



CMU SCS

**Carnegie Mellon Univ.
School of Computer Science
15-415 - Database Applications**

C. Faloutsos
Lecture #4: *Relational Algebra*



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Overview

- history
- concepts
- Formal query languages
 - relational algebra
 - rel. tuple calculus
 - rel. domain calculus

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History

- before: records, pointers, sets etc
- introduced by E.F. Codd in 1970
- revolutionary!
- first systems: 1977-8 (System R; Ingres)
- Turing award in 1981

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Concepts - reminder

- Database: a set of relations (= tables)
- rows: tuples
- columns: attributes (or keys)
- superkey, candidate key, primary key

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Example

Database:

STUDENT		
Ssn	Name	Address
123	smith	main str
234	jones	forbes ave

SSN	c-id	grade
123	15-413	A
234	15-413	B

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Example: cont'd

Database:

k-th attribute
(D_k domain)

↓

STUDENT		
Ssn	Name	Address
123	smith	main str
234	jones	forbes ave

rel. schema (attr+domains)

← tuple

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Example: cont'd

STUDENT		
Ssn	Name	Address
123	smith	main str
234	jones	forbes ave

rel. schema (attr+domains)

↕ instance

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Example: cont'd

• D_i : the domain of the i -th attribute (eg., $\text{char}(10)$)

STUDENT		
Ssn	Name	Address
123	smith	main str
234	jones	forbes ave

rel. schema (attr+domains)

↕ instance

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Overview

- history
- concepts
- **Formal query languages**
 - relational algebra
 - rel. tuple calculus
 - rel. domain calculus

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Formal query languages

- How do we collect information?
- Eg., find ssn's of people in 415
- (recall: everything is a set!)
- One solution: Rel. algebra, ie., set operators
- Q1: Which ones??
- Q2: what is a minimal set of operators?

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Relational operators

- .
- .
- .
- set union $\cup$
- set difference $-$

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Example:

- Q: find all students (part or full time)
- A: PT-STUDENT union FT-STUDENT

FT-STUDENT		
Ssn	Name	
129	peters	main str
239	lee	5th ave

PT-STUDENT		
Ssn	Name	Address
123	smith	main str
234	jones	forbes ave

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Observations:

- two tables are ‘union compatible’ if they have the same attributes (‘domains’)
- Q: how about intersection $\cap$

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Observations:

- A: redundant:
- STUDENT intersection STAFF =

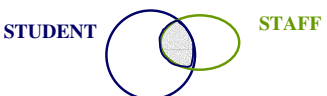


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Observations:

- A: redundant:
- STUDENT intersection STAFF =



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Observations:

- A: redundant:
- STUDENT intersection STAFF =
 $\text{STUDENT} - (\text{STUDENT} - \text{STAFF})$

STUDENT  STAFF

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Observations:

- A: redundant:
- STUDENT intersection STAFF =
 $\text{STUDENT} - (\text{STUDENT} - \text{STAFF})$

Double negation:
We'll see it again, later...

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Relational operators

- .
- .
- .
- set union $\cup$
- set difference $-$

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Other operators?

- eg, find all students on 'Main street'
- A: 'selection'

$$\sigma_{\text{address}='main\ str'}(STUDENT)$$

STUDENT		
Ssn	Name	Address
123	smith	main str
234	jones	forbes ave

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Other operators?

- Notice: selection (and rest of operators) expect tables, and produce tables (-> can be cascaded!!)
- For selection, in general:

$$\sigma_{\text{condition}}(RELATION)$$

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Selection - examples

- Find all 'Smiths' on 'Forbes Ave'

$$\sigma_{\text{name}='Smith' \wedge \text{address}='Forbes\ ave'}(STUDENT)$$

'condition' can be any boolean combination of '=', '>', '>=', ...

