Time and synchronization

(“There’s never enough time…”)

Today’s outline

- Time in distributed systems
  - A baseball example
- Synchronizing real clocks
  - Cristian’s algorithm
  - The Berkeley Algorithm
  - Network Time Protocol (NTP)
- Logical time
- Lamport logical clocks

Distributed time

- The notion of time is well-defined (and measurable) at each single location
  - But the relationship between time at different locations is unclear
  - e.g., packet-sending from HW 1 #6:

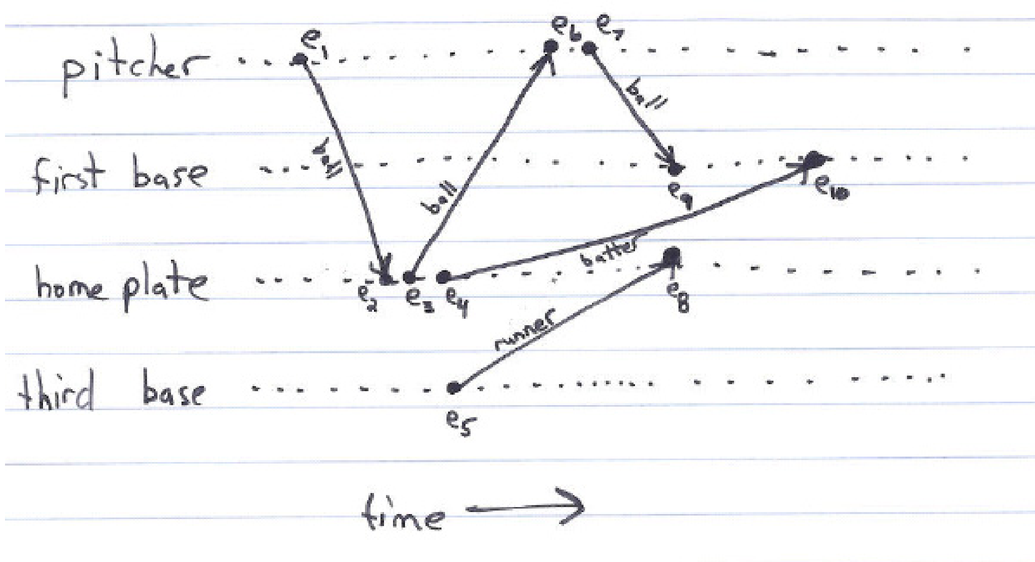
A baseball example

- Four locations: pitcher's mound, first base, home plate, and third base
- Ten events:

  e_1: pitcher throws ball to home
  e_2: ball arrives at home
  e_3: batter hits ball to pitcher
  e_4: batter runs to first base
  e_5: runner runs to home
  e_6: ball arrives at pitcher
  e_7: pitcher throws ball to first base
  e_8: runner arrives at home
  e_9: ball arrives at first base
  e_{10}: batter arrives at first base

A baseball example

- Pitcher knows e_1 happens before e_6, which happens before e_7
- Home plate umpire knows e_2 is before e_3, which is before e_4, which is before e_8, …
- Relationship between e_8 and e_9 is unclear

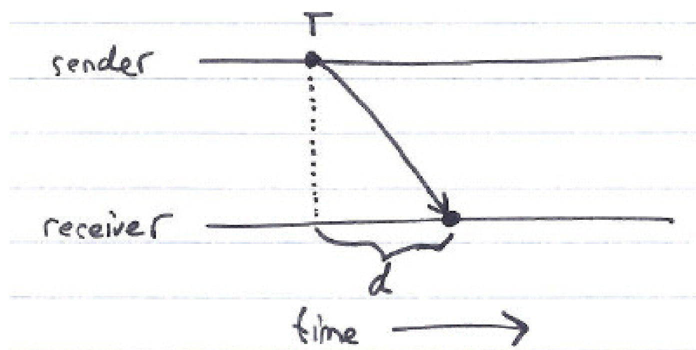


Ways to synchronize

- Send message from first base to home?
  - Or to a central timekeeper
  - How long does this message take to arrive?
- Synchronize clocks before the game?
  - Clocks drift
    - million to one => 1 second in 11 days
- Synchronize continuously during the game?
  - GPS, pulsars, etc

Perfect networks

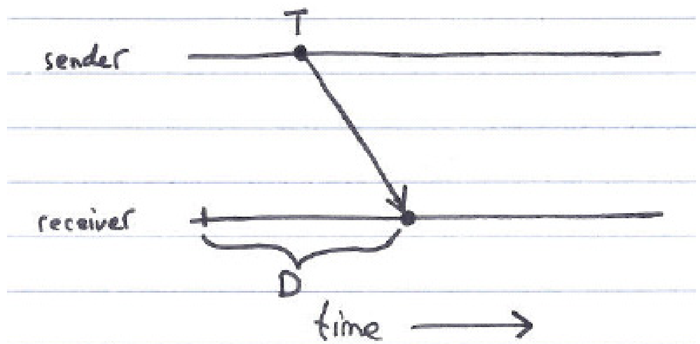
- Messages always arrive, with propagation delay exactly *d*



- Sender sends time *T* in a message
- Receiver sets clock to *T+d*
  - Synchronization is exact

