

# Multiple Sequence Alignment

CMSC 423  
Carl Kingsford

```

+ -- 41 lines: Foreword here -----
  labelColor: "white",
  labelFont: "14px sans-serif",
  backgroundColor: "black",
  dotBarStyle: 'green', // 'rgba(255, 128, 128,
  dotFillStyle: 'rgba(150,150,150,0.7)',
  dotStrokeStyle: "rgb(100, 100, 100)",
  dotHighlightFill: "rgb(255, 237, 160)",
  dotHighlightStroke: "rgba(100,100,100,0.3)",
  dotSelectFill: "coral",
  dotSelectStroke: "rgba(100,100,100,0.3)",
  dotBorderWidth: 1,
  dotBarPaddingLeft: 0, // 10
  dotBarPaddingRight: 0, // 10
  height: 50,
  topIndent: 16,
  dotSize: 50,

```

```

+ --186 lines: zoomButtonScaleFactor: 2-----

```

```

var ticks = blah.d;

```

```

var getTimelineLabel = function (d, ticks) {
  var months = new Array('Jan', 'Feb', 'Mar', 'Apr')
  var index = Math.round((d.getTime() - selected
  var y = d.getFullYear();

```

```

  if (smallStep >= MILLIS_PER_YEAR) {
    if (smallStep == 10 * MILLIS_PER_YEAR) ret
    if (smallStep == 5 * MILLIS_PER_YEAR)
      if (ticks.length > 25) return y%10==0
      else return y%5==0 ? y : '';

```

```

+ -- 10 lines: ticks per year-----

```

```

    else return '';
    else
      if (index%4 == 0) return months[d.getM
      else return '';
  }

```

```

  else if (smallStep >= MILLIS_PER_WEEK) {
    if (ticks.length < 20)
      if (index == ticks.length-1 || index ==
        return months[d.getMonth()] + ' ' +
      else
        return months[d.getMonth()] + ' ' + d.
    else

```

```

      if (index%2 == 0)
        if (index >= ticks.length-2 || index =
          return months[d.getMonth()] + ' ' +
        else
          return months[d.getMonth()] + ' ' + d
      else return '';
  }

```

```

  else if (smallStep >= MILLIS_PER_DAY) return d
  else return d.getDate();
};

```

```

/**
 * custom zoom behavior when using the scroll
 */

```

```

var myzoom = function (e, w) {
  speed = 1/7; // 1/48;
  var obj = e;
  width = w;
  function mousewheel() {
    var m = this.mouse();

```

```

+ -- 41 lines: Foreword here -----
  labelColor: "white",
  labelFont: "14px sans-serif",
  backgroundColor: "black",
  dotBarStyle: 'green', // 'rgba(255, 128, 128
  dotFillStyle: 'rgba(150,150,150,0.7)',
  dotStrokeStyle: "rgb(100, 100, 100)",
  dotHighlightFill: "rgb(255, 0, 0)",
  dotHighlightStroke: "rgba(100,100,100,0.3)",

```

```

  dotBorderWidth: 1,
  dotBarPaddingLeft: 0, // 10
  dotBarPaddingRight: 0, // 10
  height: 50,
  topIndent: 16,
  dotSize: 50,

```

```

+ --186 lines: zoomButtonScaleFactor: 2-----

```

```

var ticks = blah.d;

```

```

var getTimelineLabel = function (d, ticks) {
  var months = new Array('Jan', 'Feb', 'Mar', 'Apr')
  var index = Math.round((d.getTime() - selecte
  var y = d.getFullYear();

```

```

  if (smallStep >= MILLIS_PER_YEAR) {
    if (smallStep == 10 * MILLIS_PER_YEAR) re
    if (smallStep == 5 * MILLIS_PER_YEAR)
      if (ticks.length > 25) return y%10==0
      else return y%5==0 ? y : '';

```

```

+ -- 10 lines: ticks per year-----

```

```

    else return '';
    else
      if (index%4 == 0) return months[d.get
      else return '';
  }

```

```

  else if (smallStep >= MILLIS_PER_WEEK) {
    if (ticks.length < 20)
      return months[d.getMonth()] + ' ' + d.g

```

```

    else
      if (index%2 == 0) return months[d.getMo

```

```

    else return '';
  }
  else if (smallStep >= MILLIS_PER_DAY) return
  else return d.getDate();
};

```

```

/**
 * custom zoom behavior when using the scrol
 */

```

```

var myzoom = function (e, w) {
  speed = 1/48;
  var obj = e;
  width = w;
  function mousewheel() {
    var m = this.mouse();

```

```
+ -- 41 lines: Foreword here -----
labelColor: "white",
labelFont: "14px sans-serif",
backgroundColor: "black",
dotBarStyle: 'green',//'rgba(255, 128, 128,
dotFillStyle: 'rgba(150,150,150,0.7)',
dotStrokeStyle: "rgb(100, 100, 100)",
dotHighlightFill: "rgb(255, 237, 160)",
dotHighlightStroke: "rgba(100,100,100,0.3)"
dotSelectFill: "coral",
dotSelectStroke: "rgba(100,100,100,0.3)",
dotBorderWidth: 1,
dotBarPaddingLeft: 0, // 10
dotBarPaddingRight: 0, // 10
height: 50,
topIndent: 16,
dotSize: 50,
```

```
+ --186 lines: zoomButtonScaleFactor: 2-----
```

```
var ticks = blah.d;
```

```
var getTimelineLabel = function (d, ticks) {
  var months = new Array('Jan','Feb','Mar','Apr')
  var index = Math.round((d.getTime() - selected
  var y = d.getFullYear();
```

```
  if (smallStep >= MILLIS_PER_YEAR ) {
    if (smallStep == 10 * MILLIS_PER_YEAR) ret
    if (smallStep == 5 * MILLIS_PER_YEAR)
      if (ticks.length > 25) return y%10==0
      else return y%5==0 ? y : '';
```

```
+ -- 10 lines: ticks per year-----
```

```
    else return '';
    else
      if (index%4 == 0) return months[d.getM
      else return '';
```

```
  }
  else if (smallStep >= MILLIS_PER_WEEK) {
    if (ticks.length < 20)
      if (index == ticks.length-1 || index ==
        return months[d.getMonth()] + ' ' +
      else
        return months[d.getMonth()] + ' ' + d.
    else
```

```
      if (index%2 == 0)
        if (index >= ticks.length-2 || index =
          return months[d.getMonth()] + ' ' +
        else
          return months[d.getMonth()] + ' ' + d
      else return '';
```

```
  }
  else if (smallStep >= MILLIS_PER_DAY) return d
  else return d.getDate();
};
```

```
/**
 * custom zoom behavior when using the scroll
 */
```

```
var myzoom = function (e, w) {
  speed = 1/7;// 1/48;
  var obj = e;
  width = w;
  function mousewheel() {
    var m = this.mouse();
```

```
+ -- 41 lines: Foreword here -----
labelColor: "white",
labelFont: "14px sans-serif",
backgroundColor: "black",
dotBarStyle: 'green',//'rgba(255, 128, 128
dotFillStyle: 'rgba(150,150,150,0.7)',
dotStrokeStyle: "rgb(100, 100, 100)",
dotHighlightFill: "rgb(255, 0, 0)",
dotHighlightStroke: "rgba(100,100,100,0.3)
```

```
dotBorderWidth: 1,
dotBarPaddingLeft: 0, // 10
dotBarPaddingRight: 0, // 10
height: 50,
topIndent: 16,
dotSize: 50,
```

```
+ --186 lines: zoomButtonScaleFactor: 2-----
```

```
var ticks = blah.d;
```

```
var getTimelineLabel = function (d, ticks) {
  var months = new Array('Jan','Feb','Mar','Apr')
  var index = Math.round((d.getTime() - selecte
  var y = d.getFullYear();
```

```
  if (smallStep >= MILLIS_PER_YEAR ) {
    if (smallStep == 10 * MILLIS_PER_YEAR) re
    if (smallStep == 5 * MILLIS_PER_YEAR)
      if (ticks.length > 25) return y%10==0
      else return y%5==0 ? y : '';
```

```
+ -- 10 lines: ticks per year-----
```

```
    else return '';
    else
      if (index%4 == 0) return months[d.get
      else return '';
```

```
  }
  else if (smallStep >= MILLIS_PER_WEEK) {
    if (ticks.length < 20)
      return months[d.getMonth()] + ' ' + d.g
```

```
    else
      if (index%2 == 0) return months[d.getMo
```

```
  }
  else if (smallStep >= MILLIS_PER_DAY) return
  else return d.getDate();
};
```

