

CMU SCS

15-826: Multimedia Databases and Data Mining

Multi-key and Spatial Access Methods - II
C. Faloutsos

CMU SCS

Outline

Goal: 'Find similar / interesting things'

- Intro to DB
- ➔ • Indexing - similarity search
- Data Mining

15-826 Copyright: C. Faloutsos (2005) 2

CMU SCS

Indexing - Detailed outline

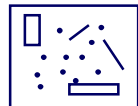
- primary key indexing
- secondary key / multi-key indexing
- ➔ • spatial access methods
 - problem dfn
 - z-ordering
 - R-trees
 - ...
- text
- ...

15-826 Copyright: C. Faloutsos (2005) 3

CMU SCS

Spatial Access Methods - problem

- Given a collection of geometric objects (points, lines, polygons, ...)
- organize them on disk, to answer spatial queries (like??)

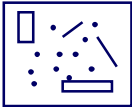


15-826 Copyright: C. Faloutsos (2005) 4

CMU SCS

Spatial Access Methods - problem

- Given a collection of geometric objects (points, lines, polygons, ...)
- organize them on disk, to answer
 - point queries
 - range queries
 - k-nn queries
 - spatial joins ('all pairs' queries)

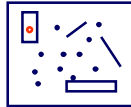


15-826 Copyright: C. Faloutsos (2005) 5

CMU SCS

Spatial Access Methods - problem

- Given a collection of geometric objects (points, lines, polygons, ...)
- organize them on disk, to answer
 - point queries
 - range queries
 - k-nn queries
 - spatial joins ('all pairs' queries)

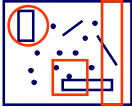


15-826 Copyright: C. Faloutsos (2005) 6


CMU SCS

Spatial Access Methods - problem

- Given a collection of geometric objects (points, lines, polygons, ...)
- organize them on disk, to answer
 - point queries
 - range queries
 - k-nn queries
 - spatial joins ('all pairs' queries)

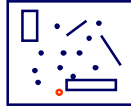


15-826
Copyright: C. Faloutsos (2005)
7


CMU SCS

Spatial Access Methods - problem

- Given a collection of geometric objects (points, lines, polygons, ...)
- organize them on disk, to answer
 - point queries
 - range queries
 - k-nn queries
 - spatial joins ('all pairs' queries)

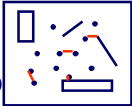


15-826
Copyright: C. Faloutsos (2005)
8


CMU SCS

Spatial Access Methods - problem

- Given a collection of geometric objects (points, lines, polygons, ...)
- organize them on disk, to answer
 - point queries
 - range queries
 - k-nn queries
 - spatial joins ('all pairs' within ϵ)



15-826
Copyright: C. Faloutsos (2005)
9


CMU SCS

SAMs - motivation

- Q: applications?

15-826
Copyright: C. Faloutsos (2005)
10

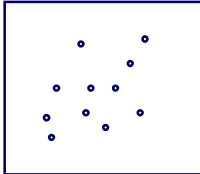

CMU SCS

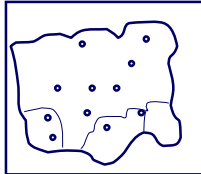
SAMs - motivation

traditional DB

GIS

age





salary

15-826
Copyright: C. Faloutsos (2005)
11

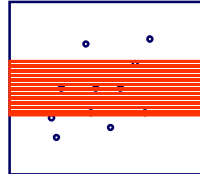

CMU SCS

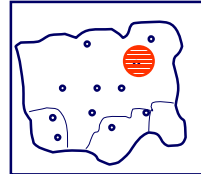
SAMs - motivation

traditional DB

GIS

age





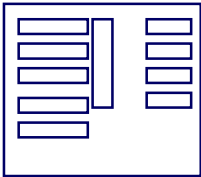
salary

15-826
Copyright: C. Faloutsos (2005)
12

CMU SCS

SAMs - motivation

CAD/CAM



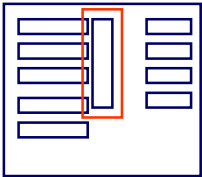
find elements too close to each other

15-826 Copyright: C. Faloutsos (2005) 13

CMU SCS

SAMs - motivation

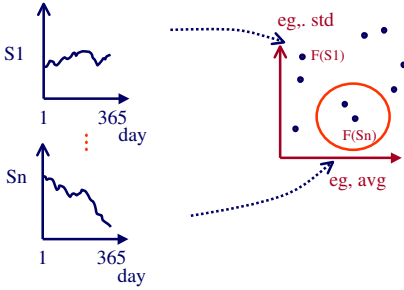
CAD/CAM



15-826 Copyright: C. Faloutsos (2005) 14

CMU SCS

SAMs - motivation



15-826 Copyright: C. Faloutsos (2005) 15

CMU SCS

Indexing - Detailed outline

- primary key indexing
- secondary key / multi-key indexing
- spatial access methods
 - problem defn
 - z-ordering
 - R-trees
 - ...
- text
- ...

