# 36-702 Homework 1 Solution

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### Problem 1

(a) Let  $n(j) = \sum_{i} I_{\{j\}}(x_i)$ ,

$$L(\theta) = \theta^{n(1)} \cdot \left(\frac{\theta}{2}\right)^{n(2)} \cdot \left(\frac{\theta}{3}\right)^{n(3)} \cdot \left(\frac{6 - 11\theta}{6}\right)^{n(4)} \propto$$

$$\propto \theta^{n(1)+n(2)+n(3)} \cdot (6-11\theta)^{n(4)}$$

Thus, there exists a constant k such that:

$$l(\theta) = k + (n - n(4))\log(\theta) + n(4)\log(6 - 11\theta)$$

$$\frac{dl(\theta)}{d\theta} = \frac{n - n(4)}{\theta} - \frac{11n(4)}{6 - 11\theta} = \frac{(n - n(4))(6 - 11\theta) - 11n(4)\theta}{\theta(6 - 11\theta)}$$

Hence,  $\frac{dl(\theta)}{d\theta} = 0$  if and only if:

$$6(n - n(4) = 11n\theta$$

$$\theta = \frac{6(n - n(4))}{11n}$$

Since L(0)=L(6/11)=0 and  $\Theta\cup\{0,6/11\}$  is a closed set, by Weierstrass's Theorem,  $\frac{6(n-n(4))}{11n}$  maximizes L.

(b) From the previous item:

$$\frac{dl(\theta)}{d\theta} = \frac{n - n(4)}{\theta} - \frac{11n(4)}{6 - 11\theta}$$

Thus,

$$\frac{d^2l(\theta)}{d\theta^2} = -\frac{n - n(4)}{\theta^2} - \frac{121n(4)}{(6 - 11\theta)^2} =$$

Finally,  $I(\theta) = -E(\frac{d^2l(\theta)}{d\theta^2}) =$ 

$$=\frac{n-\frac{n(6-11\theta)}{6}}{\theta^2}+\frac{121\frac{n(6-11\theta)}{6}}{(6-11\theta)^2}=$$

$$=\frac{11n}{6\theta}+\frac{121n}{6(6-11\theta)}=\frac{11n(6-11\theta)+121n\theta}{6\theta(6-11\theta)}=\frac{11n}{\theta(6-11\theta)}$$

(c) We know that  $\frac{\hat{\theta}-\theta}{\sqrt{I(\hat{\theta})}} \xrightarrow{\mathcal{L}} N(0,1)$ . Hence an asymptotic  $1-\alpha$  confidence interval for  $\theta$  is:

$$\left[\hat{\theta} - z_{1-\alpha/2}\sqrt{I(\hat{\theta})}, \hat{\theta} + z_{1-\alpha/2}\sqrt{I(\hat{\theta})}\right]$$

That is,

$$\left[\frac{6}{11}\left(1 - \frac{n(4)}{n}\right) - z_{1-\alpha/2} \frac{11\sqrt{n}}{6\sqrt{\frac{n(4)}{n}\left(1 - \frac{n(4)}{n}\right)}}, \frac{6}{11}\left(1 - \frac{n(4)}{n}\right) + z_{1-\alpha/2} \frac{11\sqrt{n}}{6\sqrt{\frac{n(4)}{n}\left(1 - \frac{n(4)}{n}\right)}}\right]$$

(d) From item (a),  $\hat{\theta} = \frac{6}{11} \left( 1 - \frac{n(4)}{n} \right)$ . By the LGN,  $\frac{n(4)}{n} \xrightarrow{P} \frac{6-11\theta}{6}$ . Thus, by the continuous mapping theorem:

$$\hat{\theta} \xrightarrow{P} \frac{6}{11} \left( 1 - \frac{6 - 11\theta}{6} \right) = \theta$$

## Problem 2

For any  $t \in \mathbb{R}$  define  $f_t(z) = \operatorname{sign}(\sin(tz))$ . Let  $\mathcal{F} = \{f_t : t \in \mathbb{R}\}$ . Show that  $\mathcal{F}$  has infinite VC dimension. Hint: consider a set of points like  $\{1/2, 1/4, \ldots, 1/2^n\}$ .