Synchronous networks

- Messages always arrive, with propagation delay at most *D*



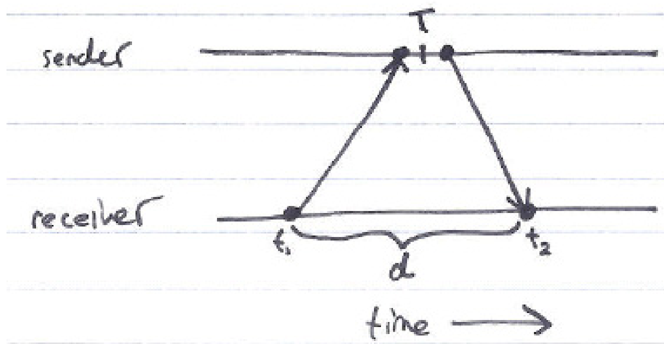
- Sender sends time *T* in a message
- Receiver sets clock to *T* + *D/2*
  - Synchronization error is at most *D/2*

Synchronization in the real world

- Real networks are asynchronous
  - Propagation delays are arbitrary
- Real networks are unreliable
  - Messages don’t always arrive

Cristian's algorithm

- Request time, get reply
  - Measure actual round-trip time *d*



- Sender's time was *T* between t_1 and t_2
- Receiver sets time to $T + d/2$
  - Synchronization error is at most $d/2$
- Can retry until we get a relatively small *d*

The Berkeley algorithm

- Master uses Cristian's algorithm to get time from many clients
  - Computes average time
  - Can discard outliers
- Sends time adjustments back to all clients

The Network Time Protocol (NTP)

- Uses a hierarchy of time servers
  - Class 1 servers have highly-accurate clocks
    - connected directly to atomic clocks, etc.
  - Class 2 servers get time from only Class 1 and Class 2 servers
  - Class 3 servers get time from any server
- Synchronization similar to Cristian's alg.
  - Modified to use multiple one-way messages instead of immediate round-trip

Real synchronization is imperfect

- Clocks never exactly synchronized
- Often inadequate for distributed systems
  - might need totally-ordered events
  - might need millionth-of-a-second precision

Logical time

- Capture just the “happens before” relationship between events
  - Discard the infinitesimal granularity of time
  - Corresponds roughly to causality
- Time at each process is well-defined
  - Definition ($\rightarrow_i$): We say $e \rightarrow_i e'$ if e happens before e' at process i

Global logical time

- Definition ($\rightarrow$): We define $e \rightarrow e'$ using the following rules:
  - Local ordering: $e \rightarrow e'$ if $e \rightarrow_i e'$ for any process i
  - Messages: send(m) $\rightarrow$ receive(m) for any message m
  - Transitivity: $e \rightarrow e''$ if $e \rightarrow e'$ and $e' \rightarrow e''$
- We say e "happens before" e' if $e \rightarrow e'$

Concurrency

- $\rightarrow$ is only a partial-order
  - Some events are unrelated
- Definition (concurrency): We say *e* is concurrent with *e'* (written $e \| e'$) if neither $e \rightarrow e'$ nor $e' \rightarrow e$

The baseball example revisited

- $e_1 \rightarrow e_2$
  - by the message rule
- $e_1 \rightarrow e_{10}$, because
  - $e_1 \rightarrow e_2$, by the message rule
  - $e_2 \rightarrow e_4$, by local ordering at home plate
  - $e_4 \rightarrow e_{10}$, by the message rule
  - Repeated transitivity of the above relations
- $e_8 \| e_9$, because
  - No application of the $\rightarrow$ rules yields either $e_8 \rightarrow e_9$ or $e_9 \rightarrow e_8$

Lamport logical clocks

- Lamport clock *L* orders events consistent with logical “happens before” ordering
  - If e $\rightarrow$ e’, then $L(e) < L(e')$
- But not the converse
  - $L(e) < L(e')$ does not imply $e \rightarrow e'$
- Similar rules for concurrency
  - $L(e) = L(e')$ implies $e \parallel e'$ (for distinct e,e')
  - $e \parallel e'$ does not imply $L(e) = L(e')$
- i.e., Lamport clocks arbitrarily order some concurrent events

Lamport's algorithm

- Each process i keeps a local clock, L_i
- Three rules:
  1. At process i, increment L_i before each event
  2. To send a message m at process i, apply rule 1 and then include the current local time in the message: i.e., $send(m,L_i)$
  3. To receive a message (m,t) at process j, set $L_j = max(L_j,t)$ and then apply rule 1 before time-stamping the receive event
- The global time $L(e)$ of an event e is just its local time
  - For an event e at process i, $L(e) = L_i(e)$

Lamport on the baseball example

- Initializing each local clock to 0, we get

$L(e_1) = 1$	(pitcher throws ball to home)
$L(e_2) = 2$	(ball arrives at home)
$L(e_3) = 3$	(batter hits ball to pitcher)
$L(e_4) = 4$	(batter runs to first base)
$L(e_5) = 1$	(runner runs to home)
$L(e_6) = 4$	(ball arrives at pitcher)
$L(e_7) = 5$	(pitcher throws ball to first base)
$L(e_8) = 5$	(runner arrives at home)
$L(e_9) = 6$	(ball arrives at first base)
$L(e_{10}) = 7$	(batter arrives at first base)

- For our example, Lamport's algorithm says that the run scores!

Total-order Lamport clocks

- Many systems require a total-ordering of events, not a partial-ordering
- Use Lamport's algorithm, but break ties using the process ID
  - $L(e) = \langle L_i(e), i \rangle$
    - $\langle L_i(e), i \rangle$ **<** $\langle L_j(e'), j \rangle$ if either
      - $L_i(e) < L_j(e')$, or
      - $L_i(e) = L_j(e')$ and $i < j$