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Relational operators

- selection $\sigma_{condition}(R)$
- .
- .
- set union $R \cup S$
- set difference $R - S$

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Relational operators

- selection picks rows - how about columns?
- A: 'projection' - eg.: $\pi_{ssn}(STUDENT)$

finds all the 'ssn' - removing duplicates

STUDENT		
Ssn	Name	Address
123	smith	main str
234	jones	forbes ave

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Relational operators

Cascading: 'find ssn of students on 'forbes ave'

$$\pi_{ssn}(\sigma_{address='forbes\ ave'}(STUDENT))$$

STUDENT		
Ssn	Name	Address
123	smith	main str
234	jones	forbes ave

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Relational operators

- selection $\sigma_{condition}(R)$
- projection $\pi_{att-list}(R)$
- .
- set union $R \cup S$
- set difference $R - S$

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Relational operators

Are we done yet?

Q: Give a query we can **not** answer yet!

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Relational operators

A: any query across **two** or more tables,
eg., 'find names of students in 15-415'

Q: what extra operator do we need??

STUDENT		
Ssn	Name	Address
123	smith	main str
234	jones	forbes ave

TAKES		
SSN	c-id	grade
123	15-413	A
234	15-413	B

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Relational operators

A: any query across **two** or more tables,
eg., 'find names of students in 15-415'

Q: what extra operator do we need??

A: surprisingly, cartesian product is enough!

STUDENT		
Ssn	Name	Address
123	smith	main str
234	jones	forbes ave

TAKES		
SSN	c-id	grade
123	15-413	A
234	15-413	B

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Cartesian product

- eg., dog-breeding: MALE x FEMALE
- gives all possible couples

MALE	
name	
spike	
spot	

x

FEMALE	
name	
lassie	
shiba	

=

M.name	F.name
spike	lassie
spike	shiba
spot	lassie
spot	shiba

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so what?

- Eg., how do we find names of students taking 415?

STUDENT		
Ssn	Name	Address
123	smith	main str
234	jones	forbes ave

SSN	c-id	grade
123	15-415	A
234	15-413	B

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Cartesian product

• A:

$\sigma_{STUDENT.ssn=TAKES.ssn}(STUDENT \times TAKES)$

Ssn	Name	Address	ssn	cid	grade
123	smith	main str	123	15-415	A
234	jones	forbes ave	123	15-415	A
123	smith	main str	234	15-413	B
234	jones	forbes ave	234	15-413	B

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Cartesian product

$\sigma_{cid=15-415}(\sigma_{STUDENT.ssn=TAKES.ssn}(STUDENT \times TAKES))$

Ssn	Name	Address	ssn	cid	grade
123	smith	main str	123	15-415	A
234	jones	forbes ave	123	15-415	A
123	smith	main str	234	15-413	B
234	jones	forbes ave	234	15-413	B

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$\pi_{name}(\sigma_{cid=15-415}(\sigma_{STUDENT.ssn=TAKES.ssn}(STUDENT \times TAKES)))$

Ssn	Name	Address	ssn	cid	grade
123	smith	main str	123	15-415	A
234	jones	forbes ave	123	15-415	A
123	smith	main str	234	15-413	B
234	jones	forbes ave	234	15-413	B

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FUNDAMENTAL Relational operators

- selection $\sigma_{condition}(R)$
- projection $\pi_{att-list}(R)$
- cartesian product MALE x FEMALE
- set union $R \cup S$
- set difference $R - S$

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Relational ops

- Surprisingly, they are enough, to help us answer almost any query we want!!
- derived/convenience operators:
 - set intersection
 - **join** (theta join, equi-join, natural join) $\bowtie$
 - ‘rename’ operator $\rho_{R'}(R)$
 - division $R \div S$

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Joins

- Equijoin: $R \bowtie_{R.a=S.b} S = \sigma_{R.a=S.b}(R \times S)$

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Cartesian product

- $A: \dots \sigma_{STUDENT.ssn=TAKES.ssn} (STUDENT \times TAKES)$

Ssn	Name	Address	ssn	cid	grade
123	smith	main str	123	15-415	A
234	jones	forbes ave	123	15-415	A
123	smith	main str	234	15-413	B
234	jones	forbes ave	234	15-413	B

