Lamport clocks

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Synchronization

- Project 4 due today
- Homework 2 due Friday
- Book report due Friday
- FCE reminder
  - I will read (and take seriously) every word of what you write

Outline

- Lamport clocks
  - Covered in 17.1, 17.2 (different focus from today)
  - Time, Clocks, and the Ordering of Events in a Distributed System
    - CACM 21:7 (1978)
- Leslie Lamport also famous for ...?

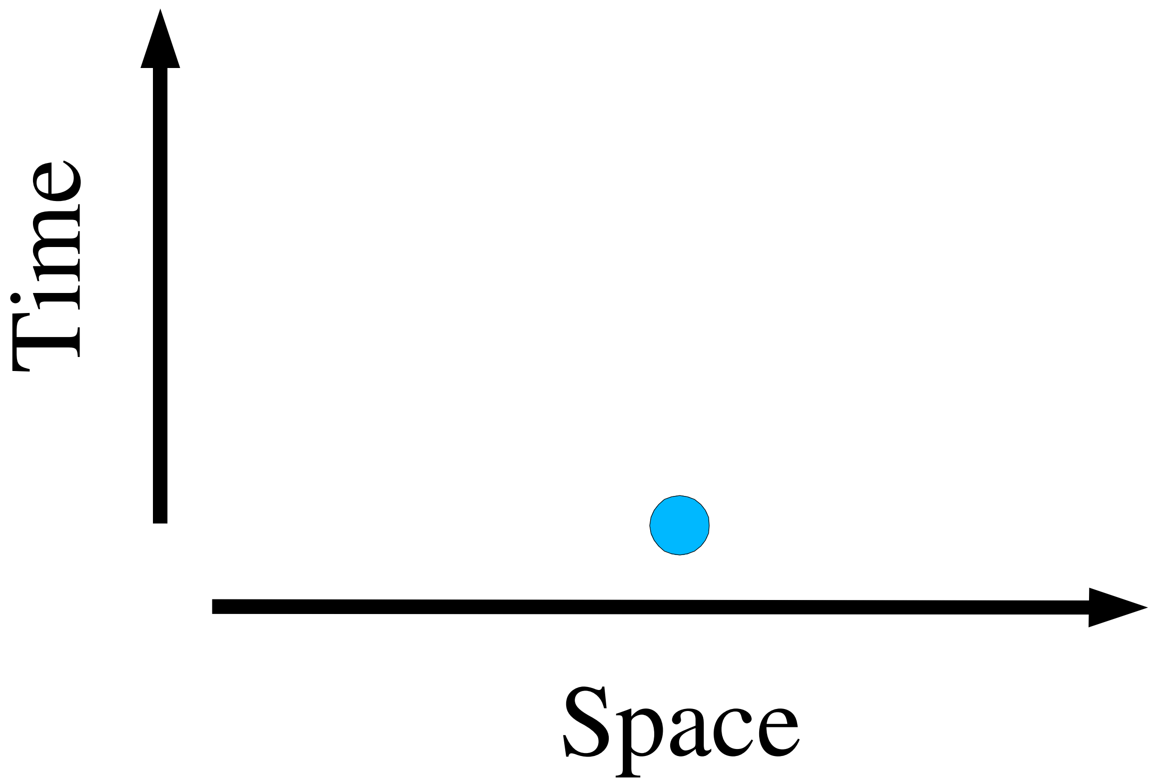
Overview

- Light cones
- Meeting for beer
- “Happened before” partial order
- Logical clocks
- Advanced techniques

Light cones

- Concept
  - Effects propagate at or below speed of light
    - Objects, light/radio/X-rays, gravity
  - ***Knowledge*** of events limited the same way
  - Event propagation modeled by expanding sphere
    - Four-dimensional “cone”

Light cones



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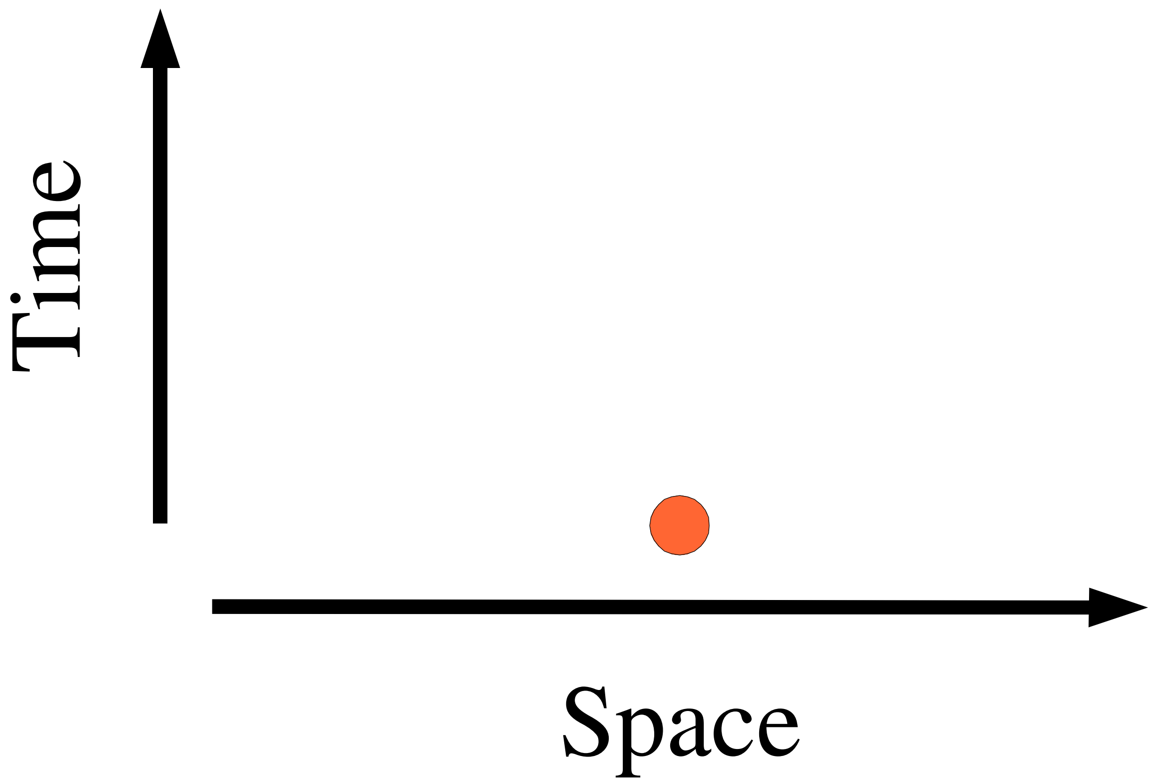

