



# Rethinking the Service Model: Scaling Ethernet to a Million Nodes

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# Vision: More Ethernet Switches Fewer IP Routers

- Today's world: IP routers + Ethernet PHY
  - Ethernet is the dominant PHY layer
  - Large number of IP routers connecting small Ethernet networks
    - E.g. CMU campus networks
- More Ethernet switches/fewer IP routers
  - ➔ large Ethernet networks
    - Enterprise/campus networks
    - Broadband access networks
    - Data center networks

# Why Large Ethernet Networks?

- Ethernet switches
  - simple, cheap, fast
  - Last fully automatic network
    - No host configuration
    - No switch configuration
  - Seamless mobility
  - Should be used to connect in the same network
- IP routers
  - Complex, expensive
  - Should be left to connect different networks

# Why Not?

## Reasons Listed In Textbooks

- Flat addressing doesn't scale
- Need to link different L2's
- Spanning tree
  - No multi-path
  - Slow fail-over
- Broadcast overhead

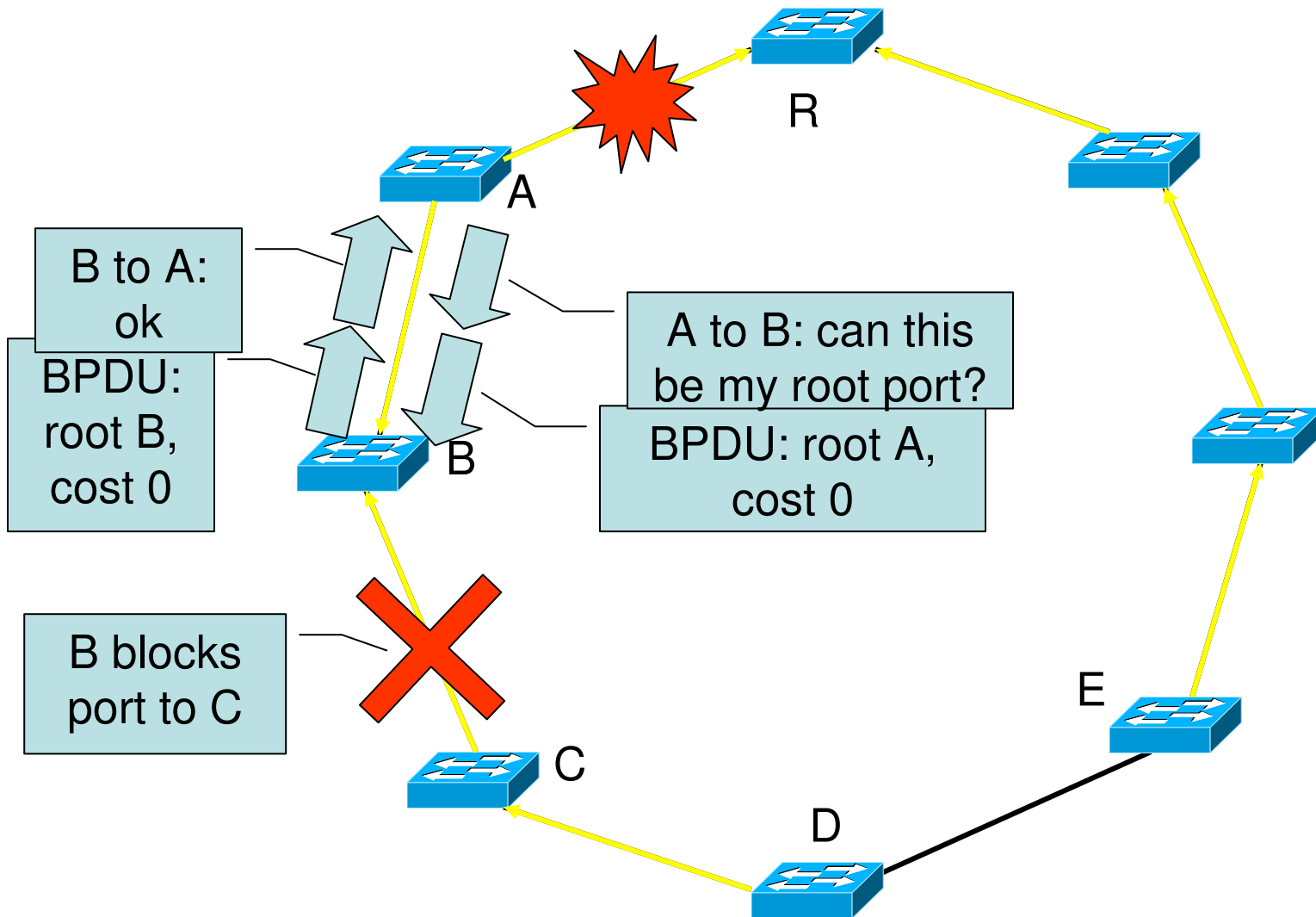
# Current Reality

- Flat addressing doesn't scale
  - Bridges with 500K-1M MAC capacity ship today
- Need to link different L2's
  - Ethernet is the only L2 left
- Spanning tree
  - ??
- Broadcast overhead
  - ??

# Outline

- Study Ethernet's flaws
  - Spanning Tree
  - Broadcast
- Identify the root cause
  - Broadcast service model
- Propose a solution
  - Turn off broadcast
  - Replace Ethernet's control plane

