15-410

"Strangers in the night..."

Synchronization #2
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Synchronization

Project 1 due tonight

- (but you knew that)

Synchronization

Register your project partner

- "Partner registration" page on Projects page
- If you know your partner today, please register today
  - You'll get your shared AFS space sooner
  - Your classmates will appreciate it

Outline

Last time

- Two building blocks
- Three requirements for mutual exclusion
- Algorithms people *don't* use for mutual exclusion

Today

- Ways to *really* do mutual exclusion

Next time

- Inside voluntary descheduling
- Project 2 – thread library

Mutual Exclusion: Reminder

Protects an atomic instruction sequence

- Do "something" to guard against
  - CPU switching to another thread
  - Thread running on another CPU

Assumptions

- Atomic instruction sequence will be “short”
- No other thread “likely” to compete

Mutual Exclusion: Goals

Typical case (no competitor) should be fast

Atypical case can be slow

- Should not be "too wasteful"

Interfering Code Sequences

Customer	*Delivery*
`cash = store->cash;`	`cash = store->cash;`
`cash += 50;`	`cash -= 2000;`
`wallet -= 50;`	`wallet += 2000;`
`store->cash = cash;`	`store->cash = cash;`

Which sequences interfere?

- "Easy": Customer interferes with Customer
- Also: Delivery interferes with Customer

Mutex aka Lock aka Latch

Specify interfering code sequences via *object*

- Data item(s) "protected by the mutex"

Object methods encapsulate entry & exit protocols

```
mutex_lock(&store->lock);
cash = store->cash
cash += 50;
personal_cash -= 50;
store->cash = cash;
mutex_unlock(&store->lock);
```

What's inside the object?

Mutual Exclusion: Atomic Exchange

Intel x86 XCHG instruction

- intel-isr.pdf page 754

xchg (%esi), %edi

```
int32 xchg(int32 *lock, int32 val) {
  register int old;
  old = *lock; /* bus is locked */
  *lock = val; /* bus is locked */
  return (old);
}
```

Inside a Mutex

Initialization

```
int lock_available = 1;
```

Try-lock

```
i_won = xchg(&lock_available, 0);
```

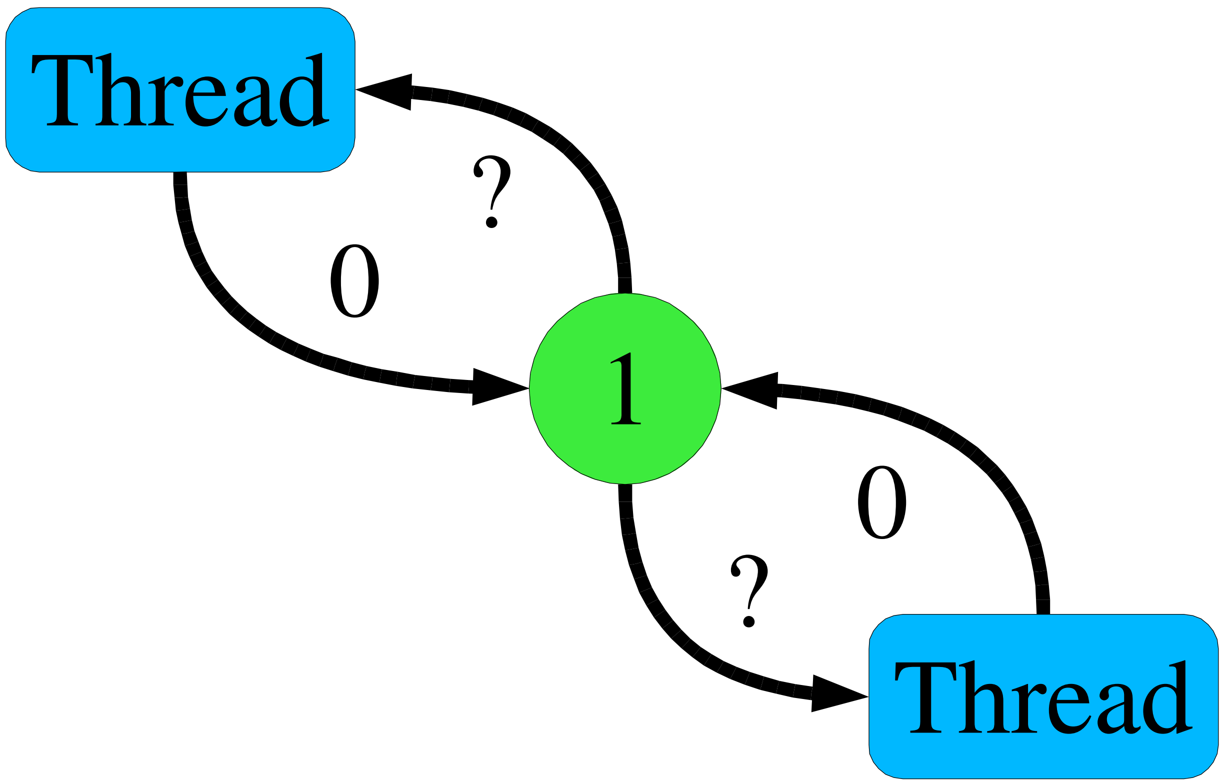
Spin-wait

```
while (!xchg(&lock_available, 0)
  /* nothing */ ;
```