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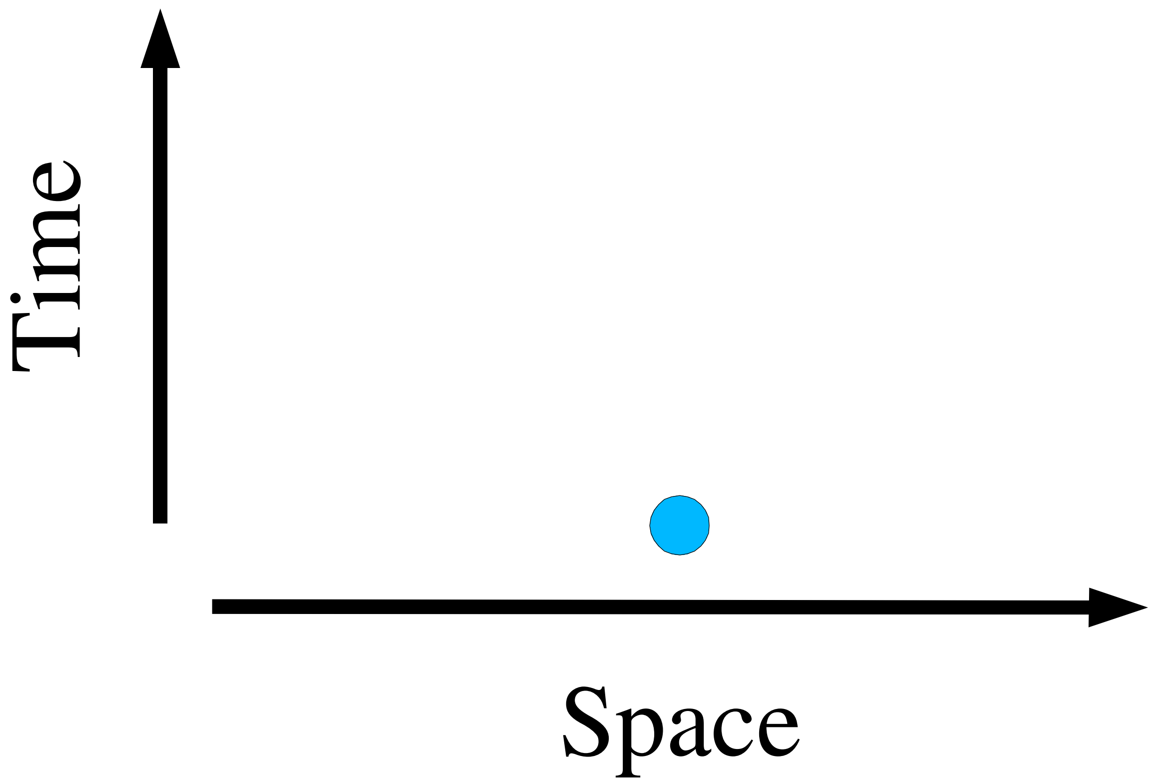

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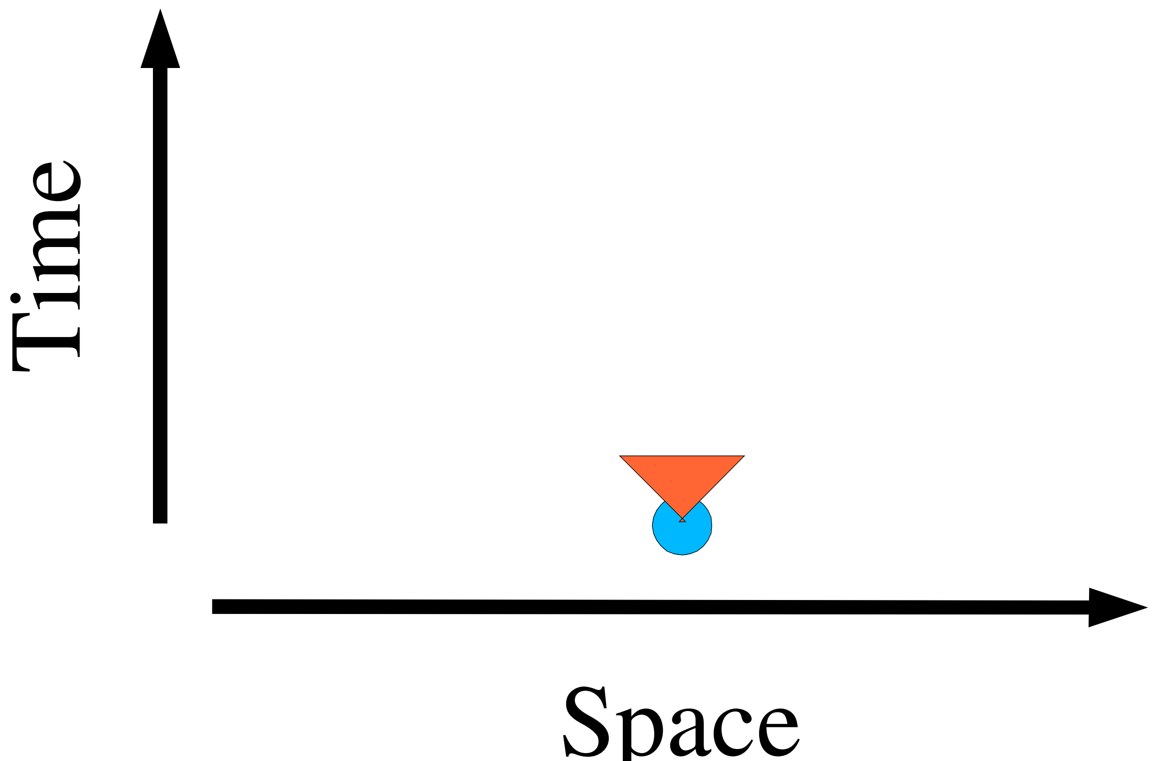

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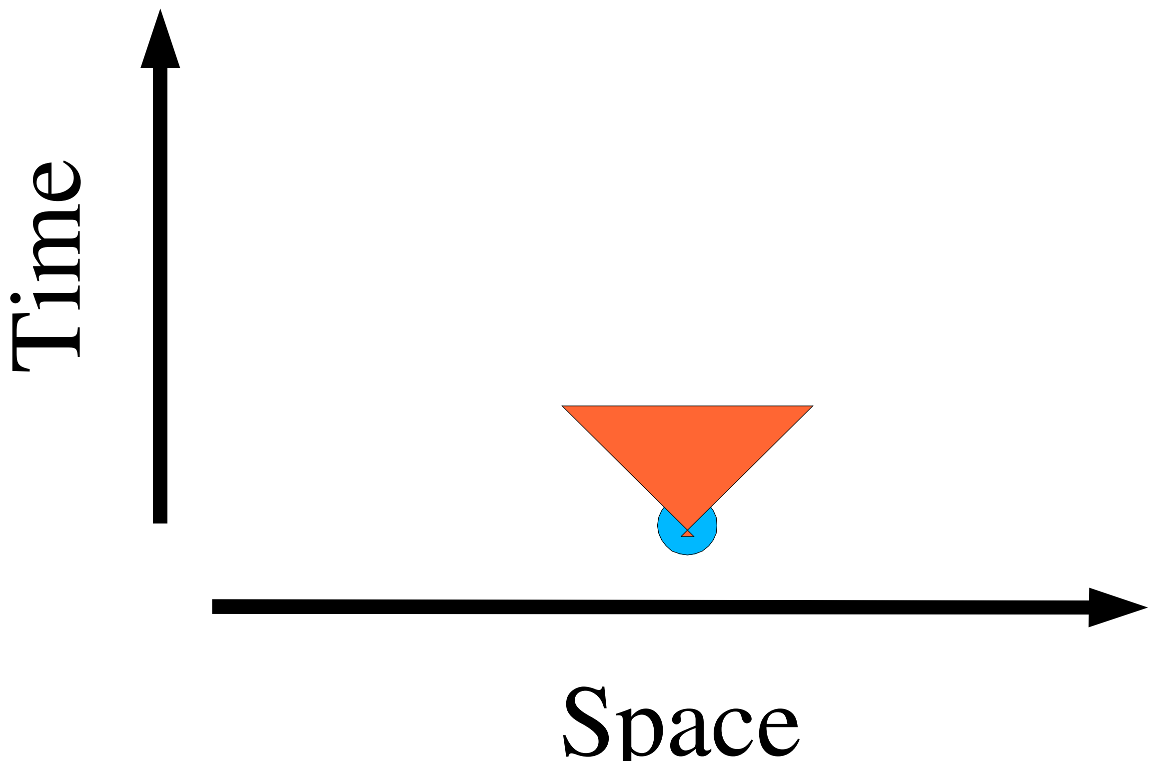