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Joins

- Equijoin: $R \bowtie_{R.a=S.b} S = \sigma_{R.a=S.b} (R \times S)$
- theta-joins: $R \bowtie_{\theta} S$
generalization of equi-join - any condition θ

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Joins

- **very** popular: natural join: $R \bowtie S$
- like equi-join, but it drops duplicate columns:
STUDENT (ssn, name, address)
TAKES (ssn, cid, grade)

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Joins

- **nat. join** has 5 attributes $STUDENT \bowtie TAKES$

Ssn	Name	Address	ssn	cid	grade
123	smith	main str	123	15-415	A
234	jones	forbes ave	123	15-415	A
123	smith	main str	234	15-413	B
234	jones	forbes ave	234	15-413	B

equi-join: 6 $STUDENT \bowtie_{STUDENT.ssn=TAKES.ssn} TAKES$

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Natural Joins - nit-picking

- if no attributes in common between R, S:
nat. join $\rightarrow$ cartesian product

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Overview - rel. algebra

- fundamental operators
- derived operators
 - joins etc
 - **rename**
 - division
- examples

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Rename op.

- Q: why? ρ_{AFTER} (BEFORE)
- A: shorthand; self-joins; ...
- for example, find the grand-parents of 'Tom', given PC (parent-id, child-id)

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Rename op.

- PC (parent-id, child-id) $PC \bowtie PC$

PC	
p-id	c-id
Mary	Tom
Peter	Mary
John	Tom

PC	
p-id	c-id
Mary	Tom
Peter	Mary
John	Tom

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Rename op.

- first, WRONG attempt:
 $PC \bowtie PC$
- (why? how many columns?)
- Second WRONG attempt:
 $PC \bowtie_{PC.c-id=PC.p-id} PC$

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Rename op.

- we clearly need two different names for the same table - hence, the 'rename' op.

$$\rho_{PC1}(PC) \bowtie_{PC1.c=id=PC.p-id} PC$$

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Overview - rel. algebra

- fundamental operators
- derived operators
 - joins etc
 - rename
 - division**
- examples

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Division

- Rarely used, but powerful.
- Example: find suspicious suppliers, ie., suppliers that supplied **all** the parts in A_BOMB

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Division

SHIPMENT	
s#	p#
s1	p1
s2	p1
s1	p2
s3	p1
s5	p3

÷

ABOMB
p#
p1
p2

=

BAD_S
s#
s1

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Division

- Observations: ~reverse of cartesian product
- It can be derived from the 5 fundamental operators (!!)
- How?

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Division

- Answer:

$$r \div s = \pi_{(R-S)}(r) - \pi_{(R-S)}[(\pi_{(R-S)}(r) \times s) - r]$$

- Observation: find ‘good’ suppliers, and subtract! (**double negation**)

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Division

SHIPMENT

s#	p#
s1	p1
s2	p1
s1	p2
s3	p1
s5	p3

÷

ABOMB

p#
p1
p2

=

BAD S

s#
s1

• Answer:

$$r \div s = \pi_{(R-S)}(r) - \pi_{(R-S)}[(\pi_{(R-S)}(r) \times s) - r]$$

• Observation: find ‘good’ suppliers, and subtract! (**double negation**)

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Division

SHIPMENT

s#	p#
s1	p1
s2	p1
s1	p2
s3	p1
s5	p3

÷

ABOMB

p#
p1
p2

=

BAD S

s#
s1

• Answer:

$$r \div s = \pi_{(R-S)}(r) - \pi_{(R-S)}[(\pi_{(R-S)}(r) \times s) - r]$$

↔

↔

All suppliers

All bad parts

• Observation: find ‘good’ suppliers, and subtract! (**double negation**)

