

## 15-453 Homework # 8

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1. Double SAT  
(20 Points)

Let  $\text{DOUBLE-SAT} = \{\phi \mid \phi \text{ has at least two satisfying assignments}\}$ . Show that  $\text{DOUBLE-SAT}$  is NP-complete.

2. HALF-CLIQUE  
(20 Points)

Let  $\text{HALF-CLIQUE} = \{G \mid G \text{ is an undirected graph having a complete subgraph with at least } n/2 \text{ nodes, where } n \text{ is the number of nodes in } G\}$ . Show that  $\text{CLIQUE} \leq_P \text{HALF-CLIQUE}$ .

3. If  $P = NP$ ...  
(40 Points)

Prove the following:

1. Show that if  $P = NP$  then every language  $A \in P$  is NP-complete, except  $A = \emptyset$  and  $A = \Sigma^*$
2. Show that if  $P = NP$ , a polynomial time algorithm exists that, given a Boolean formula  $\phi$ , actually produces a satisfying assignment for  $\phi$  if it is satisfiable. (Note: NP is a class of languages, so saying "SAT is in NP" is not a correct answer for this problem. This problem requires you to show that if you can decide whether a formula is satisfiable in polynomial time, then you can actually find a satisfying assignment for it in polynomial time.)