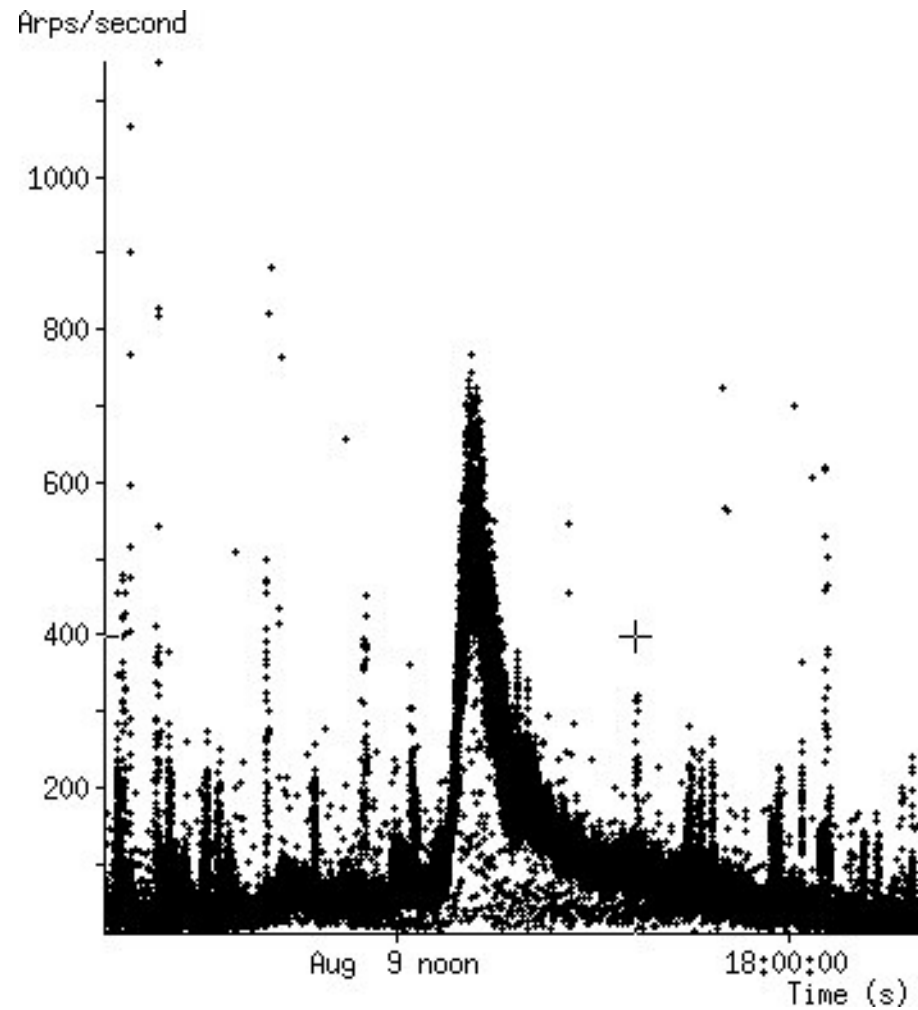
# RSTP



