

15-410

“...The cow and Zaphod...”

Virtual Memory #3

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Synchronization

First Project 3 checkpoint

- Monday during class time
- Meet in Wean 5207
 - If your group number *ends* with
 - » 0-2 try to arrive 5 minutes early
 - » 3-5 arrive at 10:42:30
 - » 6-9 arrive at 10:59:27
- Preparation
 - Your kernel should be in mygroup/p3ck1
 - It should load one program, enter user space, `gettid()`
 - » Ideally `lprintf()` the result of `gettid()`
 - We will ask you to load & run a test program we will name
 - Explain which parts are “real”, which are “demo quality”

Outline

Last time

- The mysterious TLB
- Partial memory residence (demand paging) in action
- The task of the page fault handler

Today

- Fun big speed hacks
- Sharing memory regions & files
- Page replacement policies

Demand Paging Performance

Effective access time of memory word

- $(1 - p_{\text{miss}}) * T_{\text{memory}} + p_{\text{miss}} * T_{\text{disk}}$

Textbook example (a little dated)

- $T_{\text{memory}} = 100 \text{ ns } (.1 \text{ us})$
- $T_{\text{disk}} = 25 \text{ ms } (25,000 \text{ us})$
- $p_{\text{miss}} = 1/1,000$ slows down by **factor of 250**
- slowdown of 10% needs $p_{\text{miss}} < 1/2,500,000!!!$

Speed Hacks

COW

ZFOD (Zaphod?)

Memory-mapped files

- What `msync()` is *supposed* to be used for...

Copy-on-Write

`fork()` produces two **very**-similar processes

- Same code, data, stack

Expensive to copy pages

- Many will never be modified by new process
 - Especially in `fork()`, `exec()` case

Share physical frames instead of copying?

- Easy: code pages – read-only
- Dangerous: stack pages!

Copy-on-Write

Simulated copy

- Copy page table entries to new process
- Mark PTEs read-only in old & new
- Done! (saving factor: 1024)
 - Simulation is excellent as long as process doesn't write...

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Making it real

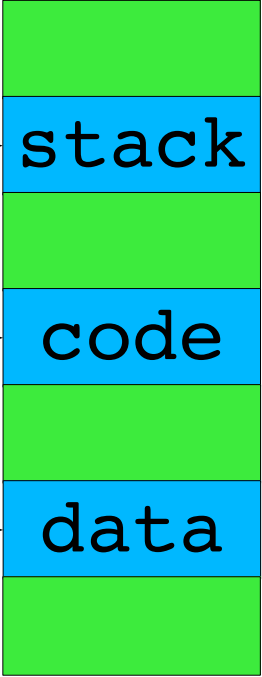
- Process writes to page (*Oops! We lied...*)
- Page fault handler responsible
 - Kernel makes a copy of the shared frame
 - Page tables adjusted
 - » ...each process points page to private frame
 - » ...page marked read-write in both PTEs

