# Assignment 1 The Untyped $\lambda$-Calculus 

15-814: Types and Programming Languages<br>Frank Pfenning

Due Thursday, September 9, 2021
60 pts

This assignment is due on the above date and it must be submitted electronically on Gradescope. Please use the attached template to typeset your assignment and make sure to include your full name and Andrew ID. For the written problems, you may also submit handwritten answers that have been scanned and are easily legible.

Please carefully read the policies on collaboration and credit on the course web pages at http://www.cs.cmu.edu/~fp/courses/15814-f20/assignments.html.

You should hand in two files

- hw01.pdf with your written solutions to the questions.
- hw01.lam with the code, where the solutions to the problems are clearly marked and auxiliary code (either from lecture or your own) is included so it passes the LAMBDA checker.


## 1 Calculating in the $\lambda$-Calculus

Task 1 (L1.1, 5 pts) Define the following functions on Booleans.

1. The "nor" operator, which yields true iff both inputs are false.
2. The conditional "if" such that

$$
\begin{aligned}
& \text { if true } e_{1} e_{2}={ }_{\beta} \quad e_{1} \\
& \text { if false } e_{1} e_{2}={ }_{\beta} \quad e_{2}
\end{aligned}
$$

3. In the solution file hw 01.1 am include the necessary definitions of nor and if and also sufficient test cases to certify their correctness.

Task 2 (L2.3, 15 pts ) One approach to representing functions defined by the schema of primitive recursion is to change the representation so that $\bar{n}$ is not an iterator but a primitive recursor.

$$
\begin{aligned}
\overline{0} & =\lambda s \cdot \lambda z \cdot z \\
\overline{n+1} & =\lambda s \cdot \lambda z \cdot s \bar{n}(\bar{n} s z)
\end{aligned}
$$

1. Define the successor function succ on this new representation (if possible) and show its correctness.
2. Define the predecessor function pred on this new representation (if possible) and show its correctness.
3. Explore if it is possible to directly represent any function $f$ specified by a schema of primitive recursion, ideally without constructing and destructing pairs. Write what you find.

Task 3 (L2.4, 10 pts) The unary representation of natural numbers requires tedious and error-prone counting to check whether your functions (such as the Lucas function in the exercise below) behave correctly on some inputs with large answers. Fortunately, you can exploit that the LAMBDA implementation counts the number or reduction steps for you and prints it in decimal form!
(i) We have

$$
\bar{n} \text { succ zero } \longrightarrow_{\beta}^{*} \bar{n}
$$

because $\bar{n}$ iterates the successor function $n$ times on 0 . Run some experiments in LAMBDA and conjecture how many leftmost-outermost reduction steps are required as a function of $n$. Note that only $\beta$-reductions are counted, and not replacing a definition (for example, zero by $\lambda s . \lambda z . z)$. We justify this because we think of the definitions as taking place at the metalevel, in our mathematical domain of discourse.
(ii) Prove your conjecture from part (i), using induction on $n$. It may be helpful to use the mathematical notation $f^{k} c$ to describe a $\lambda$-expression generated by $f^{0} c=c$ and $f^{k+1} c=$ $f\left(f^{k} c\right)$ where $f$ and $c$ are $\lambda$-expressions. For example, $\bar{n}=\lambda s . \lambda z . s^{n} z$ or succ ${ }^{3}$ zero $=$ succ (succ (succ zero)).

Task 4 (L2.5, 15 pts) Define the following functions in the $\lambda$-calculus using the LAMBDA implementation. Here we take " $=$ " to mean $={ }_{\beta}$, that is, $\beta$-conversion.

You may use all the functions in nat.lam as helper functions. Your functions should evidently reflect iteration, primitive recursion and pairs. In particular, you should avoid the use of the $Y$ combinator which will be introduced in Lecture 3.

Provide at least 3 test cases for each function and include them, together with your function definitions, in the file hw01.lam.
(i) if0 (definition by cases) satisfying the specification

$$
\begin{array}{ll}
\text { if0 } 0 \overline{0} x y & =x \\
\text { if0 } \overline{k+1} x y & =y
\end{array}
$$

(ii) even satisfying the specification

$$
\begin{aligned}
& \text { even } \overline{2 k}=\text { true } \\
& \text { even } \overline{2 k+1}=\text { false }
\end{aligned}
$$

(iii) half satisfying the specification

$$
\begin{aligned}
& \text { half } \overline{2 k}=\bar{k} \\
& \text { half } \overline{2 k+1}=\bar{k}
\end{aligned}
$$

Task 5 (L2.6, 15 pts ) The Lucas function (a variant on the Fibonacci function) is defined mathematically by

$$
\begin{array}{ll}
\text { lucas } 0 & =2 \\
\text { lucas } 1 & =1 \\
\text { lucas }(n+2) & =\text { lucas } n+\text { lucas }(n+1)
\end{array}
$$

Give an implementation of the Lucas function in the $\lambda$-calculus via the LAMBDA implementation.
You may use the functions from nat.lam as helper functions, as well as those from Task 4. Your functions should evidently reflect iteration, primitive recursion and pairs. In particular, you should avoid the use of the $Y$ combinator which will be introduced in Lecture 3.

Test your implementation on inputs $0,1,9$, and 11, expecting results 2, 1, 76, and 199. Include these tests in your code submission hw01.lam, and record the number of $\beta$-reductions used by your function in your written submission.

