

Categorical Grammar

- Basic Categorical Grammar:
 - core “5” rules
- Semantics in Categorical Grammar
- Categorical Unification Grammar
 - and how it can be used in other formalisms

Categorial Grammar

- Simplify the rules
- Move complexity from rules to lexical entries
- More tightly coupled with semantics:
 - particularly lambda calculus
- One to one relationship from
 - syntactic and semantic constituents

5 rules

- application:
 - Forward: $A/B + B = A$
 - Backward: $B + A \backslash B = A$
- composition:
 - $A/B + B/C = A/C$
- coordination:
 - $A \text{ CONJ } A' = A''$
- type raising:
 - $A = X/(X \backslash A)$

John = np

Mary = np

likes = (s\np)/np

Forward application

$X/Y Y \Rightarrow X$

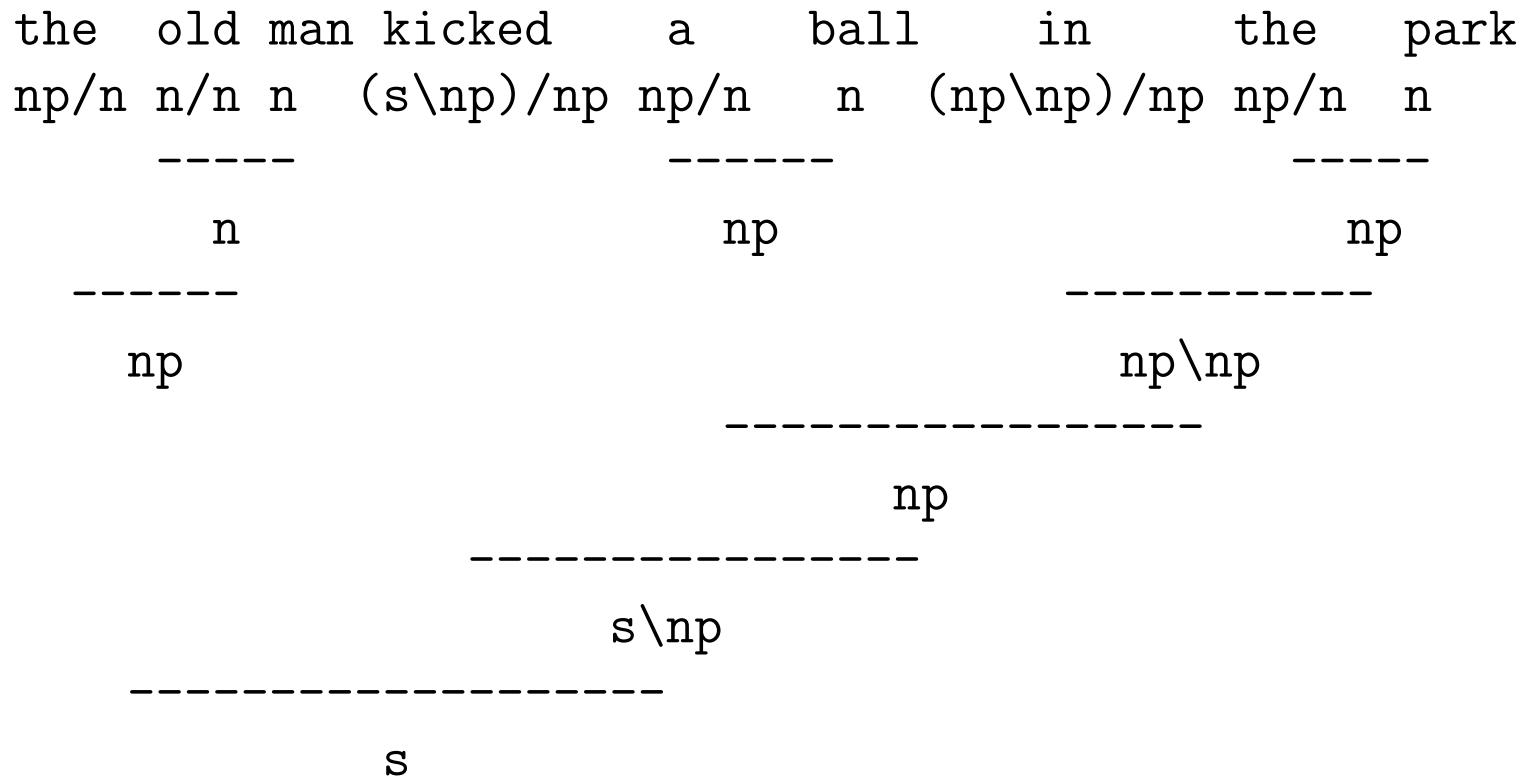
Backward application

$Y X \backslash Y \Rightarrow X$

Thus

John	likes	Mary	
np	(s\np)/np	np	
	-----		Forward
	s\np		
	-----		Backward
	s		

a, the np/n
 old n/n
 in (np\np)/np
 man, ball, park n
 kicked (s\np)/np



Coordination

- Constituent Coordination
 - John and Mary like books
 - (NP and NP) VP
 - John likes fishing and dislikes baseball.
 - NP (VP and VP)
- Non-constituent coordination
 - John likes and Mary dislikes sport.
 - ???

Type Raising

