

## Algorithms: Solutions 8

[illegible]

### Problem 1

The depth-first search algorithm may be used to identify the connected components of an undirected graph. Write a modified version of DFS for this task.

$$\text{DFS-COMPONENTS}(G)$$
 $k \leftarrow 0$      $\triangleright$  component counter

**for** each  $u \in V[G]$

$$\mathbf{do} \text{ component}[u] \leftarrow 0$$
**for** each  $u \in V[G]$ **do if**  $component[u] = 0$ 

**then**  $k \leftarrow k + 1$

$$\text{DFS-VISIT}(u, k)$$
**return**  $k$ 
$$\text{DFS-VISIT}(u, k)$$
$$component[u] \leftarrow k \quad \triangleright u \text{ has just been discovered}$$
**for** each  $v \in Adj[u]$ **do if**  $component[v] = 0$ **then** DFS-VISIT( $v, k$ )

**Problem 2**

Suppose that  $G$  is an undirected graph, and you need to check whether  $G$  has cycles. Design an algorithm that returns `TRUE` if  $G$  is acyclic, and `FALSE` if  $G$  has cycles.

The key observation is that an acyclic undirected graph has at most  $V-1$  edges. To determine whether a graph  $G$  is acyclic, we first count its edges. If the edge counter reaches  $V$ , we immediately return `FALSE`, without counting the rest of edges. On the other hand, if the number of edges is less than  $V$ , we apply DFS to search for cycles. In either case, the overall running time is  $O(V)$ .