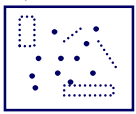
15-826 Copyright: C. Faloutsos (2005) 16

CMU SCS

SAMs: solutions

- z-ordering
- R-trees
- (grid files)

Q: how would you organize, e.g., n -dim points, on disk? (C points per disk page)



15-826 Copyright: C. Faloutsos (2005) 17

CMU SCS

z-ordering

Q: how would you organize, e.g., n -dim points, on disk? (C points per disk page)

Hint: reduce the problem to 1-d points (!!)

Q1: why?

A:

Q2: how?



15-826 Copyright: C. Faloutsos (2005) 18

z-ordering

Q: how would you organize, e.g., n -dim points, on disk? (C points per disk page)

Hint: reduce the problem to 1-d points (!!)

Q1: why?

A: B-trees!

Q2: how?



15-826

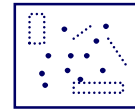
Copyright: C. Faloutsos (2005)

19

z-ordering

Q2: how?

A: assume finite granularity; z-ordering = bit-shuffling = N-trees = Morton keys = geo-coding = ...



15-826

Copyright: C. Faloutsos (2005)

20

z-ordering

Q2: how?

A: assume finite granularity (e.g., $2^{32} \times 2^{32}$; 4x4 here)

Q2.1: how to map n -d cells to 1-d cells?



15-826

Copyright: C. Faloutsos (2005)

21

z-ordering

Q2.1: how to map n -d cells to 1-d cells?



15-826

Copyright: C. Faloutsos (2005)

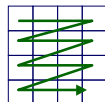
22

z-ordering

Q2.1: how to map n -d cells to 1-d cells?

A: row-wise

Q: is it good?



15-826

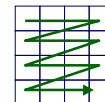
Copyright: C. Faloutsos (2005)

23

z-ordering

Q: is it good?

A: great for 'x' axis; bad for 'y' axis



15-826

Copyright: C. Faloutsos (2005)

24

CMU SCS

z-ordering

z-ordering/bit-shuffling/linear-quadtrees

Q: How to generate this curve ($z = f(x,y)$)?

A1: 'z' (or 'N') shapes, RECURSIVELY

order-1 order-2 ... order (n+1)

15-826 Copyright: C. Faloutsos (2005) 31

CMU SCS

z-ordering

Notice:

- self similar (we'll see about fractals, soon)
- method is hard to use: $z = ? f(x,y)$

order-1 order-2 ... order (n+1)

15-826 Copyright: C. Faloutsos (2005) 32

CMU SCS

z-ordering

z-ordering/bit-shuffling/linear-quadtrees

Q: How to generate this curve ($z = f(x,y)$)?

A: 3 (equivalent) answers!

Method #2?

15-826 Copyright: C. Faloutsos (2005) 33

CMU SCS

z-ordering

bit-shuffling

$z = (0101)_2 = 5$

15-826 Copyright: C. Faloutsos (2005) 34

CMU SCS

z-ordering

bit-shuffling

$z = (0101)_2 = 5$

How about the reverse:
 $(x,y) = g(z)$?

15-826 Copyright: C. Faloutsos (2005) 35

CMU SCS

z-ordering

bit-shuffling

$z = (0101)_2 = 5$

How about n -d spaces?

15-826 Copyright: C. Faloutsos (2005) 36

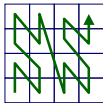
CMU SCS

z-ordering

z-ordering/bit-shuffling/**linear-quadtrees**

Q: How to generate this curve ($z = f(x,y)$)?

A: 3 (equivalent) answers!



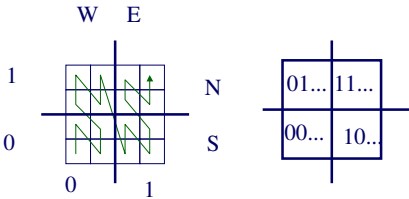
Method #3?