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Division

SHIPMENT

s#	p#
s1	p1
s2	p1
s1	p2
s3	p1
s5	p3

÷

ABOMB

p#
p1
p2

=

BAD S

s#
s1

• Answer:

$$r \div s = \pi_{(R-S)}(r) - \pi_{(R-S)}[(\pi_{(R-S)}(r) \times s) - r]$$

↔

↔

all possible suspicious shipments

• Observation: find ‘good’ suppliers, and subtract! (**double negation**)

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Division

• Answer:

SHIPMENT	
s#	p#
s1	p1
s2	p1
s1	p2
s3	p1
s5	p3

÷

ABOMB	
p#	
p1	
p2	

=

BAD S	
s#	
s1	

$$r \div s = \pi_{(R-S)}(r) - \pi_{(R-S)}[(\pi_{(R-S)}(r) \times s) - r]$$

all possible
suspicious shipments
that didn't happen

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Division

• Answer:

SHIPMENT	
s#	p#
s1	p1
s2	p1
s1	p2
s3	p1
s5	p3

÷

ABOMB	
p#	
p1	
p2	

=

BAD S	
s#	
s1	

$$r \div s = \pi_{(R-S)}(r) - \pi_{(R-S)}[(\pi_{(R-S)}(r) \times s) - r]$$

all suppliers who missed
at least one suspicious shipment,
i.e.: 'good' suppliers

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Overview - rel. algebra

- fundamental operators
- derived operators
 - joins etc
 - rename
 - division
- examples

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Sample schema

find names of students that take 15-415

STUDENT		
Ssn	Name	Address
123	smith	main str
234	jones	forbes ave

CLASS		
c-id	c-name	units
15-413	s.e.	2
15-412	o.s.	2

TAKES		
SSN	c-id	grade
123	15-413	A
234	15-413	B

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Examples

- find names of students that take 15-415

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Examples

- find names of students that take 15-415

$$\pi_{name}[\sigma_{c-id=15-415}(STUDENT \bowtie TAKES)]$$

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Sample schema

find course names of 'smith'

STUDENT		
Ssn	Name	Address
123	smith	main str
234	jones	forbes ave

CLASS		
c-id	c-name	units
15-413	s.e.	2
15-412	o.s.	2

TAKES		
SSN	c-id	grade
123	15-413	A
234	15-413	B

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Examples

- find course names of 'smith'

$$\pi_{c-name} [\sigma_{name='smith'} ($$

$STUDENT \bowtie TAKES \bowtie CLASS$

$)]$

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Examples

- find ssn of 'overworked' students, ie., that take 412, 413, 415

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Examples

- find ssn of 'overworked' students, ie., that take 412, 413, 415: almost correct answer:

$$\sigma_{c-name=412}(TAKES) \cap$$

$$\sigma_{c-name=413}(TAKES) \cap$$

$$\sigma_{c-name=415}(TAKES)$$


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Examples

- find ssn of 'overworked' students, ie., that take 412, 413, 415 - Correct answer:

$$\pi_{ssn}[\sigma_{c-name=412}(TAKES)] \cap$$

$$\pi_{ssn}[\sigma_{c-name=413}(TAKES)] \cap$$

$$\pi_{ssn}[\sigma_{c-name=415}(TAKES)]$$

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Examples

- find ssn of students that work at least as hard as ssn=123, ie., they take all the courses of ssn=123, and maybe more

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Sample schema

STUDENT		
Ssn	Name	Address
123	smith	main str
234	jones	forbes ave

CLASS		
c-id	c-name	units
15-413	s.e.	2
15-412	o.s.	2

TAKES		
SSN	c-id	grade
123	15-413	A
234	15-413	B

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Examples

- find ssn of students that work at least as hard as ssn=123 (ie., they take all the courses of ssn=123, and maybe more)

$$[\pi_{ssn, c-id}(TAKES)] \div \pi_{c-id}[\sigma_{ssn=123}(TAKES)]$$

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Conclusions

- Relational model: only tables ('relations')
- relational algebra: powerful, minimal: 5 operators can handle almost any query!

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