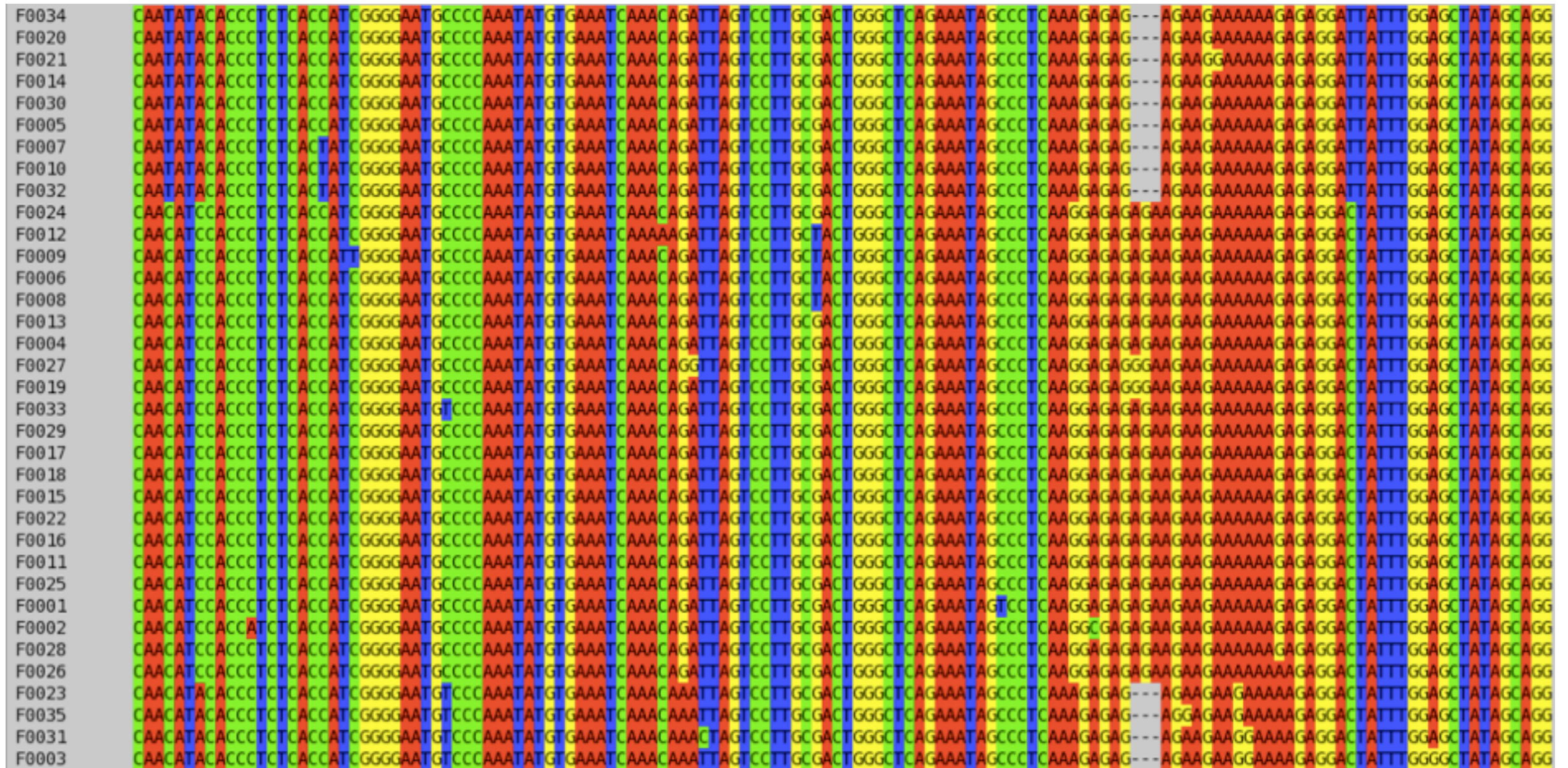
```
/**
 * custom zoom behavior when using the scrol
 */
```

```
var myzoom = function (e, w) {
  speed = 1/48;
  var obj = e;
  width = w;
  function mousewheel() {
    var m = this.mouse();
```

3rd source file here



# Multiple Sequence Alignment (MSA)



Multiple sequence alignment: find more subtle patterns & find common patterns between all sequence.



