two obstacles while still going between them

Voronoi from polygonal C-space



Voronoi from polygonal C-space

- Set of points which are equidistant from 2
 or more closest points on border of C space
- Polygonal C-space in 2d yields lines & parabolas intersecting at points
 - o lines from 2 points
 - o parabolas from line & point

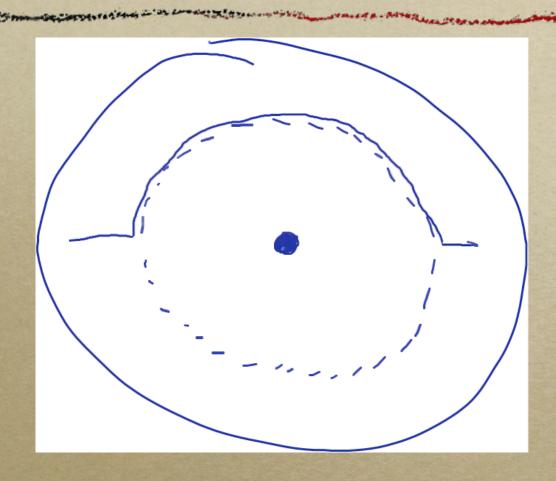
Voronoi method for planning

- o Compute Voronoi diagram of C-space
- Go straight from start to nearest point on diagram
- Plan within diagram to get near goal (guess which algorithm)
- Go straight to goal

Discussion of Voronoi

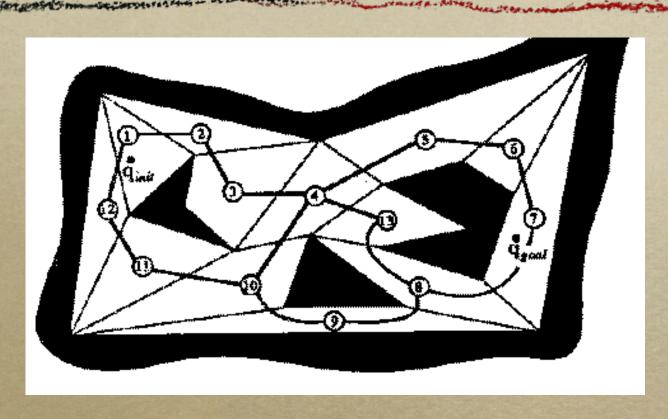
- Good: stays far away from obstacles
- Bad: assumes polygons
- Bad: assumes 2d, gets kind of hard in higher dimensions (but see http:// voronoi.sbp.ri.cmu.edu/~motion/)

Voronoi discussion



• Bad: kind of gun-shy about obstacles

Exact cell decompositions

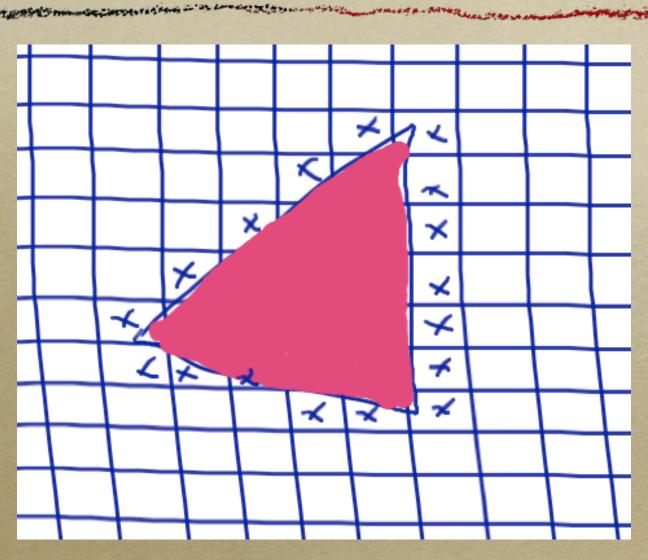


 We can try to break C-space into a bunch of convex polygons

Exact cell decompositions

- Will not discuss how to do
- Common approach for video game NPCs
- o But is also hard in higher than 2d
- And can result in wobbly paths

Approximate cell decompositions



Planning algorithm

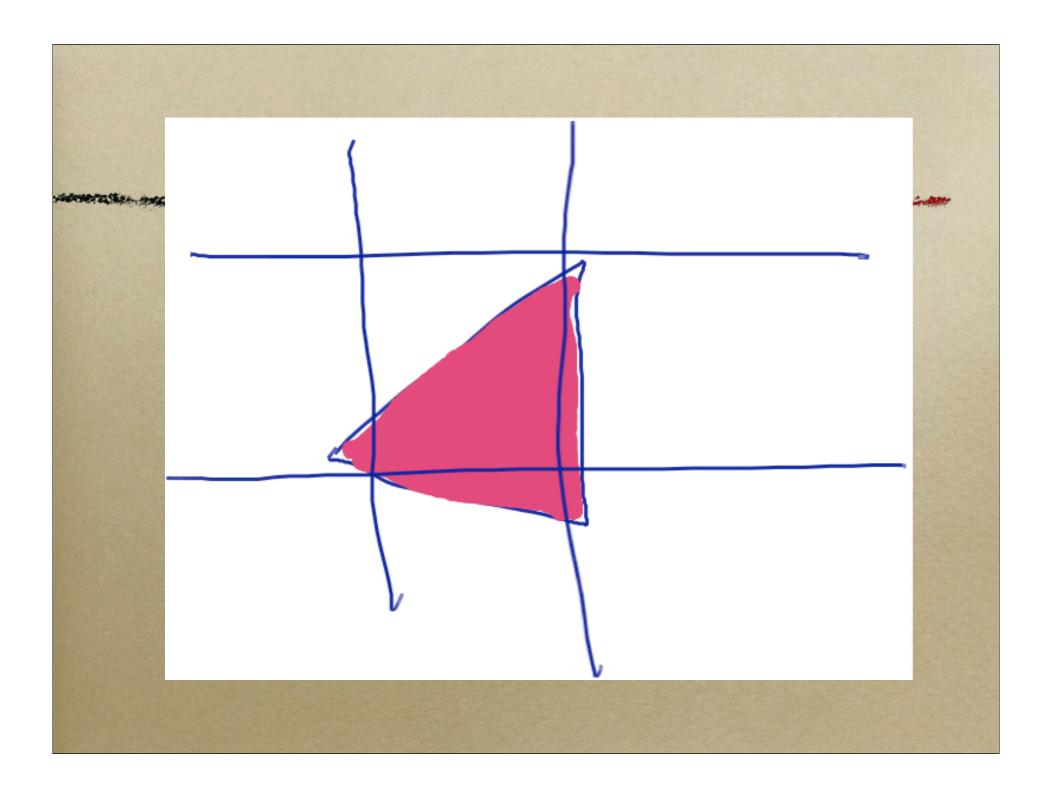
- Lay down a grid in C-space
- Delete cells that intersect obstacles
- Connect neighbors
- A* (surprise!)
- o If no path, double resolution and try again
 - o never know when we're done

