## Analysis of Algorithms: Assignment 3

Due date: September 16 (Thursday)

## **Problem 1** (4 points)

Determine asymptotic upper and lower bounds for each of the following recurrences. Make your bounds as tight as possible.

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(a) T(n) = 16T(n/4) + n

(b) T(n) = 16T(n/4) + n^2

(c) T(n) = 2T(n/4) + \sqrt{n}

(d) T(n) = T(\sqrt{n}) + 1
```

## Problem 2 (6 points)

Consider the following sorting algorithm:

```
\begin{array}{l} {\rm STOOGE\text{-}SORT}(A,i,j) \\ {\rm if} \ A[i] > A[j] \\ {\rm then} \ {\rm exchange} \ A[i] \leftrightarrow A[j] \\ {\rm if} \ i+1 \geq j \\ {\rm then} \ {\rm return} \\ k \leftarrow \lfloor (j-i+1)/3 \rfloor \\ {\rm STOOGE\text{-}SORT}(A,i,j-k) \qquad \rhd \ {\rm first} \ {\rm two\text{-}thirds} \\ {\rm STOOGE\text{-}SORT}(A,i,j-k) \qquad \rhd \ {\rm last} \ {\rm two\text{-}thirds} \\ {\rm STOOGE\text{-}SORT}(A,i,j-k) \qquad \rhd \ {\rm first} \ {\rm two\text{-}thirds} \\ {\rm STOOGE\text{-}SORT}(A,i,j-k) \qquad \rhd \ {\rm first} \ {\rm two\text{-}thirds} \\ {\rm STOOGE\text{-}SORT}(A,i,j-k) \qquad \rhd \ {\rm first} \ {\rm two\text{-}thirds} \\ {\rm STOOGE\text{-}SORT}(A,i,j-k) \qquad \rhd \ {\rm first} \ {\rm two\text{-}thirds} \\ {\rm STOOGE\text{-}SORT}(A,i,j-k) \qquad \rhd \ {\rm first} \ {\rm two\text{-}thirds} \\ {\rm STOOGE\text{-}SORT}(A,i,j-k) \qquad \rhd \ {\rm first} \ {\rm two\text{-}thirds} \\ {\rm STOOGE\text{-}SORT}(A,i,j-k) \qquad \rhd \ {\rm first} \ {\rm two\text{-}thirds} \\ {\rm STOOGE\text{-}SORT}(A,i,j-k) \qquad \rhd \ {\rm first} \ {\rm two\text{-}thirds} \\ {\rm STOOGE\text{-}SORT}(A,i,j-k) \qquad \rhd \ {\rm first} \ {\rm two\text{-}thirds} \\ {\rm STOOGE\text{-}SORT}(A,i,j-k) \qquad \rhd \ {\rm first} \ {\rm two\text{-}thirds} \\ {\rm STOOGE\text{-}SORT}(A,i,j-k) \qquad \rhd \ {\rm first} \ {\rm two\text{-}thirds} \\ {\rm STOOGE\text{-}SORT}(A,i,j-k) \qquad \rhd \ {\rm first} \ {\rm two\text{-}thirds} \\ {\rm STOOGE\text{-}SORT}(A,i,j-k) \qquad \rhd \ {\rm first} \ {\rm two\text{-}thirds} \\ {\rm STOOGE\text{-}SORT}(A,i,j-k) \qquad \rhd \ {\rm first} \ {\rm two\text{-}thirds} \\ {\rm STOOGE\text{-}SORT}(A,i,j-k) \qquad \rhd \ {\rm first} \ {\rm two\text{-}thirds} \\ {\rm STOOGE\text{-}SORT}(A,i,j-k) \qquad \rhd \ {\rm first} \ {\rm two\text{-}thirds} \\ {\rm STOOGE\text{-}SORT}(A,i,j-k) \qquad \rhd \ {\rm first} \ {\rm two\text{-}thirds} \\ {\rm STOOGE\text{-}SORT}(A,i,j-k) \qquad \rhd \ {\rm first} \ {\rm two\text{-}thirds} \\ {\rm STOOGE\text{-}SORT}(A,i,j-k) \qquad \rhd \ {\rm first} \ {\rm two\text{-}thirds} \\ {\rm STOOGE\text{-}SORT}(A,i,j-k) \qquad \rhd \ {\rm first} \ {\rm two\text{-}thirds} \\ {\rm STOOGE\text{-}SORT}(A,i,j-k) \qquad \rhd \ {\rm first} \ {\rm two\text{-}thirds} \\ {\rm STOOGE\text{-}SORT}(A,i,j-k) \qquad \rhd \ {\rm first} \ {\rm two\text{-}thirds} \\ {\rm STOOGE\text{-}SORT}(A,i,j-k) \qquad \rhd \ {\rm first} \ {\rm two\text{-}thirds} \\ {\rm STOOGE\text{-}SORT}(A,i,j-k) \qquad \rhd \ {\rm first} \ {\rm two\text{-}thirds} \\ {\rm STOOGE\text{-}SORT}(A,i,j-k) \qquad \rhd \ {\rm first} \ {\rm two\text{-}thirds} \\ {\rm STOOGE\text{-}SORT}(A,i,j-k) \qquad \rhd \ {\rm first} \ {\rm two\text{-}thirds} \\ {\rm two\text{-}thirds} \\ {\rm two\text{-}thirds}
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- (a) Argue that STOOGE-SORT(A, 1, n) correctly sorts the input array A[1..n].
- (b) Give the recurrence for the worst-case running time of Stooge-Sort and a tight asymptotic (Θ-notation) bound on the worst-case time.
- (c) Compare the worst-case time of Stooge-Sort with that of Insertion-Sort and Merge-Sort. Is it a good algorithm?

## Problem 3 (bonus)

This problem is optional, and it does not affect your grade for the homework; however, if you solve it, then you will get 2 bonus points toward your final grade for the course. You cannot submit this bonus problem after the deadline.

Consider the following algorithm, which inputs a natural number n and returns a natural number m. The algorithm uses two data structures for storing intermediate results: (1) a square matrix A[1..n, 1..n] and (2) a set S of natural numbers, which may be represented by an array or linked list. The algorithm calls a MIN-OUTSIDE(S) procedure which finds the minimal natural number that does not belong to the set S; if S is empty, this procedure returns 0.

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\begin{array}{l} \text{Slow-Counter}(n) \\ \textbf{for } i \leftarrow 1 \textbf{ to } n \\ & \textbf{do for } j \leftarrow 1 \textbf{ to } n \\ & \textbf{do } S \leftarrow \emptyset \quad \rhd \text{ make the set } S \text{ empty} \\ & \textbf{for } k \leftarrow 1 \textbf{ to } i - 1 \\ & \textbf{do } S \leftarrow S \cup \{A[k,j]\} \quad \rhd \text{ add the } A[k,j] \text{ value to } S \\ & \textbf{for } k \leftarrow 1 \textbf{ to } j - 1 \\ & \textbf{do } S \leftarrow S \cup \{A[i,k]\} \quad \rhd \text{ add the } A[i,k] \text{ value to } S \\ & A[i,j] \leftarrow \text{Min-Outside}(S) \quad \rhd \text{ find the minimal natural number } not \text{ in } S \\ & m \leftarrow A[n,n] \\ & \textbf{return } m \end{array}
```

Give a much faster algorithm that computes the same value m, without using a matrix.