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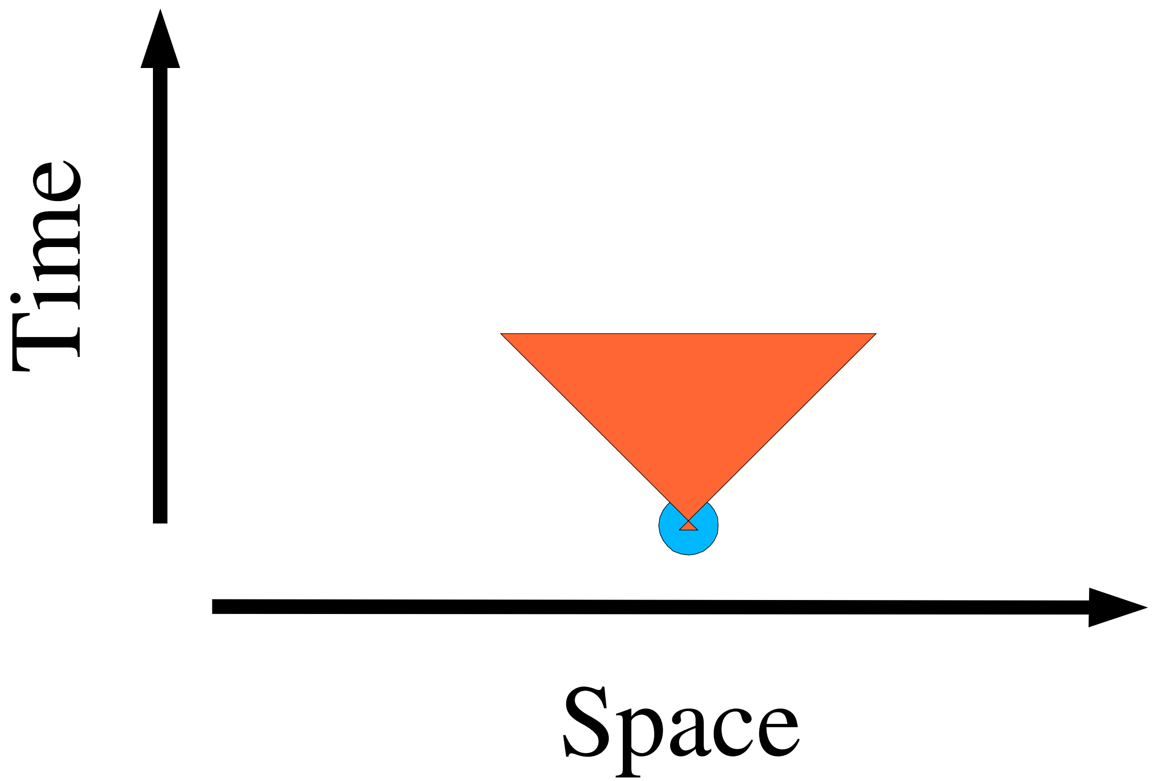

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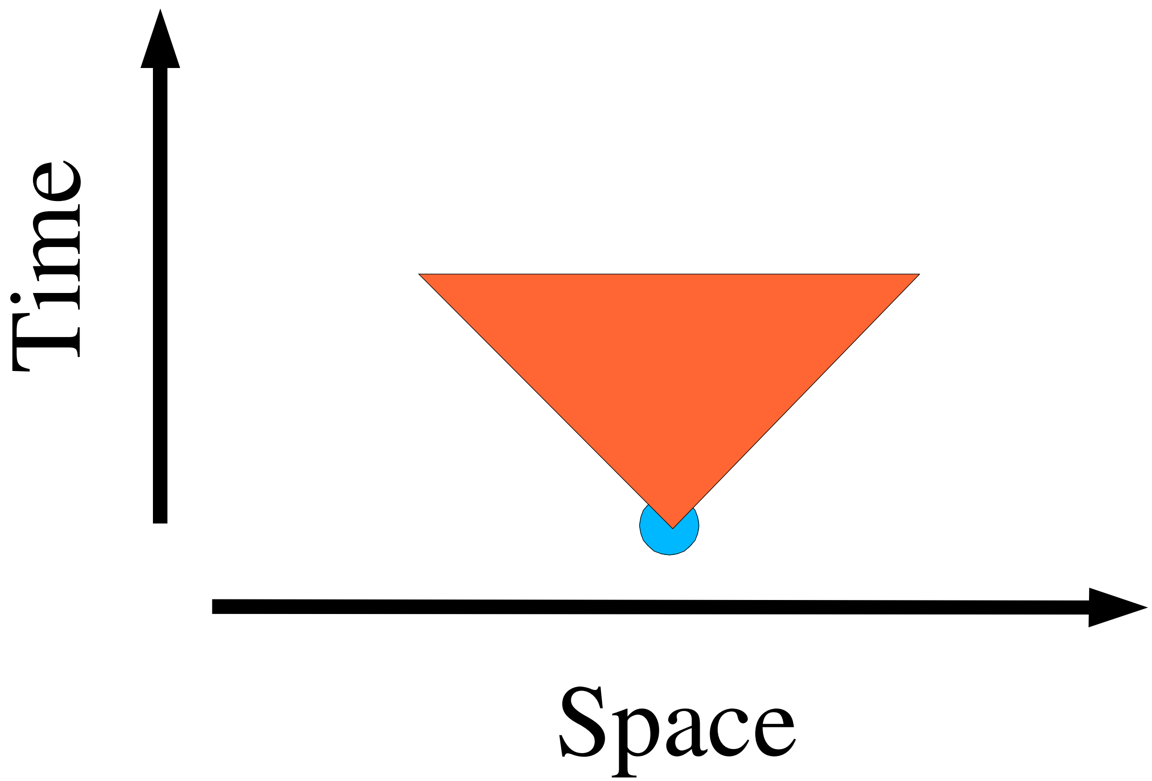