- Computationally unbounded:
 - could happen for any category
 - Makes parsing intractable

- Controlled type raising
 - needs to be guarded
 - only (some) lexical items

Categorial Grammar Semantics

- Each constituent mapped to lambda expression:
 - *Mary* → m
 - *walks* → $\text{lambda } X \text{ walks}(X)$

Lambda Calculus

- Syntax (of LC):
 - First Order Predicate logic, plus
 - lambda VAR TERM
 - where VAR scopes over all occurrences of VAR in TERM
 - lambda X walk(X)
- Lambda application:
 - lambda X walks(X) composed with m =
 - lambda X walks(X) . m
 - walks(m) by beta-reduction

A/B:S + B:T = A:S.T

B:T + A\B:S = A:S.T

John np:j

walks (s\np):lambda X walks(X)

John walks

np:j s\np:lambda X walks(X)

s : walks(j)

B:T + A\B:S = A:S . T

np:j + s\np:lambda X walks(X)

s : lambda X walks(X) . j

s : walks(j)

John np:j

Mary np:m

likes (s\np)/np: lambda Y lambda X likes(X,Y)

John likes Mary

np:j (s\np)/np: lambda Y lambda X likes(X,Y) m

s\np: lambda X likes(X,m)

s likes(j,m)

lambda Y lambda X likes(X,Y) . m

lambda X likes(X,m)

lambda X likes(X,m) . j

likes(j,m)

Conj

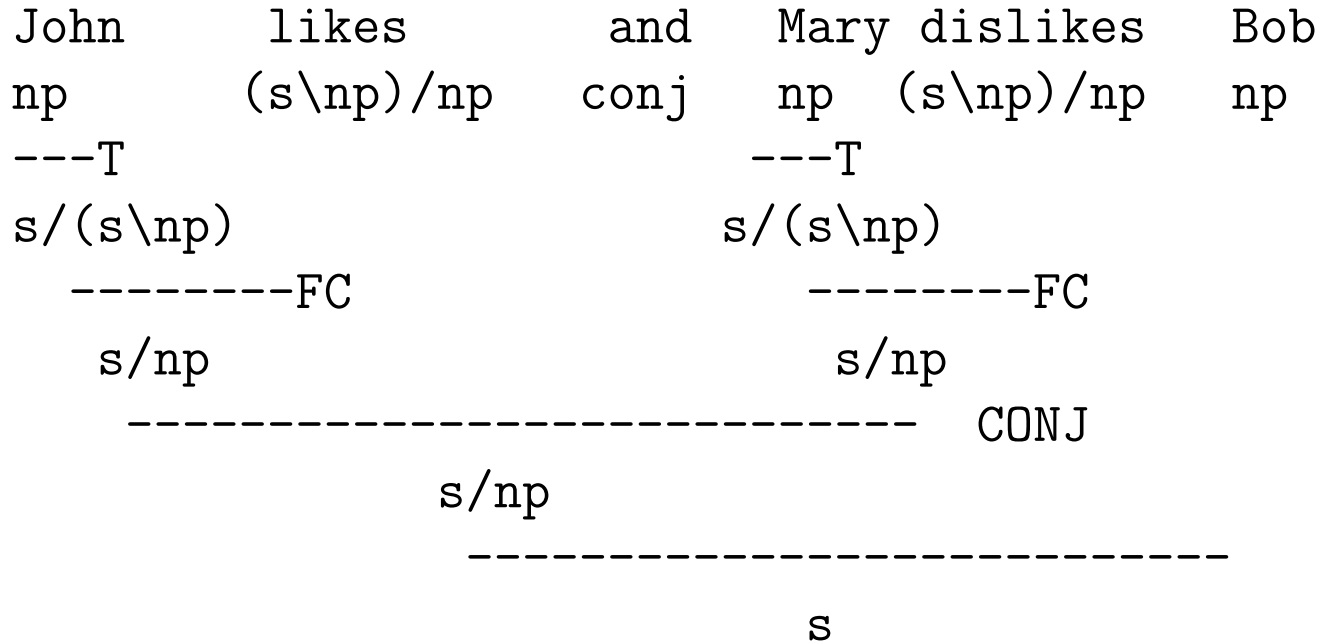
$X:A \text{ CONJ } X':A' = X'':\lambda S (A . S \& A' . S)$