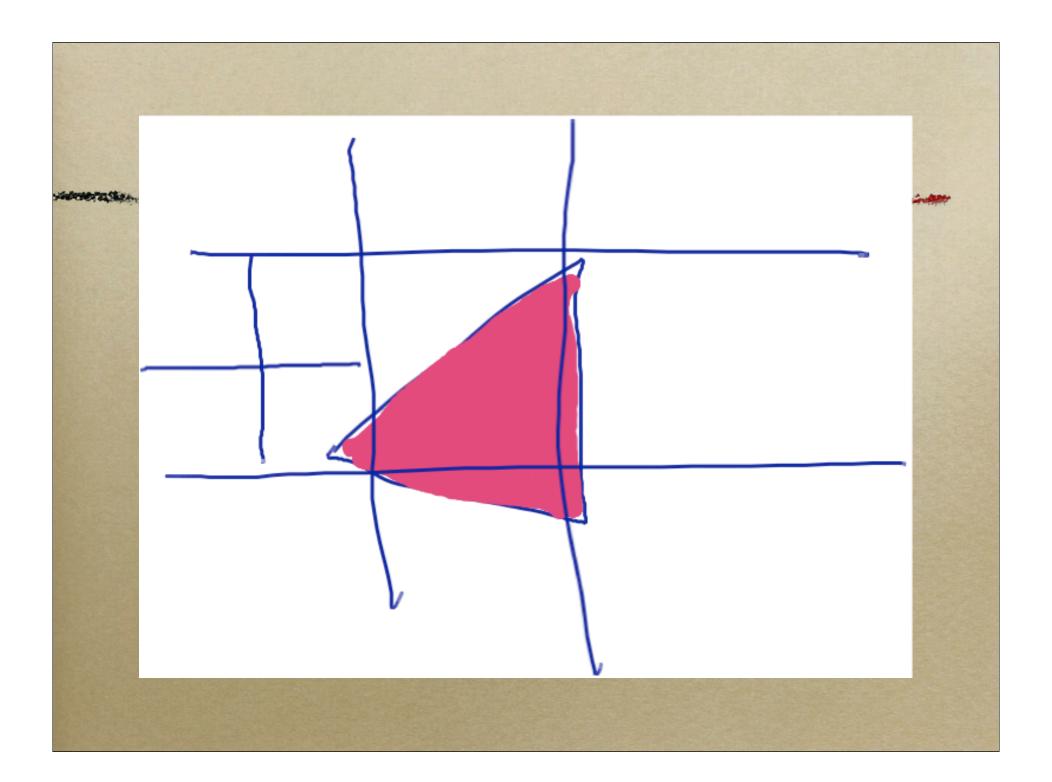
Approximate cell decomposition

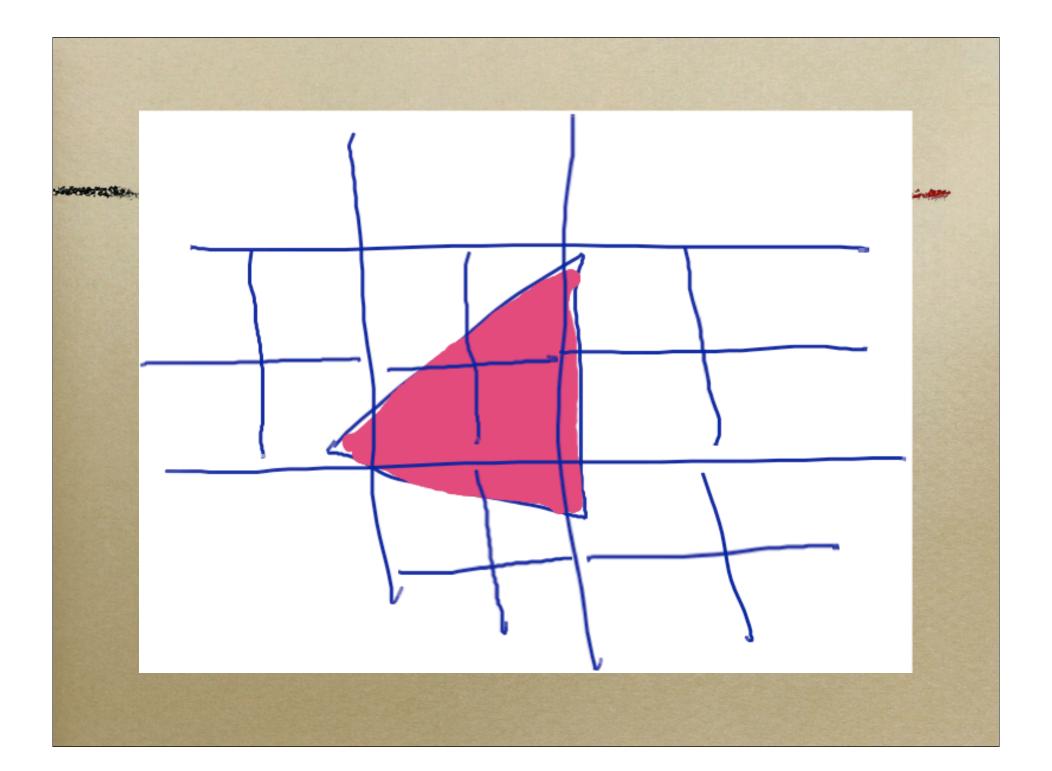
- This decomposition is what we were using for A* in examples from last class
- Works pretty well except:
 - o need high resolution near obstacles
 - want low res away from obstacles

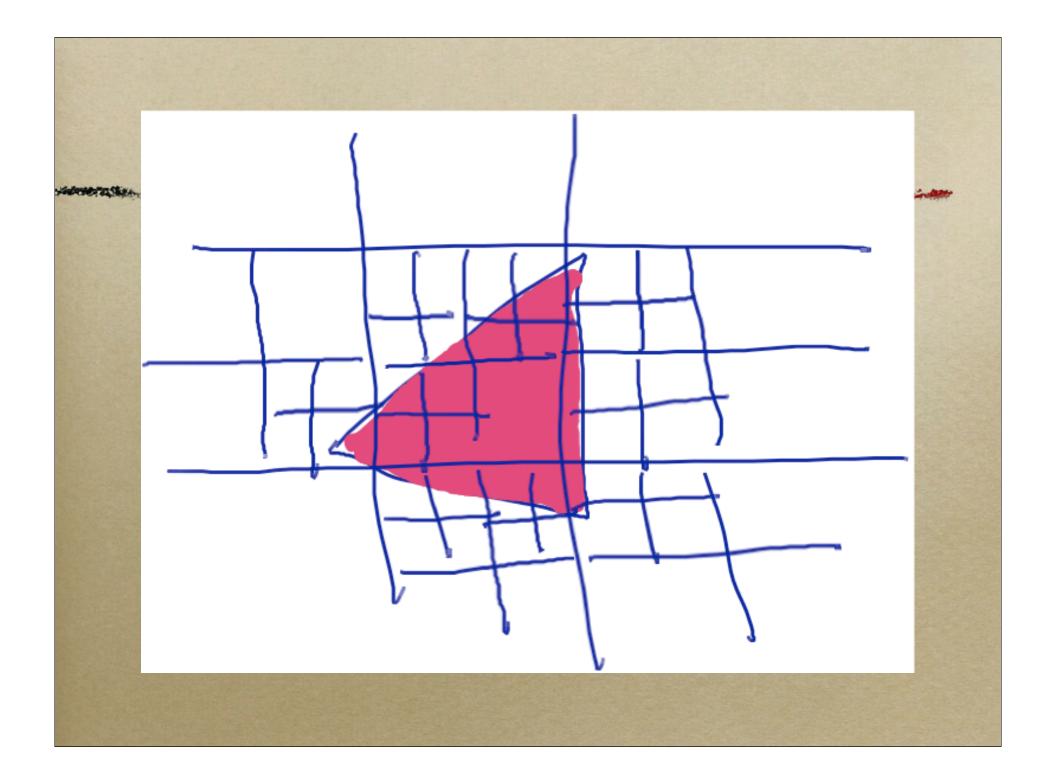
Fix: variable resolution

- Lay down a coarse grid
- Split cells that intersect obstacle borders
 - empty cells good
 - o full cells also don't need splitting
- Stop at fine resolution
- o Data structure: quadtree









