15-410

"...What goes around comes around..."

Disks
Mar. 24, 2004

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Contributions from

- **Eno Thereska**
- **15-213**
- **"How Stuff Works" web site**

Synchronization

Project 3 suggestions

- Three regular meeting times per week
  - Two hours or more at each meeting
  - Begin by asking questions about each other's code
    - » Requires having read code before meeting
    - » Requires “quiet time” between check-ins and meeting
- Source control
  - Frequent merges, not a single “big bang” at end
- Leave time at end for those multi-day bugs

Overview

Anatomy of a Hard Drive

Common Disk Scheduling Algorithms

Anatomy of a Hard Drive

On the outside, a hard drive looks like this

Taken from "How Hard Disks Work"
http://computer.howstuffworks.com/hard-disk2.htm

Anatomy of a Hard Drive

If we take the cover off, we see that there actually is a "hard disk" inside

Taken from "How Hard Disks Work"
http://computer.howstuffworks.com/hard-disk2.htm

Anatomy of a Hard Drive

A hard drive usually contains multiple disks, called *platters*

These spin at thousands of RPM (5400, 7200, etc)

Taken from "How Hard Disks Work"
http://computer.howstuffworks.com/hard-disk2.htm

Anatomy of a Hard Drive

Information is written to and read from the platters by the *read/write heads* on the *disk arm*

Taken from "How Hard Disks Work"
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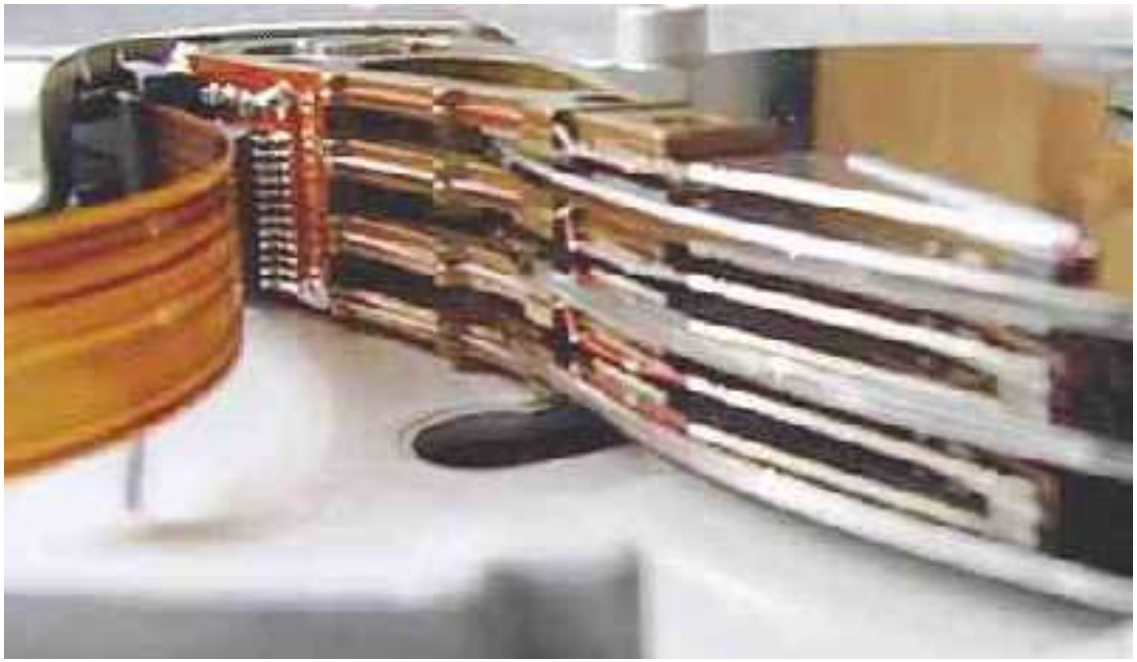
Anatomy of a Hard Drive

Both sides of each platter store information

Each side of a platter is called a *surface*

Each surface has its own read/write head



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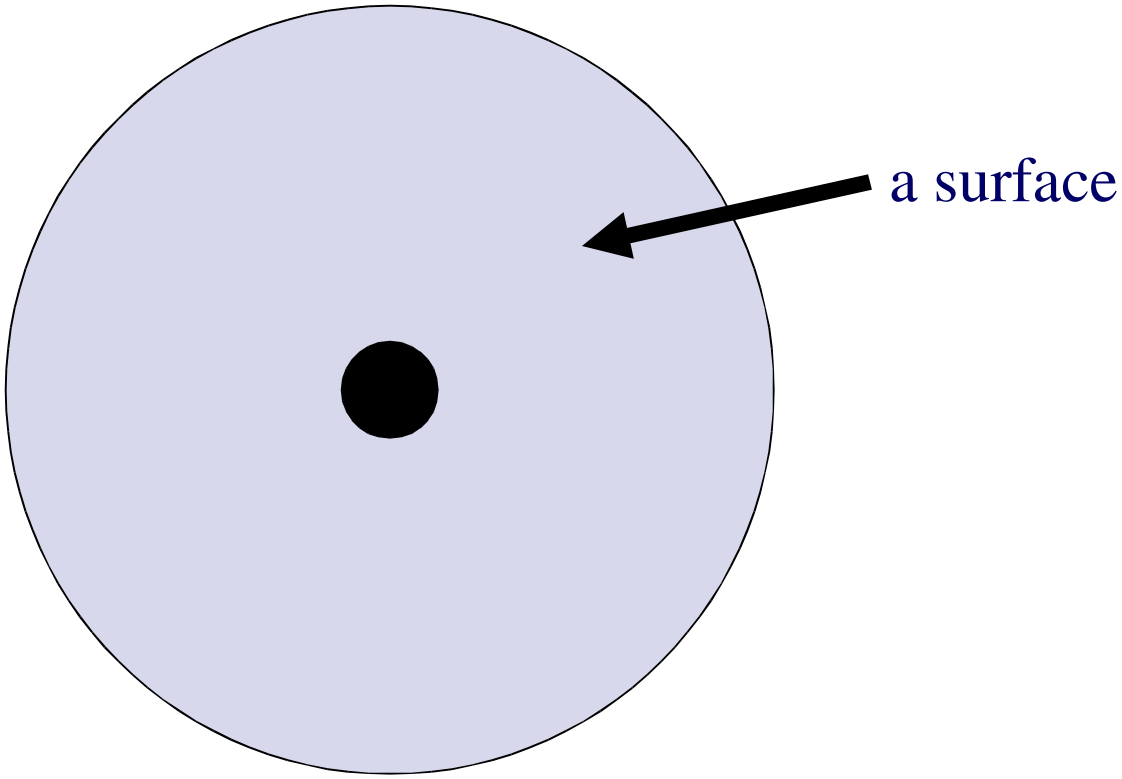
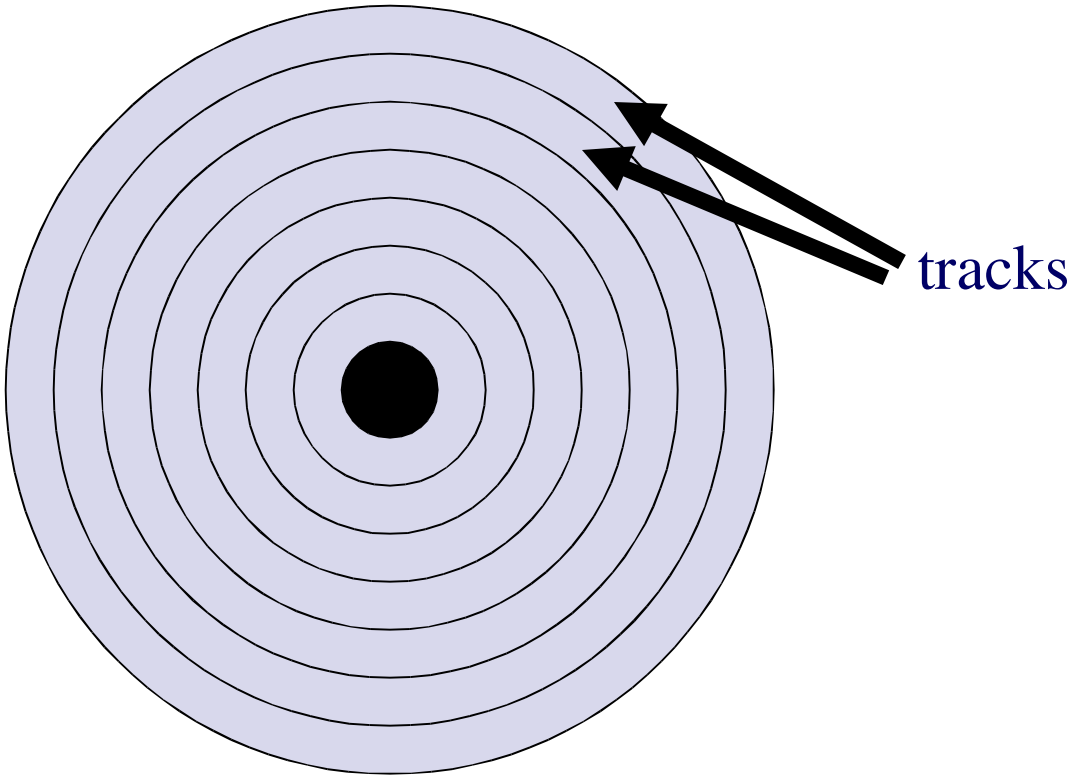
Anatomy of a Hard Drive

