



**CMU 15-381**

**LECTURE 17:**

**GAME THEORY 1**

**TEACHERS:**

**EMMA BRUNSKILL**

**ARIEL PROCACCIA (THIS TIME)**

# THE PRISONER'S DILEMMA

- Two men are charged with a crime
- They are told that:
  - If one rats out and the other does not, the rat will be freed, other jailed for nine years
  - If both rat out, both will be jailed for six years
- They also know that if neither rats out, both will be jailed for one year



# THE PRISONER'S DILEMMA

	Cooperate	Defect
Cooperate	-1,-1	-9,0
Defect	0,-9	-6,-6

Click: what would you do?



# UNDERSTANDING THE DILEMMA

- Defection is a **dominant** strategy
- But the players can do much better by cooperating
- Related to the **tragedy of the commons**



# IN REAL LIFE

- Presidential elections
  - Cooperate = positive ads
  - Defect = negative ads
- Nuclear arms race
  - Cooperate = destroy arsenal
  - Defect = build arsenal
- Climate change
  - Cooperate = curb CO<sub>2</sub> emissions
  - Defect = do not curb



# THE PROFESSOR'S DILEMMA

		Class	
		Listen	Sleep
Professor	Make effort	$10^6, 10^6$	$-10, 0$
	Slack off	$0, -10$	$0, 0$

**Click: dominant strategies?**

# NASH EQUILIBRIUM

- Each player's strategy is a **best response** to strategies of others
- Formally, a **Nash equilibrium** is a vector of strategies  $s=(s_1,\dots,s_n)$  such that  $\forall i \in N, \forall s'_i \in S_i, u_i(s) \geq u_i(s'_i, s_{-i})$
- <https://www.youtube.com/watch?v=CemLiSI5oX8>
- **Click:** what are the Nash equilibria of the professor's dilemma?





# ROCK-PAPER-SCISSORS

	R	P	S
R	0,0	-1,1	1,-1
P	1,-1	0,0	-1,1
S	-1,1	1,-1	0,0

Click: Nash equilibrium?





# MIXED STRATEGIES

- A **mixed strategy** is a probability distribution over pure strategies
- The mixed strategy of player 1 is  $x = (x_1, \dots, x_n)$ ,  $x_i = \Pr[1 \text{ plays strategy } s_i]$
- Similarly, player 2 plays  $y = (y_1, \dots, y_n)$
- The utility of player  $i$  is 
$$u_i(x, y) = \sum_{j,k} [x_j \cdot y_k \cdot u_i(s_j, s_k)]$$



# EXERCISE: MIXED NE

- $x=(1/2, 1/2, 0)$ ,  
 $y=(0, 1/2, 1/2)$ ,  
what is  $u_1(x,y)$ ?
- $x=(1/3, 1/3, 1/3)$ ,  
 $y=(1/3, 1/3, 1/3)$ ,  
what is  $u_1(x,y)$ ?
- $x=(1/2, 1/2, 0)$ ,  
 $y=(1/2, 1/2, 0)$ , is NE?
- $x=(1/3, 1/3, 1/3)$ ,  
 $y=(1/3, 1/3, 1/3)$ , is NE?

	R	P	S
R	0,0	-1,1	1,-1
P	1,-1	0,0	-1,1
S	-1,1	1,-1	0,0

# NASH'S THEOREM

- **Theorem [Nash, 1950]:** if everything is finite then there exists at least one (possibly mixed) Nash equilibrium
- However, how does one *compute* a Nash equilibrium?
- Computing a Nash equilibrium is computationally hard [Daskalakis et al., STOC 2006]



# APPLICATION: INTERDOMAIN ROUTING

- Internet composed of smaller networks called *autonomous systems* (AS)
- Owned by competing entities (Microsoft, AT&T, etc.)
- Interdomain routing = establishing routes between ASes
- Standard protocol: BGP