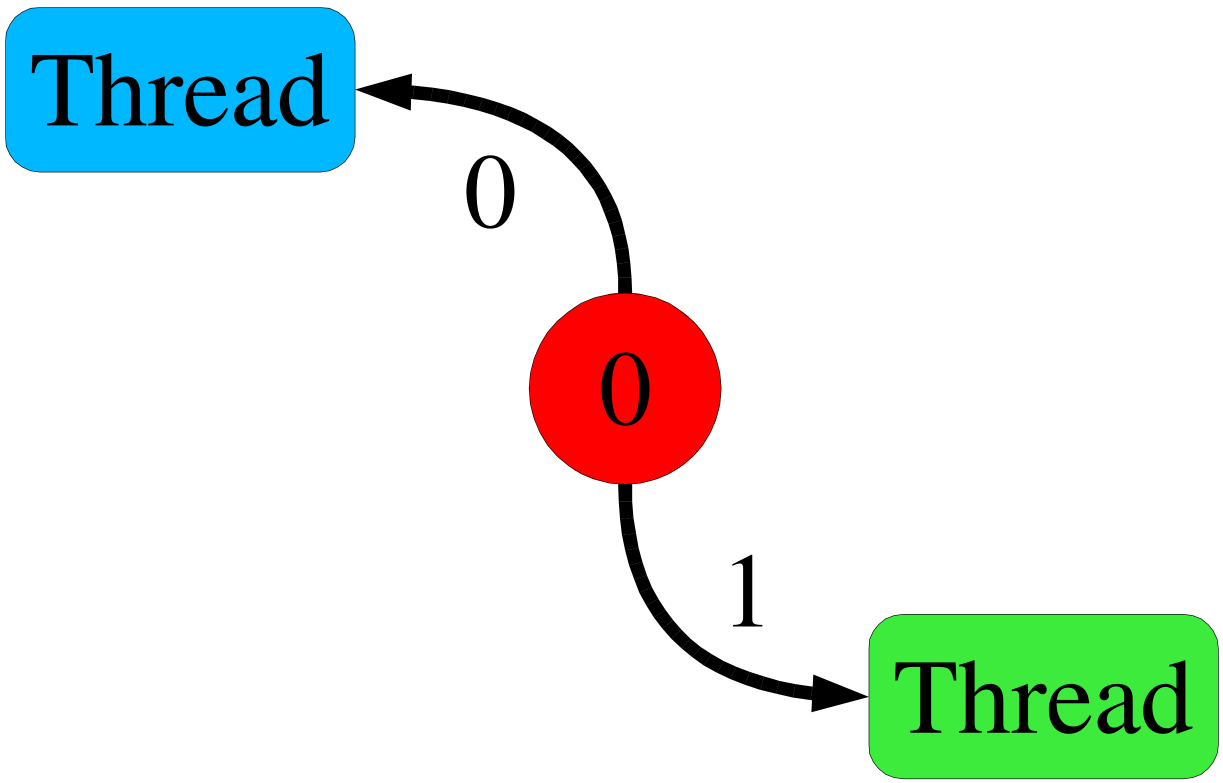
Unlock

```
xchg(&lock_available, 1); /*expect 0*/
```

Strangers in the Night, Exchanging 0's



And the winner is...



Does it work?

[What are the questions, again?]

Does it work?

Mutual Exclusion

Progress

Bounded Waiting

Does it work?

Mutual Exclusion

- There's only one 1; 1's are conserved
- Only one thread can see lock_available == 1

Progress

- Whenever lock_available == 1 a thread will get it

Bounded Waiting

- *No*
- A thread can lose *arbitrarily many times*

Ensuring Bounded Waiting

Lock

```
waiting[i] = true; /*Declare interest*/
got_it = false;
while (waiting[i] && !got_it)
  got_it = xchg(&lock_available,
                false);
waiting[i] = false;
```

Ensuring Bounded Waiting

Unlock

```
j = (i + 1) % n;
while ((j != i) && !waiting[j])
  j = (j + 1) % n;
if (j == i)
  xchg(&lock_available, true); /*W*/
else
  waiting[j] = false;
```

Ensuring Bounded Waiting

Versus (previous edition of) textbook

- Exchange vs. TestAndSet
- "Available" vs. "locked"
- Atomic release vs. normal memory write
  - Text does "blind write" at point "W"

```
lock_available = true;
```

  - This may be illegal on some machines
  - Unlocker may be *required* to use special memory access
    - Exchange, TestAndSet, etc.

