## 15-826: Multimedia Databases and

 Data MiningLecture \#26: Graph mining - patterns
Christos Faloutsos

Carnegie Nellon

## Must-read Material

- Michalis Faloutsos, Petros Faloutsos and Christos Faloutsos, On Power-Law Relationships of the Internet Topology, SIGCOMM 1999.
- R. Albert, H. Jeong, and A.-L. Barabasi, Diameter of the World Wide Web Nature, 401, 130-131 (1999).
- Reka Albert and Albert-Laszlo Barabasi Statistical mechanics of complex networks, Reviews of Modern Physics, 74, 47 (2002)
- Jure Leskovec, Jon Kleinberg, Christos Faloutsos Graphs over Time: Densification Laws, Shrinking Diameters and Possible Explanations, KDD 2005, Chicago, IL, USA


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## Must-read Material (cont'd)

- D. Chakrabarti and C. Faloutsos, Graph Mining: Laws, Generators and Algorithms, in ACM Computing Surveys, 38 (1), 2006
- J. Leskovec, D. Chakrabarti, J. Kleinberg, and C. Faloutsos, Realistic, Mathematically Tractable Graph Generation and Evolution, Using Kronecker Multiplication, in PKDD 2005, Porto, Portugal

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## Cannegievelolon

## Cannegievelelon

Graphs - why should we care?

- IR: bi-partite graphs (doc-terms)

- web: hyper-text graph $\qquad$
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Graphs - why should we care?

- 'viral' marketing
$\qquad$
- web-log ('blog') news propagation $\qquad$
- computer network security: email/IP traffic and anomaly detection $\qquad$
- ....


## Outline

- Introduction - Motivation
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- Problem\#1: Patterns in graphs $\qquad$
- Static graphs
- Weighted graphs $\qquad$
- Time evolving graphs
- Problem\#2: Tools $\qquad$
- Problem\#3: Scalability
- Conclusions
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## Graph mining

- Are real graphs random?


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## Laws and patterns

- Are real graphs random?
- A: NO!!
- Diameter
- in- and out- degree distributions
- other (surprising) patterns
- So, let's look at the data


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## Solution\# S. 1

- Power law in the degree distribution [SIGCOMM99]
internet domains



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## Solution\# S. 1

- Power law in the degree distribution [SIGCOMM99]


