

Predicate-based Test Coverage and Generation

17-654/17-765
Analysis of Software Artifacts
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These slides prepared by Thomas Ball, with additional material from M. Young, A. Memon and MSR's FSE group. Used by permission.

MSIL Unit Test Tool

→ *a hybrid helper*

- Goal
 - capture developer knowledge **ASAP** via a **strong** set of unit tests to form a **specification** of the code's behavior
- How
 - generate tests based on analysis of MSIL
 - symbolic execution + constraint satisfaction
 - runtime analysis to check complicated invariants
- Facets
 - complements specification-based test generation
 - positive feedback cycle with programmer

What criteria should guide unit test generation?

Predicate-complete Testing

- Predicates
 - relational expression such as $(x < 0)$
 - the expression $(x < 0) \parallel (y > 0)$ has two predicates
 - predicates come from program and safe runtime semantics
- Consider a program with m statements and n predicates
 - predicates partition input domain
 - $m \times 2^n$ possible *observable* states S
- Goal of Predicate-complete Testing:
 - cover all *reachable* observable states $R \subseteq S$

PCT Coverage

L2: if (A || B) S else T
L3: if (C || D) U else V

- PCT requires covering all logical combinations over $\{A, B, C, D\}$ at
 - L2 and L3
 - S, T, U and V
- Some combinations may not be reachable

PCT Coverage Subsumes Statement, Edge Coverage

- Statement coverage
 - Must cover all statements with all predicates
- Edge coverage
 - Must cover if statements with all predicates
 - Therefore touch both branches
- Path coverage
 - Must cover every path through program
 - Infinite number of them!

PCT Coverage does not imply Path Coverage

```

L1: if (x<0) (L1, x<0) (L1, !(x<0))
L2: skip; (L2, x<0)
   else
L3: x = -2; (L3, !(x<0))
L4: x = x + 1; (L4, x<0)
L5: if (x<0) (L5, x<0) (L5, !(x<0))
L6: A; (L6, x<0)
    
```

PCT Coverage does not imply Path Coverage

```

L1: if (x<0) (L1, x<0) (L1, !(x<0))
L2: skip; (L2, x<0)
   else
L3: x = -2; (L3, !(x<0))
L4: x = x + 1; (L4, x<0)
L5: if (x<0) (L5, x<0) (L5, !(x<0))
L6: A; (L6, x<0)
    
```

Three tests

$x \rightarrow -1$

PCT Coverage does not imply Path Coverage

```

L1: if (x<0) (L1, x<0) (L1, !(x<0))
L2: skip; (L2, x<0)
   else
L3: x = -2; (L3, !(x<0))
L4: x = x + 1; (L4, x<0)
L5: if (x<0) (L5, x<0) (L5, !(x<0))
L6: A; (L6, x<0)
    
```

Three tests

$x \rightarrow -1$ $x \rightarrow 2$

PCT Coverage does not imply Path Coverage

```

L1: if (x<0) (L1, x<0) (L1, !(x<0))
L2: skip; (L2, x<0)
   else
L3: x = -2; (L3, !(x<0))
L4: x = x + 1; (L4, x<0)
L5: if (x<0) (L5, x<0) (L5, !(x<0))
L6: A; (L6, x<0)
    
```

Three tests

$x \rightarrow -1$ $x \rightarrow 2$ $x \rightarrow -3$

Path Coverage does not imply PCT Coverage

```

L1: if (p) (L1, p $\bar{q}$ ) (L1,  $\bar{p}q$ ) (L1, p $\bar{q}$ ) (L1,  $\bar{p}\bar{q}$ )
L2: if (q) (L2, p $\bar{q}$ ) (L2, p $q$ )
L3: x=0; (L3, p $\bar{q}$ )
L4: y=p+q; (L4, p $\bar{q}$ ) (L4, p $q$ ) (L4,  $\bar{p}\bar{q}$ )
           ↑ path (L1, p $\bar{q}$ )
           L1 L2 L3 L4
           ↑ path (L1,  $\bar{p}q$ )
           L1 L2 L4
    
```

Path Coverage does not imply PCT Coverage

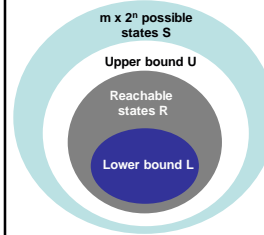
```

L1: if (p) (L1, p $\bar{q}$ ) (L1,  $\bar{p}q$ ) (L1, p $\bar{q}$ ) (L1,  $\bar{p}\bar{q}$ )
L2: if (q) (L2, p $\bar{q}$ ) (L2, p $q$ )
L3: x=0; (L3, p $\bar{q}$ )
L4: y=p+q; (L4, p $\bar{q}$ ) (L4, p $q$ ) (L4,  $\bar{p}\bar{q}$ )
           ↑ path (L1, p $\bar{q}$ )
           L1 L2 L3 L4
           ↑ path (L1,  $\bar{p}q$ )
           L1 L2 L4
           ↑ path (L1,  $\bar{p}\bar{q}$ )
           L1 L4
    
```