15-826 Copyright: C. Faloutsos (2005) 37

CMU SCS

z-ordering

linear-quadtrees : assign N->1, S->0 e.t.c.

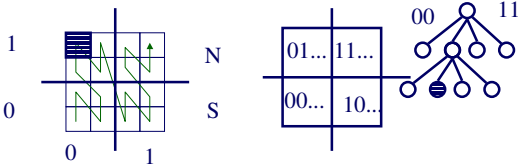


15-826 Copyright: C. Faloutsos (2005) 38

CMU SCS

z-ordering

... and repeat recursively. Eg.: $z_{\text{blue-cell}} = \text{WN;WN} = (0101)_2 = 5$

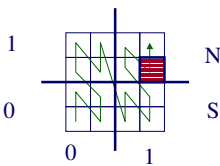


15-826 Copyright: C. Faloutsos (2005) 39

CMU SCS

z-ordering

Drill: z-value of magenta cell, with the three methods?

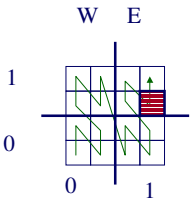


15-826 Copyright: C. Faloutsos (2005) 40

CMU SCS

z-ordering

Drill: z-value of magenta cell, with the three methods?



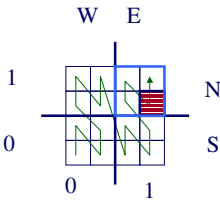
method#1: 14
method#2: shuffle(11;10)=
(1110)₂ = 14

15-826 Copyright: C. Faloutsos (2005) 41

CMU SCS

z-ordering

Drill: z-value of magenta cell, with the three methods?



method#1: 14
method#2: shuffle(11;10)=
(1110)₂ = 14
method#3: **EN**;ES = ... = 14

15-826 Copyright: C. Faloutsos (2005) 42


CMU SCS

z-ordering - Detailed outline

- spatial access methods
 - z-ordering
 - main idea - 3 methods
 - use w/ B-trees; algorithms (range, knn queries ...)
 - non-point (eg., region) data
 - analysis; variations
 - R-trees
 - ...

15-826
Copyright: C. Faloutsos (2005)
43

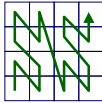

CMU SCS

z-ordering - usage & algo's

Q1: How to store on disk?

A:

Q2: How to answer range queries etc



15-826
Copyright: C. Faloutsos (2005)
44

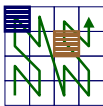

CMU SCS

z-ordering - usage & algo's

Q1: How to store on disk?

A: treat z-value as primary key; feed to B-tree

PGH





z	cname	etc
5	SF	
12	PGH	

15-826
Copyright: C. Faloutsos (2005)
45

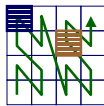

CMU SCS

z-ordering - usage & algo's

MAJOR ADVANTAGES w/ B-tree:

- already inside commercial systems (no coding/debugging!)
- concurrency & recovery is ready

PGH





z	cname	etc
5	SF	
12	PGH	

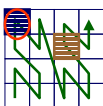
15-826
Copyright: C. Faloutsos (2005)
46


CMU SCS

z-ordering - usage & algo's

Q2: queries? (eg.: *find city at (0,3)*)?

PGH





z	cname	etc
5	SF	
12	PGH	

15-826
Copyright: C. Faloutsos (2005)
47

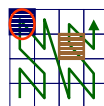

CMU SCS

z-ordering - usage & algo's

Q2: queries? (eg.: *find city at (0,3)*)?

A: find z-value; search B-tree

PGH





z	cname	etc
5	SF	
12	PGH	

15-826
Copyright: C. Faloutsos (2005)
48

CMU SCS

z-ordering - usage & algo's

Q2: range queries?

PGH

SF

z	cname	etc
5	SF	
12	PGH	

15-826 Copyright: C. Faloutsos (2005) 49

CMU SCS

z-ordering - usage & algo's

Q2: range queries?

A: compute ranges of z-values; use B-tree

PGH

SF

9,11-15

z	cname	etc
5	SF	
12	PGH	

15-826 Copyright: C. Faloutsos (2005) 50

CMU SCS

z-ordering - usage & algo's

Q2': range queries - how to reduce # of qualifying of ranges?

