## Local Alignment

CMSC 423

## Representing edits as alignments

```
prin-ciple
|||| |||xx
prinncipal
(1 gap, 2 mm)
```

misspell

mis-pell
(1 gap)
aa-bb-ccaabb

ababbbc-a-b-
(5 gaps, 1 mm )
prin-cip-le

prinncipal-
(3 gaps, 0 mm )
prehistoric

---historic
(3 gaps)
al-go-rithm-
|| xx ||x |
alKhwariz-mi
(4 gaps, 3 mm )

## Maximization vs. Minimization

## Edit distance:

$$
O P T(i, j)=\min _{\{ } \begin{cases}\operatorname{cost}\left(a_{i}, b_{j}\right)+O P T(i-1, j-1) & \text { match } a_{i}, b_{j} \\ \operatorname{gap}+O P T(i-1, j) & a_{i} \text { is not matched } \\ \operatorname{gap}+O P T(i, j-1) & b_{j} \text { is not matched }\end{cases}
$$

Sequence Similarity: replace min with a max and negate the parameters.
gap penalty $\rightarrow$ gap benefit (probably negative)
cost $\rightarrow$ score


## Local Alignment

Local alignment between $\mathbf{s}$ and $\mathbf{t}$ : Best alignment between a subsequence of $s$ and a subsequence of $t$.


Motivation:
Many genes are composed of domains, which are subsequences that perform a particular function.

## Recall: Global Alignment Matrix

$\operatorname{OPT}(i, j)$ contains the score for the best alignment between:
the first $i$ characters of string $x$ [prefix $i$ of $x$ ]
the first $j$ character of string $y$ [prefix $j$ of $y$ ]


## Local Alignment

New meaning of entry of matrix entry:
$\mathrm{A}[\mathrm{i}, \mathrm{j}]=$ best score between: some suffix of $x[1 . . . i]$ some suffix of $y[\mid \ldots j]$


## How do we fill in the local alignment matrix?

$A[i, j]=\max \left\{\begin{array}{l}A[i, j-1]+\operatorname{gap} \quad \text { (I) } \\ A[i-1, j]+\operatorname{gap} \quad(2) \\ A[i-1, j-1]+\operatorname{match}(i, j) \\ 0\end{array}\right.$
(I), (2), and (3): same cases as before: gap in $x$, gap in $y$, match $x$ and $y$

New case: 0 allows you to say the best alignment between a suffix of $x$ and a suffix of $y$ is the empty alignment.


## Local Alignment

- Initialize first row and first column to be 0 .
- The score of the best local alignment is the largest value in the entire array.
- To find the actual local alignment:
- start at an entry with the maximum score
- traceback as usual
- stop when we reach an entry with a score of 0


## Local Alignment Python Code

```
def local_align(x, y, score=ScoreParam(-7, 10, -5)):
    """Do a local alignment between x and y"""
    # create a zero-filled matrix
    A = make_matrix(len(x) + 1, len(y) + 1)
    best = 0
    optloc = (0,0)
    # fill in A in the right order
    for i in xrange(1, len(y)):
        for j in xrange(1, len(x)):
            # the local alignment recurrance rule:
            A[i][j] = max(
            A[i][j-1] + score.gap,
            A[i-1][j] + score.gap,
            A[i-1][j-1] + (score.match if x[i] == y[j] else score.mismatch),
            O
            )
            # track the cell with the largest score
            if A[i][j] >= best:
                best = A[i][j]
                optloc = (i,j)
    # return the opt score and the best location
    return best, optloc
```


## Local Alignment Python Code

```
def make_matrix(sizex, sizey):
    """Creates a sizex by sizey matrix filled with zeros."""
    return [[0]*sizey for i in xrange(sizex)]
class ScoreParam:
    """The parameters for an alignment scoring function"""
    def __init__(self, gap, match, mismatch):
    self.gap = gap
    self.match = match
    self.mismatch = mismatch
```


## Local Alignment Example \#1

| local |  |  |  | CGTA |  | CTC | C" |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | * | A | G | C | G | T | A | G |
| * | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| C | 0 | 0 | 0 | 10 | 3 | 0 | 0 | 0 |
| T | 0 | 0 | 0 | 3 | 5 | 13 | 6 | 0 |
| C | 0 | 0 | 0 | R10 | 3 | 6 | 8 | 1 |
| G | 0 | 0 | 10 | 3 | - 20 | 13 | 6 | 18 |
| T | 0 | 0 | 3 | 5 | $13^{\prime}$ | 30 | 23 | 16 |
| C | 0 | 0 | 0 | 13 | 6 | 23 | 25 | 18 |

Score(match) = 10
Score(mismatch) $=-5$
Score(gap) $=-7$

Note: this table written top-to-bottom instead of bottom-to-top

## Local Alignment Example \#2



```
Score(match) = 10
Score(mismatch) = -5
Score(gap) = -7
```

Note: this table written top-to-bottom
instead of bottom-to-top

## Local Alignment Example \#2

|  | * | b | e | s | t | $\bigcirc$ | f | t | i | m | e | s |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| * | 0 | 0 |  | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| s | 0 | 0 | 0 | 10 | $-3$ | 0 | 0 | 0 | 0 | 0 | 0 | 10 |
| - | 0 | 0 | 0 | 3 | 5 | 13 | 6 | 0 | 0 | 0 | 0 | 3 |
| f | 0 | 0 | 0 | , | 0 | 6 |  | 16 | 9 | 2 | 0 | 0 |
| t | 0 | 0 | 0 | 0 | 10 | 3 | 16 | 33 | 26 | 19 | 12 | 5 |
| e | 0 | 0 | 10 | 3 | 3 | 5 | 9 | 26 | 28 | 21 | 29 | 22 |
| n | 0 | 0 | 3 | 5 | 0 | 0 | 2 | 19 | 21 | 23 | 22 | 2 |

```
Score(match) = 10
Score(mismatch) = -5
Score(gap) = -7
```

Note: this table written top-to-bottom instead of bottom-to-top

## More Local Alignment Examples $\begin{gathered}\text { Score }(\text { match })=10 \\ \text { Score (mismatch) } \\ \text { Score }(\text { gap })=-7\end{gathered}=-5$

| loca |  |  |  |  | sh |  | " |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | * | c | a | t | d | - | g | f | i | s | h |
| * | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| d | 0 | 0 | 0 | 0 | 10 | 3 | 0 | 0 | 0 | 0 | 0 |
| $\bigcirc$ | 0 | 0 | 0 | 0 | 3 | 20 | 13 | 6 | 0 | 0 | 0 |
| g | 0 | 0 | 0 | 0 | 0 | 13 | 30 | 23 | 16 | 9 | 2 |



| local_align( "aaaa", |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
|  |  |  |  |  | a |
| * | 0 | 0 | 0 | 0 | 0 |
| a | 0 | 10 | 10 | 10 | 10 |
| a | 0 | 10 | 20 | 20 | 20 |

## Upmost and Downmost Alignments

When there are ties in the max\{\}, we have a choice about which arrow to follow.


If we prefer arrows higher in the matrix, we get the upmost alignment.

If we prefer arrows lower in the matrix, we get the downmost alignment.


## Local / Global Recap

- Alignment score sometimes called the "edit distance" between two strings.
- Edit distance is sometimes called Levenshtein distance.
- Algorithm for local alignment is sometimes called "SmithWaterman"
- Algorithm for global alignment is sometimes called "NeedlemanWunsch"
- Same basic algorithm, however.
- Underlies BLAST