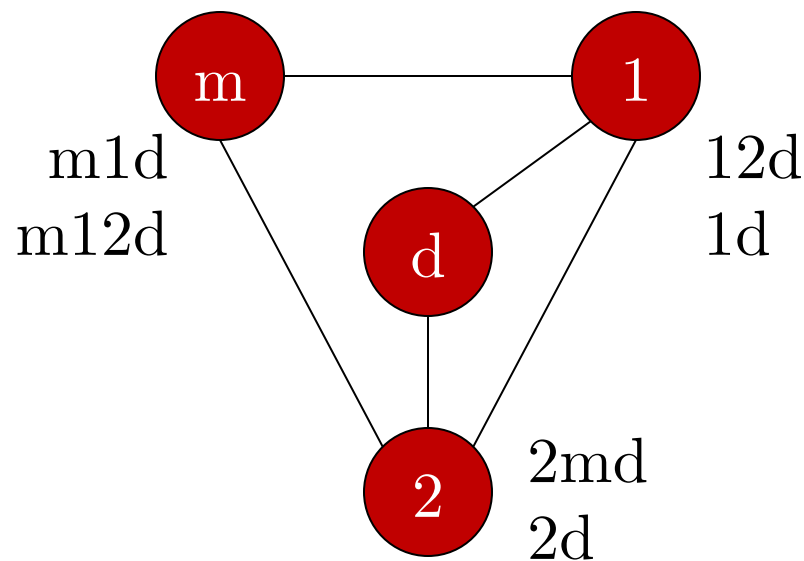
# APPLICATION: INTERDOMAIN ROUTING

- Graph with  $n$  source nodes (players) and a destination node
- Each player has preferences over routes to the destination
- Under BGP ASes continuously:
  - Receive updates about routes of neighbors
  - Choose a neighbor to send traffic to
  - Announce new route to neighboring nodes



# APPLICATION: INTERDOMAIN ROUTING

- Theorem [Levin et al, STOC 2008]: Following BGP is not an (ex-post) NE
- BGP converges to the NE (12d,2d,m12d)
- But... if m repeatedly announces to 2 the route md
- 2 would go with 2md
- 1 would go with 1d
- m gets m1d!



