

Program Representations

17-654/17-765

Analysis of Software Artifacts

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Representing Programs



- To analyze software automatically, we must be able to represent it precisely
- Some representations
 - Source code
 - Abstract syntax trees
 - Control flow graph
 - Bytecode
 - Assembly code
 - Binary code

The WHILE Language



- A simple procedural language with:
 - assignment
 - statement sequencing
 - conditionals
 - while loops
- Used in early papers (e.g. Hoare 69) as as a “sandbox” for thinking about program semantics
- We will use it to illustrate several different kinds of analysis

WHILE Syntax



- Categories of syntax
 - $S \in \mathbf{Stmt}$ statements
 - $a \in \mathbf{AExp}$ arithmetic expressions
 - $x, y \in \mathbf{Var}$ variables
 - $n \in \mathbf{Num}$ number literals
 - $b \in \mathbf{BExp}$ boolean expressions
- Syntax
 - $S ::= x := a \mid \text{skip} \mid S_1; S_2 \mid \text{if } b \text{ then } S_1 \text{ else } S_2 \mid \text{while } b \text{ do } S$
 - $a ::= x \mid n \mid a_1 \text{ op}_a a_2$
 - $\text{op}_a ::= + \mid - \mid * \mid / \mid \dots$
 - $b ::= \text{true} \mid \text{false} \mid \text{not } b \mid b_1 \text{ op}_b b_2 \mid a_1 \text{ op}_r a_2$
 - $\text{op}_b ::= \text{and} \mid \text{or} \mid \dots$
 - $\text{op}_r ::= < \mid \leq \mid = \mid > \mid \geq \mid \dots$

Example WHILE Program



```

y := x;
z := 1;
while y>1 do
  z := z * y;
  y := y - 1
    
```

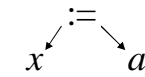
Computes the factorial function, with the input in x and the output in z

Abstract Syntax Trees

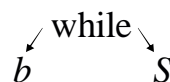


- A tree representation of source code
- Based on the language grammar
 - One type of node for each production

• $S ::= x := a$



• $S ::= \text{while } b \text{ do } S$



Parsing: Source to AST



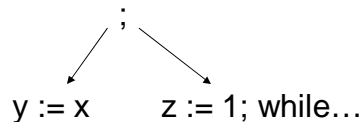
- Parsing process (top down)
 1. Determine the top-level production to use
 2. Create an AST element for that production
 3. Determine what text corresponds to each child of the AST element
 4. Recursively parse each child
- Algorithms have been studied in detail
 - For this course you only need the intuition
 - Details covered in compiler courses