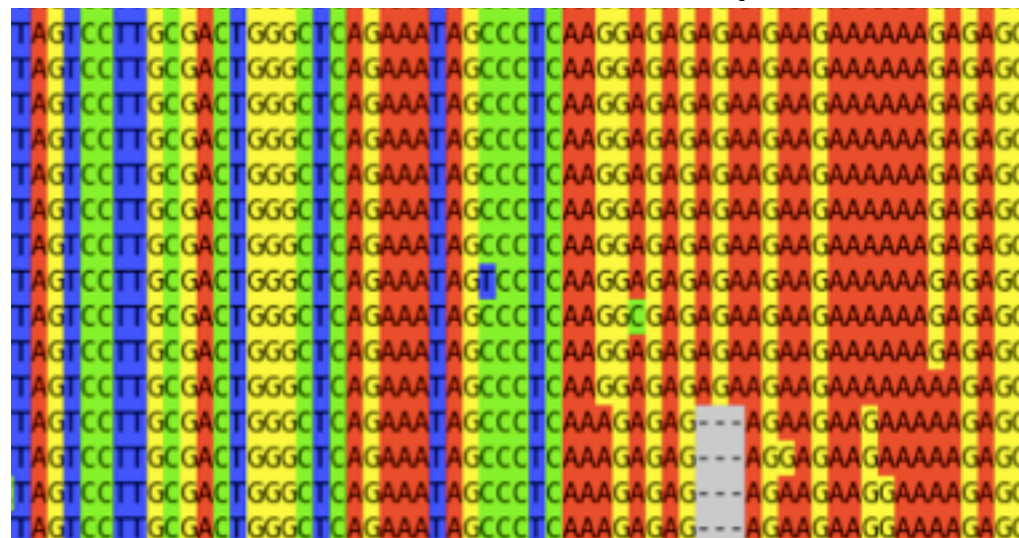
# Generalizing Alignment to $> 2$ Strings

Input: Sequences  $S_1, S_2, \dots, S_p$

Let  $\text{cost}(x_1, x_2, \dots, x_p)$  be a user-supplied function that computes the quality of a column: an alignment between characters  $x_1, x_2, \dots, x_p$ .

- **Goal:** find alignment  $M$  to **minimize**  $\sum$  cost of the columns:

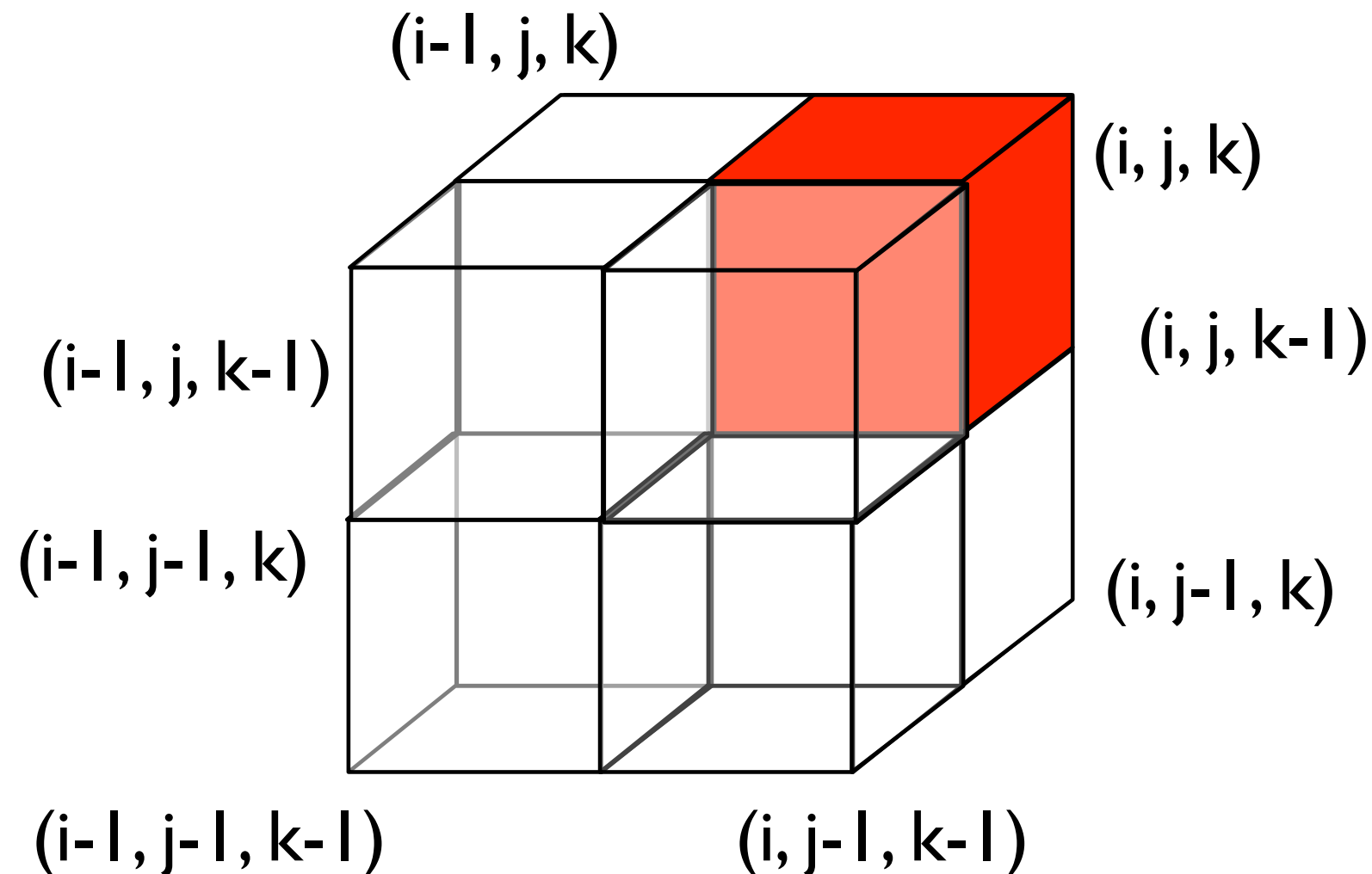
$$\text{cost}(x_1, x_2, \dots, x_p) = \text{cost}(\underbrace{\text{[red bar] [grey bar]}}_{\text{[red bar] [grey bar]}})$$



# Slow Dynamic Programming

Suppose you had just 3 sequences.

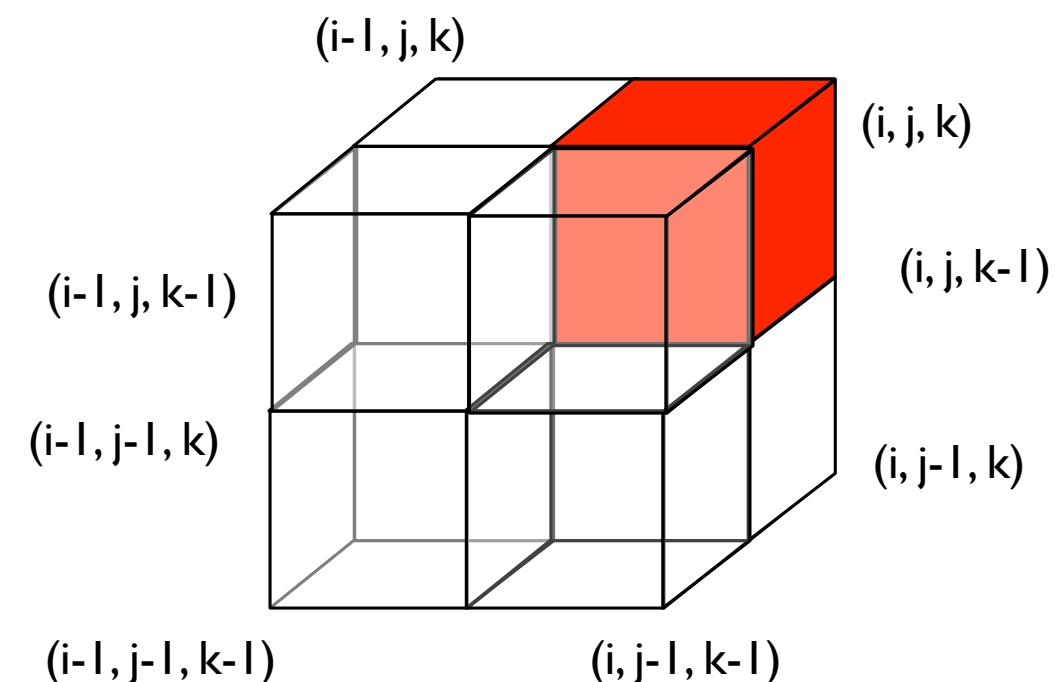
Apply the same DP idea as sequence alignment for 2 sequences, but now with a 3-dimensional matrix



