

## Algorithms: Solutions 2

- (1) Give an algorithm that determines the number of *distinct* elements in an array  $A[1..n]$ .  
(2) Estimate the worst-case running time of your algorithm.

DISTINCT-COUNTER( $A, n$ )	cost    times
$counter \leftarrow 0$	$c_1$ 1
<b>for</b> $j \leftarrow 1$ <b>to</b> $n$	$c_2$ $n + 1$
<b>do</b> $i \leftarrow j + 1$	$c_3$ $n$
<b>while</b> $i < n + 1$ and $A[j] \neq A[i]$	$c_4$ $\leq \sum_{i=1}^n (n - i + 1) = \frac{n \cdot (n+1)}{2}$
<b>do</b> $i \leftarrow i + 1$	$c_5$ $\leq \sum_{i=1}^n (n - i) = \frac{(n-1) \cdot n}{2}$
<b>if</b> $i = n + 1$	$c_6$ $n$
<b>then</b> $counter \leftarrow counter + 1$	$c_7$ $\leq n$
<b>return</b> $counter$	$c_8$ 1

$$\begin{aligned}
T(n) &\leq c_1 + c_2 \cdot (n + 1) + c_3 \cdot n + c_4 \cdot \frac{n \cdot (n + 1)}{2} + c_5 \cdot \frac{(n - 1) \cdot n}{2} + c_6 \cdot n + c_7 \cdot n + c_8 \\
&= \left(\frac{c_4}{2} + \frac{c_5}{2}\right) \cdot n^2 + \left(c_2 + c_3 + \frac{c_4}{2} - \frac{c_5}{2} + c_6 + c_7\right) \cdot n + (c_1 + c_2 + c_8) \\
&= \Theta(n^2)
\end{aligned}$$