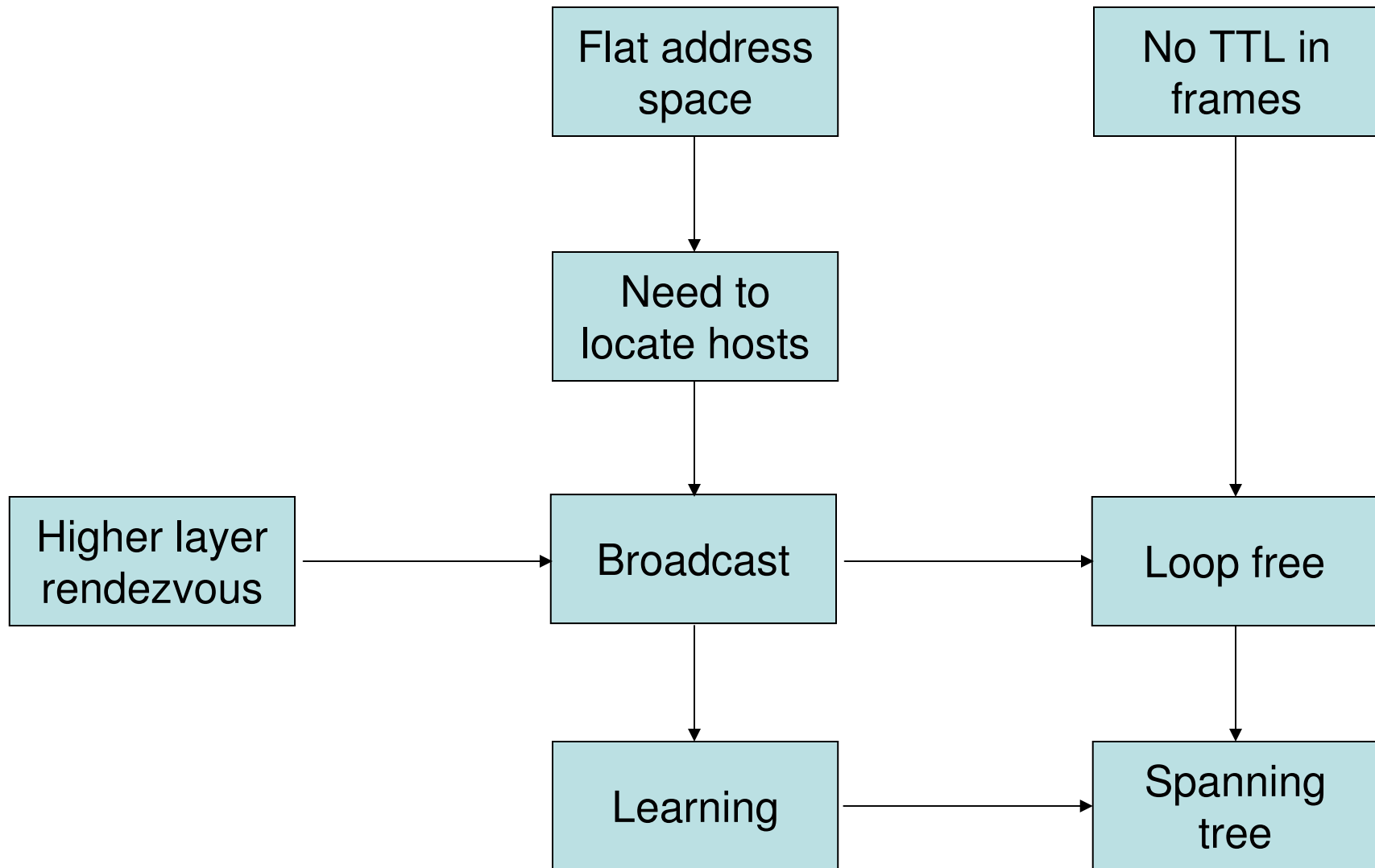