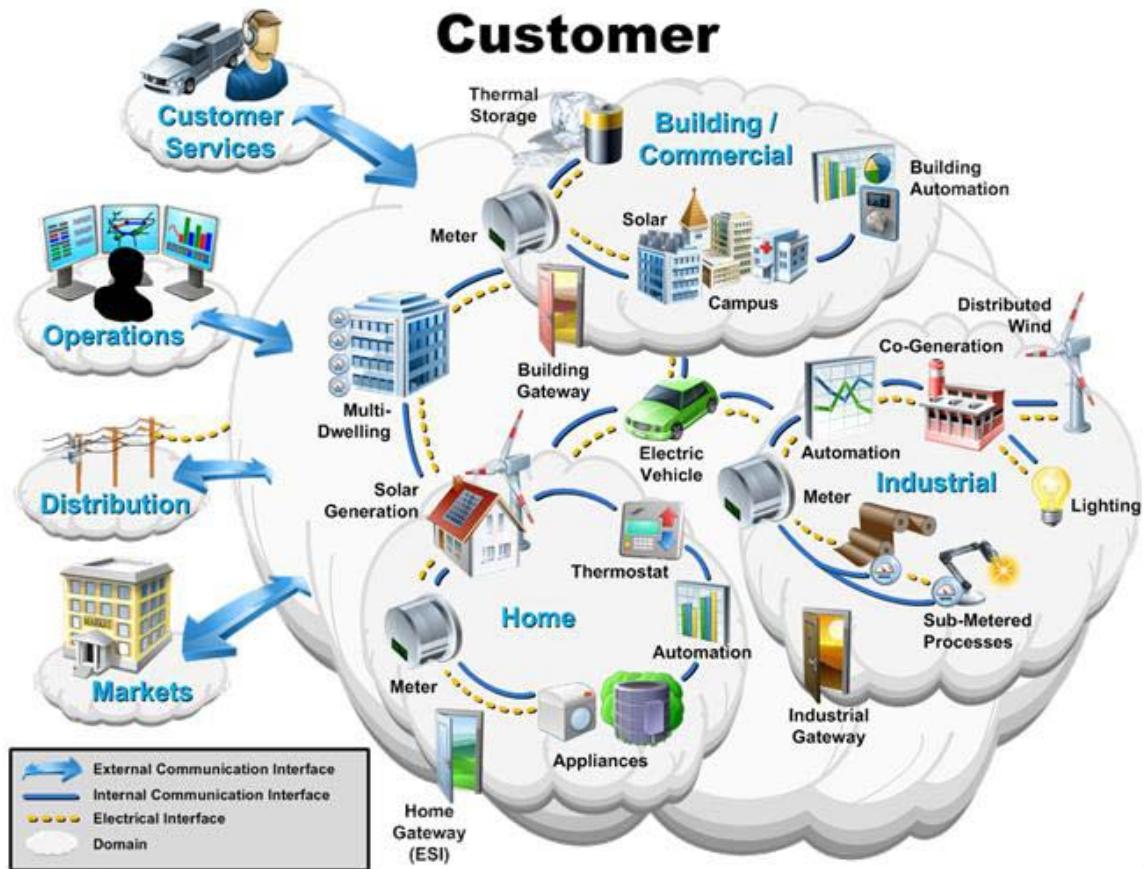
# APPLICATION: INTERDOMAIN ROUTING

- Route verification = players can verify that neighbors' declared paths actually exist
- **Theorem [Levin et al., STOC 2008]:**  
Assuming route verification (+mild technical condition), following BGP is an (ex-post) Nash equilibrium!
- Provides partial explanation for why interdomain routing functions so well!



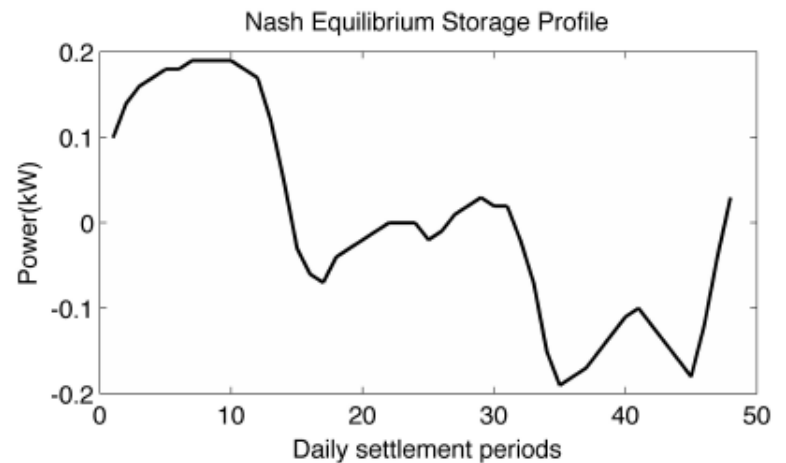
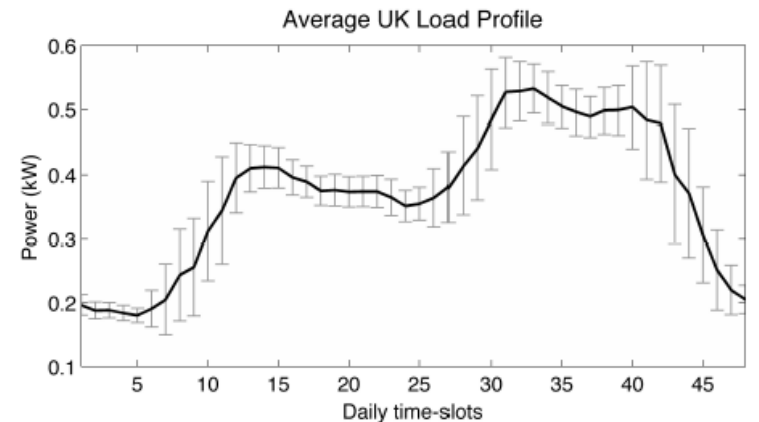


# APPLICATION: SMART GRID



# APPLICATION: SMART GRID

- Energy storage devices advocated for saving energy in future smart grid
- Bad if all are charged at the same time
- Solution: agent-based management system that allows storage devices to converge to equilibrium [Vytelingum et al., AAMAS 2010]



# APPLICATION: SMART GRID

- Strategy of an agent: how much to charge in each half hour of the day
- The behavior of electricity suppliers is specified by a supply curve
- Equilibrium can be analytically computed
- Simulations show that in eq., savings of 13% on electricity bill in UK

