

Algorithms (COT 6405): Solutions 7

Problem 1

Let $A[1..n]$ be a sorted array of n distinct integer numbers. Write an efficient algorithm $\text{INDEX-SEARCH}(A, n)$ that finds an index i such that $A[i] = i$.

The algorithm is almost identical to BINARY-SEARCH , and its time complexity is $O(\lg n)$. It works only for integer arrays, since it is based on the assumption that, for every two indices p and r (where $p \leq r$), we have $A[r] - A[p] \geq r - p$.

```
INDEX-SEARCH( $A, n$ )
 $p \leftarrow 1$ 
 $r \leftarrow n$ 
while  $p < r$ 
    do  $q = \lfloor (p + r)/2 \rfloor$ 
        if  $q \leq A[q]$ 
            then  $r \leftarrow q$ 
            else  $p \leftarrow q + 1$ 
if  $p = A[p]$ 
    then return  $p$ 
    else return 0
```

Problem 2

We consider an array $A[1..n]$ and define a segment sum from p to r :

$$\text{sum}(p, r) = \sum_{p \leq i \leq r} A[i].$$

Write a linear-time algorithm that determines the maximum over all segment sums.

```
MAX-SEGMENT( $A, n$ )
 $Local-Max \leftarrow A[1]$ 
 $Global-Max \leftarrow A[1]$ 
for  $i \leftarrow 2$  to  $n$ 
    do  $Local-Max \leftarrow \max(A[i], Local-Max + A[i])$ 
         $\triangleright$   $Local-Max$  is the maximum over the segments whose last element is  $A[i]$ 
         $Global-Max \leftarrow \max(Local-Max, Global-Max)$ 
         $\triangleright$   $Global-Max$  is the maximum over all segments in  $A[1..i]$ 
return  $Global-Max$ 
```