15-859(B) Machine Learning Theory

Lecture 14: Learning from noisy data, intro to SQ model

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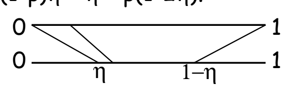
Learning when there is no perfect hypothesis

- Hoeffding/Chernoff bounds: minimizing training error will approximately minimize true error: just need $O(1/\varepsilon^2)$ samples versus $O(1/\varepsilon)$.
- What about polynomial-time algorithms? Seems harder.
  - Given data set S, finding apx best conjunction is NP-hard.
  - Can do other things, like minimize hinge-loss, maxent type loss, but not directly connected to error rate.
- One way to make progress: make assumptions on the "noise" in the data. E.g., Random Classification Noise model.

Learning from Random Classification Noise

- PAC model, target $f \in C$, but assume labels from noisy channel.
- "noisy" Oracle $EX^{\eta}(f,D)$. η is the noise rate.
  - Example x is drawn from D.
  - With probability $1-\eta$ see label $\ell(x) = f(x)$.
  - With probability η see label $\ell(x) = 1-f(x)$.
- E.g., if h has non-noisy error p, what is the noisy error rate?
  - $p(1-\eta) + (1-p)\eta = \eta + p(1-2\eta)$.

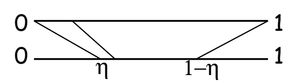


Learning from Random Classification Noise

Algorithm A PAC-learns C from random classification noise if for any $f \in C$, any distrib D, any $\eta < 1/2$, any $\varepsilon, \delta > 0$, given access to $EX^{\eta}(f,D)$, A finds a hyp h that is ε-close to f, with probability $\geq 1-\delta$.

Allowed time poly$(1/\varepsilon, 1/\delta, 1/(1-2\eta), n, size(f))$

- Q: is this a plausible goal? We are asking the learner to get closer to f than the data is.
- A: OK because noisy error rate is linear in true error rate (squashed by $1-2\eta$)



Notation

- Use "Pr[…]" for probability with respect to non-noisy distribution.
- Use "Pr_{η}[…]" for probability with respect to noisy distribution.

Learning OR-functions (assume monotone)

- Let's assume noise rate η is known. Any ideas?
- Say $p_i = Pr[f(x)=0 \wedge x_i=1]$
- Any h that includes all x_i such that $p_i=0$ and no x_i such that $p_i > \varepsilon/n$ is good.
- So, just need to estimate p_i to $\pm\, \varepsilon/2n$.
  - Rewrite as $p_i = Pr[f(x)=0|x_i=1] \times Pr[x_i=1]$.
  - 2nd part unaffected by noise (and if tiny, can ignore x_i). Define q_i as 1st part.
  - Then $Pr_{\eta}[\ell(x)=0|x_i=1] = q_i(1-\eta)+(1-q_i)\eta = \eta+q_i(1-2\eta)$.
  - So, enough to approx LHS to $O((\varepsilon/n)(1-2\eta))$.

Learning OR-functions (assume monotone)

- If noise rate not known, can estimate with smallest value of $Pr_\eta[\ell(x)=0|x_i=1]$.

Generalizing the algorithm

Basic idea of algorithm was:

- See how can learn in non-noisy model by asking about probabilities of certain events with some "slop".
- Try to learn in noisy model by breaking events into:
  - Parts predictably affected by noise.
  - Parts unaffected by noise.

Let's formalize this in notion of "statistical query" (SQ) algorithm. Will see how to convert any SQ alg to work with noise.

The Statistical Query Model

- No noise.
- Algorithm asks: "what is the probability a labeled example will have property χ? Please tell me up to additive error τ."
  - Formally, $\chi: X \times \{0,1\} \to \{0,1\}$. Must be poly-time computable. $\tau \geq 1/\text{poly}(\ldots)$.
  - Let $P_\chi = Pr[\chi(x,f(x))=1]$.
  - World responds with $P'_\chi \in [P_\chi-\tau, P_\chi+\tau]$.