Denominator Problem

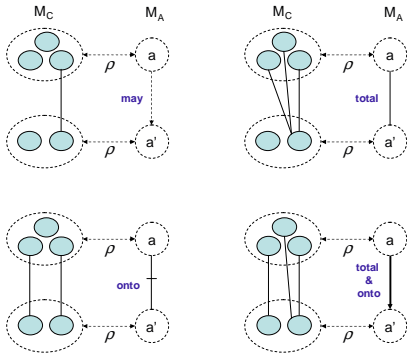
- Coverage metrics require a denominator
 - e.g. statements executed / total statements
- Easy to define for observable states
 - executed observable states / ($m \times 2^n$)
- But ($m \times 2^n$) is not a very good denominator!
 - most observable states will not be reachable
 - $R \lll S$

Upper and Lower Bounds

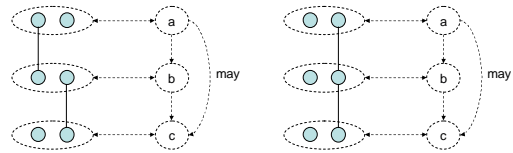


- Bound reachable observable states
 - modal transition systems and predicate abstraction
 - $|L| / |U|$ defines "goodness" of abstraction
- Test generation using lower bound L
- Refinement to increase $|L| / |U|$ ratio

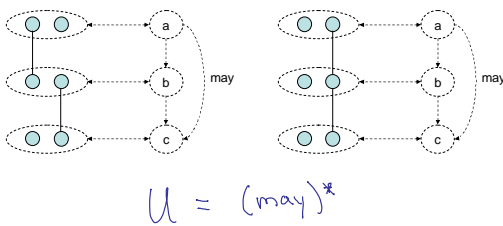
Abstraction Construction



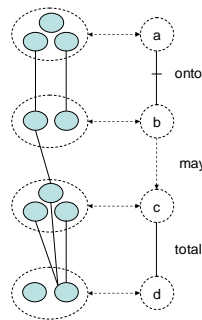
Upper Bound: May-Reachability

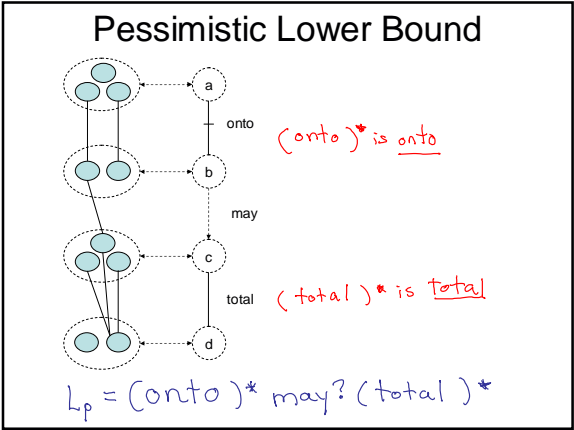
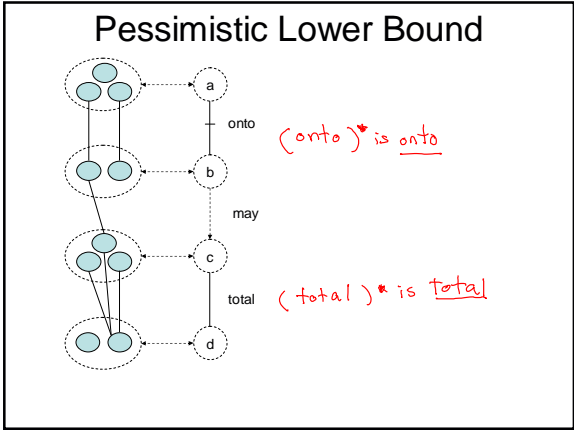


Upper Bound: May-Reachability



Pessimistic Lower Bound





Example

```

void partition(int a[]) {
  assume(a.length()>2);
  int pivot = a[0];
  int lo = 1;
  int hi = a.length()-1;
  while (lo<=hi) {
    while (a[lo]<=pivot)
      lo++;
    while (a[hi]>pivot)
      hi--;
    if (lo<hi)
      swap(a,lo,hi);
  }
}

```

Observation Vector

[lo<hi, lo<=hi, a[lo]<=pivot, a[hi]>pivot]

- lo<hi \Rightarrow lo<=hi
- $\neg \text{lo} < \text{hi} \wedge \text{lo} < = \text{hi} \Rightarrow (\text{a}[\text{lo}] < = \text{pivot} \wedge \neg \text{a}[\text{hi}] > \text{pivot}) \vee (\neg \text{a}[\text{lo}] < = \text{pivot} \wedge \text{a}[\text{hi}] > \text{pivot})$

Only 10/16 observations possible

```

void partition(int a[]) {
  assume(a.length()>2);
  int pivot = a[0];
  int lo = 1;
  int hi = a.length()-1;

L0: while (lo<=hi) {
L1:   ;
L2:   while (a[lo]<=pivot) {
L3:     lo++;
L4:   ;
L5:   while (a[hi]>pivot) {
L6:     hi--;
L7:   ;
L8:   if (lo<hi) {
L9:     swap(a,lo,hi);
LA:   ;
LB: ;}
LC: ;
}
}

```

13 labels x 10 observations = 130 observable states

But, program constrains **reachable** observable states greatly.

