

Introduction to Graphs

15-121

Introduction to Data Structures

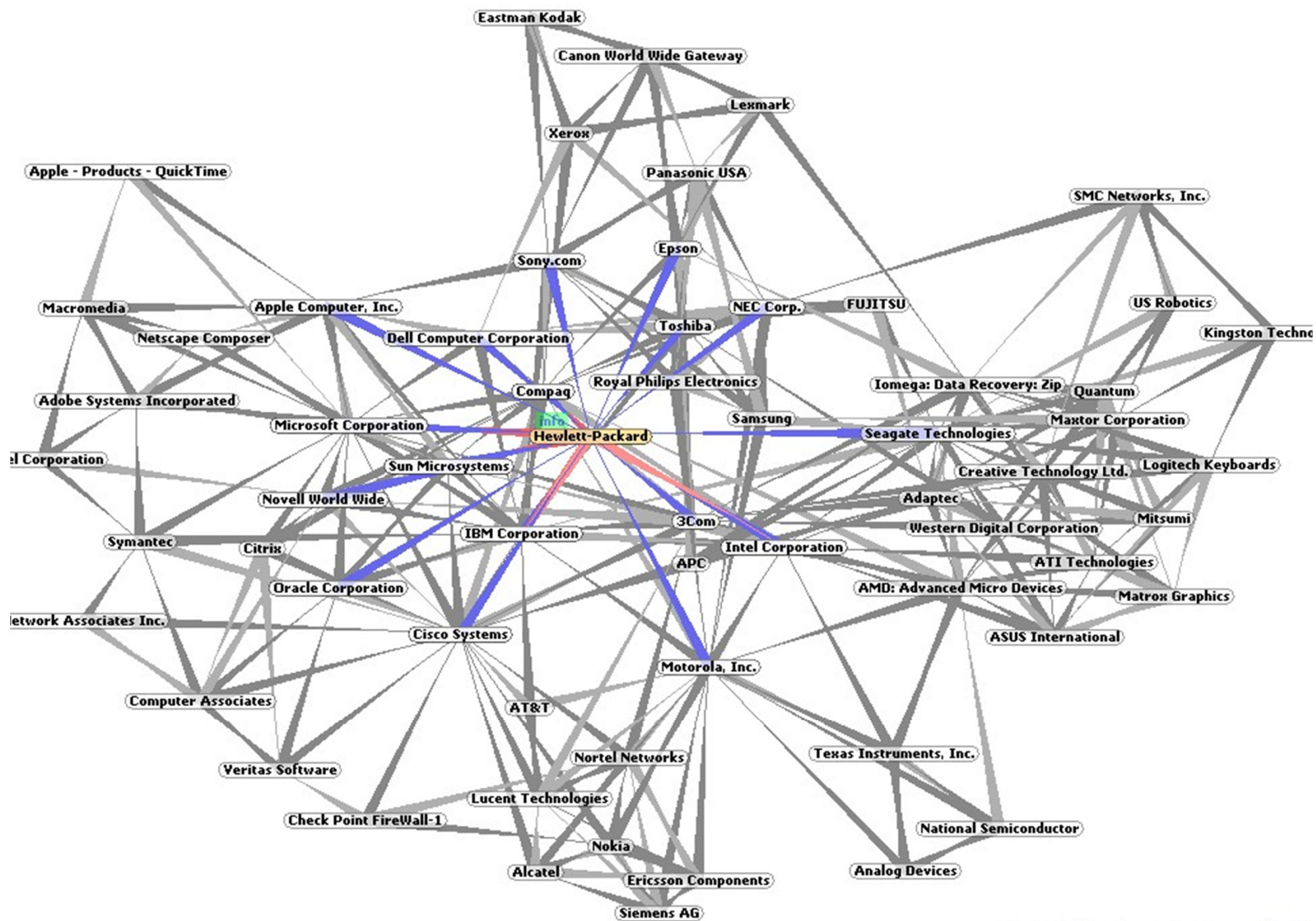
Ananda Gunawardena



Graphs are everywhere

An Airline route Map







Finding the Shortest Path

Lots of applications

Many real world problems can be modeled using graphs

- **Airline Route Map**

- What is the fastest way to get from Pittsburgh to St Louis?
- What is the cheapest way to get from Pittsburgh to St Louis?

- **Electric Circuits**

- Circuit elements - transistors, resistors, capacitors
- is everything connected together?
 - Depends on interconnections (wires)
- If this circuit is built will it work?
 - Depends on wires and objects they connect.

Graph Definitions

- Graph

- A set of vertices(nodes) $V = \{v_1, v_2, \dots, v_n\}$
- A set of edges(arcs) that connects the vertices $E = \{e_1, e_2, \dots, e_m\}$
- Each edge e_i is a pair (v, w) where v, w in V
- $|V|$ = number of vertices (cardinality)
- $|E|$ = number of edges

- Graphs can be

- directed (order (v, w) matters)
- Undirected (order of (v, w) doesn't matter)

- Edges can be

- weighted (cost associated with the edge)
- eg: Neural Network, airline route map(vanguard airlines)



Graph Representations

Graph Representation

- How do we represent a graph internally?
- Two ways
 - adjacency matrix
 - Adjacency list
- Adjacency Matrix
 - Use matrix entries to represent edges in the graph
- Adjacency List
 - Use an array of lists to represent edges in the graph (we will discuss this later)

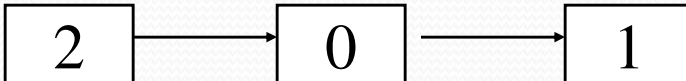
Adjacency Matrix

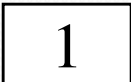
- Adjacency Matrix
 - For each edge (v,w) in E , set $A[v][w] = \text{edge_cost}$
 - Non existent edges with logical infinity
- Cost of implementation
 - $O(|V|^2)$ time for initialization
 - $O(|V|^2)$ space
 - ok for dense graphs
 - unacceptable for sparse graphs

Adjacency List

- Adjacency List
 - Ideal solution for sparse graphs
 - For each vertex keep a list of all adjacent vertices
 - Adjacent vertices are the vertices that are connected to the vertex directly by an edge.
 - Example

List 0 

List 1 

List 2 

Adjacency List

- The number of list nodes equals to number of edges
 - $O(|E|)$ space
- Space is also required to store the lists
 - $O(|V|)$ for $|V|$ lists
- Note that the number of edges is at least $\text{round}(|V|/2)$
 - assuming each vertex is in some edge
 - Therefore disregard any $O(|V|)$ term when $O(|E|)$ is present
- Adjacency list can be constructed in linear time (wrt to edges)

Breadth First Traversal

- Algorithm
 - Start from any node in the graph
 - Traverse to its neighbors (nodes that are directly connected to it) using some heuristic
 - Next traverse the neighbors of the neighbors etc.. Until some limit is reach or all the nodes in the graph are visited
 - Use a queue to perform the breadth first traversal

Depth First Traversal

- Algorithm
 - Start from any node in the graph
 - Traverse deeper and deeper until dead end
 - Back track and traverse other nodes that are not visited
 - Use a stack to perform the depth first traversal



Next: Graph Algorithms