  [can extend to [0,1]-valued or vector-valued χ]
- May repeat poly(...) times. Can also ask for unlabeled data. Must output h of error $\leq \varepsilon$. No δ in this model.

The Statistical Query Model

- Examples of queries:
  - What is the probability that $x_i=1$ and label is negative?
  - What is the error rate of my current hypothesis h? [$\chi(x,\ell)=1$ iff $h(x) \neq \ell$]
- Get back answer to $\pm\tau$. Can simulate from $\approx 1/\tau^2$ examples. [That's why need $\tau \geq 1/\text{poly}(\ldots)$.]
- To learn OR-functions, ask for $Pr[x_i=1 \wedge f(x)=0]$ with $\tau = \varepsilon/(2n)$. Produce OR of all x_i such that $P'_\chi \leq \varepsilon/(2n)$.

The Statistical Query Model

- Many algorithms can be simulated with statistical queries:
  - Perceptron: ask for $E[f(x)x : h(x)\neq f(x)]$ (formally define vector-valued $\chi = x$ if $h(x)\neq f(x)$, and 0 otherwise. Then divide by $Pr[h(x)\neq f(x)]$.)
  - Hill-climbing type algorithms: what is error rate of h? What would it be if I made this tweak?
- Properties of SQ model:
  - Can automatically convert to work in presence of classification noise.
  - Can give a nice characterization of what can and cannot be learned in it.

SQ-learnable ⇒ (PAC+Noise)-learnable

- Given query χ, need to estimate from noisy data. Idea:
  - Break into part predictably affected by noise, and part unaffected.
  - Estimate these parts separately.
  - Can draw fresh examples for each query or estimate many queries from same sample if VCDim of query space is small.
- Running example: $\chi(x,\ell)=1$ iff $x_i=1 \wedge \ell=0$.

How to estimate $Pr[\chi(x,f(x))=1]$?

- Let CLEAN = $\{x : \chi(x,0) = \chi(x,1)\}$
- Let NOISY = $\{x : \chi(x,0) \neq \chi(x,1)\}$
  - What are these for $\chi(x,\ell)=1$ iff $x_i=1 \wedge \ell=0$?
- Now we can write:
  - $Pr[\chi(x,f(x))=1] = Pr[\chi(x,f(x))=1 \wedge x\in CLEAN] + Pr[\chi(x,f(x))=1 \wedge x\in NOISY]$.
- Step 1: first part is easy to estimate from noisy data (easy to tell if $x \in$ CLEAN).
- What about the 2nd part?

How to estimate $Pr[\chi(x,f(x))=1]$?

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  - What are these for $\chi(x,\ell)=1$ iff $x_i=1 \wedge \ell=0$?
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- Can estimate $Pr[x\in NOISY]$.
- Also estimate $P_\eta \equiv Pr_\eta[\chi(x,\ell)=1 \mid x\in NOISY]$.
- Want $P \equiv Pr[\chi(x,f(x))=1 \mid x\in NOISY]$.
- Write $P_\eta = P(1-\eta) + (1-P)\eta = \eta + P(1-2\eta)$.
- So, $P = (P_\eta - \eta)/(1-2\eta)$.
  - Just need to estimate P_η to additive error $\tau(1-2\eta)$.
  - If don't know η, can have "guess and check" wrapper around entire algorithm.

Characterizing what's learnable using SQ algorithms

- Key tool: Fourier analysis of boolean functions.
- Sounds scary but it's a cool idea!
- Let's think of functions from $\{0,1\}^n \rightarrow \{-1,1\}$.
- View function f as a vector of 2^n entries:

$(D[000]^{1/2}f(000), D[001]^{1/2}f(001), \ldots, D[x]^{1/2}f(x), \ldots)$

- What is $\langle f, f\rangle$? What is $\langle f, g\rangle$?
- What is an orthonormal basis? Will see connection to SQ algs next time…