Boolean Program

```

void partition() {
  decl lt, le, al, ah;
  enforce ( (lt>le) &
    ( (!lt&le) => (al&!ah) | (!al&ah) ) );
  lt,le,al,ah := T,T,*,*;
L0: while (le) {
L1:   ;
L2:   while (al) {
L3:     lt,le,al := (!lt ? F:*), lt, *;
L4:   ;}
L5:   while (ah) {
L6:     lt,le,ah := (!lt ? F:*), lt, *;
L7:   ;}
L8:   if (lt) {
L9:     al,ah := !ah,!al;
LA: ;}
LB: ;}
LC: ;
}

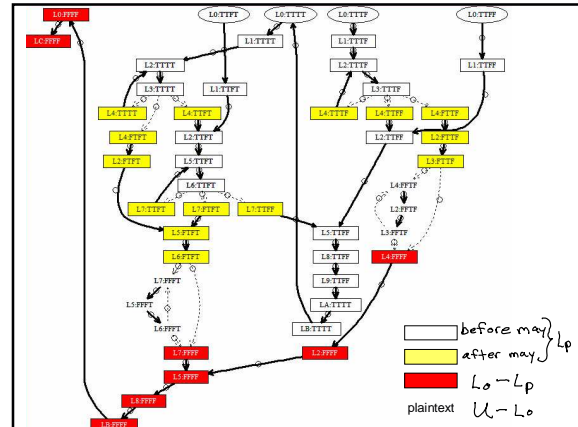
```

State Space of Boolean Program

[$lo < hi$, $lo \leq hi$, $a[lo] \leq pivot$, $a[hi] > pivot$]

	TTTT	TTTF	FTTF	FTTF	TFTF	FFTF	FTTF	FTTF	FTTF
L0	x	x			x			x	x
L1	x	x	x		x			x	
L2	x	x	x	x	x	x		x	x
L3	x	x	x	x					
L4	x	x	x	x	x	x		x	x
L5					x	x	x	x	x
L6					x	x	x	x	x
L7					x	x	x	x	x
L8								x	x
L9								x	
LA	x								
LB	x								x
LC									x

Upper Bound = 49 states



Test Generation

- DFS of L_p generates covering set of paths
- Symbolically execute paths to generate tests
- Run program on tests to find errors and compute coverage of observable states

Generated Inputs

```
void partition(int a[]) {
    assume(a.length()>2);
    int pivot = a[0];
    int lo = 1;
    int hi = a.length()-1;

    L0: while (lo<=hi) {
    L1: ;
    L2: while (a[lo]<=pivot) {
    L3:   lo++;
    L4:   ;
    L5: while (a[hi]>pivot) {
    L6:   hi--;
    L7:   ;
    L8: if (lo<hi) {
    L9:   swap(a,lo,hi);
    LA: ;
    LB: ;
    LC: ;
    }
```

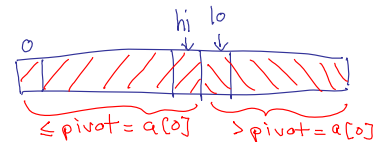
Array bounds violations

(L0:TTTT,L4:FTTF)	{0,-8,1}
(L0:TTTT,L4:TTTF)	{0,-8,2,1}
(L0:TTTT,L4:TTTT)	{0,-8,-8,1}
(L0:TTTF,L4:TTTF)	{1,-7,3,0}
(L0:TTTF,L4:FTTF)	{0,-7,-8}
(L0:TTTF,L4:TTTT)	{1,-7,-7,0}
(L0:TTTF,L7:TTTF)	{0,2,-8,1}
(L0:TTTF,L7:FTTF)	{0,1,2}
(L0:TTTF,L7:TTTT)	{0,3,1,2}
(L0:TTTF,L0:TTTT)	{1,2,-1,0}

Results

- Buggy partition function
 - U=49, L=43, Tested=42
- Fixed partition function
 - U=56, L=37, Tested=43
- What about the remaining 13 states?

Refinement



Existing predicates do not precisely track boundary condition
 $lo < hi$, $lo \leq hi$, $a[lo] \leq pivot$, $a[hi] > pivot$

New Observation Vector

```
[ lo<hi, lo<=hi, lo=hi+1,  
  a[lo]<=pivot, a[hi]>pivot,  
  a[lo-1]<=pivot, a[hi+1]>pivot  
]
```

Only 48/128 observations possible

For this set of predicates, $L_p = U$

Conclusions

- PCT coverage
 - new form of state-based coverage
 - similar to path coverage but finite
- Upper and lower bounds
 - computed using predicate abstraction and modal transitions
 - use lower bound to guide test generation
 - refine bounds