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Solution\# S.2: Eigen Exponent $E$

```
    Eigenvalue
```



```
- A2: power law in the eigenvalues of the adjacency matrix
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```


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Solution\# S.2: Eigen Exponent $\boldsymbol{E}$
Eigenvalue


- [Mihail, Papadimitriou '02]: slope is $1 / 2$ of rank exponent
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## And numerous more

- \# of sexual contacts
- Income [Pareto] -'80-20 distribution'
- Duration of downloads [Bestavros+]
- Duration of UNIX jobs ('mice and elephants’)
- Size of files of a user
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- ...
- 'Black swans'


## Outline

- Introduction - Motivation
- Problem\#1: Patterns in graphs
- Static graphs
- degree, diameter, eigen,
- triangles
- cliques
- Weighted graphs
- Time evolving graphs
- Problem\#2: Tools

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Solution\# S.3: Triangle 'Laws'


- Real social networks have a lot of triangles $\qquad$
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Solution\# S.3: Triangle 'Laws'


- Real social networks have a lot of triangles
- Friends of friends are friends
- Any patterns?

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Triangle Law: Computations
[Tsourakakis ICDM 2008]

But: triangles are expensive to compute
(3-way join; several approx. algos)
Q: Can we do that quickly?

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Triangle Law: Computations
[Tsourakakis ICDM 2008] $\qquad$
$\qquad$
But: triangles are expensive to compute
(3-way join; several approx. algos) $\qquad$
Q: Can we do that quickly?
A: Yes!
\#triangles $=\mathbf{1 / 6 ~ S u m ~ ( ~} \lambda_{\mathrm{i}}{ }^{3}$ )
(and, because of skewness (S2), we only need the top few eigenvalues!

# Triangle Law: Computations 

[Tsourakakis ICDM 2008] Wikipedia graph 2006-Nov-04
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$1000 x+$ speed - up, $>90 \%$ accuracy
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Triangle counting for large graphs?

Anomalous nodes in Twitter(~ 3 billion edges)
[U Kang, Brendan Meeder, +, PAKDD'11]
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## Carnegie Mellon

Triangle counting for large graphs?

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Anomalous nodes in Twitter(~ 3 billion edges)
[U Kang, Brendan Meeder, +, PAKDD'11]
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## Any other 'laws'?

Yes!
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## Any other 'laws'?

Yes!

- Small diameter ( $\sim$ constant!) -
- six degrees of separation / 'Kevin Bacon' - small worlds [Watts and Strogatz]


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## Any other 'laws'?

- Bow-tie, for the web [Kumar+ '99]
- IN, SCC, OUT, 'tendrils'
- disconnected components

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## EigenSpokes

B. Aditya Prakash, Mukund Seshadri, Ashwin Sridharan, Sridhar Machiraju and Christos Faloutsos: EigenSpokes: Surprising
$\qquad$ Patterns and Scalable Community Chipping in Large Graphs, PAKDD 2010, Hyderabad, India, 21-24 June 2010.

## Faloutsos

## Carnegievelelon

## EigenSpokes

- Eigenvectors of adjacency matrix $\qquad$
- equivalent to singular vectors
(symmetric, undirected graph) $\qquad$
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## EigenSpokes

- Eigenvectors of adjacency matrix $\qquad$
- equivalent to singular vectors
(symmetric, undirected graph) $\qquad$


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## EigenSpokes

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## EigenSpokes

- Eigenvectors of adjacency matrix $\qquad$
- equivalent to singular vectors
(symmetric, undirected graph) $\qquad$


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## EigenSpokes

| - EE plot: <br> $2^{\text {nd }}$ Principal component <br> - Scatter plot of u2 scores of u1 vs u2 |  |
| :---: | :---: |
|  |  |
| - One would expect | $\stackrel{\circ}{\circ}$ |
| - Many points @ origin |  |
| - A few scattered $\sim$ randomly | u1 |
|  | $1^{\text {st }}$ Principal component |
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## EigenSpokes

- EE plot:
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- Scatter plot of scores of u1 vs u2
- One would expect
- Many points @ origin
- Af $\quad$ tered
$\sim 1$ ic $n$ ly

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## EigenSpokes - pervasiveness

- Present in mobile social graph
- across time and space
- Patent citation graph ${ }^{3_{02}^{02}}{ }^{02}$ $\qquad$
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## EigenSpokes - explanation

Near-cliques, or near-

bipartite-cores, loosely
$\qquad$ connected


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## EigenSpokes - explanation

Near-cliques, or near-
 bipartite-cores, loosely connected

So what?

- Extract nodes with high scores
- high connectivity
- Good "communities"

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## Outline

- Introduction - Motivation
- Problem\#1: Patterns in graphs $\qquad$
- Static graphs
- degree, diameter, eigen,
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- cliques
$\Rightarrow \quad-$ Weighted graphs
- Time evolving graphs
- Problem\#2: Tools