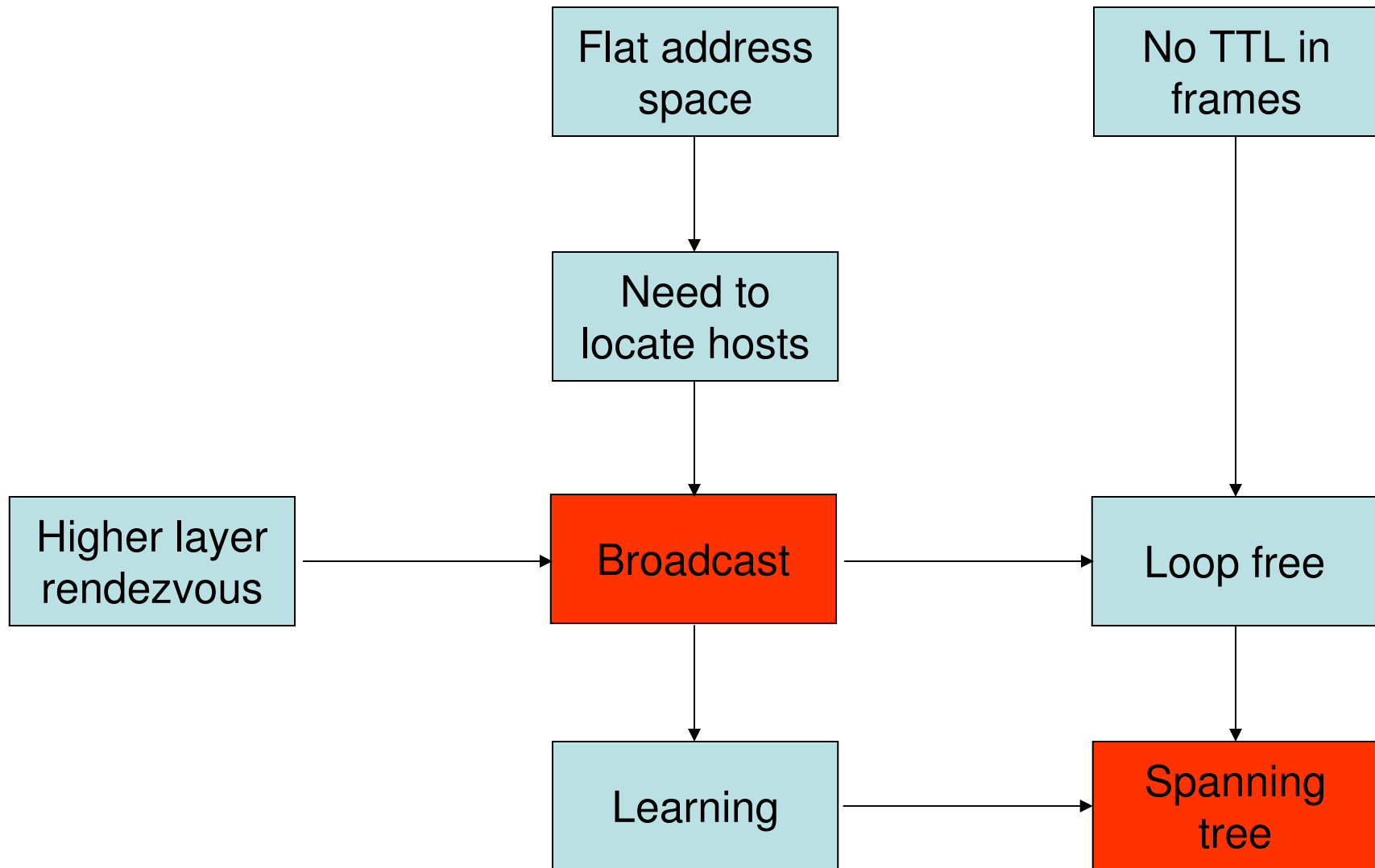
# Broadcast (ARP)