How are the surfaces organized?



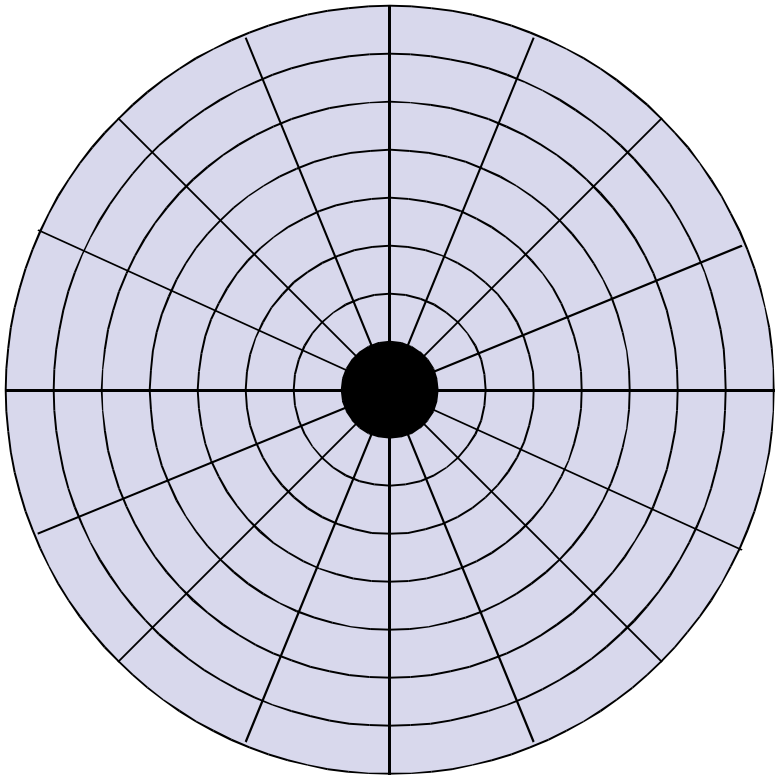
Anatomy of a Hard Drive

Each surface is divided by concentric circles, creating *tracks*



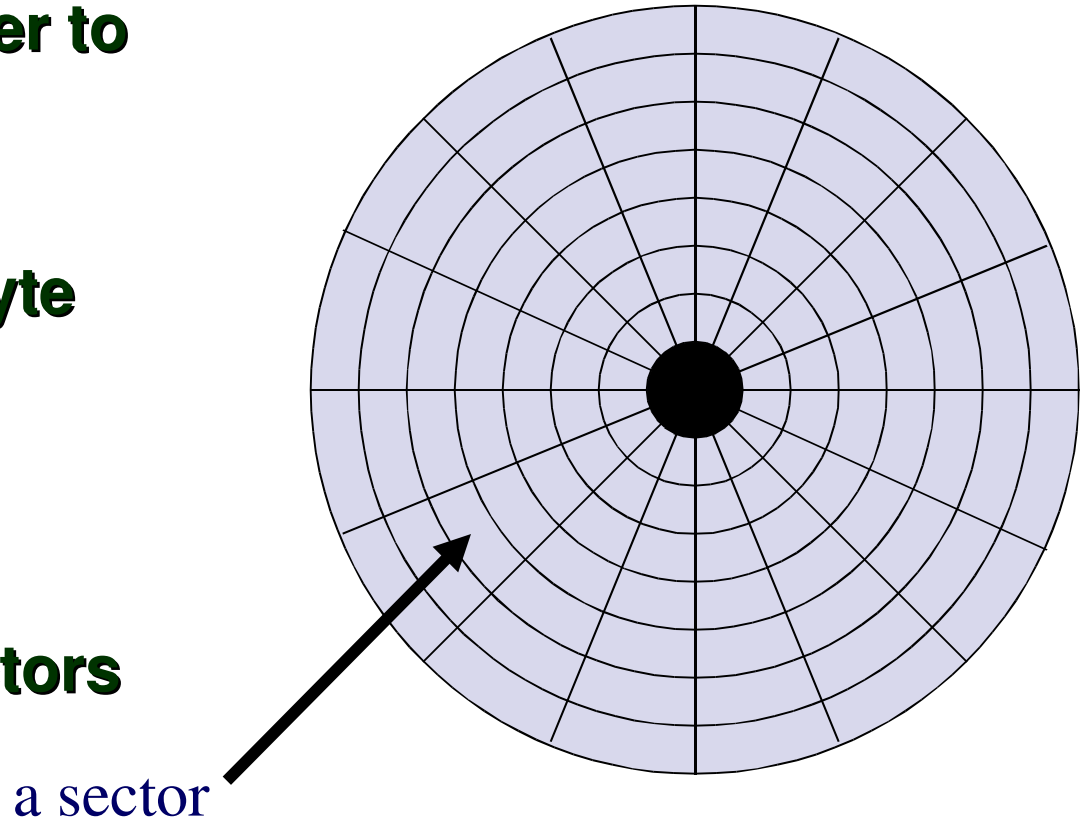
Anatomy of a Hard Drive

These tracks are further divided into *sectors*



Anatomy of a Hard Drive

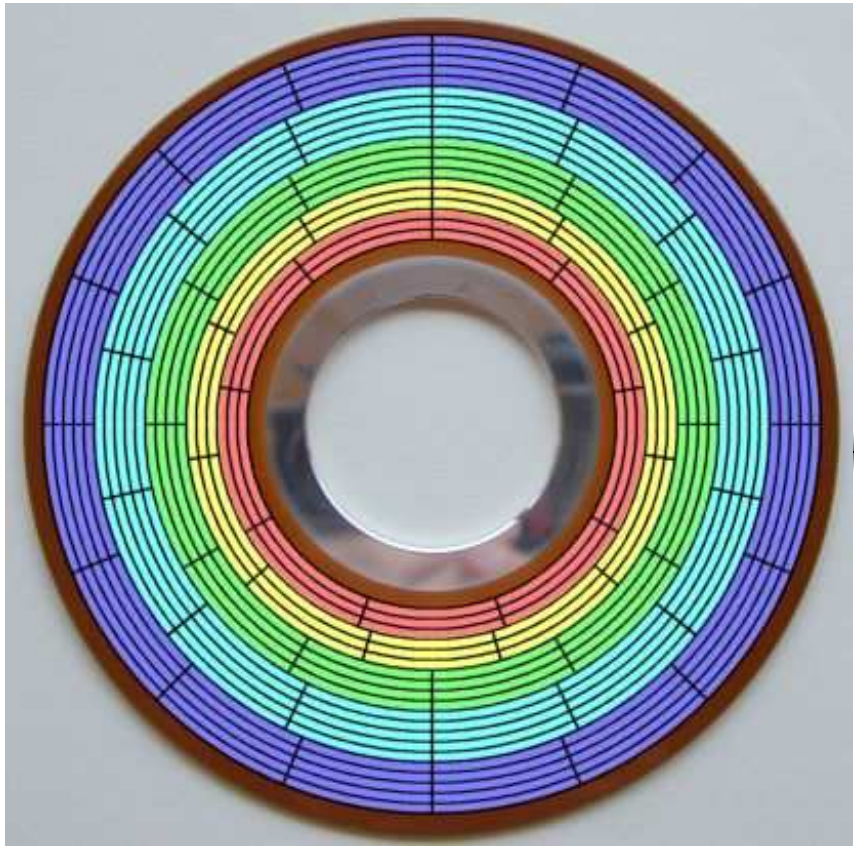
