

# 15-354: Midterm

October 11, 2007

Name:

Andrew ID:

## Instructions

- Fill in the box above with your name and your Andrew ID. **Do it, now!**
- Clearly mark your answers in the allocated space. If need be, use the back of a page for scratch space. If you have made a mess, cross out the invalid parts of your solution, and circle the ones that should be graded.
- Scan the test first to make sure that none of the 7 pages are missing. The problems are of varying difficulty and are not necessarily sorted according to increasing difficulty. You might wish to pick off the easy ones first.
- You have 80 minutes. Good luck.

1	20	
2	20	
3	20	
4	20	
5	20	
Total	100	

**Problem 1: Counting and Decidability** (20 pts.)

Suppose  $A \subseteq \mathbb{N}$ . Often it is important to understand the behavior of the corresponding counting function  $f_A : \mathbb{N} \rightarrow \mathbb{N}$  defined by

$$f(n) = \text{cardinality of } (A \cap \{1, 2, \dots, n\})$$

For example, when  $A$  is the set of primes the counting function has been studied in great detail in number theory.

- A. Show that  $f_A$  is computable when  $A$  is decidable.
- B. Assume that  $A$  is semi-decidable. Show that  $A$  must be decidable when  $f_A$  is computable.

**Problem 2: Iteration and Primitive Recursion** (20 pts.)

Primitive recursive functions are obtained from simple basic functions by composition and primitive recursion. As we have seen, iteration is another important operation that produces useful functions. For  $f : \mathbb{N} \rightarrow \mathbb{N}$  define its iterate

$$f^\dagger(n) = f^n(n)$$

- A. Determine  $f^\dagger$  when  $f$  is the successor function,  $f(x) = x + 1$ .
  
- B. Show that  $f^\dagger$  is primitive recursive whenever  $f$  is.

**Problem 3: Selector Functions and Equivalence Relations** (20 pts.)

The forward state merging algorithm is usually implemented using selector functions to represent equivalence relations on the state set  $Q = \{1, 2, \dots, n-1, n\}$ : relation  $\rho$  is represented by an array  $R$  such that

$$R(x) = \min(z \mid z \rho x).$$

If you write pseudo-code to describe your algorithms below, make sure to provide ample comments; no credit otherwise. You can use any auxiliary data structure you like.

- A. Given  $R$ , explain how to compute the number  $r$  of equivalence classes of  $\rho$  in time linear in  $n$  using  $O(1)$  extra memory.
- B. Given  $R$  and  $r$ , explain how to compute the size of a smallest equivalence class of  $\rho$  in time linear in  $n$  (there may be several smallest classes).

**Problem 4: A Minimal Automaton** (20 pts.)

For this problem, the input alphabet is  $\Sigma = \{a, b\}$ .

Consider the regular language  $L = (bb)^*aaa$ .

- A. Using the table method used in class, compute all the quotients of  $L$ . For your convenience, the table below has the first column filled in – but note that there may well be fewer than 8 quotients.

$-$	$(bb)^*aaa$	$L_1$
$a^{-1}L_1$		
$b^{-1}L_1$		
$a^{-1}L_2$		
$b^{-1}L_2$		
$a^{-1}L_3$		
$b^{-1}L_3$		
$a^{-1}L_4$		
$b^{-1}L_4$		
$a^{-1}L_5$		
$b^{-1}L_5$		
$a^{-1}L_6$		
$b^{-1}L_6$		
$a^{-1}L_7$		
$b^{-1}L_7$		
$a^{-1}L_8$		
$b^{-1}L_8$		

- B. Draw a *nice* picture of the resulting minimal automaton (make sure to indicate the initial state

and all final states).

**Problem 5: Alternative Union of Regular Languages** (20 pts.)

Given two DFAs  $M_1$  and  $M_2$  for regular languages  $L_1$  and  $L_2$  we know how to construct a DFA  $M_1 \times M_2$  for  $L = L_1 \cup L_2$  using a product machine. Here is an alternative approach, exploiting non-determinism. Assume the two state sets are disjoint and define a new machine

$$M = \langle Q_1 \cup Q_2, \Sigma, \delta_1 \cup \delta_2, \{q_{01}, q_{02}\}, F_1 \cup F_2 \rangle$$

Note  $M$  has size  $n_1 + n_2$  and is thus potentially quite a bit smaller than the product automaton. Of course,  $M$  is not a DFA.

- A. Prove that  $M$  accepts  $L$ : you must argue about all possible runs of  $M$  on input  $x$ .
- B. Describe the machine  $M'$  obtained by performing deterministic simulation on  $M$ . More precisely,
  - Describe the state set  $Q'$  in terms of  $Q_1$  and  $Q_2$ .
  - Describe the transition function  $\delta'$  in terms of  $\delta_1$  and  $\delta_2$ .
  - Describe the initial state.
  - Describe the final states.