*Proof.* Define  $z_i = \frac{1}{2^i}$  and define  $\mathcal{Z}_n = \{z_i : i \in [1, 2, \dots n]\}$ . Note that if  $f_t(z_i) = 1$  it must be that  $\operatorname{sign}(\sin(tz_i)) = 1$ , which implies that:

$$\sin(tz_i) > 0$$
$$tz_i > \sin^{-1}(0)$$

This implies that

$$0 < tz_i < \pi \text{ or } 2\pi < tz_i < 3\pi \text{ or } 4\pi < tz_i < 5\pi \dots$$

Written succinctly, it must be that  $0 + a_i 2\pi < tz_i < \pi + a_i 2\pi, a_i \in [0, 1, 2, ..., \infty)$ .

Now, consider the case when  $f_t z_i = -1$ . In this case it must be that:

$$\sin(tz_i) < 0$$
$$tz_i < \sin^{-1}(0)$$

This implies that:

$$\pi < tz_i < 2\pi$$
 or  $3\pi < tz_i < 4\pi$  or  $5\pi < tz_i < 6\pi$ ...

Written succinctly, it must be that  $\pi + a_i 2\pi < tz_i < 2\pi + a_i 2\pi, a_i \in [0, 1, 2, \dots, \infty)$ .

Now, let  $y_i = f_t(z_i)$  and let  $\mathcal{Y}_n = \{y_i : i \in [1, \dots n]\}$ . For a given  $\mathcal{Y}_n$  we can write a set of n inequalities:

$$\begin{cases} 0 + a_i 2\pi < tz_i < \pi + 2a_i \pi, & y_i = 1\\ \pi + a_i 2\pi < tz_i < 2\pi + 2a_i \pi, & y_i = -1 \end{cases}$$

Substituting for  $z_i$ , this becomes:

$$\begin{cases} 0 + a_i 2\pi < t \frac{1}{2^i} < \pi + 2a_i \pi, & y_i = 1\\ \pi + a_i 2\pi < t \frac{1}{2^i} < 2\pi + 2a_i \pi, & y_i = -1 \end{cases}$$

Let  $t = k\pi$ , we can then write the set of equations as:

$$\begin{cases} 0 + a_i 2\pi < k\pi \frac{1}{2^i} < \pi + 2a_i \pi, & y_i = 1\\ \pi + a_i 2\pi < k\pi \frac{1}{2^i} < 2\pi + 2a_i \pi, & y_i = -1 \end{cases}$$

and this is equal to:

$$\begin{cases} 0 + a_i 2 < k \frac{1}{2^i} < 1 + 2a_i, & y_i = 1 \\ 1 + a_i 2 < k \frac{1}{2^i} < 2 + 2a_i, & y_i = -1 \end{cases}$$

We add 1 to the top equation to get:

$$\begin{cases} 1 + a_i 2 < 1 + \frac{k}{2^i} < 2 + 2a_i, & y_i = 1\\ 1 + a_i 2 < \frac{k}{2^i} < 2 + 2a_i, & y_i = -1 \end{cases}$$

Finally, for convenience, subtract  $2a_i$  from both equations to get:

$$\begin{cases} 1 < 1 + \frac{k}{2^i} - 2a_i < 2, & y_i = 1 \\ 1 < \frac{k}{2^i} - 2a_i < 2, & y_i = -1 \end{cases}$$

Now, let  $\mathcal{Y} = \{y_1, \dots, y_n\}$  be a set of any arbitrary labelings for the points in  $\mathcal{Z}_n$ . Note, that there are  $2^n$  such labelings. I will now prove that it is possible to find a solution to the set of equations defined directly above for all such labelings.

In general, if  $y_i = 1$ , we can say the following:

$$1 < 1 + \frac{k}{2^{i}} - 2a_{1} < 2$$

$$2^{i} < 2^{i} + k - 2^{i+1}a_{1} < 2^{i+1}$$

$$0 < k - a_{1}2^{i+1} < 2^{i+1} - 2^{i}$$

And if  $y_i = -2$ , we can say:

$$1 < \frac{k}{2^{i}} - 2a_{1} < 2$$
$$2^{i} < k - 2^{i+1}a_{1} < 2^{i+1}$$

This is a very important pattern. For a fixed i, we have an alternating set of ranges for which k can be in if  $y_i = 1$ , ie:  $(0, 2^{i+1} - 2^i), (1 \times 2^{i+1}, 1 \times 2^{i+1} + 2^{i+1} - 2^i), (2 \times 2^{i+1}, 2 \times 2^{i+1} + 2^{i+1} - 2^i), \ldots$  We also have an alternating set of ranges for which k can be in if  $y_i = 1$ . Notice that these ranges are situated between those for  $y_i = 1$ , ie:  $(2^i, 2^{i+1}), (1 \times 2^{i+1} + 2^i, 1 \times 2^{i+1} + 2^{i+1}), (2 \times 2^{i+1} + 2^i, 2 \times 2^{i+1} + 2^{i+1}), \ldots$ 