Discussion

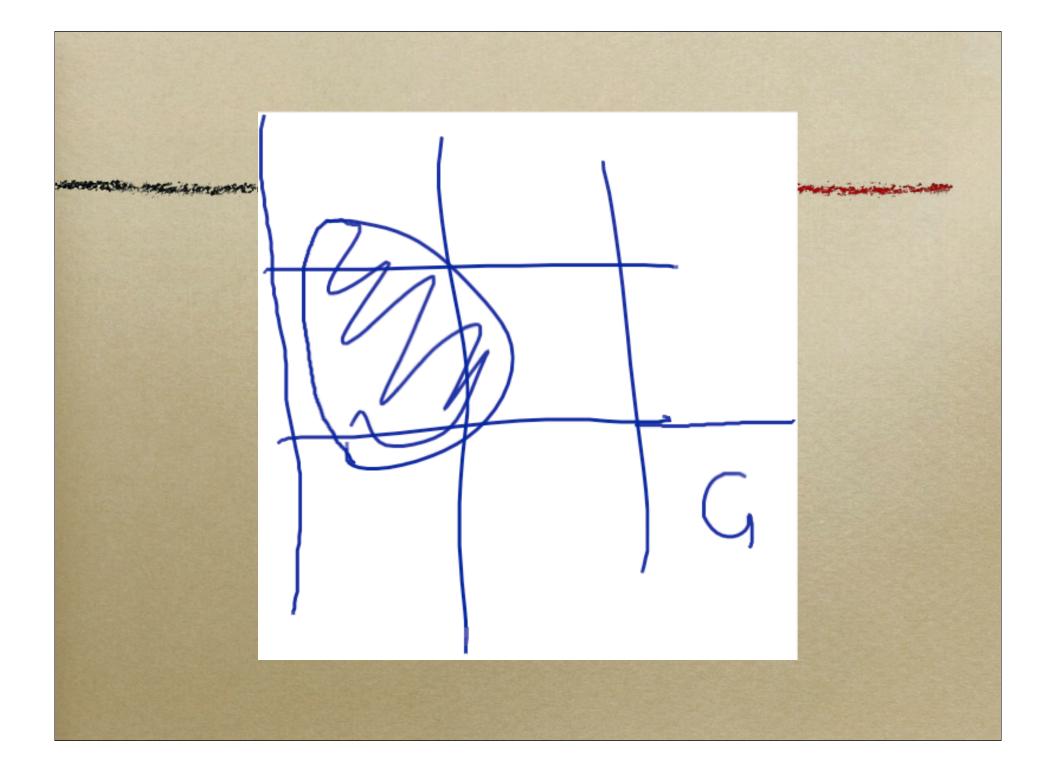
- Works pretty well, except:
 - Still don't know when to stop
 - Won't find shortest path
 - o Still doesn't really scale to high-d

Better yet

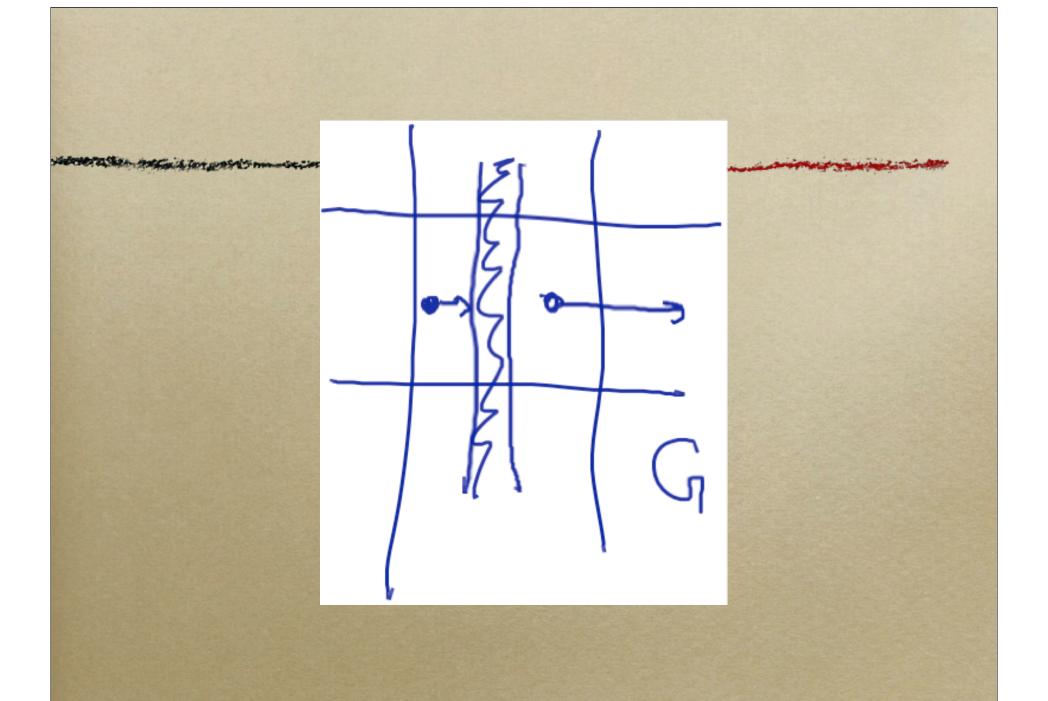
- Adaptive decomposition
- Split only cells that actually make a difference
 - o are on path from start
 - make a difference to our policy

Parti-game paper

- Andrew Moore and Chris Atkeson. The Parti-game Algorithm for Variable Resolution Reinforcement Learning in Multidimensional State-spaces
- http://www.autonlab.org/autonweb/14699.html



Advertising the second second second

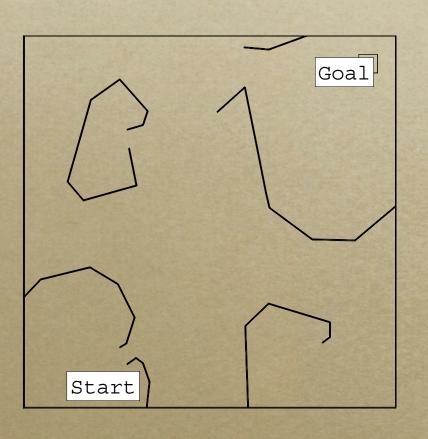


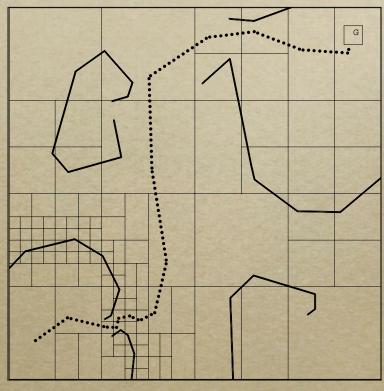
same asker and a forest mesers on series

Parti-game algorithm

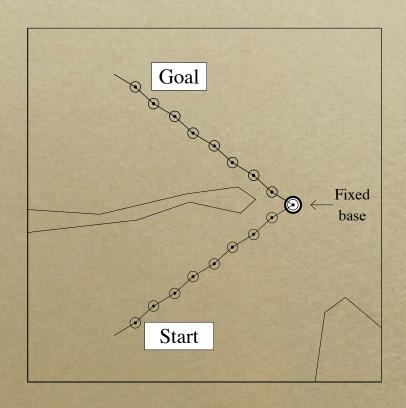
- Try actions from several points per cell
- o Try to plan a path from start to goal
- On the way, pretend an opponent gets to choose which outcome happens (out of all that have been observed in this cell)
- o If we can get to goal, we win
- o Otherwise we can split a cell

