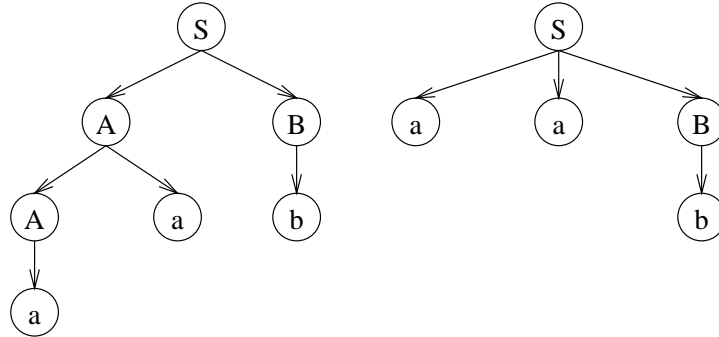


(a) Show that this grammar is ambiguous.

We may derive aab using either $S \rightarrow AB$ or $S \rightarrow aaB$, which give rise to different trees:



(b) Give a regular expression that describes the same language.

The language comprises all words that consist of one or more a 's followed by exactly one b ; thus, the corresponding regular expression is a^+b .

(c) Construct an unambiguous grammar that describes the same language.

We can remove the production $S \rightarrow aaB$ to eliminate the ambiguity, which leads to the following grammar:

$$\begin{aligned} S &\rightarrow AB \\ A &\rightarrow a \mid Aa \\ B &\rightarrow b \end{aligned}$$

We can further simplify this grammar, by eliminating the variables A and B :

$$S \rightarrow aS \mid ab$$

Problem 2

For each of the following two grammars, construct an equivalent grammar that has no λ -productions and no unit-productions.

(a)
$$\begin{aligned} S &\rightarrow A \mid B \\ A &\rightarrow a \mid aA \\ B &\rightarrow b \mid bB \end{aligned}$$

Simplified grammar:

$$\begin{aligned} S &\rightarrow a \mid aA \mid b \mid bB \\ A &\rightarrow a \mid aA \\ B &\rightarrow b \mid bB \end{aligned}$$

(b)
$$\begin{aligned} S &\rightarrow aAb \mid bAa \mid aSb \mid bSa \\ A &\rightarrow aAa \mid \lambda \end{aligned}$$

Simplified grammar:

$$\begin{aligned} S &\rightarrow aAb \mid bAa \mid aSb \mid bSa \mid ab \mid ba \\ A &\rightarrow aAa \mid aa \end{aligned}$$