

Program Representations

17-654/17-765

Analysis of Software Artifacts

Jonathan Aldrich



Representing Programs



- To analyze software automatically, we must be able to represent it precisely
- Some representations
 - Source code
 - Abstract syntax trees
 - Control flow graph
 - Bytecode
 - Assembly code
 - Binary code

The WHILE Language



- A simple procedural language with:
 - assignment
 - statement sequencing
 - conditionals
 - while loops
- Used in early papers (e.g. Hoare 69) as as a “sandbox” for thinking about program semantics
- We will use it to illustrate several different kinds of analysis

WHILE Syntax



- Categories of syntax
 - $S \in \mathbf{Stmt}$ statements
 - $a \in \mathbf{AExp}$ arithmetic expressions
 - $x, y \in \mathbf{Var}$ variables
 - $n \in \mathbf{Num}$ number literals
 - $b \in \mathbf{BExp}$ boolean expressions
- Syntax
 - $S ::= x := a \mid \text{skip} \mid S_1; S_2 \mid \text{if } b \text{ then } S_1 \text{ else } S_2 \mid \text{while } b \text{ do } S$
 - $a ::= x \mid n \mid a_1 \text{ op}_a a_2$
 - $\text{op}_a ::= + \mid - \mid * \mid / \mid \dots$
 - $b ::= \text{true} \mid \text{false} \mid \text{not } b \mid b_1 \text{ op}_b b_2 \mid a_1 \text{ op}_r a_2$
 - $\text{op}_b ::= \text{and} \mid \text{or} \mid \dots$
 - $\text{op}_r ::= < \mid \leq \mid = \mid > \mid \geq \mid \dots$

Example WHILE Program



```

y := x;
z := 1;
while y > 1 do
  z := z * y;
  y := y - 1
    
```

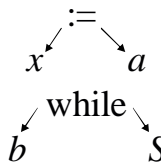
Computes the factorial function, with the input in x and the output in z

Abstract Syntax Trees

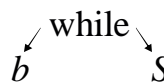


- A tree representation of source code
- Based on the language grammar
 - One type of node for each production

• $S ::= x := a \quad \rightarrow$



• $S ::= \text{while } b \text{ do } S \quad \rightarrow$



Parsing: Source to AST



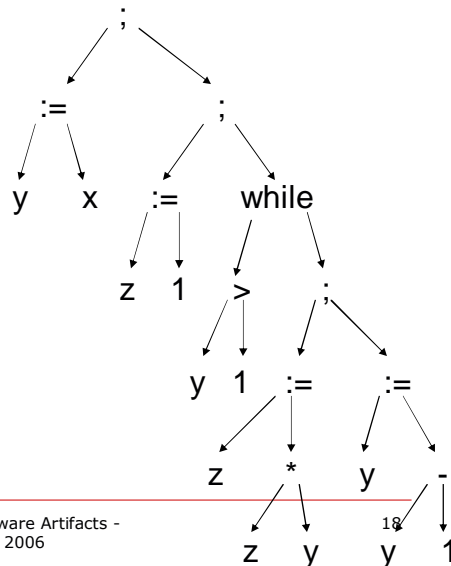
- Parsing process (top down)
 1. Determine the top-level production to use
 2. Create an AST element for that production
 3. Determine what text corresponds to each child of the AST element
 4. Recursively parse each child
- Algorithms have been studied in detail
 - For this course you only need the intuition
 - Details covered in compiler courses

Parsing Example



```

y := x;
z := 1;
while y>1 do
  z := z * y;
  y := y - 1
    
```



- Top-level production?
 - $S_1; S_2$
- What are the parts?
 - $y := x$
 - $z := 1; \text{while } \dots$

WHILE ASTs in Java



- Java data structures mirror grammar

- $S ::= x := a$
 | skip
 | $S_1; S_2$
 | if b then S_1 else S_2
 | while b do S

```
class AST { ... }
class Stmt extends AST { ... }
class Assign extends Stmt {
    Var var;
    AExpr expr;
}
class Skip extends Stmt { }
class Seq extends Stmt {
    Stmt left;
    Stmt right;
}
class If extends Stmt {
    BExpr cond;
    Stmt thenStmt;
    Stmt elseStmt;
}
class While extends Stmt {
    BExpr cond;
    Stmt body;
}
```

Course Analysis Toolkit



- Eclipse
 - Open-source Java integrated development environment
 - Extensible through plugins
 - Exposes Java AST to plugins
- Crystal
 - Plugin for Eclipse
 - Provides a basic analysis framework
 - Supports interaction with end user

Extending Crystal



- Download and install Java 5
- Download and install Eclipse 3.1
- Download and install Crystal
- Implement a class that extends:
 - `ICrystalAnalysis` for global analyses
 - `AbstractCrystalMethodAnalysis` for method-at-a-time analyses
 - This will usually be the case
- Register your new analysis with Crystal
 - It can then be run from the Crystal menu

AbstractCrystalMethodAnalysis



- `public void beforeAllMethods() { }`
 - Called at the beginning of an analysis cycle
 - Use for analysis setup
- `public abstract void analyzeMethod(MethodDeclaration d);`
 - Invoked by the framework for each method in the system
 - You must override this to perform your analysis task for each method
- `public void afterAllMethods() { }`
 - Called at the end of an analysis cycle
 - Use for analysis cleanup and any reporting that's still left

Example: PrintMethods



```
Crystal crystal = Crystal.getInstance();
public void beforeAllMethods() {
    crystal.userOut().println("Printing methods:");
}
public void analyzeMethod(MethodDeclaration md) {
    crystal.userOut().println(md.getName());
}
public void afterAllMethods() {
    crystal.userOut().println("Done.");
}
```

Registering the Analysis



- In CrystalPlugin.java:

```
public void setupCrystalAnalyses(
    Crystal crystal) {
    PrintMethods pm = new PrintMethods();
    crystal.registerAnalysis(pm);
}
```

The Eclipse AST



- View Tree
- Browse javadoc for:
 - MethodDeclaration
 - Block
 - Statement
 - VariableDeclarationStatement
 - VariableDeclaration
 - ExpressionStatement
 - Assignment
 - Name
 - IVariableBinding

ASTNodes and Bindings



- ASTNode
 - The AST representation of Java source
 - There will be an ASTNode for each occurrence of a variable in the source
- Binding
 - A single canonical object representing the variable
 - Eclipse doesn't provide a way to get from the Binding to the ASTNode
 - Efficiency choice
 - Crystal provides a convenient shortcut
 - `ASTNode Utilities.getASTNode(IBinding b)`

Demo



- Installing Crystal
- Run Assignment 0
- Look at Assignment 0 code
- Look at Visitor
- Results of Assignment 1

The Visitor Pattern



```
class Visitor {
    // called before visit
    void preVisit(Node n) {}
    // if return true, children visited
    boolean visit(Element e) {
        return true; }
    // called after child visits
    void endVisit(Element e) {
        return true; }
    // called after visit
    void postVisit(Node n) {}
}

class Node {
    abstract void accept(Visitor
    v);
}
class Element extends Node {
    void accept(Visitor v) {
        v.preVisit(this);
        boolean c = v.visit(this);
        if (c)
            children.accept(v);
        v.endVisit(this);
        v.postVisit(this);
    }
}
```