Example Page Table

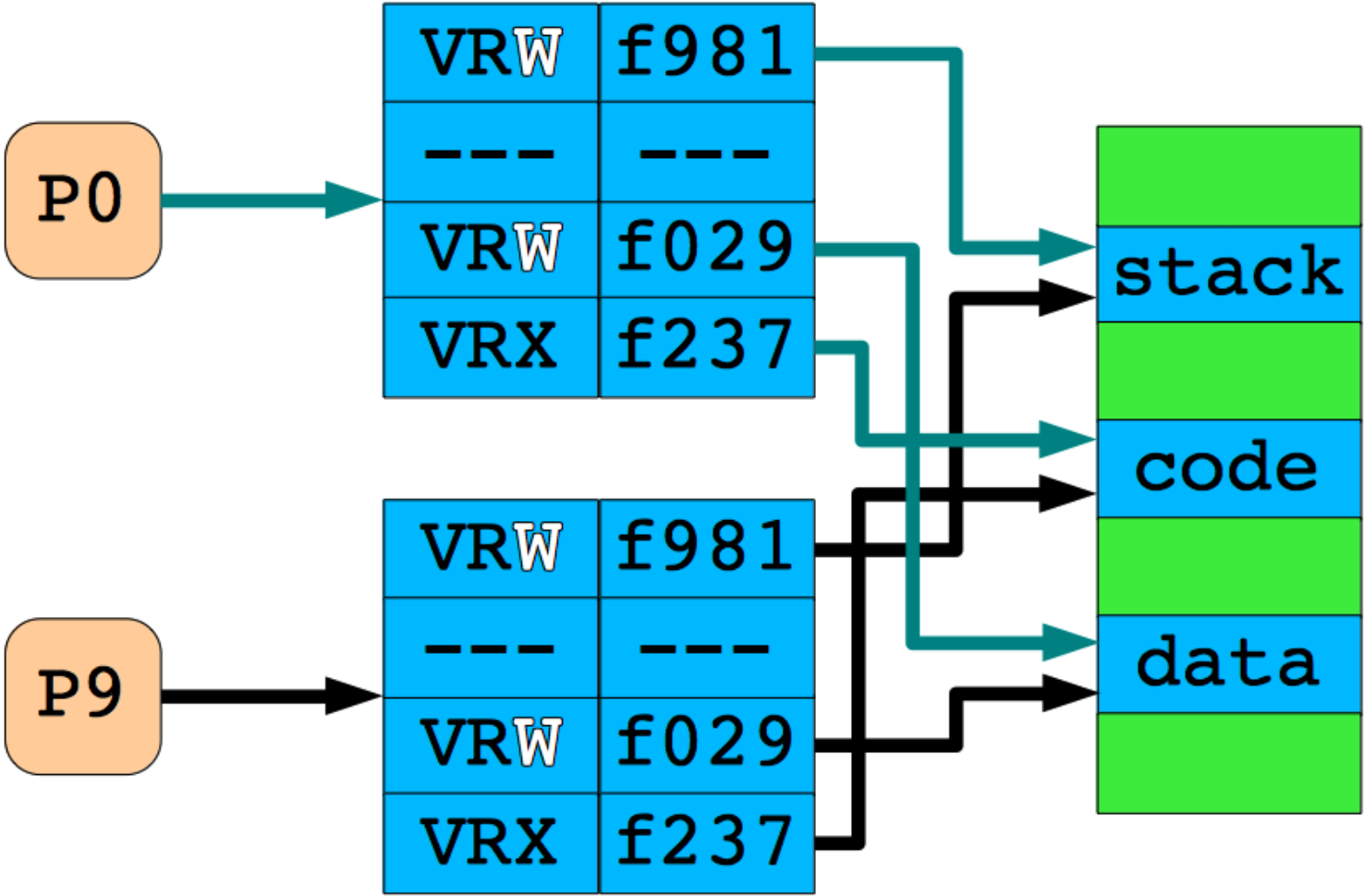
Virtual Address

VRW	f981
---	---
VRW	f029
VRX	f237

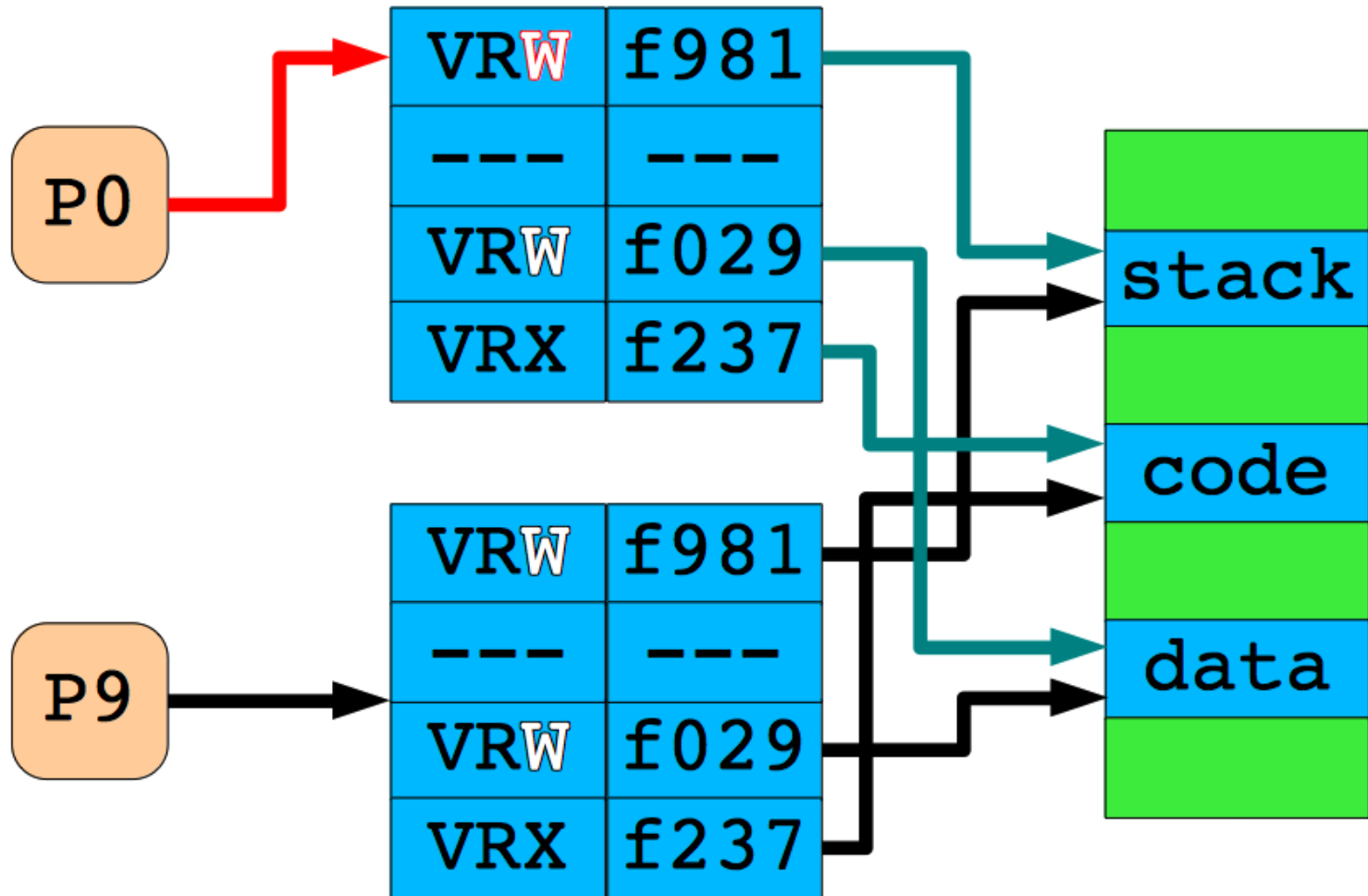
Page table



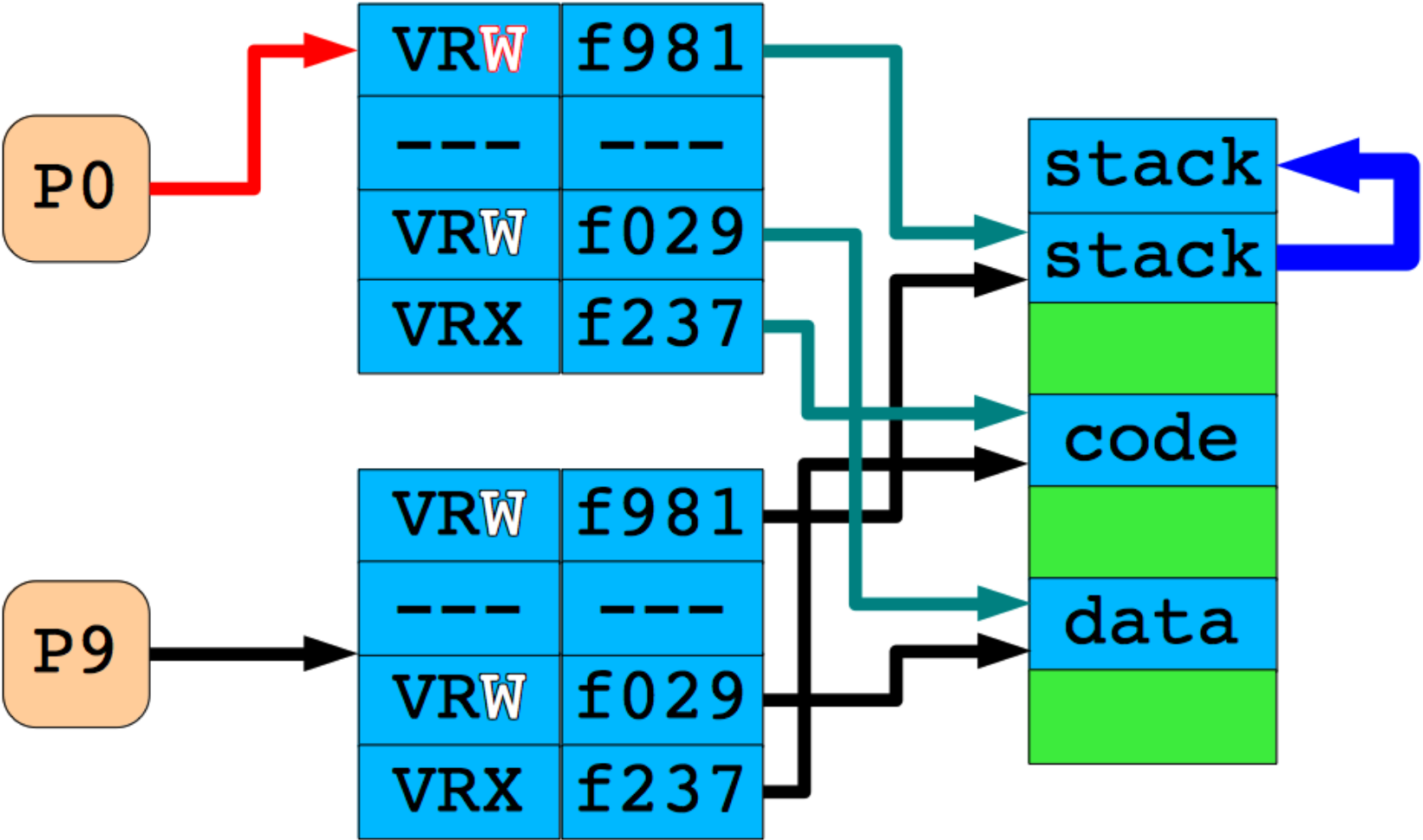
Copy-on-Write of Address Space



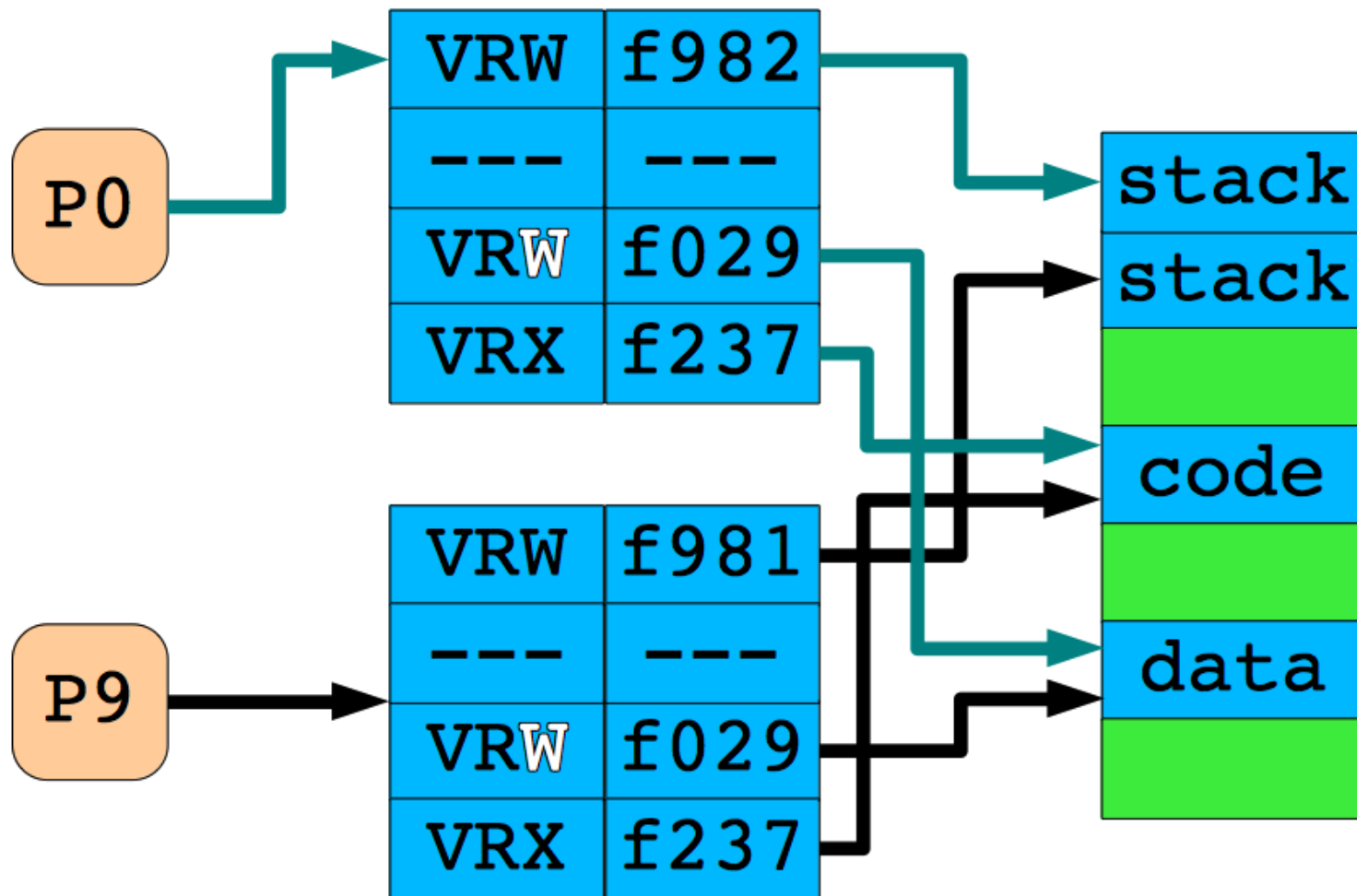
Memory Write \Rightarrow Permission Fault



Copy Into Blank Frame



Adjust PTE frame pointer, access



Zero Pages

Very special case of copy-on-write

- ZFOD = “Zero-fill on demand”

Many process pages are “blank”

- All of bss
- New heap pages
- New stack pages

Have one *system-wide* all-zero frame

- Everybody points to it
- Logically read-write, physically read-only
- Reads of zeros are free
- Writes cause page faults & cloning

Memory-Mapped Files

Alternative interface to `read()`, `write()`

- `mmap(addr, len, prot, flags, fd, offset)`
- new memory region presents file contents
- write-back policy typically unspecified
 - unless you `msync()` ...

Benefits

- Avoid serializing pointer-based data structures
- Reads and writes may be much cheaper
 - Look, Ma, no syscalls!