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## Observations on weighted graphs?

- A: yes - even more 'laws'!

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$\qquad$
M. McGlohon, L. Akoglu, and C. Faloutsos
$\qquad$ Components: Patterns and a Generator. SIG-KDD 2008

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Observation W.1: Fortification
Q: How do the weights
of nodes relate to degree?

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Observation W.1: Fortification

More donors,
more \$ ?


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Observation W.1: fortification:

## Snapshot Power Law

- Weight: super-linear on in-degree
- exponent 'iw': $1.01<\mathrm{iw}<1.26$

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## Problem: Time evolution

- with Jure Leskovec (CMU -> Stanford)

- and Jon Kleinberg (Cornell sabb. @ CMU)


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## T. 1 Evolution of the Diameter

- Prior work on Power Law graphs hints at slowly growing diameter:
- diameter $\sim \mathrm{O}(\log \mathrm{N})$
- diameter $\sim \mathrm{O}(\log \log \mathrm{N})$$\stackrel{\sim}{3}$
- What is happening in real data?


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## T. 1 Evolution of the Diameter

- Prior work on Power Law graphs hints
$\qquad$ at slowly growing diameter:

$\qquad$
- What is happening in real data?
- Diameter shrinks over time


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## T. 1 Diameter - "Patents"

- Patent citation network
- 25 years of data
- @1999
- 2.9 M nodes
- 16.5 M edges

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## T. 2 Temporal Evolution of the Graphs

- $\mathrm{N}(\mathrm{t})$... nodes at time t
- $\mathrm{E}(\mathrm{t}) \ldots$ edges at time t
- Suppose that

$$
\mathrm{N}(\mathrm{t}+1)=2 * \mathrm{~N}(\mathrm{t})
$$

- Q: what is your guess for $\qquad$ $\mathrm{E}(\mathrm{t}+1)=? 2$ * $\mathrm{E}(\mathrm{t})$


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## T. 2 Temporal Evolution of the Graphs

- $\mathrm{N}(\mathrm{t})$... nodes at time t
- $\mathrm{E}(\mathrm{t})$... edges at time t
- Suppose that
$\mathrm{N}(\mathrm{t}+1)=2 * \mathrm{~N}(\mathrm{t})$
- Q: what is your guess for $\mathrm{E}(\mathrm{t}+1)={ }^{2}$ * $\mathrm{E}(\mathrm{t})$
- A: over-doubled!
- But obeying the "'Densification Power Law’

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## T. 2 Densification - Patent Citations

| - Citations among patents granted <br> @1999 <br> - 2.9 M nodes <br> - 16.5 M edges <br> - Each year is a datapoint |  |  | $\begin{aligned} & 1.66 \\ & \underbrace{1999} \\ & x^{166} R^{2}=0 \end{aligned}$ |
| :---: | :---: | :---: | :---: |
|  | ${ }^{10^{5}} 10^{\text {b }}$ | $\frac{10^{\circ}}{\substack{100 \\ \text { Number of nodes }}}$ |  |
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More on Time-evolving graphs

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M. McGlohon, L. Akoglu, and C. Faloutsos

| Weighted Graphs and Disconnected |
| :--- |
| Components: Patterns and a Generator. |
| SIG-KDD 2008 |
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## [ Gelling Point ]

- Most real graphs display a gelling point
- After gelling point, they exhibit typical behavior. This is marked by a spike in diameter.



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Observation T.3: NLCC behavior
Q: How do NLCC's emerge and join with
$\qquad$ the GCC?
(" ${ }^{\prime}$ NLCC" ${ }^{\prime}=$ non-largest conn. components)
$\qquad$

- Do they continue to grow in size?
- or do they shrink? $\qquad$
- or stabilize?



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Observation T.3: NLCC behavior
Q: How do NLCC's emerge and join with the GCC?
(" ${ }^{\prime} \mathrm{NLCC} "$ " $=$ non-largest conn. components)

- Do they continue to grow in size?
- or do they shrink?
- or stabilize?


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Observation T.3: NLCC behavior
Q: How do NLCC's emerge and join with
$\qquad$ the GCC?
(" ${ }^{\prime}$ NLCC" $"=$ non-largest conn. components)
YES - Do they continue to grow in size?
YES - or do they shrink? $\qquad$ YES - or stabilize?

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Observation T.3: NLCC behavior

- After the gelling point, the GCC takes off, but NLCC's remain $\sim$ constant (actually, oscillate).
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## Timing for Blogs

- with Mary McGlohon (CMU->Google)
- Jure Leskovec (CMU->Stanford)
$\qquad$
- Natalie Glance (now at Google)
- Mat Hurst (now at MSR)
[SDM'07]

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## Cannegievelelon

## T.5: duration of phonecalls

Surprising Patterns for the Call Duration Distribution of Mobile Phone Users

Pedro O. S. Vaz de Melo, Leman
Akoglu, Christos Faloutsos, Antonio $\qquad$
A. F. Loureiro

PKDD 2010 $\qquad$

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## 'TLaC: Lazy Contractor'

- The longer a task (phonecall) has taken,
- The even longer it will take

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## Data Description

- Data from a private mobile operator of a large
$\qquad$ city
- 4 months of data
- 3.1 million users
- more than 1 billion phone records
- Over $96 \%$ of 'talkative' users obeyed a TLAC
$\qquad$ distribution ('talkative': >30 calls)


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## Outline