- A sector is the smallest unit of data transfer to or from the disk
- Most modern hard drives have 512 byte sectors
- CD-ROM sectors are 2048 bytes
- Gee, those outer sectors look bigger...



Anatomy of a Hard Drive

Does this mean that sectors on the outside of a surface are larger than those on the inside?

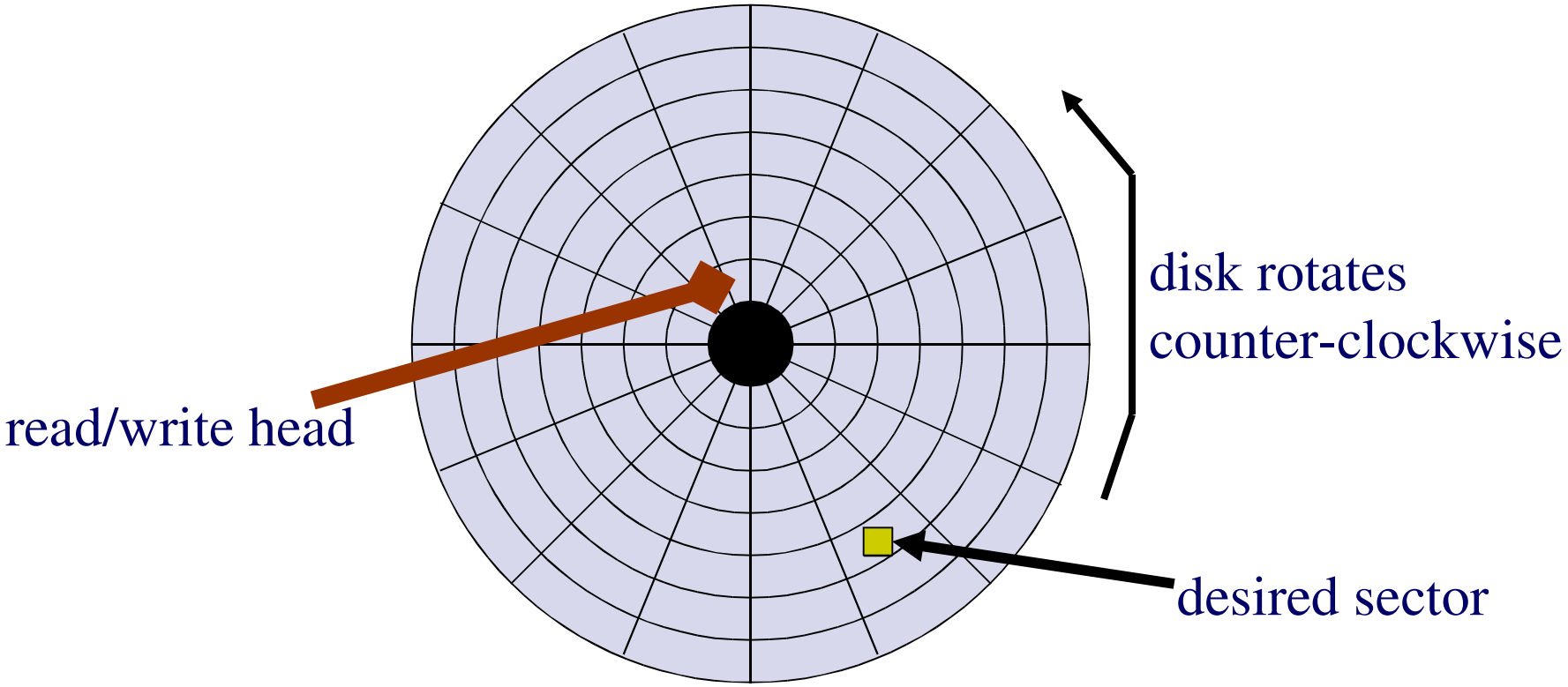
Modern hard drives fix this with *zoned bit recording*