# Ethernet's Features



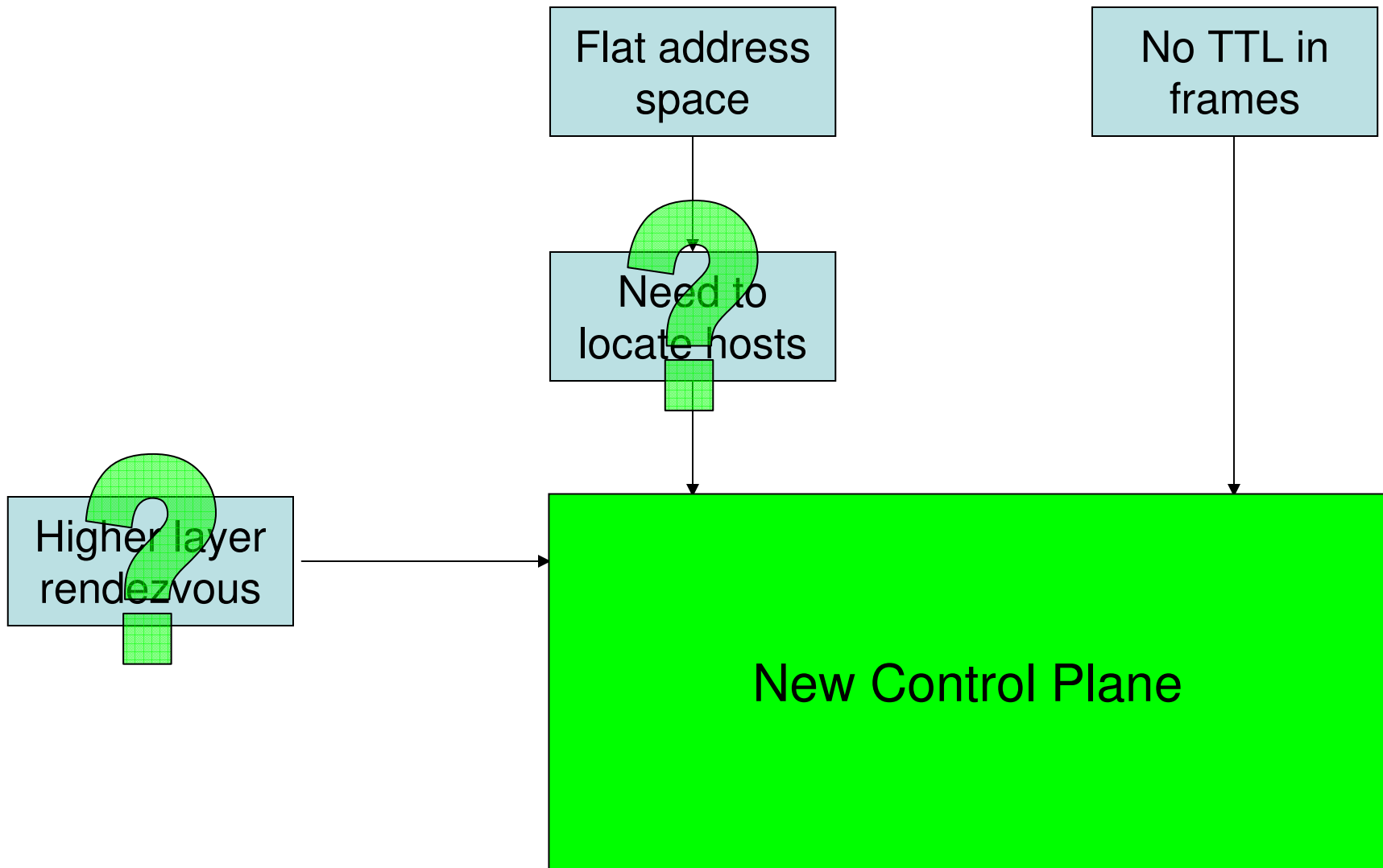
# Ethernet's Features



# Breaking the Broadcast/RSTP Dependency

- Change the service model: Turn off broadcast
  - Eliminates security risk
  - Improves scalability
  - Removes exponential packet copying
- Can eliminate RSTP
  - Unicast packets may loop, but no blowup
  - Network doesn't overload during transient loops