- Introduction - Motivation
$\qquad$
- Problem\#1: Patterns in graphs
- Problem\#2: Tools
$\Rightarrow \quad$ - OddBall (anomaly detection) $\qquad$
- Belief Propagation
- Immunization $\qquad$
- Problem\#3: Scalability
- Conclusions $\qquad$


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## OddBall: Spotting Anomalies in Weighted Graphs <br> 

Leman Akoglu, Mary McGlohon, Christos
Faloutsos
Carnegie Mellon University
School of Computer Science

PAKDD 2010, Hyderabad, India

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## Main idea

For each node,

- extract 'ego-net' (=1-step-away neighbors)
- Extract features (\#edges, total weight, etc etc)
- Compare with the rest of the population


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What is an egonet?


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## Near-Clique/Star


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## Near-Clique/Star


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Near-Clique/Star


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## Near-Clique/Star



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## Popular press



And less desirable attention:

- E-mail from 'Belgium police' ('copy of your code?')

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## Outline

- Introduction - Motivation
$\qquad$
- Problem\#1: Patterns in graphs
- Problem\#2: Tools
- OddBall (anomaly detection) $\qquad$
$\Rightarrow$ - Belief Propagation - antivirus app
- Immunization $\qquad$
- Problem\#3: Scalability
- Conclusions $\qquad$
$\qquad$


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Polonium：Tera－Scale Graph Mining and Inference for Malware Detection
$\qquad$ SDM 2011，Mesa，Arizona Software Engineer

## Polonium：The Data

| 晨的部部 | 60＋terabytes of data anonymously |
| :---: | :---: |
|  | contributed by participants of worldwide |
|  | Norton Community Watch program |
|  | 50＋million machines |
|  | 900＋million executable files |
|  | Constructed a machine－file bipartite graph (0.2 TB+) |
|  | 1 billion nodes（machines and files） |
| $\square^{\cdots}$ 号 | 37 billion edges |
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## Polonium：Key Ideas

－Use Belief Propagation to propagate domain knowledge in machine－file graph to detect malware
－Use＂guilt－by－association＂（i．e．，homophily） $\qquad$
－E．g．，files that appear on machines with many bad files are more likely to be bad
－Scalability：handles 37 billion－edge graph

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Polonium: One-Interaction Results


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## Outline

- Introduction - Motivation
- Problem\#1: Patterns in graphs
- Problem\#2: Tools
- OddBall (anomaly detection) $\qquad$
- Belief propagation
$\Rightarrow$ - Immunization $\qquad$
- Problem\#3: Scalability -PEGASUS
- Conclusions
$\qquad$

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## Immunization and epidemic thresholds

- Q1: which nodes to immunize?
- Q2: will a virus vanish, or will it create an epidemic?


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## Q1: Immunization:

-Given
-a network,
-k vaccines, and
-the virus details
-Which nodes to immunize?


## Carnegie Nellor

## -Given

## Q1: Immunization:

-a network,
-k vaccines, and
-the virus details
-Which nodes to immunize?


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## Carnegie. Nellon <br>  <br> Q2: will a virus take over?

- Flu-like virus (no immunity, 'SIS')
- Mumps (life-time immunity, 'SIR’)
- Pertussis (finite-length immunity, ‘SIRS')



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## Q2: will a virus take over?

- Flu-like virus (no immunity, 'SIS')
- Mumps (life-time immunity, 'SIR')
- Pertussis (finite-length immunity, 'SIRS')
$\beta$ : attack prob
$\delta$ : heal prob
A: depends on connectivity (avg degree? Max degree? variance? Something else?
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Epidemic threshold $\tau$

What should $\tau$ depend on?

- avg. degree? and/or highest degree?
- and/or variance of degree?
- and/or third moment of degree?
- and/or diameter?


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## Epidemic threshold

- [Theorem] We have no epidemic, if

$$
\beta / \delta<\tau=1 / \lambda_{1, A}
$$

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## Epidemic threshold

- [Theorem] We have no epidemic, if
$\qquad$
$\qquad$
$\qquad$


Proof: [Wang+03] (for SIS=flu only)

## A2: will a virus take over?

- For all typical virus propagation models (flu, mumps, pertussis, HIV, etc)
- The only connectivity measure that matters, is


## $1 / \lambda_{1}$

the first eigenvalue of the adj. matrix
[Prakash+, '10, arxiv]


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## Thresholds for some models

- $s=$ effective strength
- $s<1$ : below threshold


| Models | Effective Strength <br> $(\mathrm{s})$ |
| :--- | :--- |
| Threshold (tipping <br> point) |  |

$$
\begin{array}{ll}
\text { SIS, SIR, SIRS, } & s=\lambda \cdot\left(\frac{\beta}{\delta}\right) \\
\text { SEIR } & s=1 \\
\text { SIV, SEIV } & s=\lambda \cdot\left(\frac{\beta \gamma}{\delta(\gamma+\theta)}\right)
\end{array}
$$

$$
\mathrm{SI}_{1} \mathrm{I}_{2} \mathrm{~V}_{1} \mathrm{~V}_{2} \text { (H.I.V.) } \quad s=\lambda \cdot\left(\frac{\beta_{1} v_{2}+\beta_{2} \varepsilon}{v_{2}\left(\varepsilon+v_{1}\right)}\right)