Taken from "Reference Guide – Hard Disk Drives"
http://www.storagereview.com/map/lm.cgi/zone

Anatomy of a Hard Drive

Let's read in a sector from the disk



Anatomy of a Hard Drive

We need to do two things to transfer a sector

1. **Move the read/write head to the appropriate track (seek)**

2. **Wait until the desired sector spins around**

Anatomy of a Hard Drive

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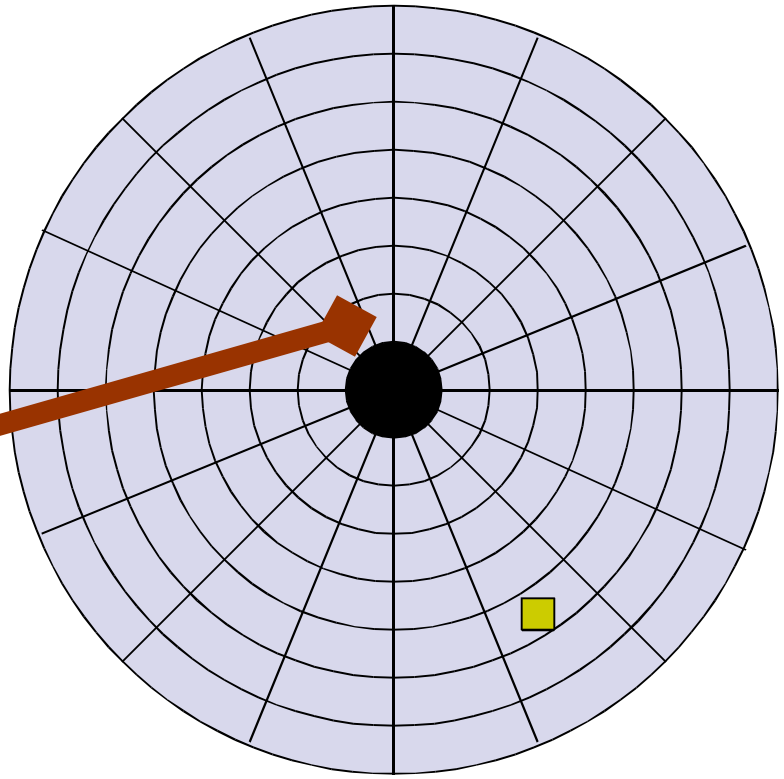

read/write head

Anatomy of a Hard Drive

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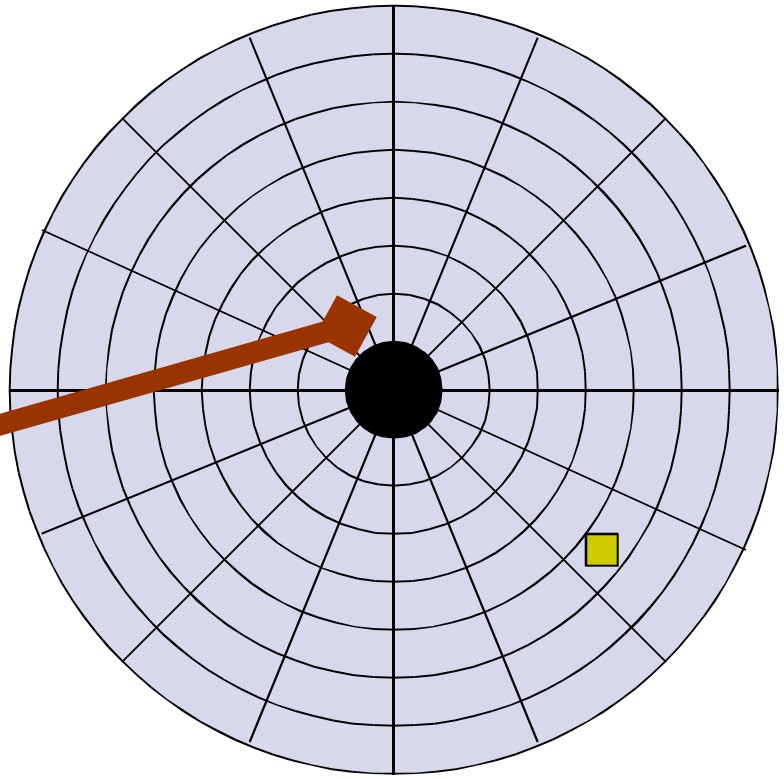

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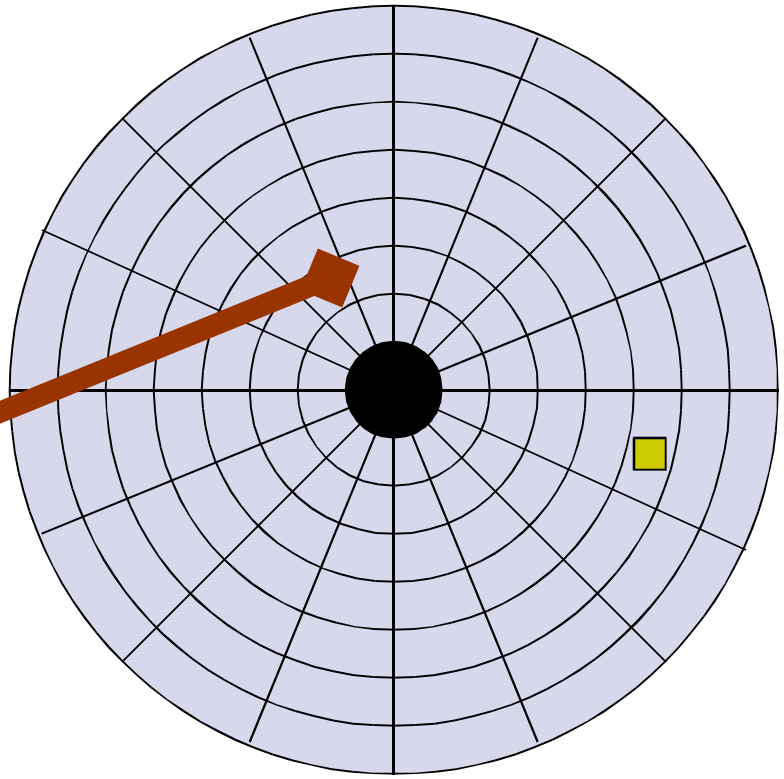