# Fixing Ethernet



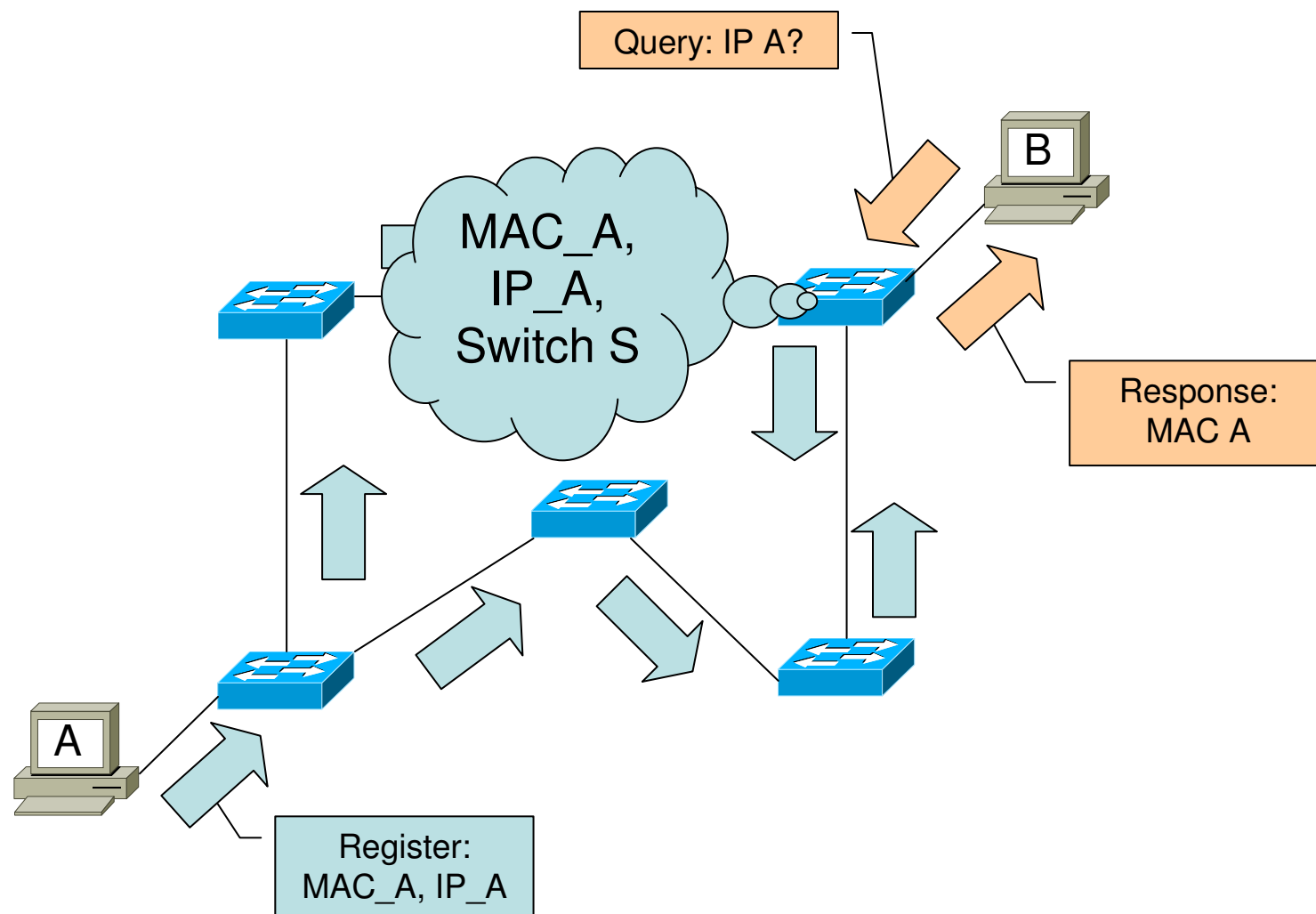
# Why Replace the Control Plane?

- Fix what's broken
- Enable extensibility
  - Faster convergence (MAN)
  - Traffic engineering (SAN)
  - Isolation (Access net)
- Two control planes to consider
  - Fully distributed
  - Thin control plane