Parsing Example



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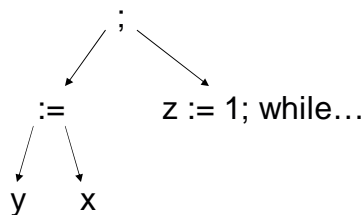
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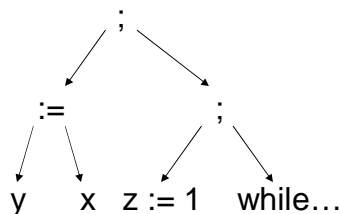
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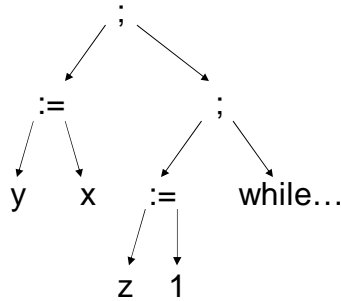
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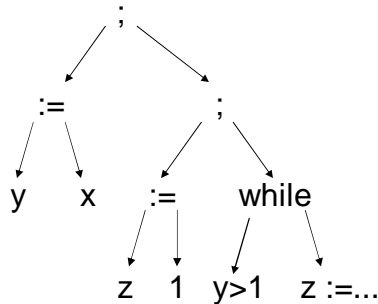
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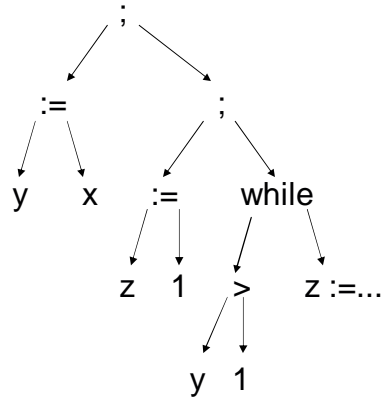
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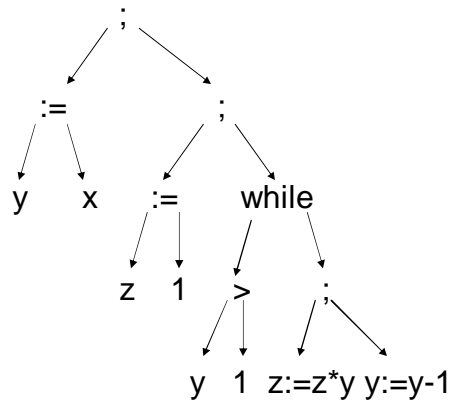
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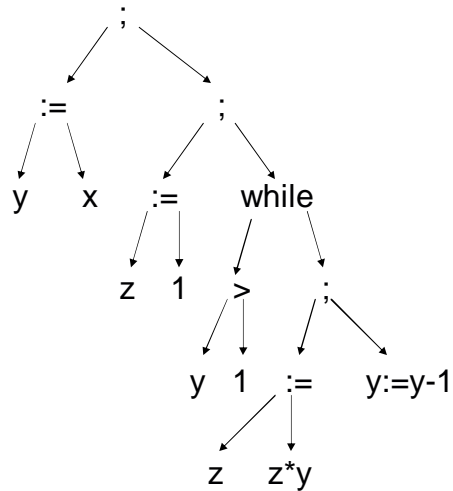
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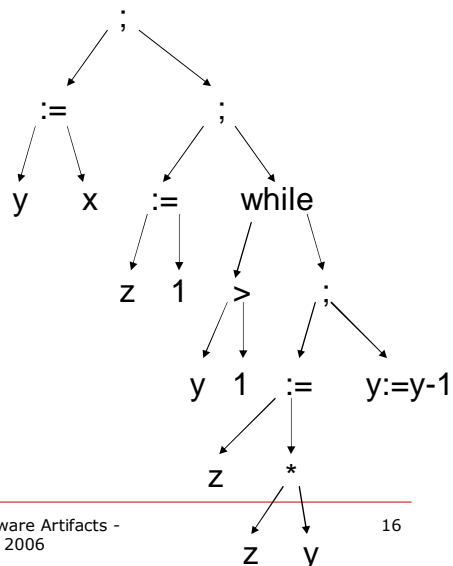
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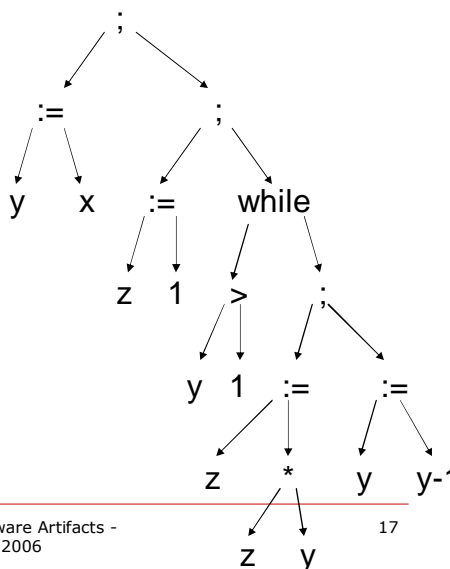
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Spring 2006

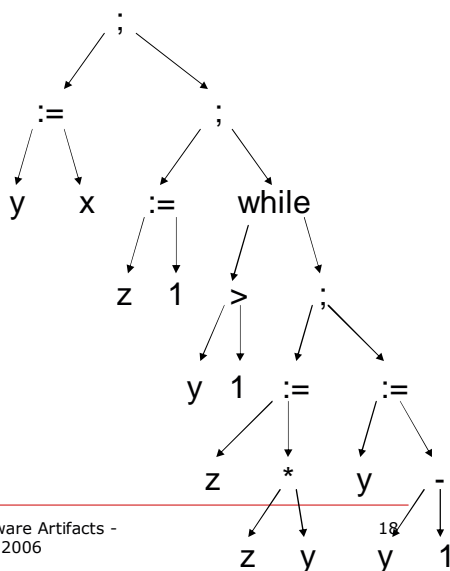
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WHILE ASTs in Java



- Java data structures mirror grammar

- $S ::= x := a$
 - | skip
 - | $S_1; S_2$
 - | if b then S_1 else S_2
 - | while b do S

```
class AST { ... }
class Stmt extends AST { ... }
class Assign extends Stmt {
    Var var;
    AExpr expr;
}
class Skip extends Stmt { }
class Seq extends Stmt {
    Stmt left;
    Stmt right;
}
class If extends Stmt {
    BExpr cond;
    Stmt thenStmt;
    Stmt elseStmt;
}
class While extends Stmt {
    BExpr cond;
    Stmt body;
}
```