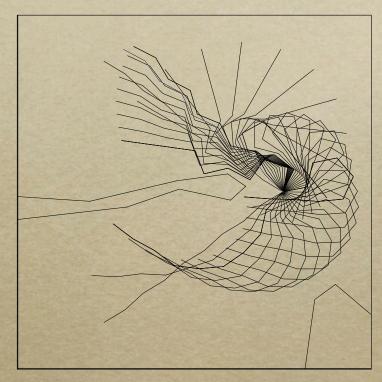
Parti-game example





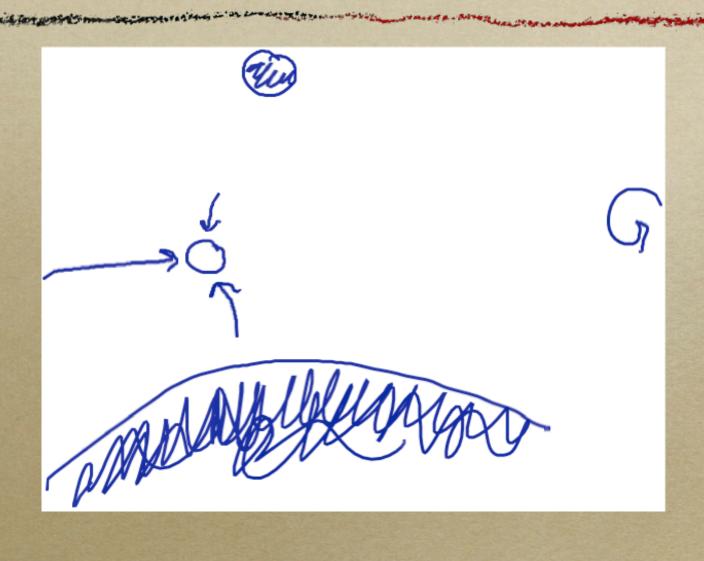
9dof planar arm





85 partitions total

Potential-based algorithms

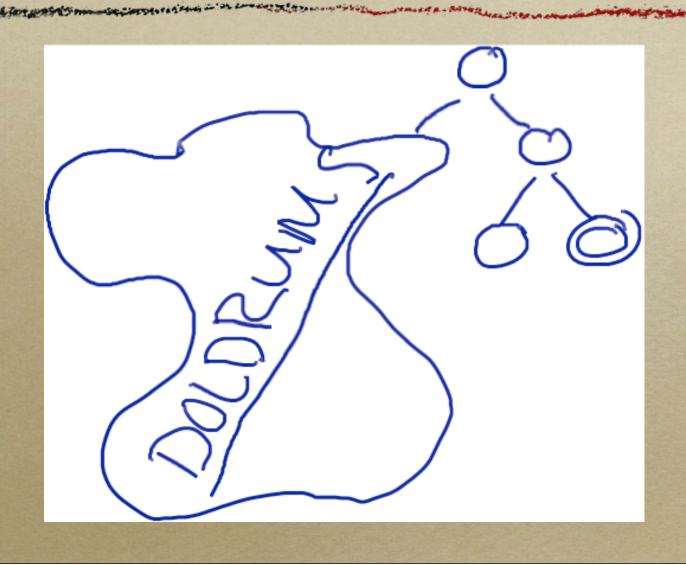


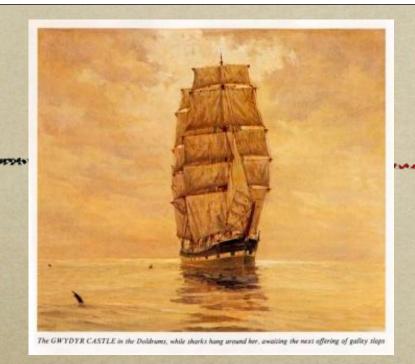
Potential-based algorithms

- Add a fictitious force that moves us away from obstacles
 - o stronger when closer
- Add a force towards goal
- Local minima galore ...
- Or, expensive but cool ways to calculate potentials that don't have local minima

Randomness in search

We can be very lucky or unlucky





Doldrums: One Of Murphy's Yarns
http://oldpoetry.com/opoem/56157 Cicely Fox Smith

"I heard onst of a barque," said Murphy.

"Becalmed, that couldn't get a breath,

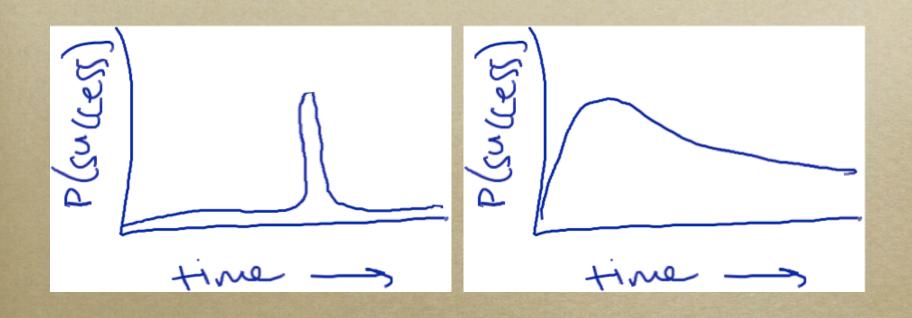
Till all the crowd was sick with scurvy

An' the skipper drunk himself to death."

Simple idea

- Try multiple starting points, random seeds for order of expanding neighbors
- Interleave computation (or iterative lengthening)
- When does this work?

Randomization cont'd



 Randomization works well if search times are sometimes short but have heavy tail

RRTs

- We will come back to randomness for more planning algorithms later
- For now, here's a randomized way of dividing up C-space that seems to work quite well in high-dimensions
- Rapidly-exploring Random Trees

RRTs

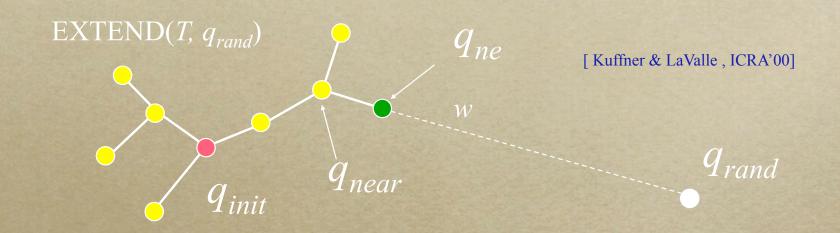
- Put landmarks into C-space
- Break up C-space into Voronoi regions around landmarks
- Put landmarks densely only if high resolution is needed to find a path
- Will not guarantee optimal path

RRT assumptions

- RANDOM_CONFIG
 - samples from some distribution on Cspace; can use to bias search
- EXTEND(q, q')
 - uses a local controller to head towards
 q'from q
 - stops before hitting obstacle
- FIND_NEAREST(q, Q)