read/write head

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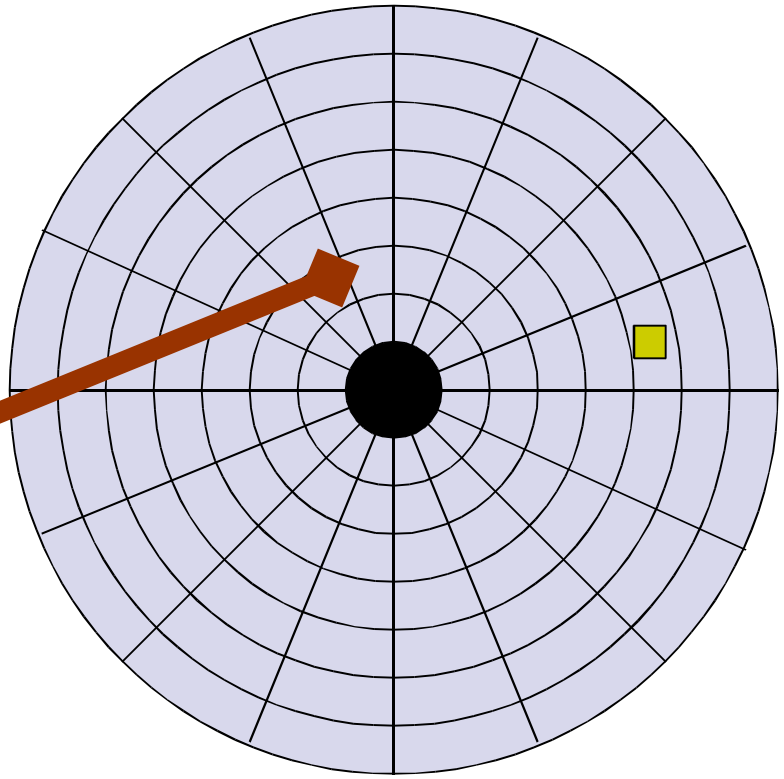

read/write head

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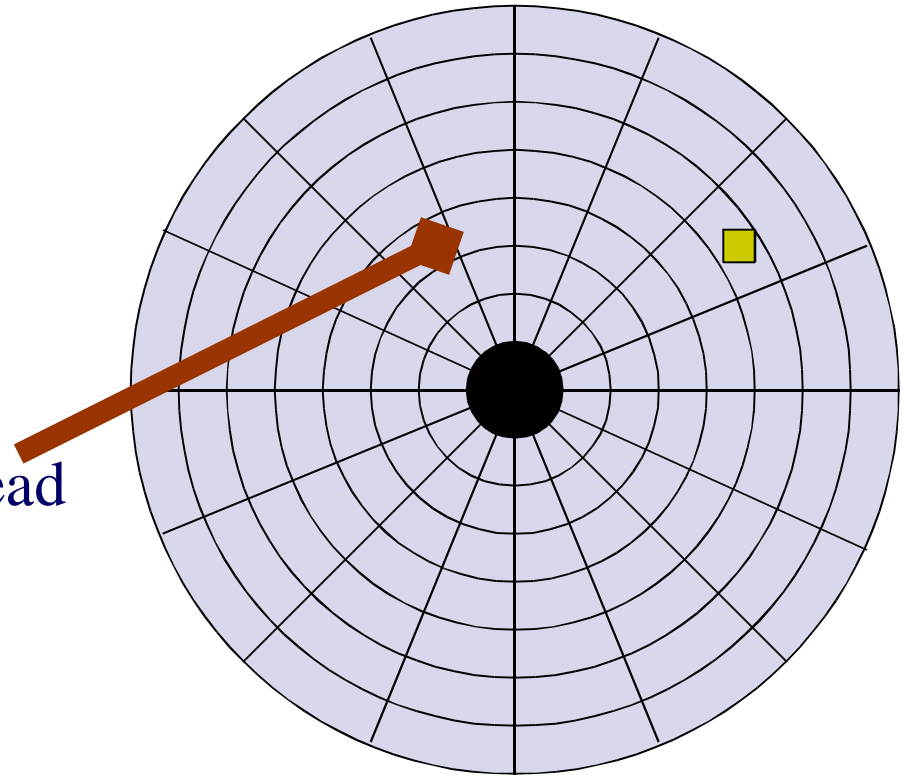

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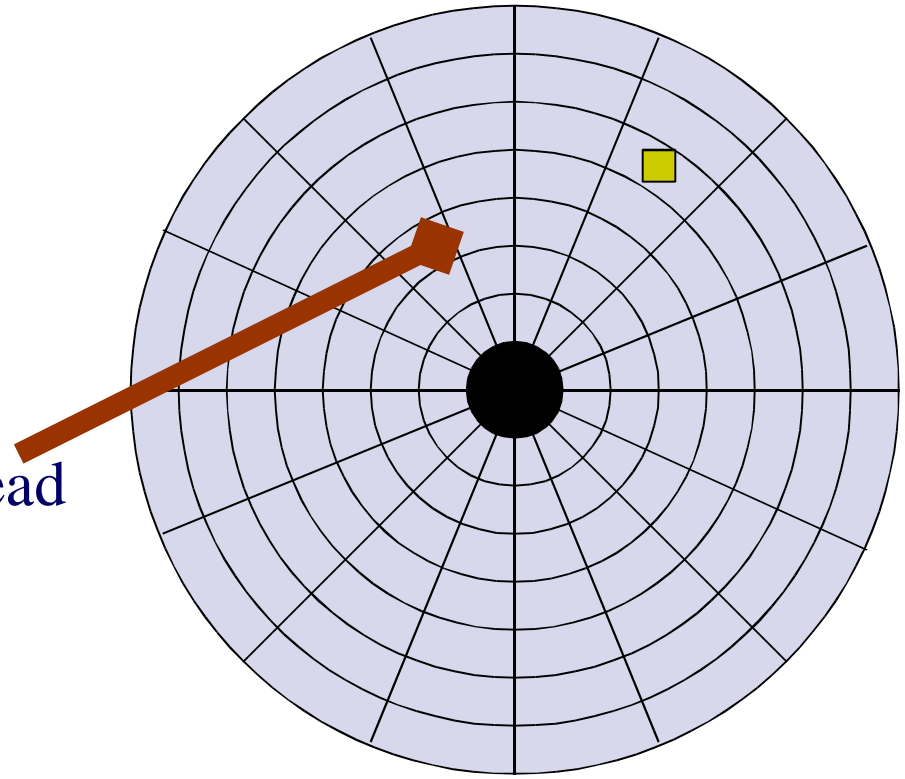