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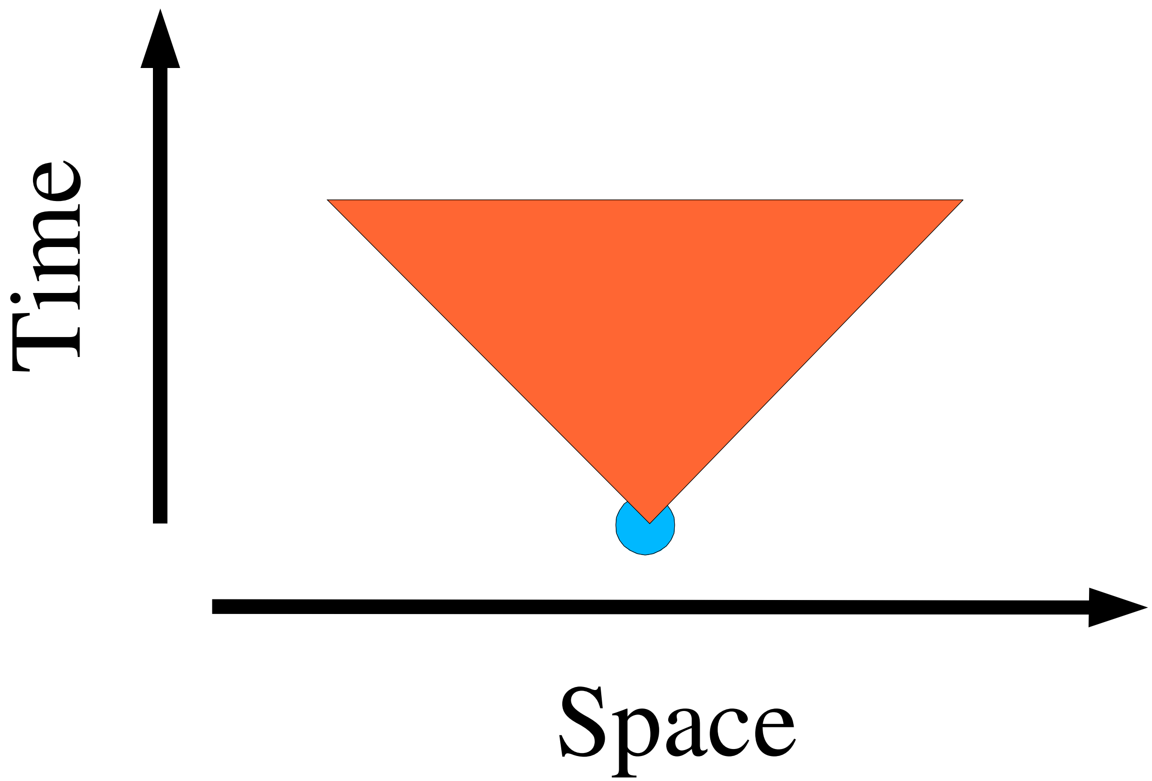

Light cones

- Future light cone
  - The part of spacetime potentially influenced by an event
- Past light cone
  - The part of spacetime that could have influenced an event

Meeting for Beer

- P1 transmits "Panther Hollow Inn" to blackboard

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  - Hey, P2, let's go have a beer.
  - I have transmitted the bar's name to the blackboard.
  - See you there!

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- It says “Squirrel Cage” - ***how???***

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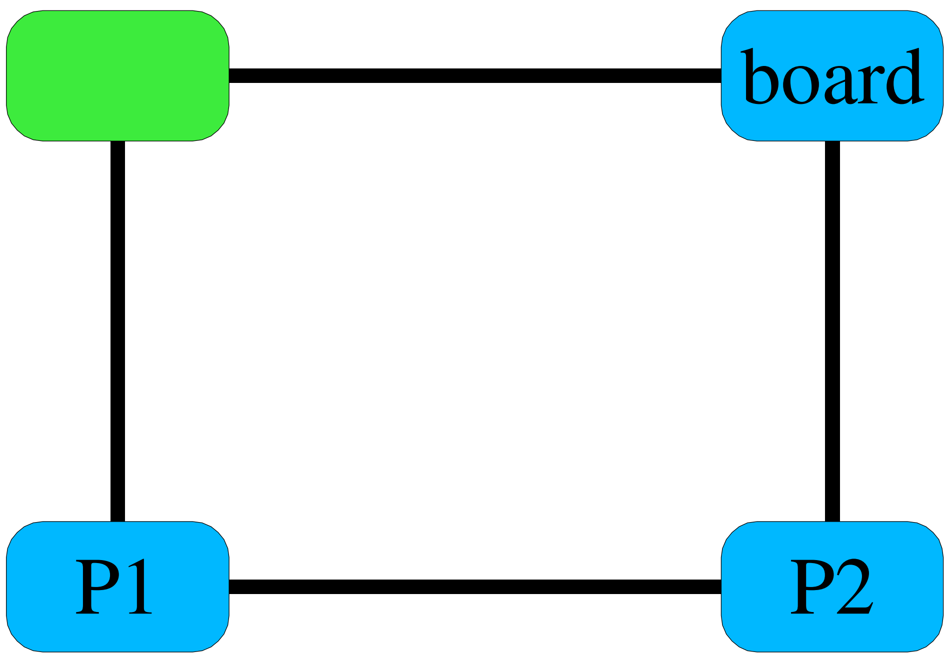

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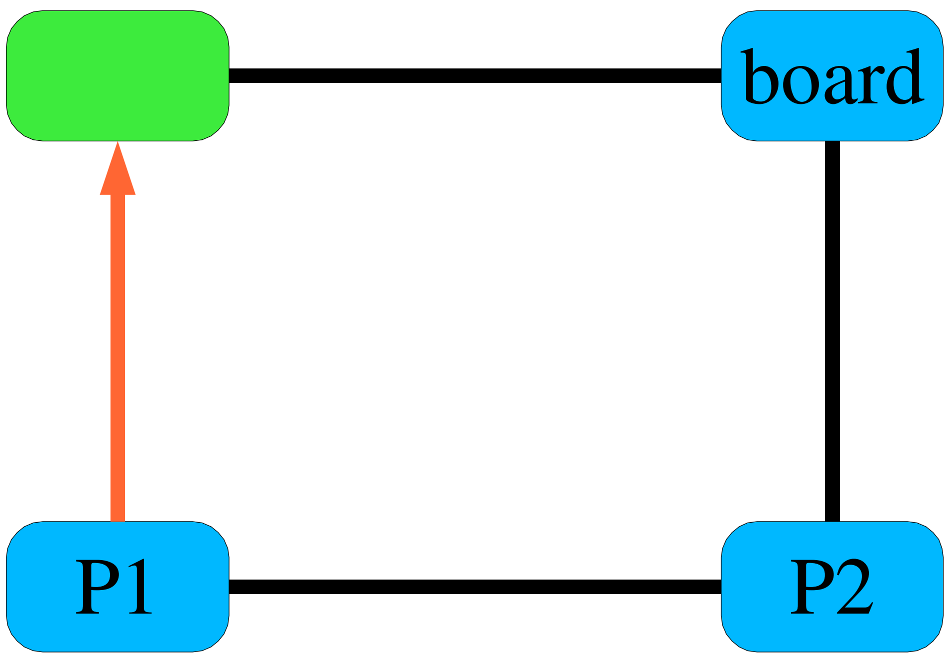