Composition

$X/Y:A \ Y/Z:B \Rightarrow X/Z: \lambda Q (A . (B . Q))$

Type raising

$NP:a \rightarrow T/(T\backslash NP): \lambda R (R . a)$



Conj

X:A CONJ X':A' = X'':lambda S (A . S & A' . S)

Composition

X/Y:A Y/Z:B => X/Z: lambda Q (A . (B . Q))

Type raising

NP:a -> T/(T\NP): lambda R (R . a)

John likes ...

np:j (s\np)/np:lambda Y lambda X likes(X,Y)

---T

s/(s\np): lambda R (R . j)

-----FC

s/np:

lambda Q (A . (B . Q))

lambda Q ((lambda R (R . j)) . (lambda Y lambda X likes(X,Y) . Q))

lambda Q ((lambda R (R . j)) . (lambda X likes(X,Q)))

lambda Q (lambda X likes(X,Q) . j)

lambda Q (likes(j,Q))

Conj

X:A CONJ X':A' = X'':lambda S (A . S & A' . S)

Composition

X/Y:A Y/Z:B => X/Z: lambda Q (A . (B . Q))

Type raising

NP:a -> T/(T\NP): lambda R (R . a)

... Mary dislikes ...

np:m (s\np)/np:lambda Y lambda X dislikes(X,Y)

---T

s/(s\np): lambda R (R . m)

-----FC

s/np:

lambda X (A . (B . Q))

lambda Q ((lambda R (R . m)) . (lambda Y lambda X dislikes(X,Y) . Q))

lambda Q ((lambda R (R . m)) . (lambda X dislikes(X,Q)))

lambda Q (lambda X dislikes(X,Q) . m)

lambda Q (dislikes(m,Q))

Conj

X:A CONJ X':A' = X'':lambda S (A . S & A'. S)

Composition

X/Y:A Y/Z:B => X/Z: lambda Q (A . (B . Q))

Type raising

NP:a -> T/(T\NP): lambda R (R . a)

John	likes	and	Mary dislikes	Bob
....		CONJ	np:b
-----FC			-----FC	

s/np: lambda Q (likes(j,Q))

s/np: lambda Q (dislikes(m,Q))

----- CONJ

s/np: lambda S (lambda Q (likes(j,Q)) . S & lambda Q (dislikes(m,Q)) . S)

s/np: lambda S (likes(j,S) & dislikes(m,S))

s : lambda S (likes(j,S) & dislikes(m,S)) . b

s : likes(j,b) & dislikes(m,b)

Compositionality and Incrementality

- Compositionality:
 - all constituents have a denotation
- Incrementality:
 - all initial substrings have a denotation
 - all substrings have a denotation (stronger)

Categorical Unification Grammar

- Extending the formalism to allow features:
 - agreement, grammatical relations
- Embedding CCG techniques in other formalisms
 - SUBCAT, predicate/arguments

the np/n

boy n

boys n

walk s\np

walks s\np

Forward application

$X/Y \ Y \Rightarrow X$

Backward application

$Y \ X \backslash Y \Rightarrow X$

Thus

the	boy	walks
np/n	n	s\np

----- FA

np

----- BA

s

the [cat: np]/[cat: n]
 boy [cat: n]
 boys [cat: n]
 walk [cat: s]\[cat: np]
 walks [cat: s]\[cat: np]

