Homework 7: Axiomatic semantics and Hoare-style verification

17-355/17-665/17-819: Program Analysis
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Due: Tuesday, March 26 11:59 pm

100 points total

Assignment Objectives:

- Demonstrate understanding of verification condition generation and use in verifying programs.
- Write new Hoare Rules/axiomatic semantics for a particular language.
- Reason about soundness/completeness of a system for axiomatic semantics.

Handin Instructions. Please submit your assignment on Canvas as a **PDF** by the due date. Name it **[your-andrew-id]-hw7.pdf**.

Question 1, VCGen, *(18 points).* Consider the following rules for VCGen. We saw the first two in class, the third is a new proposed rule for `let`:

$$\begin{array}{lcl} VC(S_1; S_2, Q) & = & VC(S_1, VC(S_2, Q)) \\ VC(x := e, Q) & = & [e/x]Q \\ VC(\texttt{let}\ x = e\ \texttt{in}\ S, Q) & = & [e/x]VC(S, Q) \end{array}$$

The rule for `let` is incorrect. (a) Explain why in English prose, and then (b) give a correct rule for `let`.

Question 2, Let rule soundness, *(12 points).* Given $\{P\}\ S\ \{Q\}$, we desire that $P \Rightarrow VC(c, Q) \Rightarrow WP(c, Q)$. We say that our VC rules are *sound* if $\vDash \{VC(S, Q)\}\ S\ \{Q\}$. Demonstrate the unsoundness of the buggy `let` rule above by giving/showing the following six things:

1. a statement S and
2. a post-condition Q and
3. a state E, all such that
4. $E \vDash VC(S, Q)$ and
5. $\langle S, E\rangle \Downarrow E'$ but
6. $E' \nvDash Q$

Question 3, Do-while, *(16 points).* Write a sound and complete Hoare rule for `do` S `while` b. The statement has the standard semantics (i.e., S is executed at least once, before b is tested). You do not need to formally prove soundness/completeness, but make sure the rule makes sense.

*This homework was developed together with Claire Le Goues

Question 4, Loop proof obligations, *(18 points).* Consider the following program:

```
{N > 0 }
i := 0;
{ i <= N }
while (i < N) do
    { i <= N }
     i := i + 1
    { i <= N }
{i=N}
```

Assuming the loop invariant $i \leq N$, write the proof obligations for the `while` loop. Don't solve/prove them, just write them out. The form of your answer should be three mathematical implications, one for each of the following proof obligations:

- Invariant is initially true:
- Invariant is preserved by the loop body:
- Invariant and exit condition imply postcondition:

Question 5, Soundness/Completeness, *(24 points).* Consider the following three Hoare rules:

$$\frac{\vdash \{X\}\ S\ \{b \Rightarrow X \wedge \neg b \Rightarrow Y\}}{\vdash \{b \Rightarrow X \wedge \neg b \Rightarrow Y\}\ \texttt{while}\ b\ \texttt{do}\ S\ \{Y\}}\ \text{rule1} \qquad \frac{\vdash \{X \wedge b\}\ S\ \{X\}}{\vdash \{X\}\ \texttt{while}\ b\ \texttt{do}\ S\ \{X\}}\ \text{rule2}$$

$$\frac{\vdash \{X\}\ S\ \{X\}}{\{X\}\ \texttt{while}\ b\ \texttt{d}\ S\ \{X \wedge \neg b\}}\ \text{rule3}$$

Recall that a system of axiomatic semantics is *sound* if everything we can prove is also true: if $\vdash \{P\}\ S\ \{Q\}$ then $\vDash \{P\}\ S\ \{Q\}$. A system of axiomatic semantics is *complete* if we can prove all true things: if $\vDash \{P\}\ S\ \{Q\}$ then $\vdash \{P\}\ S\ \{Q\}$. All three of the rules above are sound, but none are complete.

For two of the identified incomplete rules, give/show example P, Q, E, S, and E' such that $\langle S, E\rangle \Downarrow E'$ and both $E \vDash P$ and and $E' \vDash Q$, but it is not possible given the rule to prove $\vdash \{P\}\ S\ \{Q\}$. Be clear about which rules you have chosen.

(Educational) note: An incomplete system cannot prove all possible properties or handle all possible programs. Incompleteness in an axiomatic semantics or type system is typically not as dire as unsoundness. Many research results that claim to work for the C language, for example, are actually incomplete because they fail to address, say, setjmp/longjmp or bitfields. (Many of them are also unsound because they do not correctly model various language features like unsafe casts, pointer arithmetic, or integer overflow.)

Acknowledgement: Thanks to Selva Samuel for finding a bug in the original formulation of this question back in 2016.