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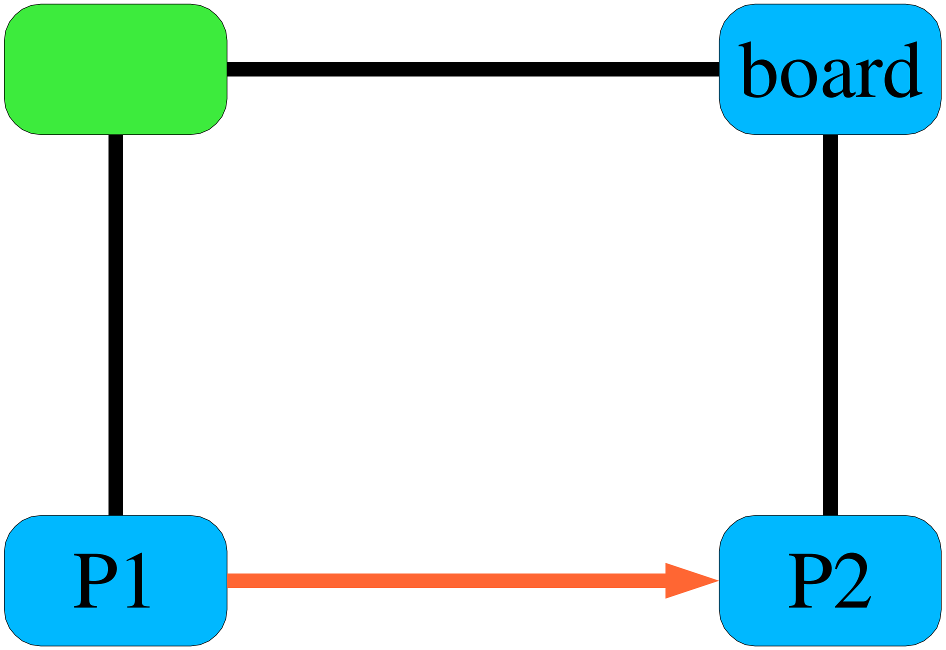

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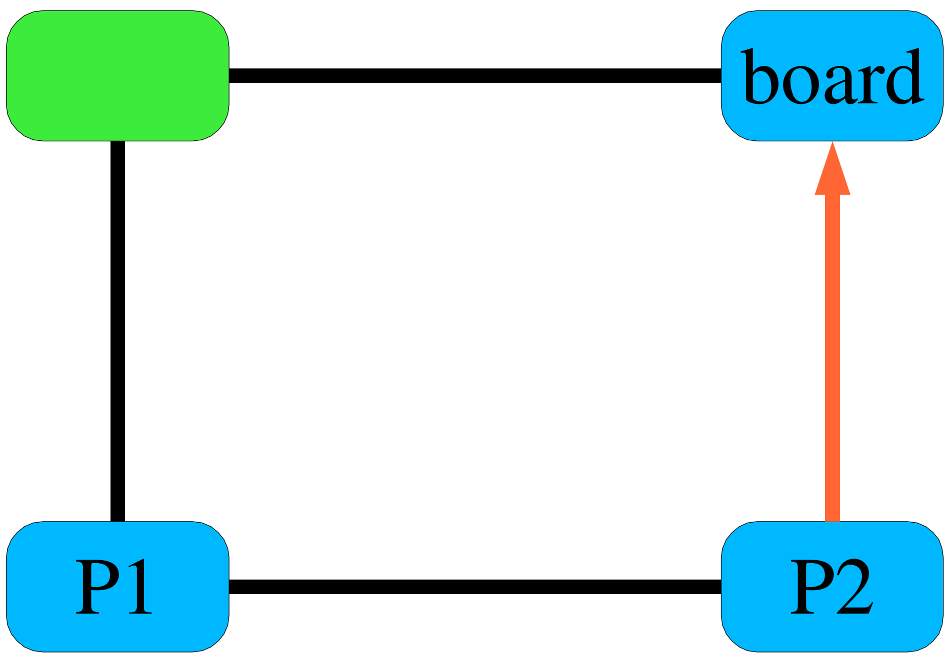

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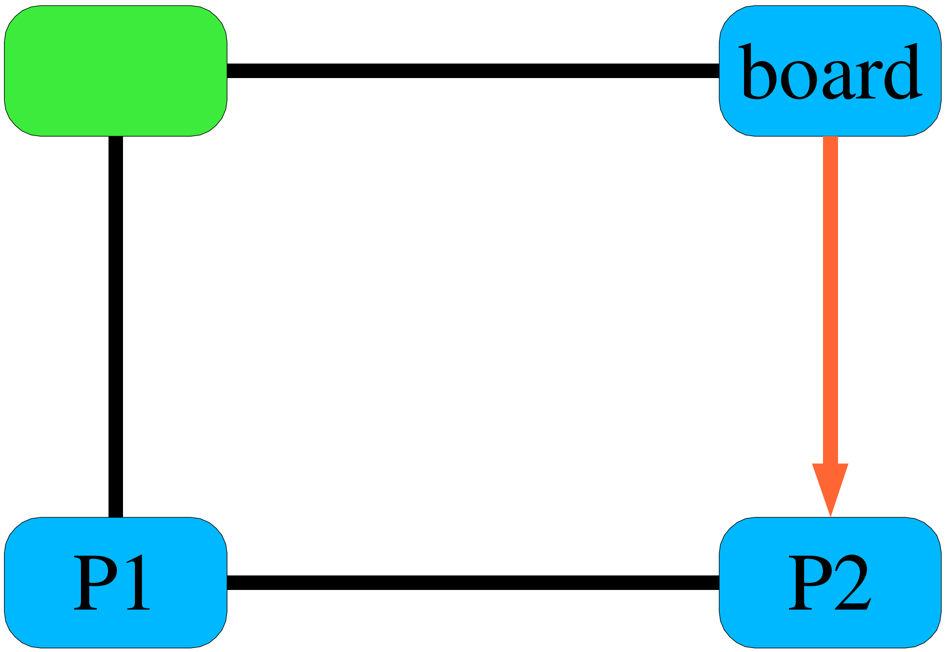