Path Planning with RRTs

RRT = Rapidly-Exploring Random Tree



```
BUILD_RRT (q_{init}) {

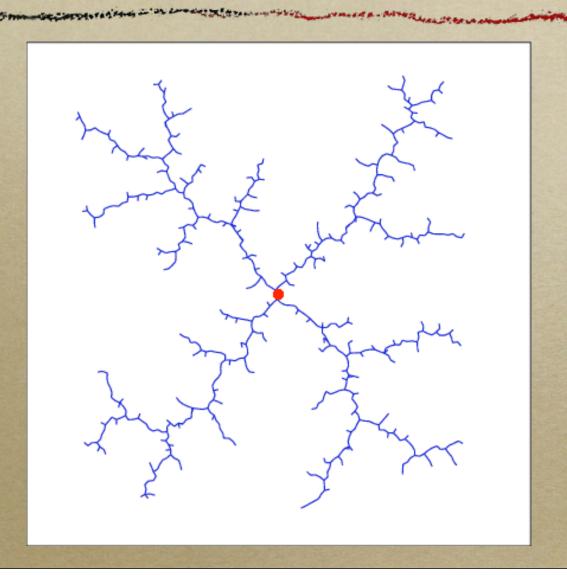
T.init(q_{init});

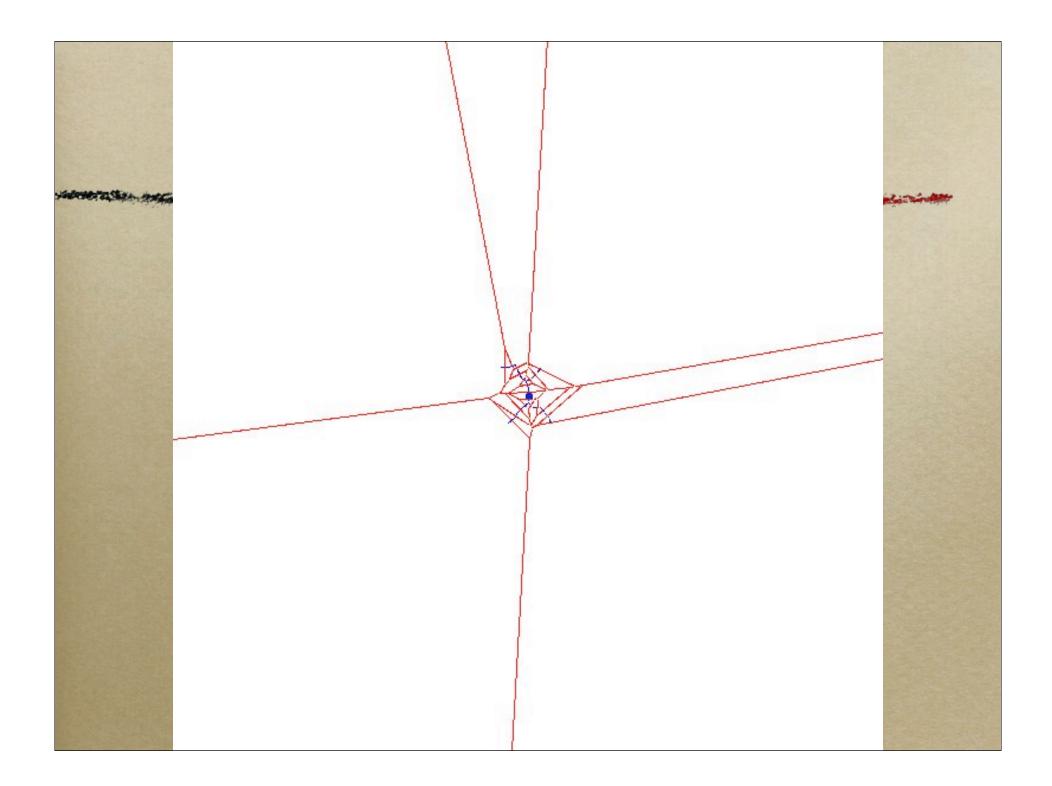
for k = 1 to K do

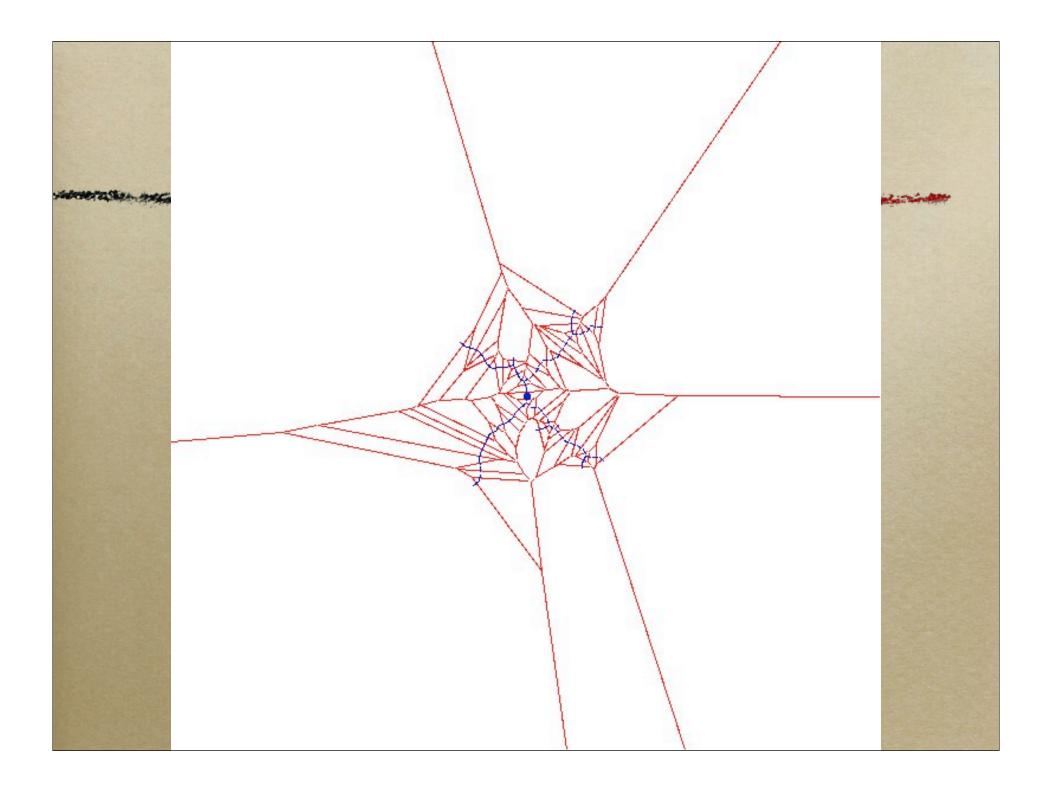
q_{rand} = \text{RANDOM\_CONFIG}();

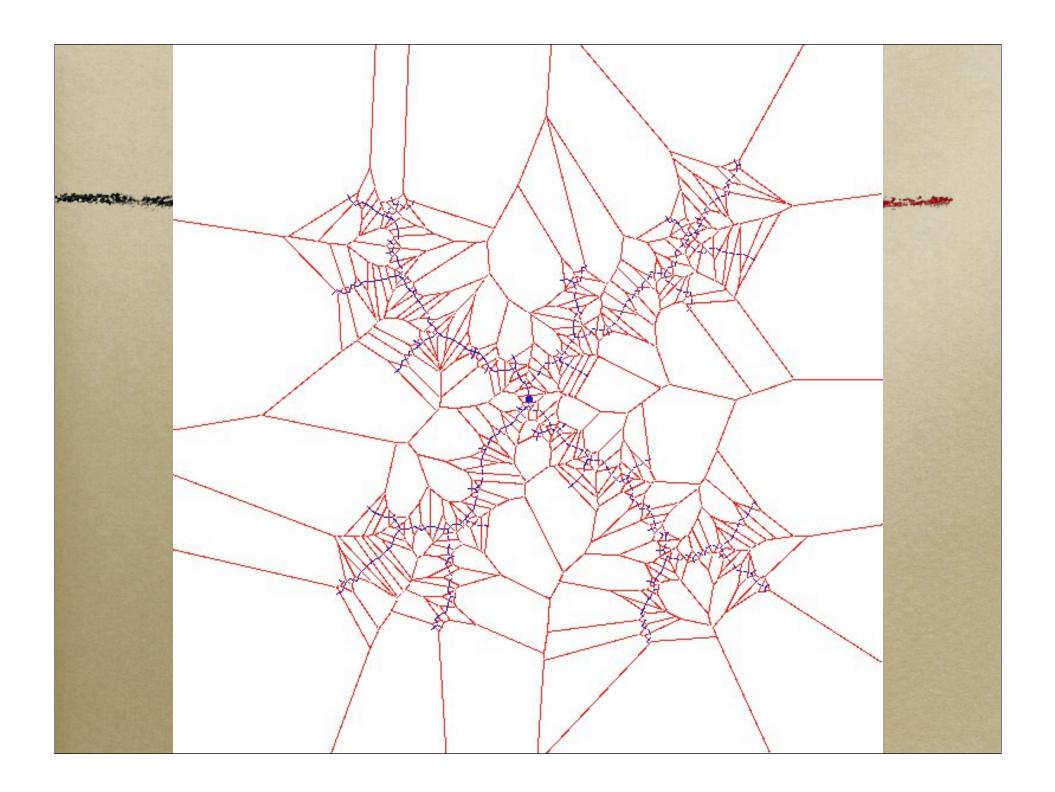
EXTEND(T, q_{rand})
}
```