Evaluation

One awkward requirement

One unfortunate behavior

Evaluation

One awkward requirement

- Everybody knows size of thread population
  - Always & instantly!
  - Or uses an upper bound

One unfortunate behavior

- Recall: expect *zero* competitors
- Algorithm: O(n) in *maximum possible* competitors

Is this criticism too harsh?

- After all, Baker's Algorithm has these misfeatures...

Looking Deeper

Look beyond abstract semantics

- Mutual exclusion, progress, bounded waiting

Consider

- *Typical* access pattern
- *Particular* runtime environments

Environment

- Uniprocessor vs. Multiprocessor
  - Who is doing what when we are trying to lock/unlock?
- Threads aren't mysteriously “running” or “not running”
  - Decision made by scheduling algorithm with properties

Uniprocessor Environment

Lock

- What if xchg() didn't work the first time?
- Some other process has the lock
  - That process isn't running (because we are)
  - *xchg() loop is a waste of time*
  - We should let the lock-holder run instead of us

Unlock

- What about bounded waiting?
- When we mark mutex available, who wins next?
  - Whoever *runs* next..only one at a time! ("Fake competition")
  - How unfair are real OS kernel thread schedulers?
  - If scheduler is vastly unfair, the right thread will *never* run!

Multiprocessor Environment

Lock

- Spin-waiting probably justified
  - (why?)

Unlock

- Next xchg() winner "chosen" by memory hardware
- How unfair are real memory controllers?

Test&Set

```
boolean testandset(int32 *lock) {
register boolean old;
  old = *lock;  /* bus is locked */
  *lock = true; /* bus is locked */
  return (old);
}
```

Conceptually simpler than XCHG?

- Or not

Load-linked, Store-conditional

For multiprocessors

- "Bus locking considered harmful"

Split XCHG into halves

- *Load-linked(addr)* fetches old value from memory
- *Store-conditional(addr,val)* stores new value back
  - If nobody else stored to that address in between
  - If so, instruction "fails" (sets an error code)

Load-linked, Store-conditional

```
loop:  LL   $3, mutex_addr
       BEQ  $3, $0, loop     # avail == 0
       LI   $3, 0            # prep. 0
       SC   $3, mutex_addr   # write 0?
       BEQ  $3, $0, loop     # aborted...
```

Your cache “snoops” the shared memory bus

- Locking would shut down *all* memory traffic
- Snooping allows all traffic, watches for *conflicting* traffic
- Are aborts “ok”? *When* are they “ok”?

Intel i860 magic lock bit

Instruction sets processor in "lock mode"

- Locks bus
- Disables interrupts

Isn't that dangerous?

- 32-instruction countdown timer triggers exception
- Any exceptions (page fault, zero divide, ...) unlock bus

Why would you want this?

- Implement test&set, compare&swap, semaphore – you choose

Mutual Exclusion: Inscrutable Software

Lamport's “Fast Mutual Exclusion” algorithm

- 5 writes, 2 reads (if no contention)
- Not bounded-waiting (in theory, i.e., if contention)
- http://www.hpl.hp.com/techreports/Compaq-DEC/SRC-RR-7.html

Why not use it?

- What *kind* of memory writes/reads?
- Remember, the computer is “modern”...

Passing the Buck?

Q: Why not ask the OS for mutex_lock() *system call*?

Easy on a uniprocessor...

- Kernel *automatically* excludes other threads
- Kernel can easily disable interrupts
- No need for messy unbounded loop, weird XCHG...

Kernel has special power on a multiprocessor

- Can issue “remote interrupt” to other CPUs
- No need for messy unbounded loop...

So why *not* rely on OS?

Passing the Buck

A: Too expensive

- Because... (you know this song!)

Mutual Exclusion: *Tricky* Software

Fast Mutual Exclusion for Uniprocessors

- **Bershad, Redell, Ellis: ASPLOS V (1992)**

Want uninterruptable instruction sequences?

- **Pretend!**

```
scash = store->cash;
scash += 10;
wallet -= 10;
store->cash = scash;
```

- **Uniprocessor: interleaving requires thread switch...**
- **Short sequence *almost always* won't be interrupted...**

How can that work?

Kernel *detects* "context switch in atomic sequence"

- Maybe a small set of instructions
- Maybe particular memory areas
- Maybe a flag

```
no_interruption_please = 1;
```

Kernel *handles* unusual case

- Hand out another time slice? (Is that ok?)
- Hand-simulate unfinished instructions (yuck?)
- "Idempotent sequence": slide PC back to start

Summary

Atomic instruction sequence

- Nobody else may interleave same/"related" sequence

Specify interfering sequences via *mutex object*

Inside a mutex

- Last time: race-condition memory algorithms
- Atomic-exchange, Compare&Swap, Test&Set, ...
- Load-linked/Store-conditional
- Tricky software, weird software

Mutex strategy

- How should you behave given runtime environment?