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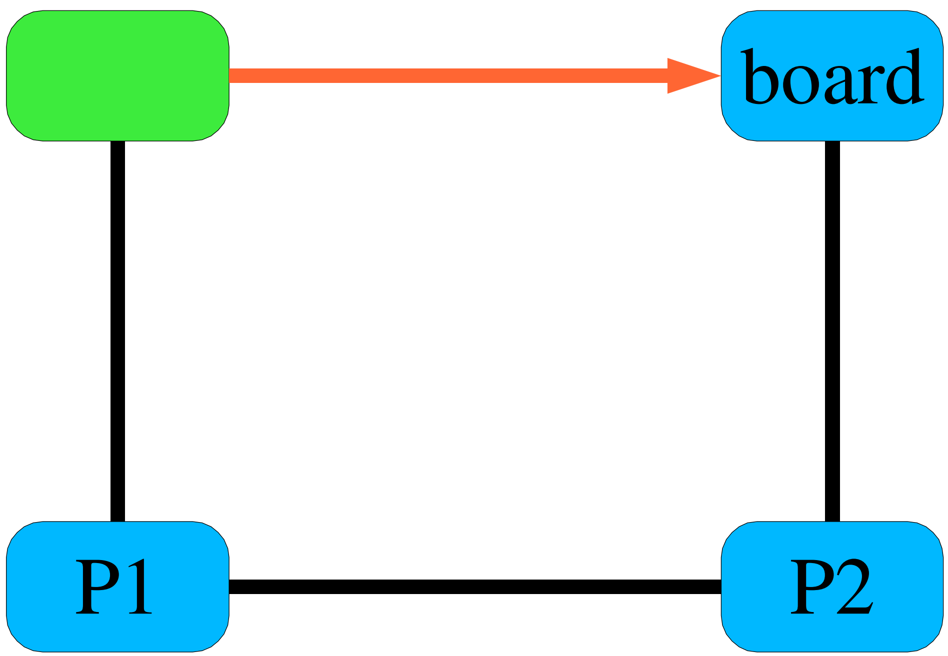


What went wrong?

- P1 thought
  - Blackboard update ***happened before*** invitation
- P2 thought
  - Invitation ***happened before*** blackboard update
- When does an event "happen"?
  - When its effects propagate "everywhere relevant"
- What does "happen before" mean?
- Could that green node really be so slow?

Universe Model

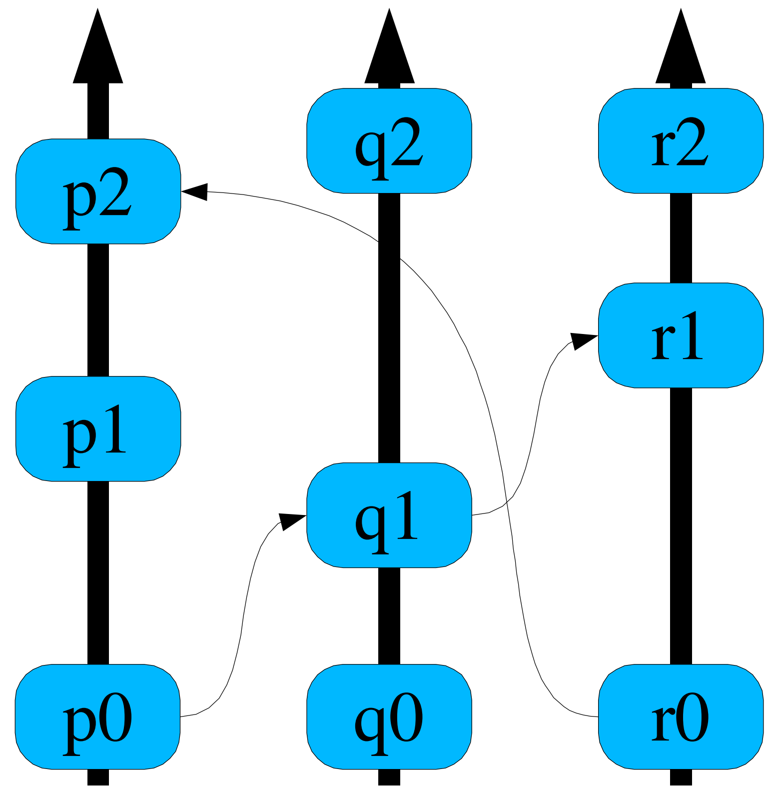
- System = set of processes
- Process = sequence of events
- Event
  - Internal: ++x;
  - Message transmission
  - Message reception

"Happened before" partial order

- A ***happens before*** B ($A \rightarrow B$)
  - If A and B happen inside a process, in (A, B) order
  - If A = transmission, B = reception, of same message
  - If $A \rightarrow B$ and $B \rightarrow C$, then $A \rightarrow C$
- A and B are ***concurrent*** when
  - $A \,!\!\rightarrow B$ and $B \,!\!\rightarrow A$
- Observe: $A \,!\!\rightarrow A$

Space-time Diagram

- $\rightarrow$
  - inside a process, or
  - follow a message
- $p0 \rightarrow r2$
- concurrent
  - p0, q0, r0
  - p1, q1
  - q1, r0
  - p1, r0



→ means "possibly causes"

- p0 possibly causes p1
  - ...by storing something in P's memory
- p0 possibly causes q1
  - Message could trigger q1
- Concurrent events
  - ...cannot cause each other

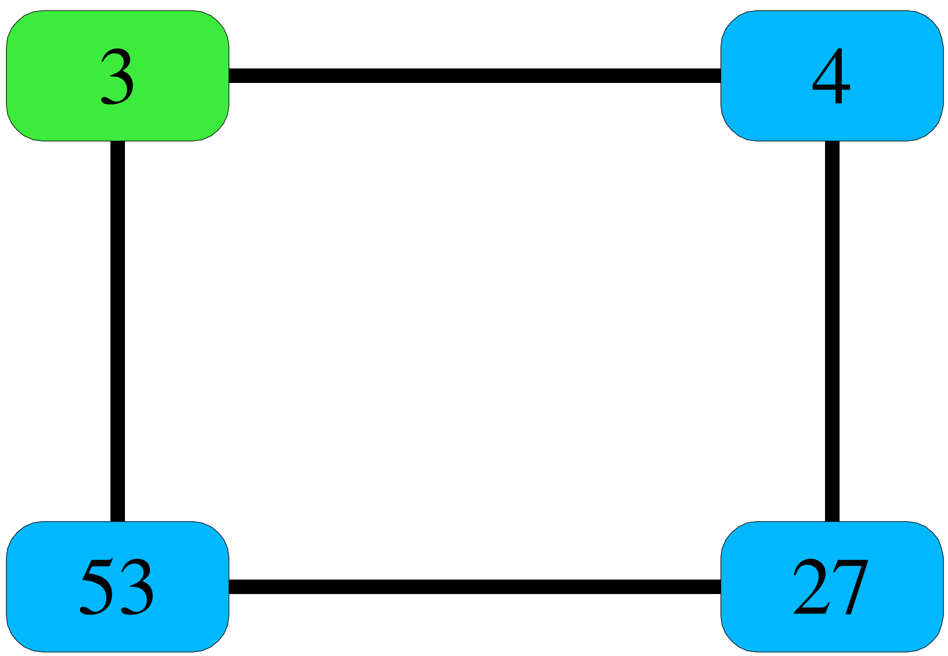
Logical clocks

- Can we assign timestamps to events?
- Want
  - If $A \rightarrow B$ then $C(A) < C(B)$
- Events inside P_i
  - $a \rightarrow b \Rightarrow C_i(a) < C_i(b)$
- Message from P_i to P_j
  - $a = P_i$'s send, $b = P_j$'s receive $\Rightarrow C_i(a) < C_j(b)$