# DP Recurrence for 3 sequences

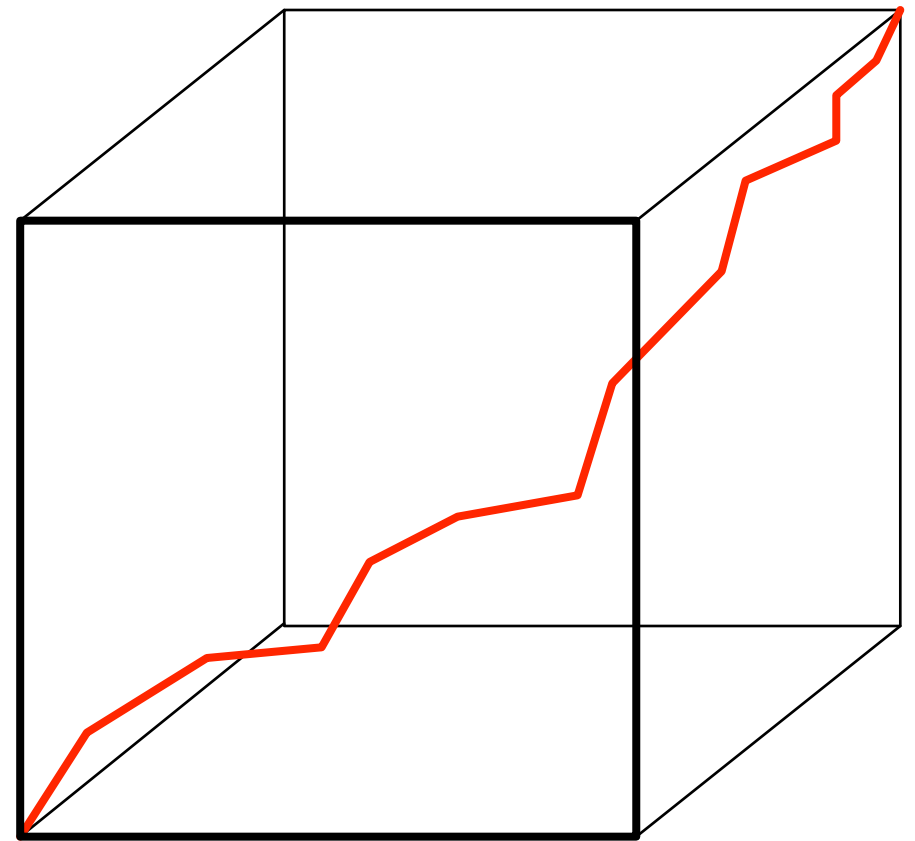
$$A[i, j, k] = \min \begin{cases} \text{cost}(x_i, y_j, z_k) + A[i-1, j-1, k-1] \\ \text{cost}(x_i, -, -) + A[i-1, j, k] \\ \text{cost}(x_i, y_j, -) + A[i-1, j-1, k] \\ \text{cost}(-, y_j, z_k) + A[i, j-1, k-1] \\ \text{cost}(-, y_j, -) + A[i, j-1, k] \\ \text{cost}(x_i, -, z_k) + A[i-1, j, k-1] \\ \text{cost}(-, -, z_k) + A[i, j, k-1] \end{cases}$$

Every possible pattern for the gaps.



# Running time

- $n^3$  subproblems, each takes  $2^3$  time  
 $\Rightarrow O(n^3)$  time.
- For  $p$  sequences:  $n^p$  subproblems,  
each takes  $2^p$  time for the max and  
 $p^2$  to compute `cost()`  $\Rightarrow O(p^2 n^p 2^p)$
- Even  $O(n^3)$  is often too slow for the  
length of sequences encountered in  
practice.
- One solution: approximation  
algorithm.