PGH

SF

9,11-15

z	cname	etc
5	SF	
12	PGH	

15-826 Copyright: C. Faloutsos (2005) 51

CMU SCS

z-ordering - usage & algo's

Q2': range queries - how to reduce # of qualifying of ranges?

A: Augment the query!

PGH

SF

9,11-15 -> 8-15

z	cname	etc
5	SF	
12	PGH	

15-826 Copyright: C. Faloutsos (2005) 52

CMU SCS

z-ordering - usage & algo's

Q2'': range queries - how to break a query into ranges?

9,11-15

15-826 Copyright: C. Faloutsos (2005) 53

CMU SCS

z-ordering - usage & algo's

Q2'': range queries - how to break a query into ranges?

A: recursively, quadtree-style; decompose only non-full quadrants

12-15

9,11-15

15-826 Copyright: C. Faloutsos (2005) 54

CMU SCS

z-ordering - usage & algo's

Q2'': range queries - how to break a query into ranges?

A: recursively, quadtree-style; decompose only non-full quadrants

15-826 Copyright: C. Faloutsos (2005) 55

CMU SCS

z-ordering - Detailed outline

- spatial access methods
 - z-ordering
 - main idea - 3 methods
 - use w/ B-trees; algorithms (range, knn queries ...)
 - non-point (eg., region) data
 - analysis; variations
 - R-trees
 - ...

15-826 Copyright: C. Faloutsos (2005) 56

CMU SCS

z-ordering - usage & algo's

Q3: k-nn queries? (say, 1-nn)?

A: traverse B-tree; find nn wrt z-values and ...

z	cname	etc
5	SF	
12	PGH	

15-826 Copyright: C. Faloutsos (2005) 57

CMU SCS

z-ordering - usage & algo's

Q3: k-nn queries? (say, 1-nn)?

A: traverse B-tree; find nn wrt z-values and ...

z	cname	etc
5	SF	
12	PGH	

15-826 Copyright: C. Faloutsos (2005) 58

CMU SCS

z-ordering - usage & algo's

... ask a range query.

z	cname	etc
5	SF	
12	PGH	

15-826 Copyright: C. Faloutsos (2005) 59

CMU SCS

z-ordering - usage & algo's

... ask a range query.

z	cname	etc
5	SF	
12	PGH	

15-826 Copyright: C. Faloutsos (2005) 60

CMU SCS

z-ordering - regions

Q: How to store in B-tree?
Q: How to search (range etc queries)

15-826 Copyright: C. Faloutsos (2005) 67

CMU SCS

z-ordering - regions

Q: How to store in B-tree? A: sort ($* < 0 < 1$)
Q: How to search (range etc queries)

z	obj-id	etc
0010	C	
0101	A	
1000	C	
11**	B	

15-826 Copyright: C. Faloutsos (2005) 68

CMU SCS

z-ordering - regions

Q: How to search (range etc queries) - eg 'red' range query

z	obj-id	etc
0010	C	
0101	A	
1000	C	
11**	B	

15-826 Copyright: C. Faloutsos (2005) 69

CMU SCS

z-ordering - regions

Q: How to search (range etc queries) - eg 'red' range query
A: break query in z-values; check B-tree

z	obj-id	etc
0010	C	
0101	A	
1000	C	
11**	B	

15-826 Copyright: C. Faloutsos (2005) 70

CMU SCS

z-ordering - regions

Almost identical to range queries for point data, except for the "don't cares" - i.e.,

1100 ?? 11**

z	obj-id	etc
0010	C	
0101	A	
1000	C	
11**	B	

15-826 Copyright: C. Faloutsos (2005) 71

CMU SCS

z-ordering - regions

Almost identical to range queries for point data, except for the "don't cares" - i.e.,

$z1 = 1100 \text{ ?? } 11** = z2$

Specifically: does $z1$ contain/avoid/intersect $z2$?
Q: what is the criterion to decide?

15-826 Copyright: C. Faloutsos (2005) 72

z-ordering - regions

$z1 = 1100$?? $11^{**} = z2$

Specifically: does $z1$ contain/avoid/intersect $z2$?

Q: what is the criterion to decide?