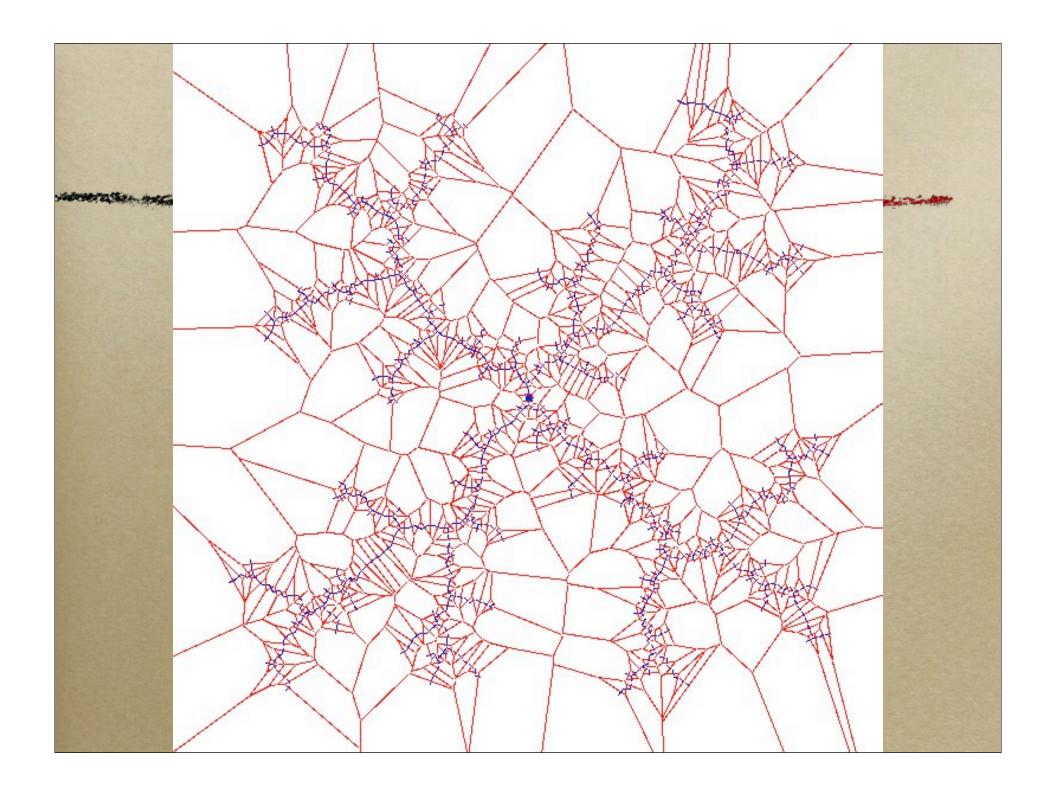
RRT example

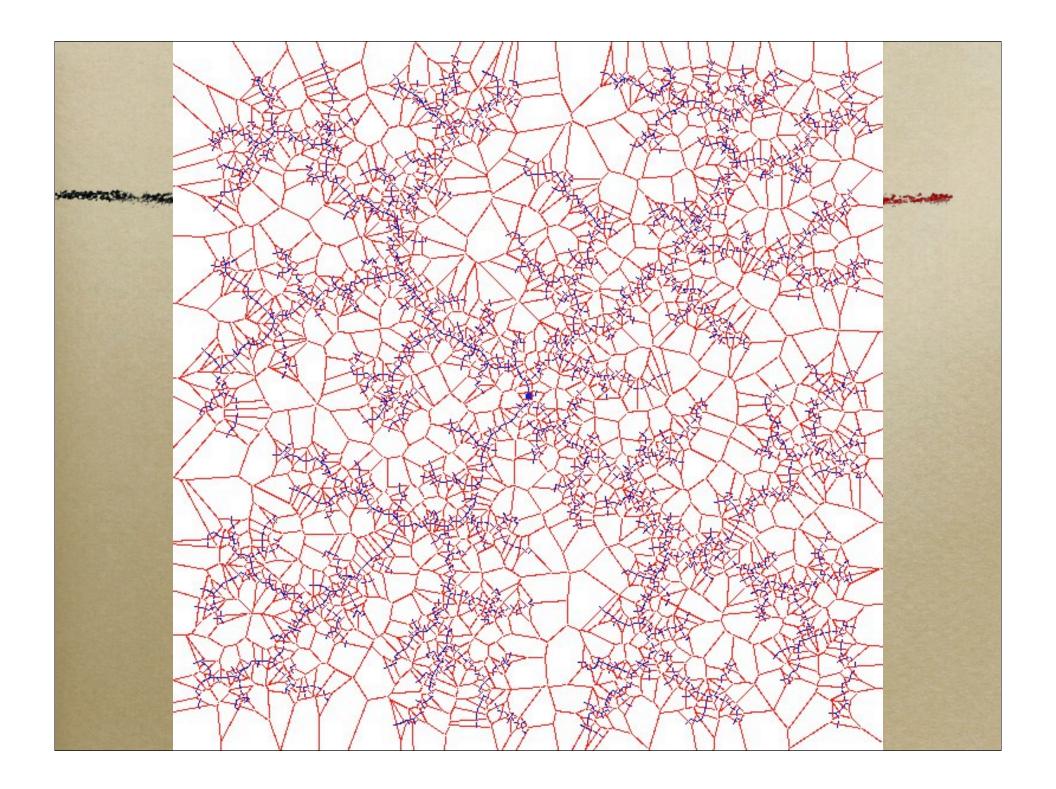




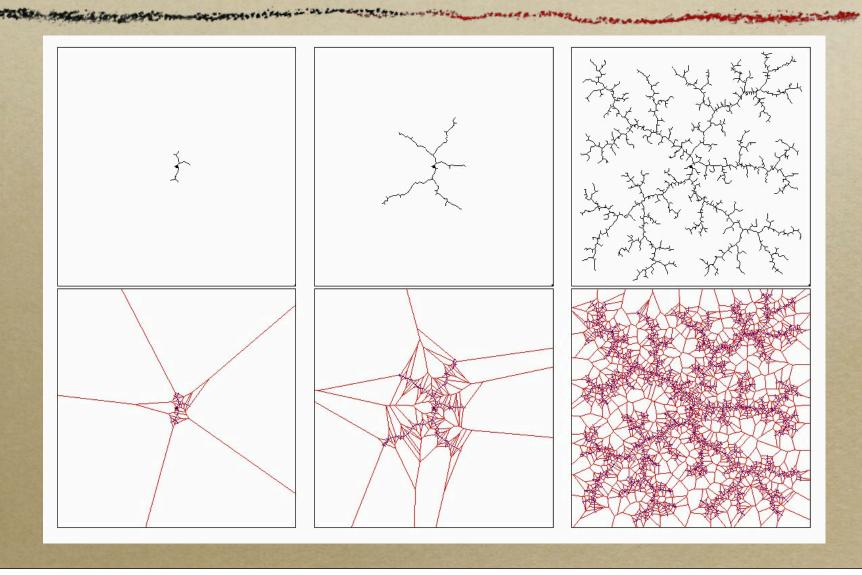




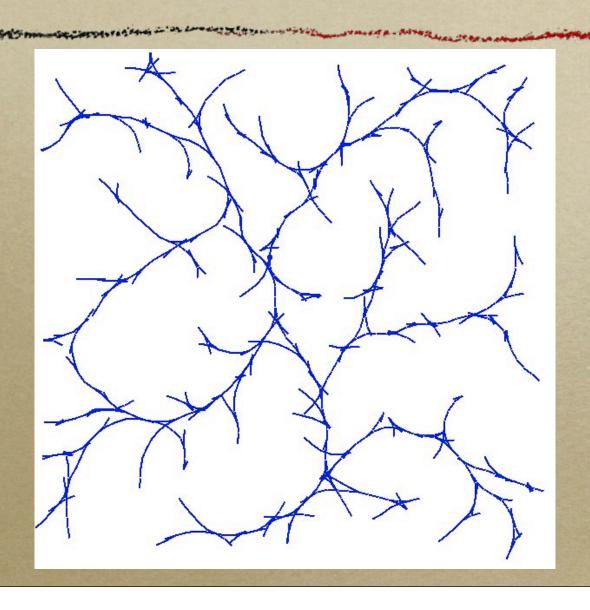




RRT example



RRT for a car (3 dof)