Anatomy of a Hard Drive

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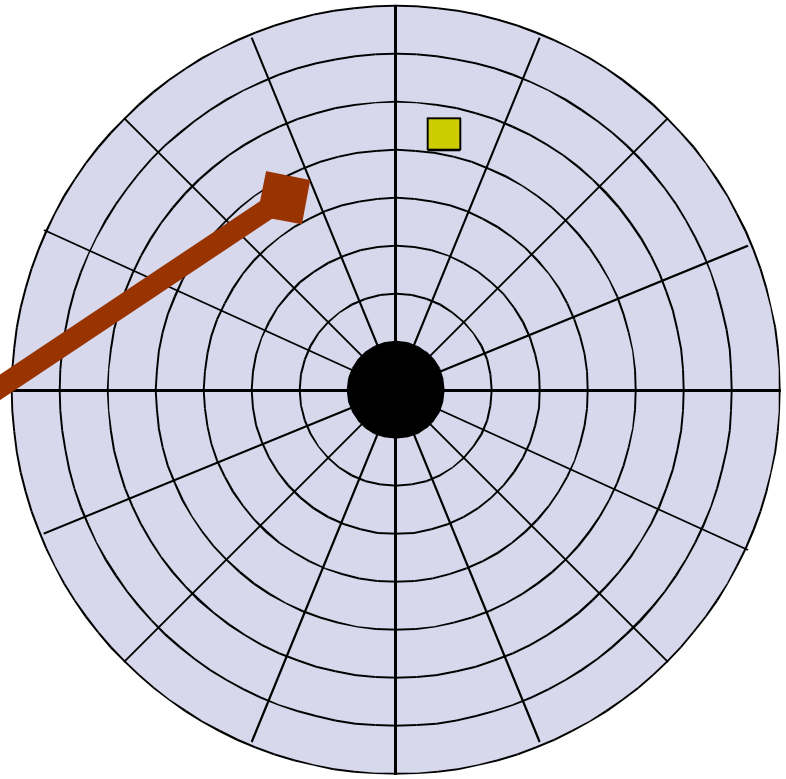

read/write head

Anatomy of a Hard Drive

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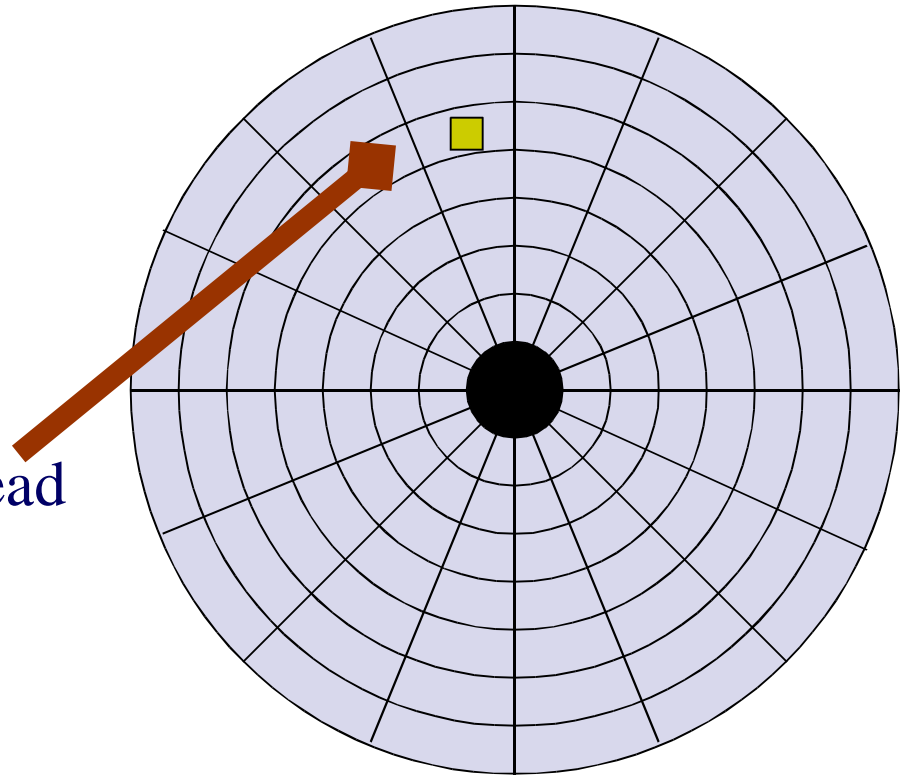

Anatomy of a Hard Drive

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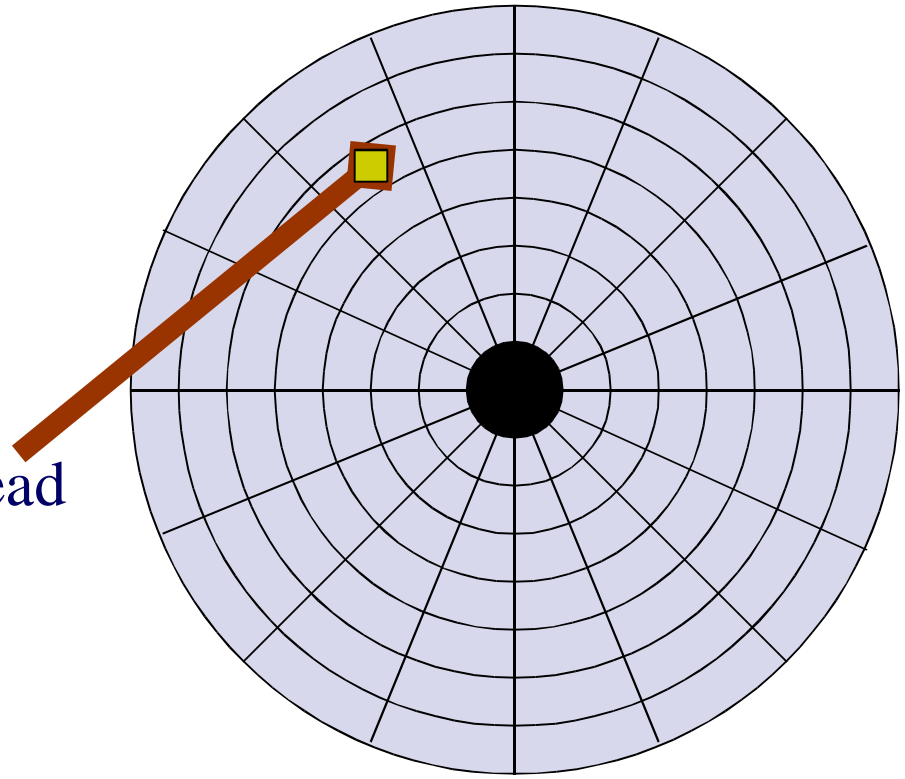
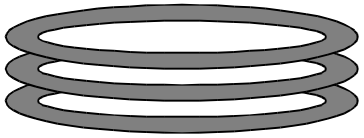