Forward application

$X/Y Y \Rightarrow X$

Backward application

$Y X \backslash Y \Rightarrow X$

Thus

the	boy	walks
[cat: np]/[cat: n]	[cat: n]	[cat: s]\[cat: np]
FA		
[cat: np]		
	BA	
	[cat: s]	

the [cat: np num: !X]/
 [cat: n num: !X]
 boy [cat: n num: sg]
 boys [cat: n num: pl]
 walk [cat: s]\
 [cat: np num: pl]
 walks [cat: s]\
 [cat: np num: sg]

Thus

the	boy	walks
[cat: np	[cat: n	[cat: s]\
num: !X]/	num: sg]	[cat: np
[cat: n		num: sg]
num: !X]		

----- FA

[cat: np
num: sg]

----- BA

[cat: s]

the [cat: np num: !X]/
 [cat: n num: !X]
 boy [cat: n num: sg]
 boys [cat: n num: pl]
 walk [cat: s]\
 [cat: np num: pl]
 walks [cat: s]\
 [cat: np num: sg]

Thus

the	boys	walk
[cat: np	[cat: n	[cat: s]\
num: !X]/	num: pl]	[cat: np
[cat: n		num: pl]
num: !X]		

----- FA

[cat: np
num: pl]

----- BA

[cat: s]

the [cat: np num: !X]/
 [cat: n num: !X]
 boy [cat: n num: sg]
 boys [cat: n num: pl]
 walk [cat: s]\
 [cat: np num: pl]
 walks [cat: s]\
 [cat: np num: sg]

Thus

the	boys	walks
[cat: np	[cat: n	[cat: s]\
num: !X]/	num: pl]	[cat: np
[cat: n		num: sg]
num: !X]		

----- FA

[cat: np
 num: pl]

----- ****

the [becomes [cat: np num: !X]
direction: forward
wants: [cat: n num: !X]]
boy [cat: n num: sg]
walks [becomes: [cat: s]
direction: backward
wants: [cat: np num: sg]]

Forward Application

[becomes: !X
direction: forward + !Y = !X
wants: !Y]

Backward Application

[becomes: !X
!Y + direction: backward = !X
wants: !Y]

the [becomes [cat: np num: !X]
 direction: forward
 wants: [cat: n num: !X]]
 boy [cat: n num: sg]
 walks [becomes: [cat: s]
 direction: backward
 wants: [cat: np num: sg]]

Thus

the		boy	walks
[becomes [cat: np num: !X]			...
direction: forward		[cat: n num: sg]	
wants: [cat: n num: !X]]			
----- FA			
	[cat: np		
	num: sg]		
	----- BA		
	[cat: s]		

SUBCAT feature

- In GPSG and HPSG:
 - SUBCAT feature identifies features of argument(s)
 - [SUBCAT [NP]]
- This is actually CG-like:
 - $S \backslash NP$ is a verb looking for one argument
 - $(S \backslash NP) / NP$ is a verb looking for two arguments
- Typing the arguments:
 - number/case etc

CG Parsing

□ CG:

forward application	$A/B + B = A$
backward application	$B + A \setminus B = A$
composition	$A/B + B/C = A/C$
conjunctions	$A \text{ CONJ } A' = A''$

□ Simple bottom-up algorithm:

- basic chart technique
- Insert lexical entries in chart
- **do**
- **for each** v vertice
- **for each** e_1 edge from v
- **for each** e_2 edge from e_1 's end
- **if** $e_1 + e_2$ match rule
- apply and insert new edge
- **until** new new edges inserted

CCG Parsing

forward application	$A/B + B = A$
backward application	$B + A \backslash B = A$
composition	$A/B + B/C = A/C$
conjunctions	$A \text{ CONJ } A' = A''$
type raising	$A = X/(X \backslash A)$

- Simple bottom-up algorithm:
 - basic chart technique
 - Insert lexical entries in chart
 - apply restricted type raising
 - **do**
 - **for each** v vertice
 - **for each** $e1$ edge from v
 - **for each** $e2$ edge from $e1$'s end
 - **if** $e1 + e2$ match rule
 - apply and insert new edge
 - **until** new new edges inserted

Some other examples

- gapping:
 - Harry will buy bread and Barry, potatoes

Some Properties

- mildly context sensitive
- weakly equivalent to LTAG:
 - Lexicalized Tree Adjoining Grammars
- Relationship to intonation:
 - Steedman 1991
- Lambek calculus
 - change argument order
- Complexity:
 - unrestricted type raising: unbounded
 - restricted type raising: $O(n^3)$
 - on CCG minus CONJ which can be $(S \setminus S) / S$