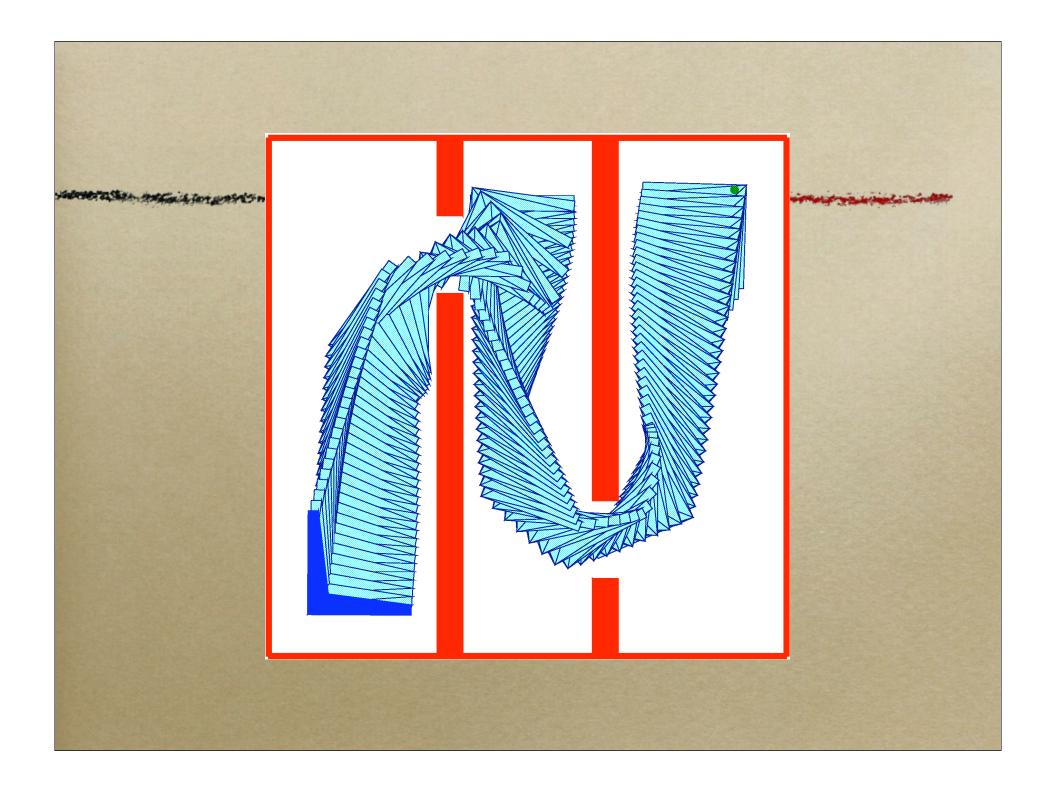


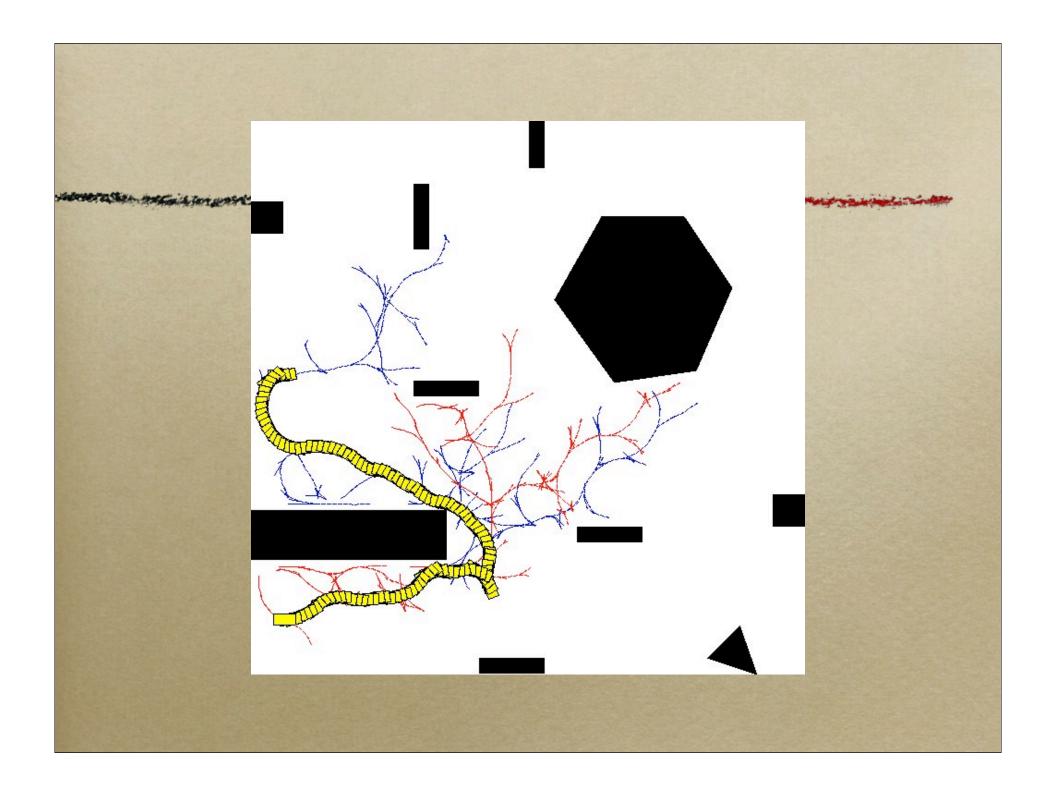
RRTs explore coarse to fine

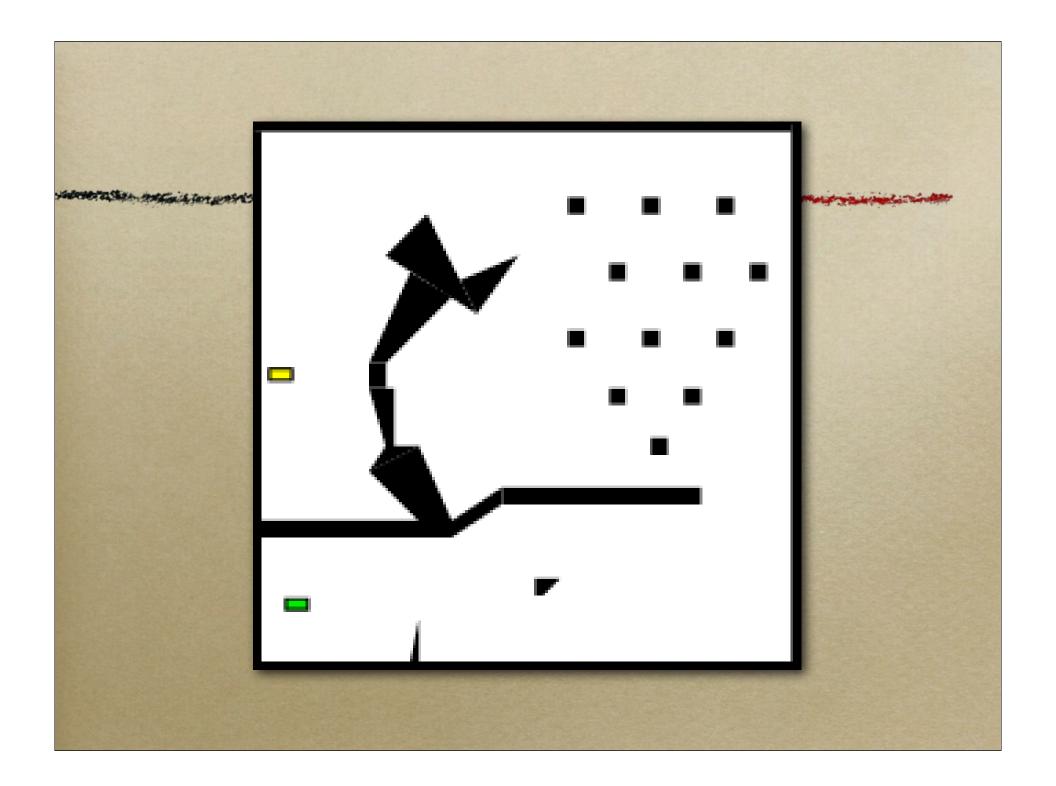
- Tend to break up large Voronoi regions
- Limiting distribution of vertices is RANDOM_CONFIG
 - Key idea in proof: as RRT grows, probability that qrand is reachable with local controller (and so immediately becomes a new vertex) approaches 1
- o In limit, we get a Voronoi cell decomposition

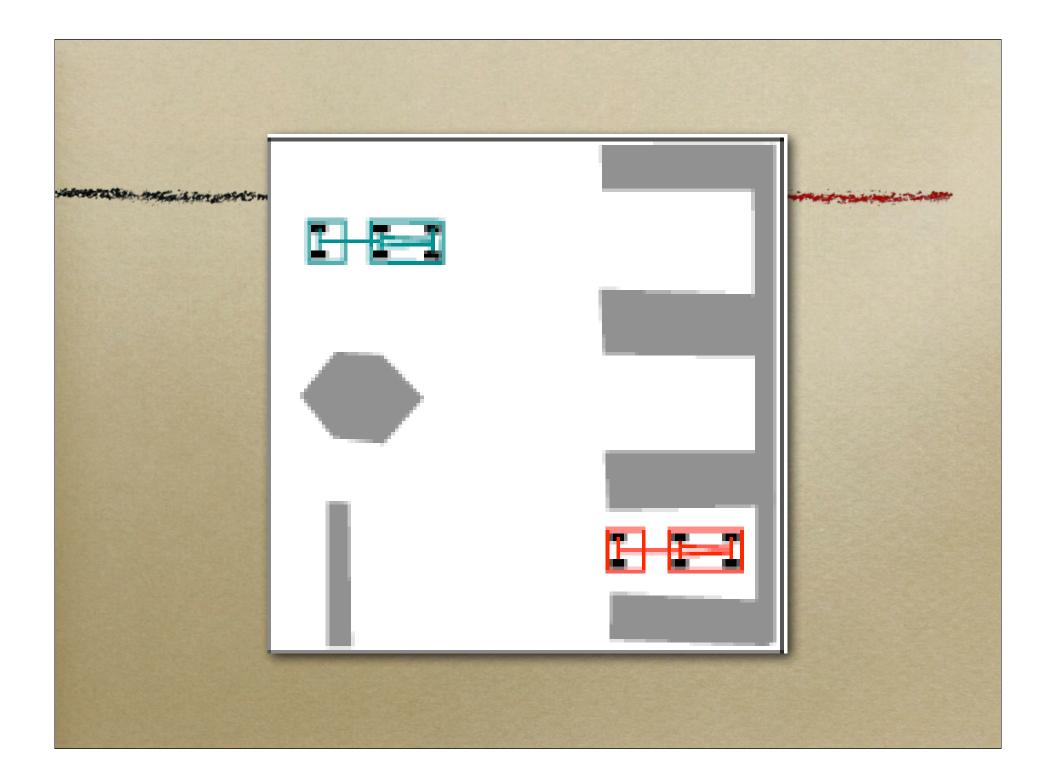
Planning with RRTs