Further, notice what happens as we move from ranges for i and ranges for i+1. The ranges for i+1 are twice as big as the ranges for i. Additionally, the ranges for k when  $y_{i+1} = 1$  will cover the ranges for k when  $y_i = 1$  and  $y_i = -1$ . This also holds for the ranges when  $y_{i+1} = -1$ . Thus, no matter what the label for  $y_i$ , it will always be possible to find a value of k that correctly predicts the label for  $y_i$  and  $y_{i+1}$ . This relationship holds for all i, so it must be that for any arbitrary set  $\mathcal{Y}_n$ , we can find a  $k = t\pi$ , such that  $f_t$  correctly maps all  $z \in \mathcal{Z}_n$  to  $y \in \mathcal{Y}_n$ . This is equivalent to saying  $s(\mathcal{F}, n) = 2^n$  for all n, which implies the VC dimension for  $\mathcal{F}$  is infinity.

### Problem 3

(a) We wish to find an  $\epsilon$  such that  $P(|\bar{X} - \mu| > \epsilon) \leq \delta$ . By Bernstein's Inequality, it is sufficient to take an  $\epsilon$  such that:

$$2exp\left(-\frac{n\epsilon^2}{2\sigma^2 + 2c\epsilon/3}\right) \le \delta$$

$$\frac{n\epsilon^2}{2\sigma^2 + 2c\epsilon/3} \ge \log(2/\delta)$$

$$\epsilon^2 - \frac{2c\log(2/\delta)\epsilon}{3n} - \frac{2\sigma^2\log(2/\delta)}{n} \ge 0$$

Hence, it is sufficient to take  $\epsilon$  such that:

$$\epsilon \ge \frac{c\log(2/\delta)}{3n} + \sqrt{\frac{\left(\frac{2c\log(2/\delta)}{3n}\right)^2 + \frac{8\sigma^2\log(2/\delta)}{n}}{4}}$$

Remembering that  $\sqrt{x+y} \leq \sqrt{x} + \sqrt{y}$ , it is sufficient to take  $\epsilon$  such that:

$$\epsilon \ge \frac{2c\log(2/\delta)}{3n} + \sqrt{\frac{2\sigma^2\log(2/\delta)}{n}}$$

(b) We wish to show that:

$$\forall \epsilon > 0, \exists C > 0 : P(n|\bar{X}_n - \mu_n| > C) \le \epsilon$$

First, observe that since  $X_i$  are dicotomic,  $|X_i| \leq 1$ . Next, if we call  $P(Y_i \in A_n) = p_n$ ,  $\sigma^2 \leq p_n(1-p_n) \leq p_n$ . Also call  $M = \sup f$  and observe that,  $p_n \leq \frac{M}{n}$  and, thus,  $\sigma^2 \leq \frac{M}{n}$ . Hence, by Bernstein's Inequality:

$$P(|\bar{X}_n - \mu_n| > \frac{C}{n}) \le 2exp\left(-\frac{\frac{C^2}{n}}{\frac{2M}{n} + \frac{2C}{3n}}\right)$$

$$P(|\bar{X}_n - \mu_n| > \frac{C}{n}) \le 2exp\left(-\frac{C}{2M/C + 2/3}\right)$$

The result follows observing that  $2exp\left(-\frac{C}{2M/C+2/3}\right)\xrightarrow{C\mapsto\infty}0$ .

### Problem 4

(Rademacher Complexity). Let  $\mathcal{F} = \{f_1, \dots, f_N\}$  where each f is a binary function:  $f(x) \in \{0, 1\}$ . Show that

$$\mathcal{R}_n(\mathcal{F}) \le 2\sqrt{\frac{\log N}{n}}$$

*Proof.* From theorem 42.33 in the notes, we know that:

$$\mathcal{R}_n(\mathcal{F}) \le \sqrt{\frac{2\log(s(\mathcal{F},n))}{n}}$$

For a finite set of functions  $\mathcal{F} = \{f_1, \dots, f_N\}$  it must be that  $s(\mathcal{F}, n) \leq N$  as it is impossible to produce more unique labelings than there are functions in the set.

Thus, it immediately follows that:

$$\mathcal{R}_n(\mathcal{F}) \le \sqrt{\frac{2\log(N)}{n}} \le 2\sqrt{\frac{\log N}{n}}$$