Logical clocks

- Events inside P_i
  - Increment $C_i()$ between successive events
- Message from P_i to P_j
  - Sender: place ***timestamp*** T in message: $C_i(\text{send})$
  - Receiver: ensure $C_j(\text{receive}) > T$

Meeting for Beer



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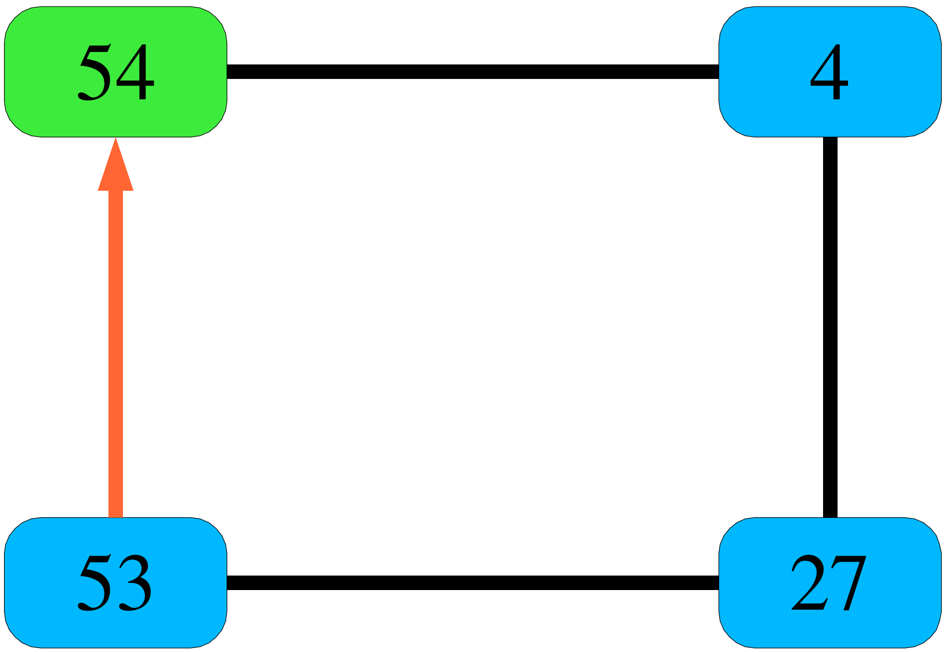

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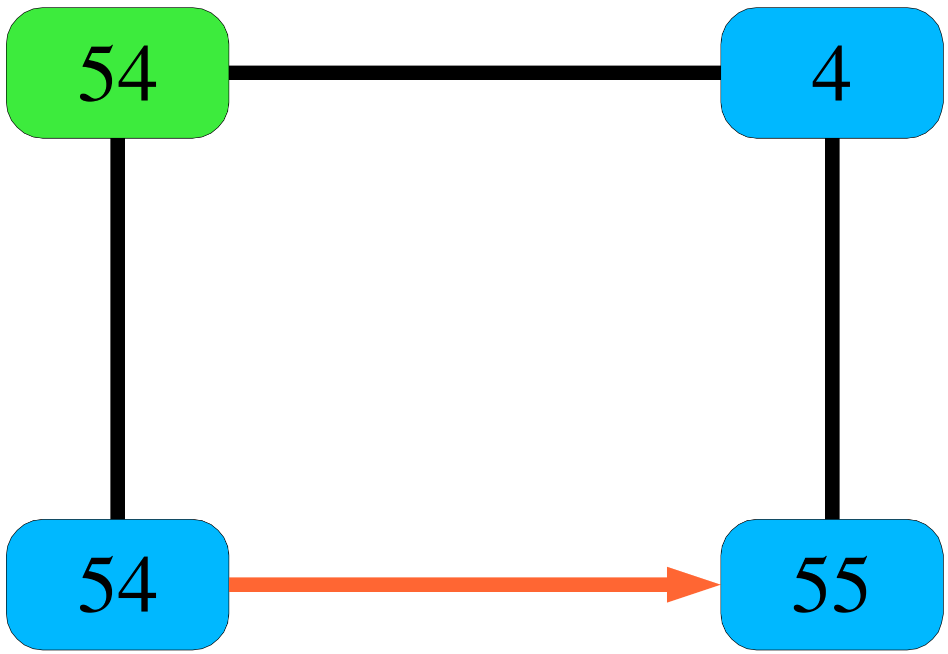

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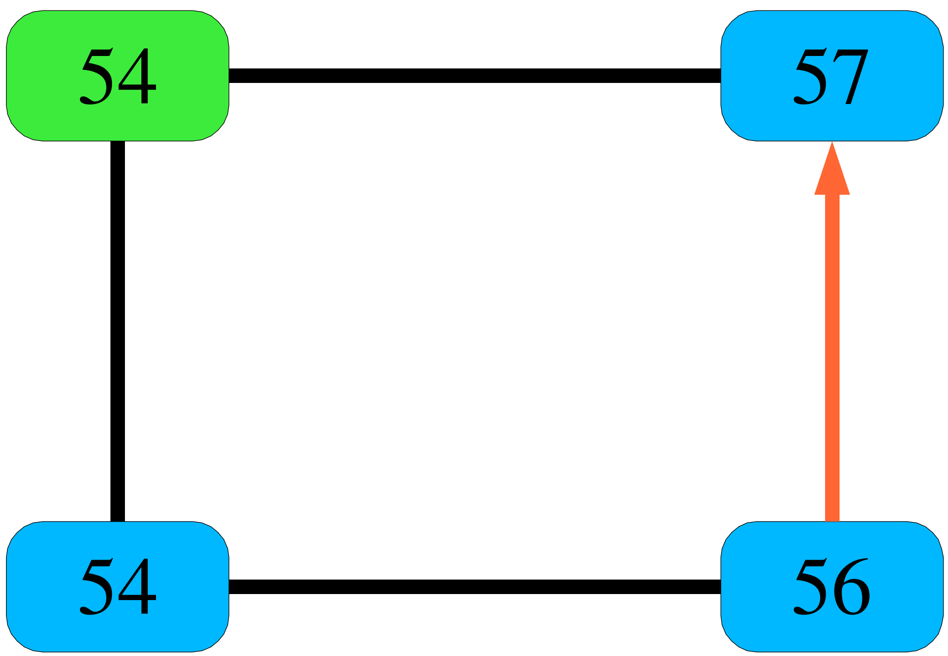