Memory-Mapped Files

Implementation

- Memory region remembers `mmap()` parameters
- Page faults trigger `read()` calls
- Pages stored back via `write()` to file

Shared memory

- Two processes `mmap()` “the same way”
- Point to same memory region

Page Replacement/Page Eviction

Process always want *more* memory frames

- Explicit deallocation is rare
- Page faults are implicit allocations

System inevitably runs out of frames

Solution outline

- Pick a frame, store contents to disk
- Transfer ownership to new process
- Service fault using this frame

Pick a Frame

Two-level approach

- Determine # frames each process “deserves”
- “Process” chooses which frame is least-valuable
 - Most OS's: kernel actually does the choosing

System-wide approach

- Determine globally-least-useful frame

Store Contents to Disk

Where does it belong?

- Allocate backing store for each page
 - What if we run out?

Must we *really* store it?

- Read-only code/data: no!
 - Can re-fetch from executable
 - Saves paging space & disk-write delay
 - But file-system read() may be slower than paging-disk read
- Not modified since last page-in: no!
 - Hardware typically provides “page-dirty” bit in PTE
 - Cheap to “store” a page with dirty==0

Page Eviction Policies

Don't try these at home

- FIFO
- Optimal
- LRU

Practical

- LRU approximation

Current Research

- ARC (Adaptive Replacement Cache)
- CAR (Clock with Adaptive Replacement)
- CART (CAR with Temporal Filtering)

Page Eviction Policies

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Practical

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- CARTHAGE (CART with Hilarious AppendaGE)

FIFO Page Replacement

Concept

- Queue of all pages – named as (task id, virtual address)
- Page added to tail of queue when first given a frame
- Always evict oldest page (head of queue)

Evaluation

- Fast to “pick a page”
- Stupid
 - Will indeed evict old unused startup-code page
 - But *guaranteed* to eventually evict process's favorite page too!

Optimal Page Replacement

Concept

- Evict whichever page will be referenced *latest*
 - “Buy the most time” until next page fault

Evaluation

- Requires perfect prediction of program execution
- Impossible to implement

So?

- Used as upper bound in simulation studies

LRU Page Replacement

Concept

- Evict Least-Recently-Used page
- “Past performance *may* not predict future results”
 - ...but it's an important hint!

Evaluation

- Would probably be reasonably accurate
- LRU is computable without a fortune teller
- Bookkeeping *very* expensive
 - (right?)

LRU Page Replacement

Concept

- Evict Least-Recently-Used page
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Evaluation

- Would probably be reasonably accurate
- LRU is computable without a fortune teller
- Bookkeeping *very* expensive
 - Hardware must sequence-number every page reference
 - » Evictor must scan every page's sequence number
 - Or you can “just” do a doubly-linked-list operation per ref

Approximating LRU

Hybrid hardware/software approach

- 1 reference bit per page table entry
- OS sets reference = 0 for all pages
- Hardware sets reference=1 when PTE is used in lookup
- OS periodically scans
 - (reference == 1) ⇒ “recently used”
- Result:
 - Hardware sloppily partitions memory into “recent” vs. “old”
 - Software periodically samples, makes decisions

Approximating LRU

“Second-chance” algorithm

- Use stupid FIFO queue to choose victim candidate page
- reference == 0?
 - not “recently” used, evict page, steal its frame
- reference == 1?
 - “somewhat-recently used” - don't evict page this time
 - append page to rear of queue (“second chance”)
 - set reference = 0
 - » Process must use page again “soon” for it to be skipped

Approximation

- Observe that queue is randomly sorted
 - We are evicting not-recently-used, not *least*-recently-used

Approximating LRU

“Clock” algorithm

- **Observe: “Page queue” requires linked list**
 - Extra memory traffic to update pointers
- **Observe: Page queue's order is essentially random**
 - Doesn't add anything to accuracy
- **Revision**
 - Don't have a queue of pages
 - Just treat memory as a circular array

Clock Algorithm

```
static int nextpage = 0;
boolean reference[NPAGES];

int choose_victim() {
    while (reference[nextpage]) {
        reference[nextpage] = false;
        nextpage = (nextpage+1) % NPAGES;
    }
    return (nextpage) ;
}
```

“Page Buffering”

Problem

- Don't want to evict pages only *after* a fault needs a frame
- Must wait for disk write before launching disk read (slow!)

“Assume a blank page...”