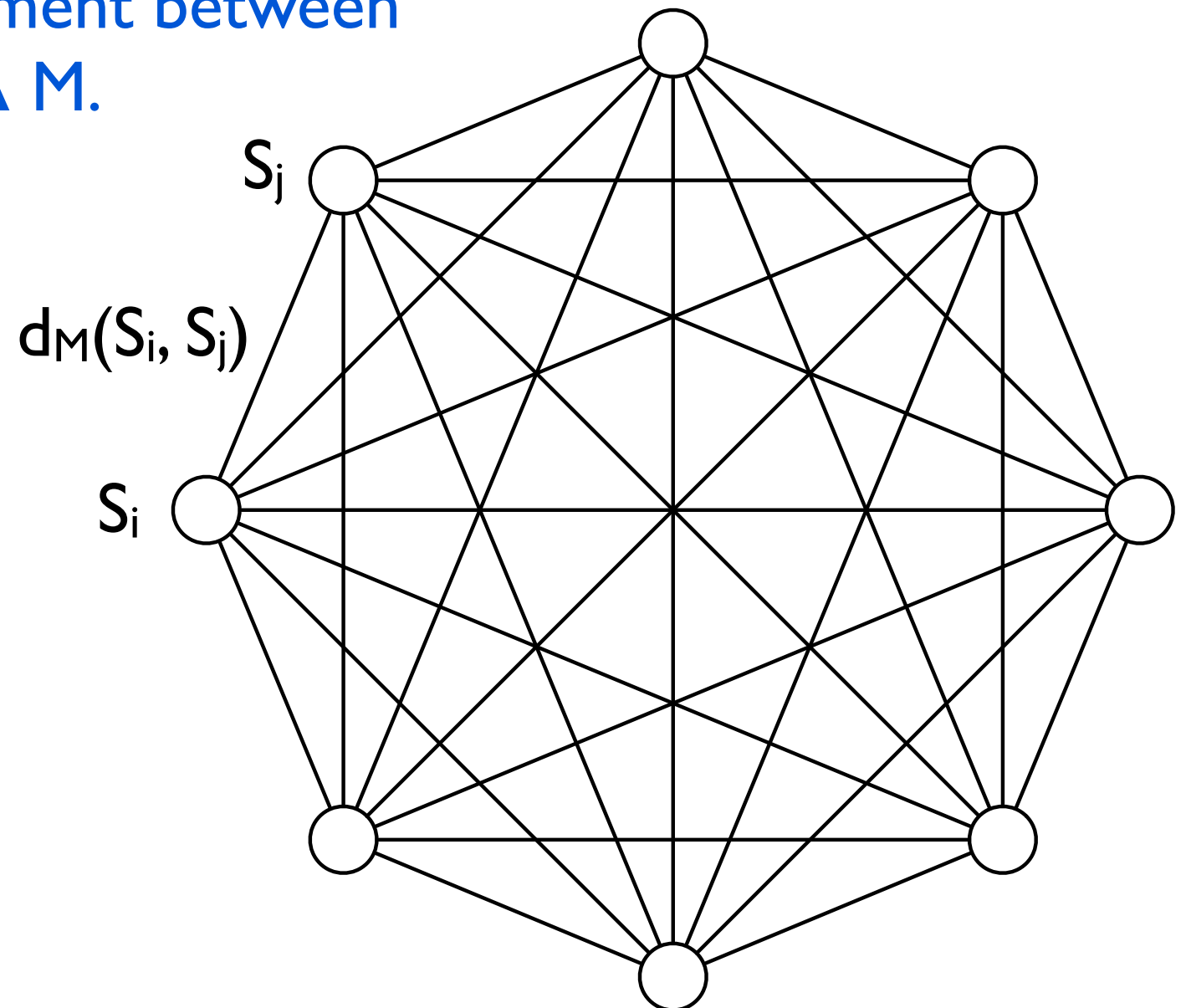


# SP-Score

A particular cost() function, the SP-Score, is commonly used and allows us to design an approximation algorithm for the MSA problem.

$d_M(S_i, S_j)$  = the cost of the alignment between  $S_i$  and  $S_j$  **as implied by MSA M**.

$SP\text{-Score}(M) = \sum_{i < j} d_M(S_i, S_j)$   
= sum of all the scores of the pairwise alignments implied by M.



# MSA

- A multiple sequence alignment (MSA) implies a pairwise alignment between every pair of sequences.
- This implied alignment need not be optimal, however:

match = -1, a mismatch = 1, gap = 2

Sequences: AT, A, T, AT, AT

Optimal MSA:

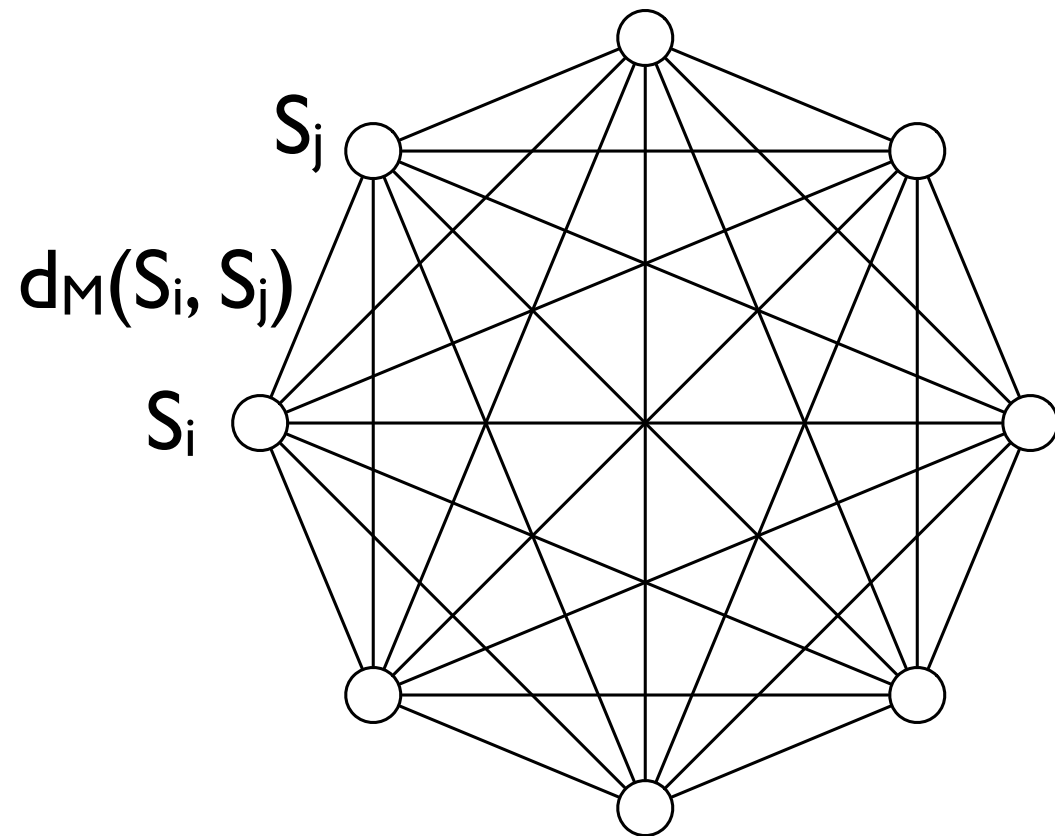
AT
A-
-T
AT
AT

+2 +2 = 4

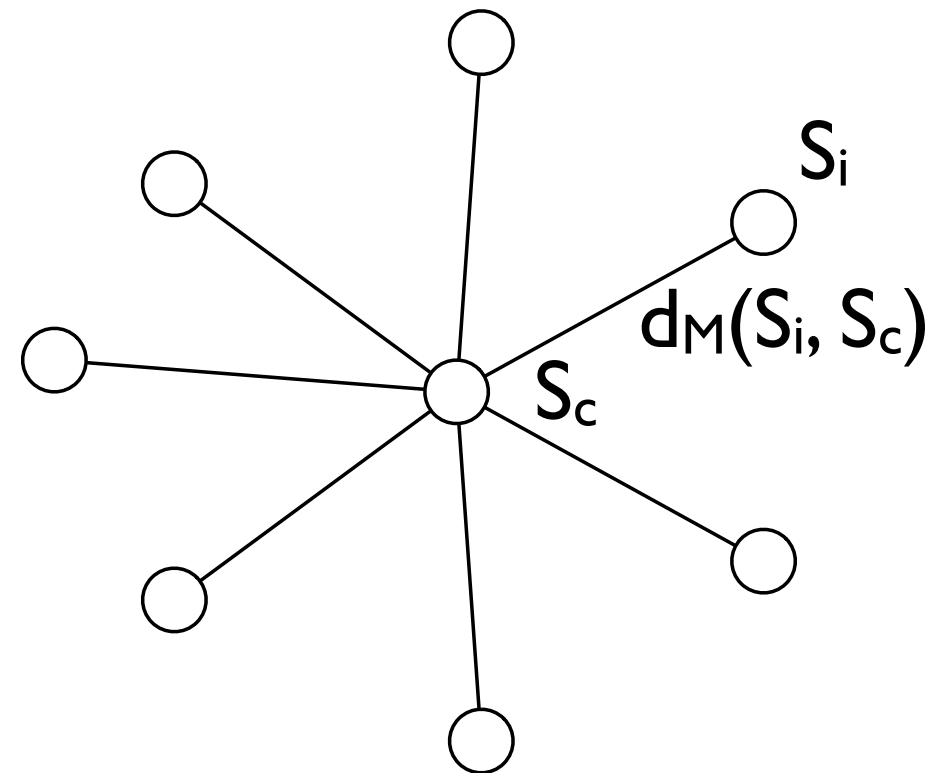
Optimal  
Alignment  
between  
A and T: A  
T  
+1