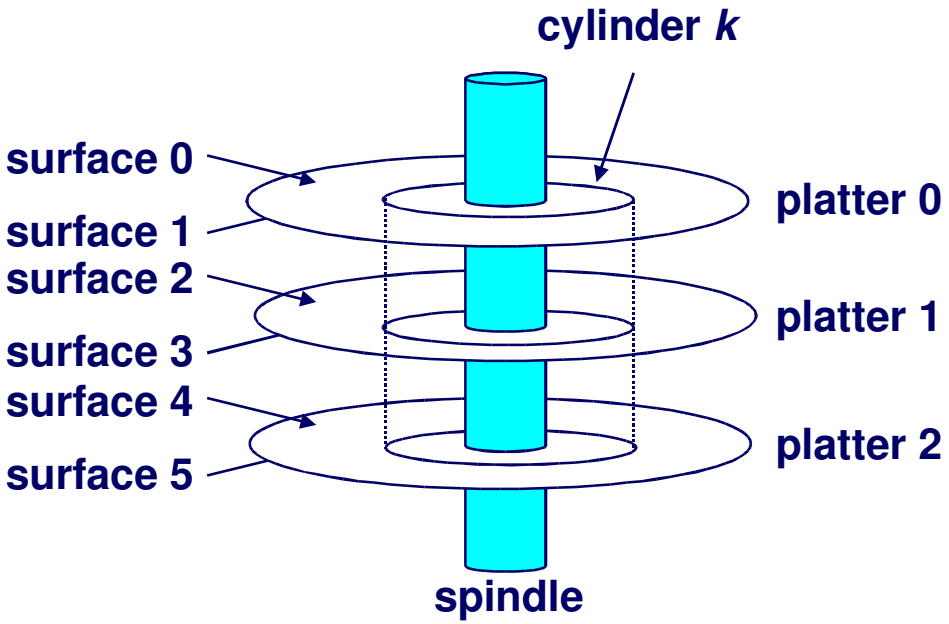
Disk Cylinder

Matching tracks across surfaces are collectively called a *cylinder*



Disk Cylinder

Matching tracks form a cylinder.



Cheap Access Within A Cylinder

Heads on single arm

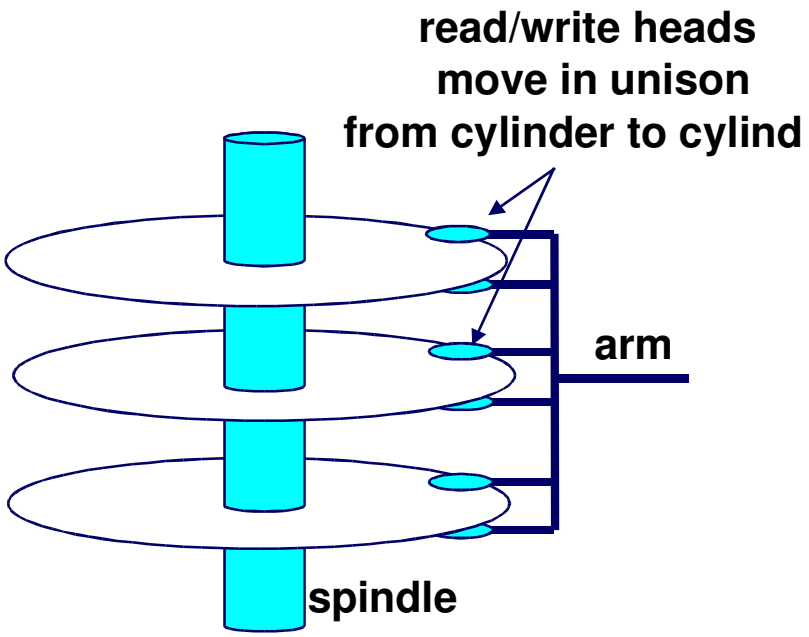
- All heads always on same cylinder

Switching heads is "cheap"

- Deactive head 3
- Activate head 4
- Wait for 1st sector header

Optimal transfer rate

- Transfer all sectors on a track
- Transfer all tracks on a cylinder
- *Then* move elsewhere



Anatomy of a Hard Drive

On average, we will have to move the read/write head over half the tracks

The time to do this is the “average seek time”, and is ~10ms for a 5400 rpm disk

We will also must wait half a rotation

The time to do this is rotational latency, and on a 5400 rpm drive is ~5.5ms

Anatomy of a Hard Drive

Other factors influence overall disk access time including

- Settle time, the time to stabilize the read/write head after a seek
- Command overhead, the time for the disk to process a command and start doing something

Minor compared to seek time and rotational latency

Anatomy of a Hard Drive

Total drive random access time is on the order of 15 to 20 milliseconds

- 50 ½-kilobyte transfers per second = 25 Kbyte/sec
- Oh man, disks are slow
  - But wait! Disk transfer rates are quoted at tens of Mbytes/sec!

What can we, as operating system programmers, do about this?

- Read more per seek (multi-sector transfers)
- Don't seek so randomly (“disk scheduling”)

Disk Scheduling Algorithms

The goal of a disk scheduling algorithm is to be nice to the disk

We can help the disk by giving it requests that are located close to each other

This minimizes seek time, and possibly rotational latency

There exist a variety of ways to do this

Addressing Disks

What the OS knows about the disk?

- Interface type (IDE/SCSI), unit number, number of sectors

What happened to sectors, tracks, etc?

- Old disks were addressed by cylinder/head/sector (CHS)
- Modern disks are addressed by abstract sector number
  - LBA = logical block addressing

Who uses sector numbers?

- File systems assign logical blocks to files

Terminology

- To disk people, “block” and “sector” are the same
- To file system people, a “block” is some number of sectors

Disk Addresses vs. Scheduling

Goal of OS disk-scheduling algorithm

- Understand which sectors are close to each other
- Maintain queue of requests
- When disk finishes one request, give it the “best” request

Goal of disk's logical addressing

- Hide messy details of which sectors are located where

Oh, well

- Older OS's used to try to understand disk layout
- Modern OS's just assume nearby sector numbers are close
- Experimental OS's try to understand disk layout again
- Next few slides assume “modern”, not “old”/“experimental?

First Come First Served (FCFS)

Send requests to disk as they are generated by the OS

Trivial to implement – FIFO queue in device driver

Fair

- **What could be more fair?**

"Unacceptably high mean response time"

- **File "abc" in sectors 1, 2, 3, ...**
- **File "def" in sectors 16384, 16385, 16386, ...**
- **Sequential reads: 1, 16384, 2, 16385, 3, 16386**

"Fair, but cruel"

Shortest Seek Time First (SSTF)

Maintain "queue" of disk requests

Serve the request nearest to the disk arm

- **Easy to estimate by difference in block numbers**

Generates very fast response time

Intolerable response time *variance*, however

Why?

What Went Wrong?

FCFS - "fair, but cruel"

- Ignores position of disk arm, very slow

SSTF – good throughput, very unfair

- Ignores necessity of eventually scanning entire disk

"Scan entire disk" - now that's an idea!