- Page fault handler can be much faster

“page-out daemon”

- Scans system for dirty pages
 - Write to disk
 - Clear dirty bit
 - Page can be instantly evicted later
- When to scan, how many to store? Indeed...

Frame Allocation

How many frames should a process have?

Minimum allocation

- Examine worst-case instruction
 - Can multi-byte instruction cross page boundary?
 - Can memory parameter cross page boundary?
 - How many memory parameters?
 - Indirect pointers?

“Fair” Frame Allocation

Equal allocation

- Every process gets same *number of frames*
 - “Fair” - in a sense
 - Probably wasteful

Proportional allocation

- Every process gets same *percentage of residence*
 - (Everybody 83% resident, larger processes get more frames)
 - “Fair” - in a different sense
 - Probably the right approach
 - » Theoretically, encourages greediness

Thrashing

Problem

- Process *needs* N frames...
 - Repeatedly rendering image to video memory
 - Must be able to have all “world data” resident 20x/second
- ...but OS provides N-1, N/2, etc.

Result

- Every page OS evicts generates “immediate” fault
- More time spent paging than executing
- Paging disk constantly busy
 - Denial of “paging service” to other processes
- Widespread unhappiness

“Working-Set” Allocation Model

Approach

- Determine necessary # frames for each process
 - “Working set” - size of frame set you need to get work done
- If unavailable, swap entire process out
 - (later, swap some *other* process entirely out)

How to measure working set?

- Periodically scan all reference bits of process's pages
- Combine multiple scans (see text)

Evaluation

- Expensive
- Can we approximate it?

Page-Fault Frequency Approach

Approach

- Recall, “thrashing” == “excessive” paging
- Adjust per-process frame quotas to balance fault rates
 - System-wide “average page-fault rate” (10 faults/second)
 - Process A fault rate “too high”: increase frame quota
 - Process A fault rate “too low”: reduce frame quota

What if quota increase doesn't help?

- If giving you *some* more frames didn't help, maybe you need *a lot* more frames than you have...
 - Swap you out entirely for a while

Program Optimizations

Is paging an “OS problem”?

- Can a programmer reduce working-set size?

Locality depends on data structures

- Arrays encourage sequential accesses
 - Many references to same page
 - Predictable access to next page
- Random pointer data structures scatter references

Compiler & linker can help too

- Don't split a routine across two pages
- Place helper functions on same page as main routine

Effects can be *dramatic*

Summary

Speed hacks

Page-replacement policies

- The eviction problem
- Sample policies
 - For real: LRU approximation with hardware support
- Page buffering
- Frame Allocation (process page quotas)

Definition & use of

- Dirty bit, reference bit

Virtual-memory usage optimizations

Synchronization

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- Meet in GHC 3000
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 - » Ideally `lprintf()` the result of `gettid()`
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Synchronization

Second Project 3 checkpoint

- Wednesday during class time
- Attendance is *mandatory* except by prior arrangement
- Meet in Wean 5207
 - If your group number *ends* with
 - » 0-2 arrive at 10:42:30
 - » 3-5 arrive at 10:59:27
 - » 6-9 try to arrive 5 minutes early
- Preparation
 - Your kernel should be in mygroup/p3ck2
 - Depending on whether you are demo'ing fork() or exec() we will ask you to run different test programs
 - Either way we hope to observe context switching

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Synchronization

Project 2 due tonight

- Check you can write your mygroup/p2 directory *early*
- Please put your files in mygroup/p2
 - Not p2/p2, p2/our_project_2, p2/p2.tar
- Please don't mail us files
- Don't forget about the late-day form if you need it
 - Remember to balance against P3

Upcoming

- HW1 out soon, due sometime Wednesday
- Exam – Thursday
- Project 3 (including one checkpoint before spring break)

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Synchronization

Upcoming

- HW1 out today, due Tuesday evening
- P2 – due Friday evening
- Exam review - Monday
- Exam – Wednesday evening (watch for e-mail!)
- Project 3 (including one checkpoint before spring break)