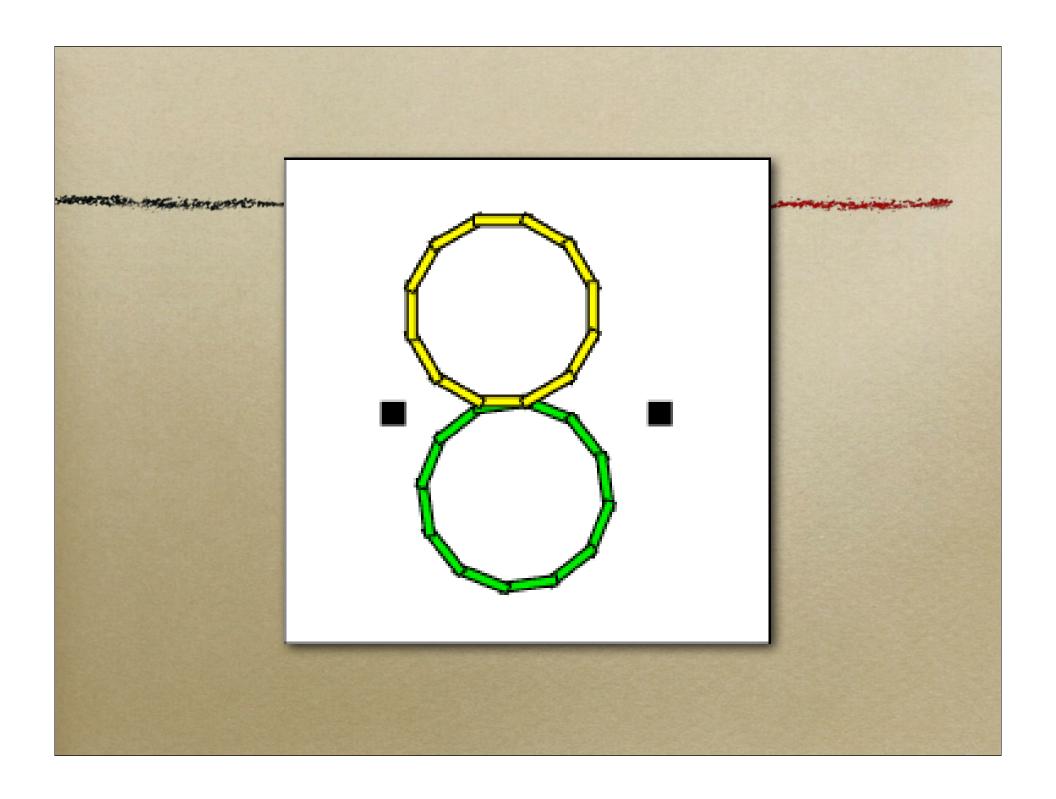
- Build RRT from start until we add a node that can reach goal using local controller
- ∘ (Unique) path: root → last node → goal
- Optional: cross-link tree by testing local controller, search within tree using A*
- Optional: grow forward and backward











What you should know

- C-space
- Ways of splitting up C-space
 - Visibility graph
 - Voronoi
 - Exact, approximate cell decomposition
 - Adaptive cells (quadtree, parti-game)
- Potential fields
- RRTs