(A,A), (A,-), (A,A), (A,A), (A, -), (A,A), (A,A) (-,A), (-,A), (A,A)  
-1 + 2 -1 -1 +2 -1 -1 +2 +2 -1 = +2

# Star Alignment Approximation



SP-Score



Star-Score =

$$\sum_i d_M(S_i, S_c)$$



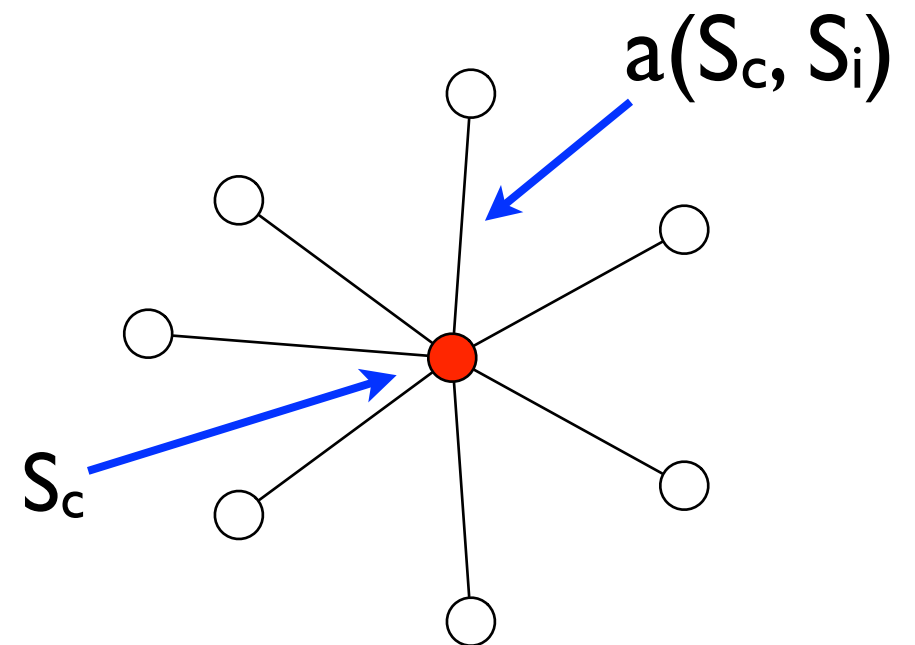
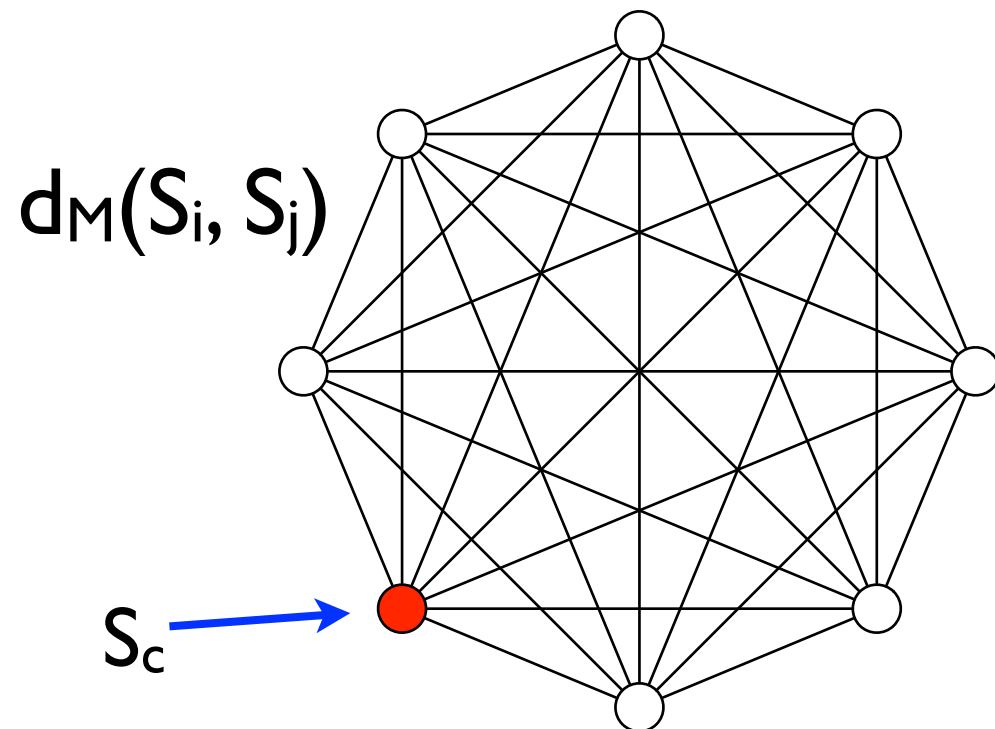
# Star Alignment Algorithm

**Input:** sequences  $S_1, S_2, \dots, S_p$

- Build all  $O(p^2)$  pairwise alignments.
- Let  $S_c$  = the sequence in  $S_1, S_2, \dots, S_p$  that is closest to the others.  
That is, choose  $S_c$  to minimize:

$$\sum_{i \neq c} a(S_c, S_i)$$

- *Progressively align* all other sequences to  $S_c$ .



# Progressive Alignment

- Build a multiple sequence alignment up from pairwise alignments.

Start with an alignment between  $S_c$  and some other sequence:

```
SC  YFPHFDSLHGSAQVKAHGKKVGDALTAVGHLDDLPGAL
S1  YFPHFDSLHG-AQVKG--KKVADALTNAVAHVDDMPNAL
```

Add 3rd sequence, say  $S_2$ , and use the  $SC - S_2$  alignment as a guide, adding spaces into the MSA as needed.