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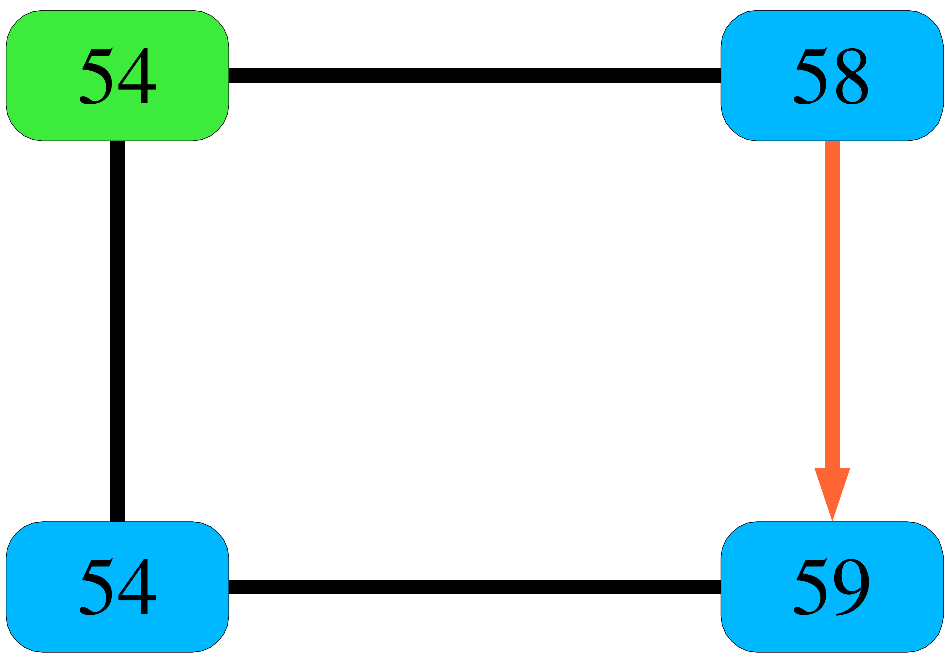

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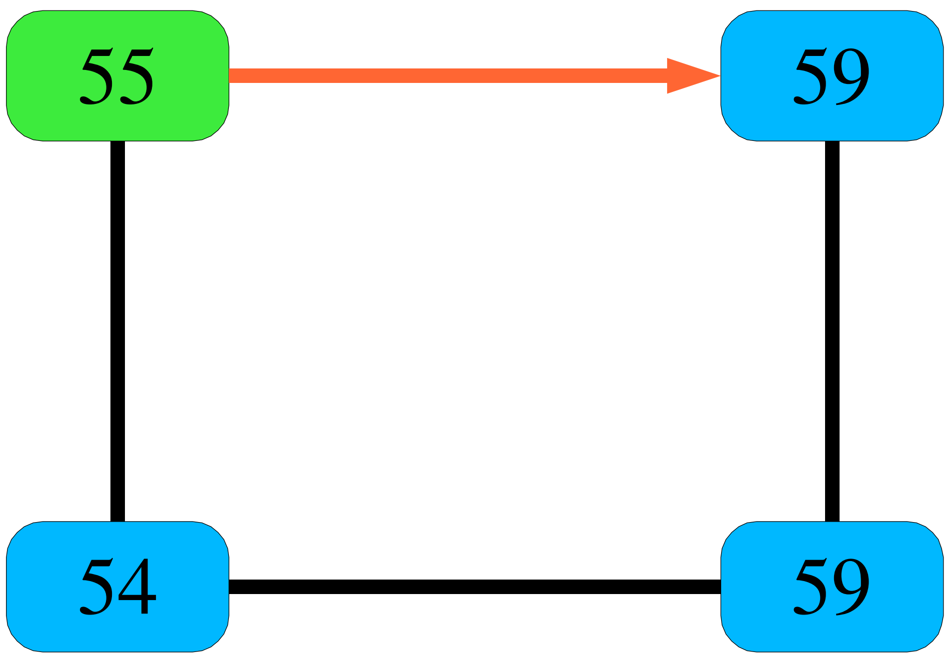
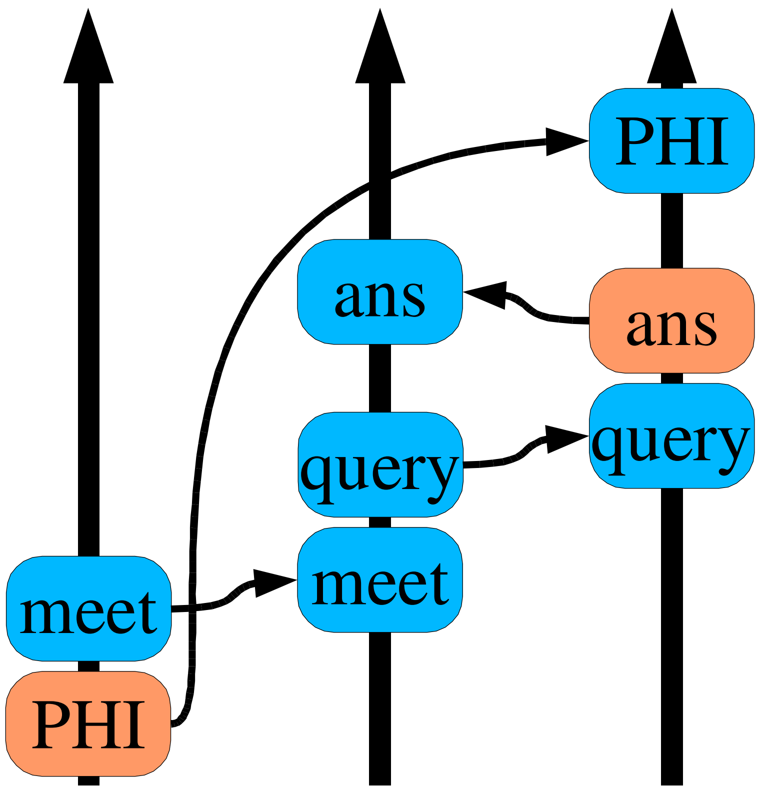


What this means

- P1 wants
  - <"PHI" written> ***happened before*** <read by P2>
- Equivalent to "59 < 57" (oops)
- The events were ***concurrent***
- "PHI" ***could not cause*** P2's bar trip

Space-time Diagram

- P1 wants
  - <“PHI” declared> ***happened before*** <P2 decided>
- Equivalent to “59 < 57” (oops)
- The events were ***concurrent***
- “PHI” ***could not cause*** P2's bar trip



Fixing the problem

- P1 should wait for board to acknowledge
- “PHI” causes ACK
- ACK causes “Meet me at...”
- “Meet me at...” causes bar trip
- Then: “PHI” causes bar trip

Extensions

- Define ***total ordering*** of system events
  - Typical (timestamp, process #) tuple comparison
    - Process # used to break timestamp ties
- Distributed agreement algorithms
  - Such as "***fair*** distributed mutual exclusion"
    - Requests must be granted "in order"
    - See text: 17.2
- Adding physical (real-time) clocks

Summary

- Light cones
- "Happened before" partial order
- Potential causality
- Another definition of concurrency
  - You've dealt with single-clock race conditions
    - (one memory bus provides one global clock)
  - In distributed systems there is no global clock
- Timestamps track message causality