$$

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## Q1: Immunization:

-Given

- a network,
-k vaccines, and
-the virus details •the virus details
•Which nodes to immunize? for any virus!


A: immunize the ones that
Max eigen-drop $\Delta \lambda$

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- Belief propagation
- Immunization $\qquad$
- Problem\#3: Scalability -PEGASUS
- Conclusions


## Scalability

- Google: $>450,000$ processors in clusters of $\sim 2000$ processors each [Barroso, Dean, Hölzle, "Web Search for a Planet: The Google Cluster Architecture" IEEE Micro $\qquad$ 2003]
- Yahoo: 5Pb of data [Fayyad, KDD’07] $\qquad$
- Problem: machine failures, on a daily basis
- How to parallelize data mining tasks, then?
- A: map/reduce - hadoop (open-source clone) http://hadoop.apache.org/

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## Outline - Algorithms \& results

$\Rightarrow$|  | Centralized | Hadoop/ <br> PEGASUS |
| :--- | :---: | :---: |
| Degree Distr. | old | old |
| Pagerank | old | old |
| Diameter/ANF | old | HERE |
| Conn. Comp | old | HERE |
| Triangles | done | HERE |
| Visualization | started |  |

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HADI for diameter estimation

- Radius Plots for Mining Tera-byte Scale Graphs U Kang, Charalampos Tsourakakis, Ana Paula Appel, Christos Faloutsos, Jure Leskovec, SDM'10
- Naively: diameter needs $\mathbf{O}(\mathbf{N} * * 2)$ space and up to $\mathrm{O}\left(\mathrm{N}^{* *} 3\right)$ time - prohibitive ( $\mathrm{N} \sim 1 \mathrm{~B}$ )
- Our HADI: linear on E ( $\sim 10 \mathrm{~B}$ )
- Near-linear scalability wrt \# machines $\qquad$
- Several optimizations -> 5x faster

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- Largest publicly available graph ever studied.

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YahooWeb graph (120Gb, 1.4B nodes, 6.6 B edges)

- Largest publicly available graph ever studied.
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## Faloutsos


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YahooWeb graph (120Gb, 1.4B nodes, 6.6 B edges)

- effective diameter: surprisingly small.
- Multi-modality (?!)
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Radius Plot of GCC of YahooWeb.

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YahooWeb graph (120Gb, 1.4B nodes, 6.6 B edges)

- effective diameter: surprisingly small.
- Multi-modality: probably mixture of cores .
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YahooWeb graph (120Gb, 1.4B nodes, 6.6 B edges) - effective diameter: surprisingly small.

- Multi-modality: probably mixture of cores . 15:826 (c) 2011 C . Faloutos
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## Faloutsos


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Running time - Kronecker and Erdos-Renyi $\qquad$ Graphs with billions edges.

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## Outline - Algorithms \& results

$\Rightarrow$|  | Centralized | Hadoop/ <br> PEGASUS |
| :--- | :---: | :---: |
| Degree Distr. | old | old |
| Pagerank | old | old |
| Diameter/ANF | old | HERE |
| Conn. Comp | old | HERE |
| Triangles |  | HERE |
| Visualization | started |  |

## Generalized Iterated Matrix Vector Multiplication (GIMV)

PEGASUS: A Peta-Scale Graph Mining System - Implementation and Observations. U Kang, Charalampos E. Tsourakakis, and Christos Faloutsos.
(ICDM) 2009, Miami, Florida, USA.
Best Application Paper (runner-up).

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## Generalized Iterated Matrix details $<$

Vector Multiplication (GIMV)

- PageRank
- proximity (RWR)
- Diameter
- Connected components
- (eigenvectors,
- Belief Prop.
- ...)

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Matrix - vector
Multiplication (iterated)
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## Example: GIM-V At Work

- Connected Components - 4 observations:

Count

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## Example: GIM-V At Work

- Connected Components

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## Example: GIM-V At Work

- Connected Components


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## Example: GIM-V At Work

- Connected Components

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## Example: GIM-V At Work

- Connected Components



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## Example: GIM-V At Work

- Connected Components
Count


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## GIM-V At Work

- Connected Components over Time
- LinkedIn: 7.5M nodes and 58M edges

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## Outline

- Introduction - Motivation
- Problem\#1: Patterns in graphs
- Problem\#2: Tools
- Problem\#3: Scalability
$\Rightarrow$ - Conclusions


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## OVERALL CONCLUSIONS low level:

- Several new patterns (fortification, shrinking diameter, triangle-laws, conn. components, etc)
- New tools:
- anomaly detection (OddBall), belief propagation, immunization
- Scalability: PEGASUS / hadoop
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OVERALL CONCLUSIONS high level

- BIG DATA: Large datasets reveal patterns/ outliers that are invisible otherwise $\qquad$
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