$SC - S_2$  alignment:

```
SC  YFPHF-DLS-----HGSAQVKAHGKKVGDALTAVGHL-----DDLPGAL
S2  FFPKFKGLTTADQLKKSADVVRWHAERII-----NAVNDAVASMDDEKMS
```

New  $\{SC, S_1, S_2\}$  alignment (**red** gaps added in  $S_1$ ):

```
SC  YFPHF-DLS-----HGSAQVKAHGKKVGDALTAVGHL-----DDLPGAL
S1  YFPHF-DLS-----HG-AQVKG--KKVADALTNAVAHV-----DDMPNAL
S2  FFPKFKGLTTADQLKKSADVVRWHAERII-----NAVNDAVASMDDEKMS
```

Continue with  $S_3, S_4, \dots$

# Performance

Assume the cost function satisfies the triangle inequality:

$$\text{cost}(x,y) \leq \text{cost}(x, z) + \text{cost}(z,y)$$

Example:  $\text{cost}(A, C) \leq \text{cost}(A, T) + \text{cost}(T, C)$

$\underbrace{\text{cost}(A, C)}_{\text{cost of 1 mutation from } A \rightarrow C} \leq \underbrace{\text{cost}(A, T) + \text{cost}(T, C)}_{\text{cost of a mutation from } A \rightarrow T \text{ and then from } T \rightarrow C}$

STAR = cost of result of star algorithm under SP-score

OPT = cost of optimal multiple sequence alignment (under SP-score)

**Theorem.** If cost satisfies the triangle inequality, then  $\text{STAR} \leq 2 \times \text{OPT}$ .

Example: if optimal alignment has cost 10, the star alignment will have cost  $\leq 20$ .



# Proof (1)

**Theorem.** If cost satisfies the triangle inequality, then  $\text{STAR} \leq 2\text{OPT}$ .

$$\frac{\text{STAR}}{\text{OPT}} \leq 2$$

For some  $B$  we will  
prove the 2 statements:

$$\begin{array}{l} \text{STAR} \leq 2B \\ \text{OPT} \geq B \end{array}$$

This will imply:

$$\Rightarrow \frac{\text{STAR}}{\text{OPT}} \leq \frac{2B}{B} = 2$$

## Proof (2)

**Theorem.** If cost satisfies the triangle inequality, then  $\text{STAR} \leq 2\text{OPT}$ .

$$2 \cdot \text{STAR} = \sum_{ij} d_{\text{STAR}}(S_i, S_j) \quad \text{defn of SP-score}$$

$$\begin{array}{l} \text{by triangle} \\ \text{inequality} \end{array} \leq \sum_{ij} (d_{\text{STAR}}(S_i, S_c) + d_{\text{STAR}}(S_c, S_j))$$

$$\begin{array}{l} \text{because STAR} \\ \text{alignment is optimal} \\ \text{for pairs involving } S_c \end{array} = \sum_{ij} (\mathbf{a}(S_i, S_c) + \mathbf{a}(S_c, S_j))$$

$$\begin{array}{l} \text{distribute } \sum \end{array} = \sum_{ij} \mathbf{a}(S_i, S_c) + \sum_{ij} \mathbf{a}(S_c, S_j)$$

$$\leq 2p \sum_i \mathbf{a}(S_i, S_c) \quad \begin{array}{l} \text{sums are the same} \\ \text{and each term appears} \\ \leq p \text{ (\# of sequences)} \\ \text{times.} \end{array}$$

# Proof (3)

**Theorem.** If cost satisfies the triangle inequality, then  $\text{STAR} \leq 2\text{OPT}$ .

$$2 \cdot \text{OPT} = \sum_{i,j} d_{\text{OPT}}(S_i, S_j) \quad \text{defn of SP-score}$$

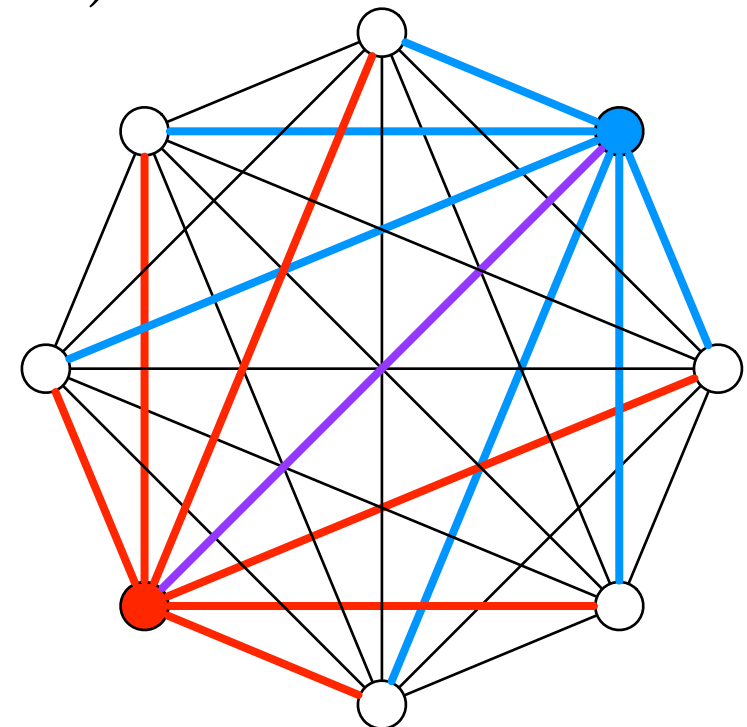
optimal pairwise alignment  
is  $\leq$  pairwise alignment  
induced by any MSA

$$\geq \sum_{i,j} a(S_i, S_j)$$

sum of cost of all pairwise  
alignments is = the sum of  $p$   
different stars.

$$\geq p \sum_i a(S_i, S_c)$$

We chose  $S_c$  because it was  
the lowest-cost star.





# End of Proof

For some  $B$  we will  
prove the 2 statements:

$$\begin{array}{l} \text{STAR} \leq 2B \\ \text{OPT} \geq B \end{array}$$

This will imply:

$$\Rightarrow \frac{\text{STAR}}{\text{OPT}} \leq \frac{2B}{B} = 2$$


$$2 \cdot \text{STAR} \leq 2p \sum_i \mathbf{a}(S_i, S_c)$$

$$2 \cdot \text{OPT} \geq p \sum_i \mathbf{a}(S_i, S_c)$$

$$\Rightarrow \frac{\text{STAR}}{\text{OPT}} \leq \frac{2p \sum_i \mathbf{a}(S_i, S_c)}{p \sum_i \mathbf{a}(S_i, S_c)} = 2$$

# Consensus Sequence

For every column  $j$ ,  
choose  $c \in \Sigma$  that  
minimizes  $\sum_i \text{cost}(c, S_i[j])$