# Fully Distributed Control Plane

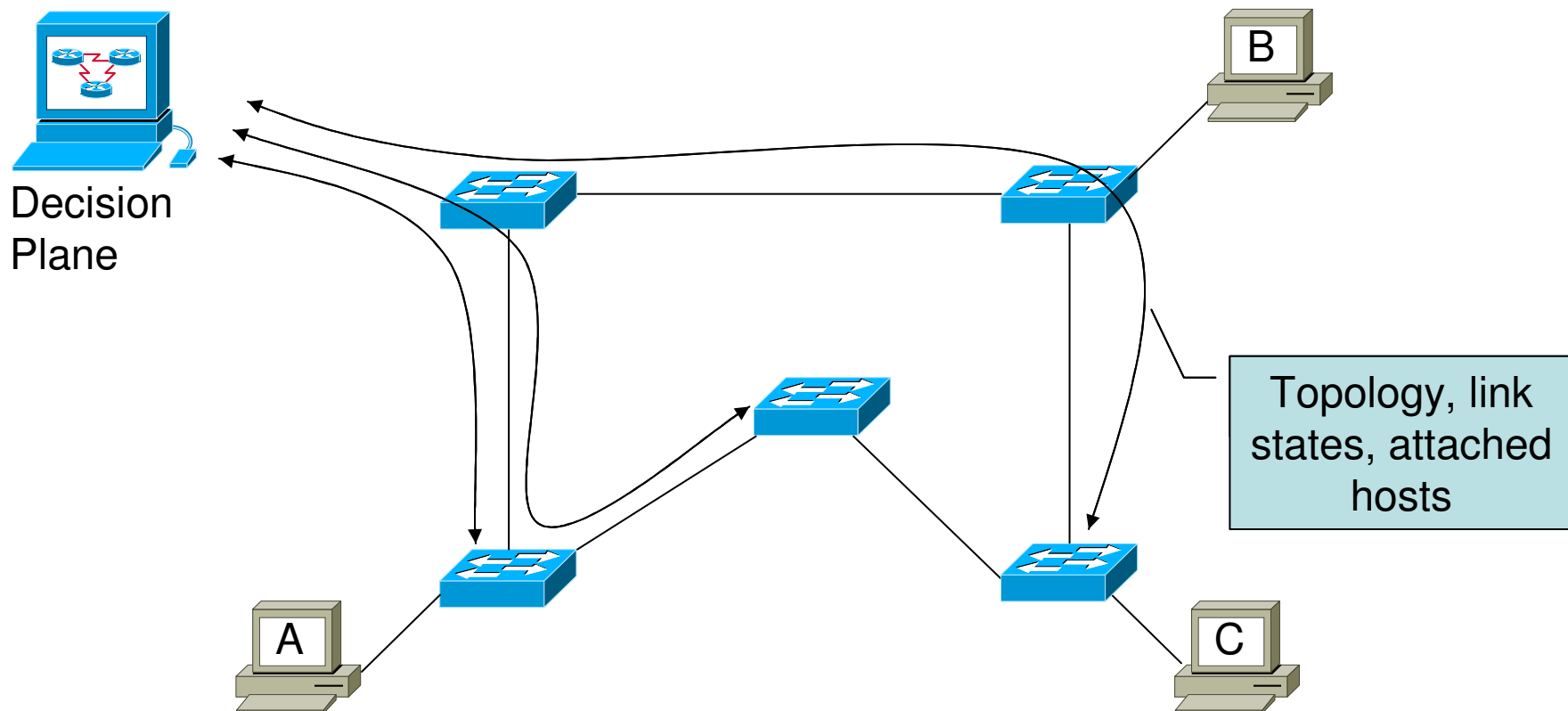
- Link state computation of forwarding paths
  - Fast convergence
  - Multiple paths, not just a spanning tree
- Distributed directory replicated at all bridges
  - Provides IP to MAC mapping
  - Also used for service location
- Hosts register with local switch

# Distributed Directory Example





# Thin Control Plane



# Thin Control Plane Advantages

- Switches remain simple
- Decisions made with global view of network
  - Multi-path forwarding
  - Directory service
- Can introduce new services
  - Traffic engineering
  - Pre-planned failure response

# Related Work

- Control plane
  - OSI's CLNP/ISIS
  - Rexford04's Thin Control Plane
- Multi-path forwarding with [R]STP
  - SmartBridge00, STAR02, Pellegrini04, Viking04
- Replacing spanning tree with link state
  - Garcia03 ("LSOM")
  - Perlman04 ("RBridges")
    - Adds header with TTL for links between bridges
    - No host registration needed

# Summary

- Vision: More switches, fewer routers
  - Ethernet switches are cheaper, less complex than IP routers
  - Leads to larger Ethernet networks
  - Many potential application scenarios
- To realize
  - Eliminate broadcast
  - New control plane to enable practical L2