A: **Prefix property**: let $r1, r2$ be the corresponding regions, and let $r1$ be the smallest ($\Rightarrow z1$ has fewest '*'s). Then:

15-826
 Copyright: C. Faloutsos (2005)
 73

z-ordering - regions

- $r2$ will either contain completely, or avoid completely $r1$.
- it will contain $r1$, if $z2$ is the prefix of $z1$

1100 ?? 11^{**}

region of $z1$:
completely contained in
region of $z2$

15-826
 Copyright: C. Faloutsos (2005)
 74

z-ordering - regions

Drill (True/False). Given:

- $z1 = 011001^{**}$
- $z2 = 01^{*****}$
- $z3 = 0100^{****}$

T/F $r2$ contains $r1$

T/F $r3$ contains $r1$

T/F $r3$ contains $r2$

15-826
 Copyright: C. Faloutsos (2005)
 75

z-ordering - regions

Drill (True/False). Given:

- $z1 = 011001^{**}$
- $z2 = 01^{*****}$
- $z3 = 0100^{****}$

T/F $r2$ contains $r1$ - TRUE (prefix property)

T/F $r3$ contains $r1$ - FALSE (disjoint)

T/F $r3$ contains $r2$ - FALSE ($r2$ contains $r3$)

15-826
 Copyright: C. Faloutsos (2005)
 76

z-ordering - regions

Drill (True/False). Given:

- $z1 = 011001^{**}$
- $z2 = 01^{*****}$
- $z3 = 0100^{****}$

15-826
 Copyright: C. Faloutsos (2005)
 77

z-ordering - regions

Drill (True/False). Given:

- $z1 = 011001^{**}$
- $z2 = 01^{*****}$
- $z3 = 0100^{****}$

T/F $r2$ contains $r1$ - TRUE (prefix property)

T/F $r3$ contains $r1$ - FALSE (disjoint)

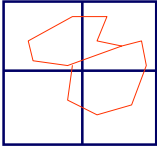
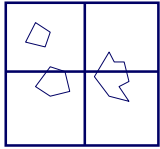
T/F $r3$ contains $r2$ - FALSE ($r2$ contains $r3$)

15-826
 Copyright: C. Faloutsos (2005)
 78


CMU SCS

z-ordering - regions

Spatial joins: find (quickly) all
counties intersecting lakes

15-826
Copyright: C. Faloutsos (2005)
79


CMU SCS

z-ordering - regions

Spatial joins: find (quickly) all
counties intersecting lakes

Naive algorithm: $O(N * M)$
 Something faster?

15-826
Copyright: C. Faloutsos (2005)
80


CMU SCS

z-ordering - regions

Spatial joins: find (quickly) all
counties intersecting lakes

z	obj-id	etc
0010	ALG	
...	...	
1000	WAS	
11**	ALG	

z	obj-id	etc
0011	Erie	
0101	Erie	
...		
10**	Ont.	

15-826
Copyright: C. Faloutsos (2005)
81


CMU SCS

z-ordering - regions

Spatial joins: find (quickly) all
counties intersecting lakes

Solution: merge the lists of (sorted) z-values,
 looking for the prefix property

footnote#1: '*' needs careful treatment
 footnote#2: need dup. elimination

15-826
Copyright: C. Faloutsos (2005)
82


CMU SCS

z-ordering - Detailed outline

- spatial access methods
 - z-ordering
 - main idea - 3 methods
 - use w/ B-trees; algorithms (range, knn queries ...)
 - non-point (eg., region) data
 - analysis; variations
 - R-trees
 - ...

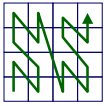
➔

15-826
Copyright: C. Faloutsos (2005)
83


CMU SCS

z-ordering - variations

Q: is z-ordering the best we can do?

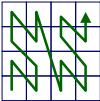


15-826
Copyright: C. Faloutsos (2005)
84

CMU SCS

z-ordering - variations

Q: is z-ordering the best we can do?
 A: probably not - occasional long 'jumps'
 Q: then?



15-826 Copyright: C. Faloutsos (2005) 85

CMU SCS

z-ordering - variations

Q: is z-ordering the best we can do?
 A: probably not - occasional long 'jumps'
 Q: then? A1: Gray codes



15-826 Copyright: C. Faloutsos (2005) 86

CMU SCS

z-ordering - variations

A2: Hilbert curve! (a.k.a. Hilbert-Peano curve)



15-826 Copyright: C. Faloutsos (2005) 87

CMU SCS

z-ordering - variations

'Looks' better (never long jumps). How to derive it?