Course Analysis Toolkit



- Eclipse
 - Open-source Java integrated development environment
 - Extensible through plugins
- Crystal
 - Plugin for Eclipse
 - Provides a Java AST for analysis
 - Eclipse has its own AST, but it's not ideal for analysis
 - Provides a basic analysis framework
 - Supports interaction with end user

Extending Crystal



- Download and install Java version 5 or 6
- Download and install Eclipse 3.2
- Download and install Crystal
- Implement a class that extends:
 - `ICrystalAnalysis` for global analyses
 - `AbstractCrystalMethodAnalysis` for method-at-a-time analyses
 - This will usually be the case
- Register your new analysis with Crystal
 - It can then be run from the Crystal menu

AbstractCrystalMethodAnalysis



- `public void beforeAllMethods() { }`
 - Called at the beginning of an analysis cycle
 - Use for analysis setup
- `public abstract void analyzeMethod(IMethodDeclarationNode d);`
 - Invoked by the framework for each method in the system
 - You must override this to perform your analysis task for each method
- `public void afterAllMethods() { }`
 - Called at the end of an analysis cycle
 - Use for analysis cleanup and any reporting that's still left

Example: PrintMethods



```
Crystal crystal = Crystal.getInstance();
public void beforeAllMethods() {
    crystal.userOut().println("Printing methods:");
}
public void analyzeMethod(IMethodDeclarationNode md) {
    crystal.userOut().println(md.getId());
}
public void afterAllMethods() {
    crystal.userOut().println("Done.");
}
```

Registering the Analysis



- In CrystalPlugin.java:

```
public void setupCrystalAnalyses(
    Crystal crystal) {
    PrintMethods pm = new PrintMethods();
    crystal.registerAnalysis(pm);
}
```

The Crystal AST



- View Tree in Eclipse
 - Interface package: com.surelogic.ast.java.operator
- Browse hierarchy
 - IJavaOperatorNode – root of tree
 - IDeclarationNode – class, field, methods
 - IClassDeclarationNode - classes
 - IMethodDeclarationNode - methods
 - IFieldDeclarationNode and DeclaratorNode – “int i,j=6,k;” vs “j=6”
 - IStatementNode
 - IBlockStatementNode – blocks { }
 - IDeclStatementNode – variables
 - If, For, Return...
 - IExprStatementNode – expression statements (incl. assignments)
 - IReturnTypeNameNode – void, other types
 - IPrimitiveTypeNode: boolean, int, float...
 - IReferenceTypeNode: int[], String, Foo<T>

The Crystal AST



- Browse hierarchy (continued)
 - IInitializerNode → IExpressionNode – expressions
 - IBinopExpressionNode: binary expressions
 - IArithBinopExpressionNode: + - * / etc.
 - IAssignmentExpressionNode: assignment
 - IAssignExpressionNode (regular)
 - IOpAssignExpressionNode (+=, -=, *=, etc.)
 - IUnopExpressionNode: !b, x++
 - ILiteralExpressionNode: 5, 7.3, true, “hello”, null
 - IPrimaryExpressionNode
 - IAllocationExpressionNode: new X()
 - IFieldRefNode: this.f
 - ISomeFunctionCallNode: x.m(5)
 - ISomeThisExpression: this, X.this
 - ITypeExpressionNode: used as the “receiver” of static methods (artificial)
 - IVariableUseExpressionNode: x, y
 - IBinding – binds a node to a declaration
 - example: call.resolveBinding(), type.resolveType()
 - IDeclarationNode: declarations are bindings (getNode() returns this)
 - IFieldDeclarationNode
 - IFunctionBinding: gets the function declaration
 - IType
 - IPrimitiveType – binds to IPrimitiveTypeNode
 - IDerivedRefType – arrays, etc.
 - IDeclaredType – a class or interface
 - ITypeFormal – a type parameter

Demo



- Installing Crystal
- Run Assignment 0
- Look at Assignment 0 code
- Look at Visitor
- Results of Assignment 1

The Visitor Pattern



```
interface IVisitor<T> {  
    // one for each element type  
    T visit(Element e);  
}  
  
class DescendingVisitor<T>  
    implements IVisitor<T> {  
    T visit(Element e) {  
        T rv = null;  
        for(Element child : e.getChildren()) {  
            rv = child.accept(this);  
        }  
        return rv;  
    }  
}  
  
interface INode {  
    <T> T accept(Visitor<T> v);  
}  
class Element implements INode {  
    <T> T accept(Visitor<T> v) {  
        return visitor.visit(this);  
    }  
}
```