(typically this means the  
most common letter)

```
S1 YFPHF-DLS-----HGSAQVKAHGKKVG-----DALTLAVAHLLDDLPGAL
S2 YFPHF-DLS-----HG-AQVKG-GKKVA-----DALTNAVAVHVD DMPNAL
S3 FFPKFKGLTTADQLKKSADV RWH AERII-----NAVND AVAS MDDTEKMS
S4 LFSFLKGTSEVP--QNNPELQAHAGKVFKLVYEAAIQ LQVTGVVVVTDATL
CO YFPHFKDLS-----HGSAQVKAHGKKVG-----DALTLAVAHVDDTPGAL
```

- Consensus is a summarization of the whole alignment.
- Consensus sequence is sometimes used as an estimate for the ancestral sequence.
- Sometimes the MSA problem is formulated as: find MSA  $M$  that minimizes:  $\sum_i d_M(\text{CO}, S_i)$

# Profiles

- Another way to summarize an MSA:

S1 ACG-TT-GA  
S2 ATC-GTCGA  
S3 ACGCGA-CC  
S4 ACGCGT-TA

Column in the alignment

Character

	1	2	3	4	5	6	7	8	9
A	1	0	0	0	0	0.25	0	0	0.75
C	0	0.75	0.25	0.5	0	0	0.25	0.25	0.25
G	0	0	0.75	0	0.75	0	0	0.5	0
T	0	0.25	0	0	0.25	0.75	0	0.25	0
-	0	0	0	0.5	0	0	0.75	0	0

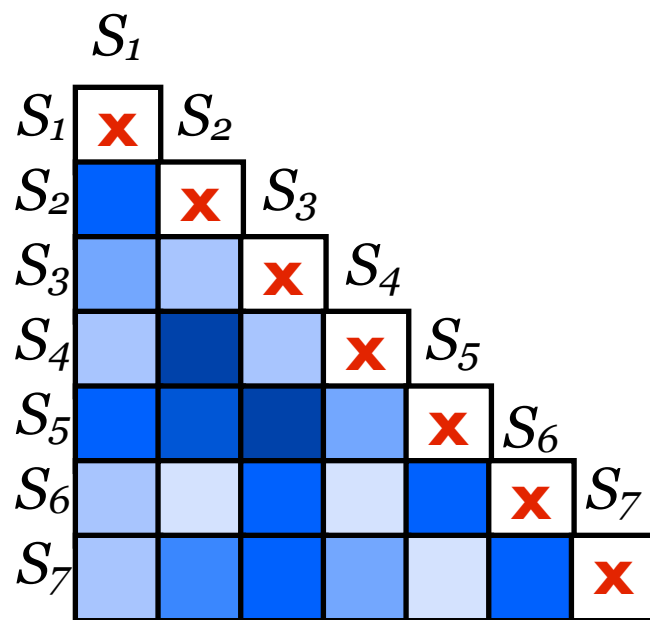
Call this profile  
matrix R

Fraction of time  
given column had  
the given character

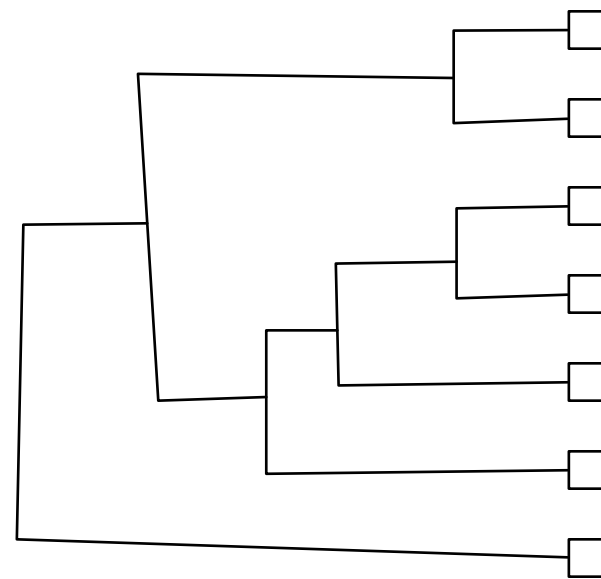
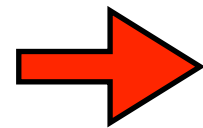


# CLUSTLW

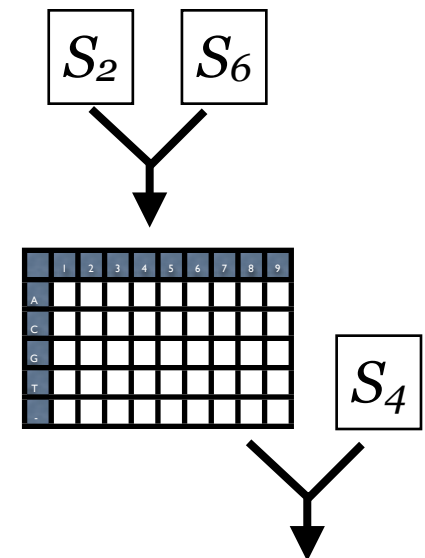
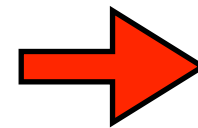
- CLUSTLW is a widely used, “classical” heuristic multiple aligner.
- Not the fastest, not the most accurate, but pretty good.
- Large # of heuristic tricks included in the software, but basic idea is straightforward:



Step (1): Build  
pairwise distance  
matrix



Step (2): Build guide  
tree



Step (3): Align  
sequences / **sets of  
sequences** from  
the most similar to  
least similar

# Profile-based Alignment

gap in profile  
introduced to  
better fit sequence

R =

	1	2	3	4
A	1	0	0	0
C	0	0.75	0.25	0.5
G	0	0	0.75	0
T	0	0.25	0	0
-	0	0	0	0.5

5	6	7	8	9
0	0.25	0	0	0.75
0	0	0.25	0.25	0.25
0.75	0	0	0.5	0
0.25	0.75	0	0.25	0
0	0	0.75	0	0

A C C - A G A C G A

Score of matching character  $x$  with  
column  $j$  of the profile:

$$P(x, j) = \sum_{c \in \Sigma} \text{sim}(x, c) \times R[c, j]$$

$\text{sim}(x, c)$  = how similar character  $x$  is  
to character  $c$ .

$$A[i, j] = \max \begin{cases} A[i-1, j-1] + P(x_i, j) & \text{align } x_i \text{ to column } j \\ A[i-1, j] + \text{gap} & \text{introduce gap into profile} \\ A[i, j-1] + P("-", j) & \text{introduce gap into } x \end{cases}$$

# Recap

- Multiple sequence alignments (MSAs) are a fundamental tool. They help reveal subtle patterns, compute consistent distances between sequences, etc.
- Quality of MSAs often measured using the SP-score: sum of the scores of the pairwise alignments implied by the MSA.
- Same DP idea as pairwise alignment leads to exponentially slow algorithm for MSA for general  $p$ .
- 2-approximation obtainable via star alignments.
- MSAs often used to create profiles summarizing a family of sequences.