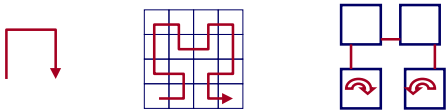


15-826 Copyright: C. Faloutsos (2005) 88

CMU SCS

z-ordering - variations

'Looks' better (never long jumps). How to derive it?



order-1 order-2 ... order (n+1)

15-826 Copyright: C. Faloutsos (2005) 89

CMU SCS

z-ordering - variations

Q: function for the Hilbert curve ($h = f(x,y)$)?
 A: bit-shuffling, followed by post-processing, to account for rotations. Linear on # bits.
 See textbook, for pointers to code/algorithms (eg., [Jagadish, 90])

15-826 Copyright: C. Faloutsos (2005) 90

CMU SCS

z-ordering - variations

Q: how about Hilbert curve in 3-d? n-d?

A: Exists (and is not unique!). Eg., 3-d, order-1 Hilbert curves (Hamiltonian paths on cube)



#1



#2

15-826 Copyright: C. Faloutsos (2005) 91

CMU SCS

z-ordering - Detailed outline

- spatial access methods
 - z-ordering
 - main idea - 3 methods
 - use w/ B-trees; algorithms (range, knn queries ...)
 - non-point (eg., region) data
 - analysis; variations
 - R-trees
 - ...

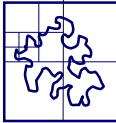
15-826 Copyright: C. Faloutsos (2005) 92

CMU SCS

z-ordering - analysis

Q: How many pieces ('quad-tree blocks') per region?

A: proportional to perimeter (surface etc)

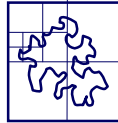
15-826 Copyright: C. Faloutsos (2005) 93

CMU SCS

z-ordering - analysis

(How long is the coastline, say, of England?

Paradox: The answer changes with the yardstick -> fractals ...)

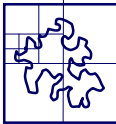



15-826 Copyright: C. Faloutsos (2005) 94

CMU SCS

z-ordering - analysis

Q: Should we decompose a region to full detail (and store in B-tree)?

15-826 Copyright: C. Faloutsos (2005) 95

CMU SCS

z-ordering - analysis

Q: Should we decompose a region to full detail (and store in B-tree)?

A: NO! approximation with 1-3 pieces/z-values is best [Orenstein90]

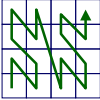
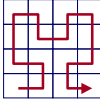



15-826 Copyright: C. Faloutsos (2005) 96

CMU SCS

z-ordering - analysis

Q: how to measure the 'goodness' of a curve?

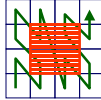
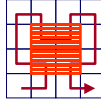



15-826 Copyright: C. Faloutsos (2005) 97

CMU SCS

z-ordering - analysis

Q: how to measure the 'goodness' of a curve?
A: e.g., avg. # of runs, for range queries

4 runs
(#runs ~ #disk accesses on B-tree)

3 runs

15-826 Copyright: C. Faloutsos (2005) 98

CMU SCS

z-ordering - analysis

Q: So, is Hilbert really better?
A: 27% fewer runs, for 2-d (similar for 3-d)

Q: are there formulas for #runs, #of quadtree blocks etc?
A: Yes ([Jagadish; Moon+ etc] see textbook)

15-826 Copyright: C. Faloutsos (2005) 99

CMU SCS

z-ordering - fun observations

Hilbert and z-ordering curves: "space filling curves": eventually, they visit every point in n-d space - therefore:





order-1 order-2 ... order (n+1)

15-826 Copyright: C. Faloutsos (2005) 100

CMU SCS

z-ordering - fun observations

... they show that the plane has as many points as a line (-> headaches for 1900's mathematics/topology). (fractals, again!)





order-1 order-2 ... order (n+1)

15-826 Copyright: C. Faloutsos (2005) 101

CMU SCS

z-ordering - fun observations

Observation #2: Hilbert (like) curve for video encoding [Y. Matias+, CRYPTO '87]:
Given a frame, visit its pixels in randomized hilbert order; compress; and transmit




15-826 Copyright: C. Faloutsos (2005) 102

z-ordering - fun observations

In general, Hilbert curve is great for preserving distances, clustering, vector quantization etc

Indexing - Detailed outline

- primary key indexing
- secondary key / multi-key indexing
- spatial access methods
 - problem defn
 - z-ordering
 - R-trees
 - ...



- text

- ...

Conclusions

- z-ordering is a great idea (n-d points -> 1-d points; feed to B-trees)
- used by TIGER system and (most probably) by other GIS products
- works great with low-dim points