- Start disk arm moving in one direction
- Serve requests as the arm moves past them
  - No matter what order they arrived in
- When arm bangs into stop, reverse direction

SCAN

Doubly-linked ordered list of requests

- Insert according to order

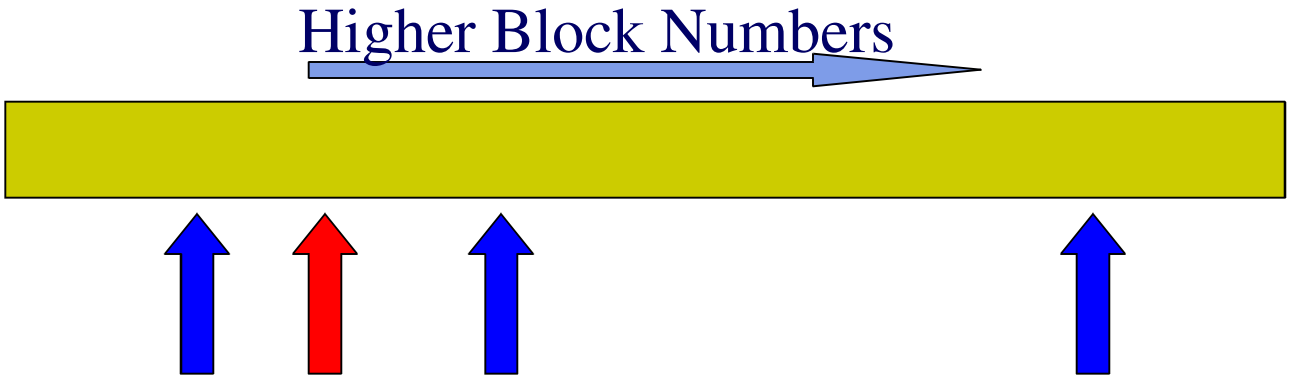
Bi-directional scanning

- Direction = +1 or -1
- Tell disk: “seek to cylinder X=current+direction”
- Examine list for requests in cylinder X
- If X == 0 or X == max, direction = -direction
- Else current = X

SCAN

Blue are requests

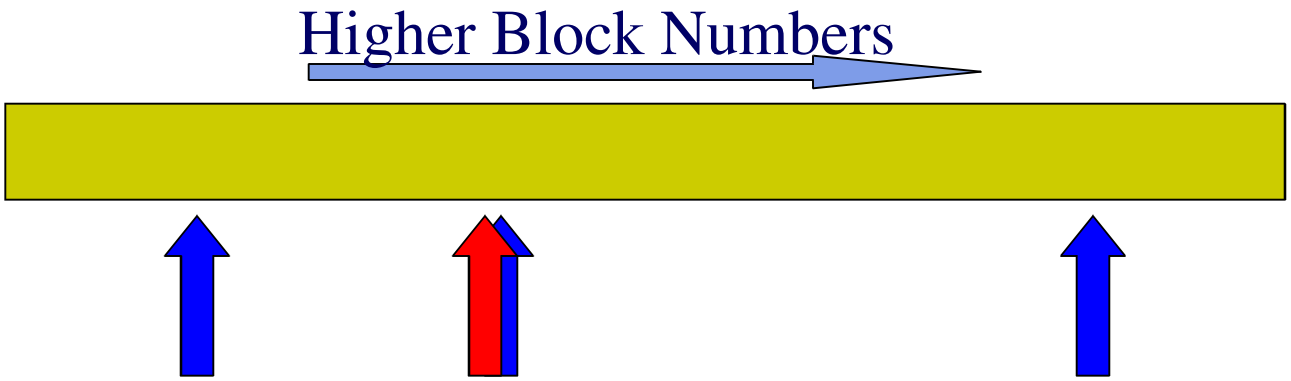
Yellow is disk



Red is disk head

Green is completed requests

SCAN



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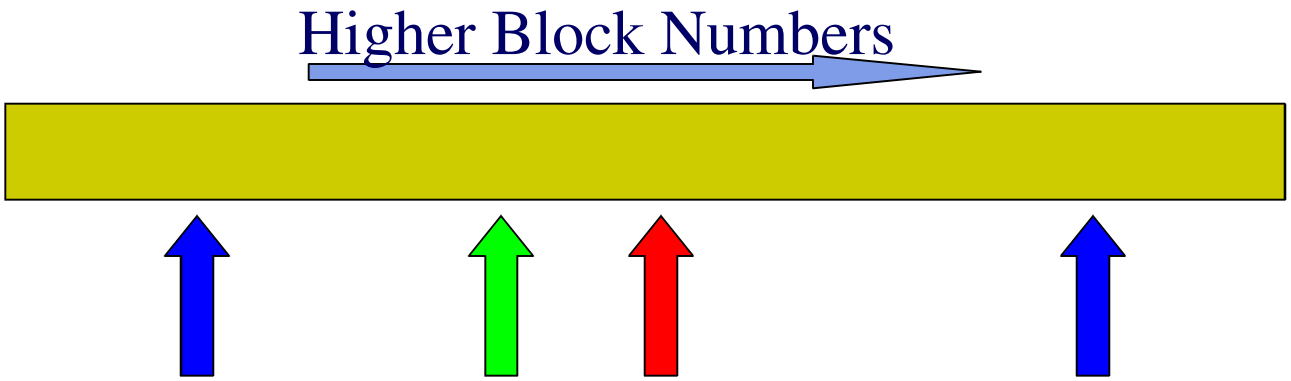

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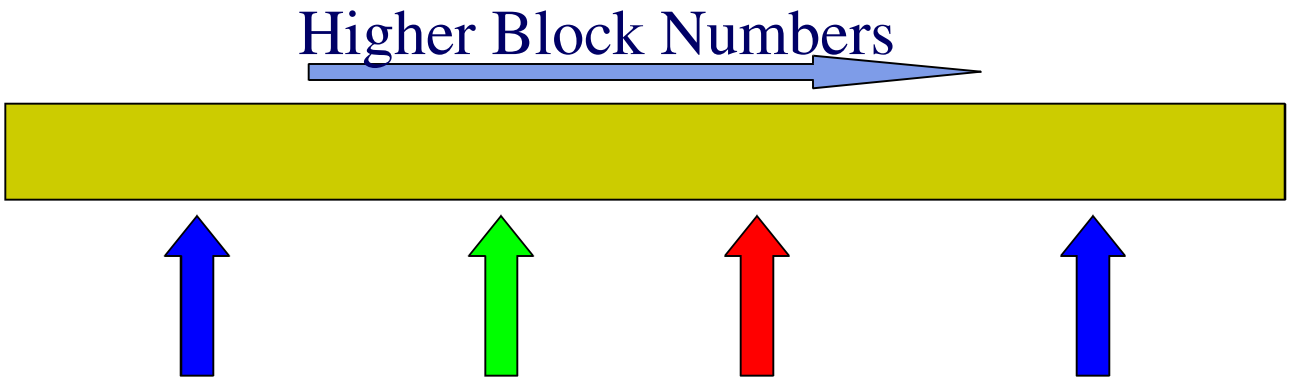

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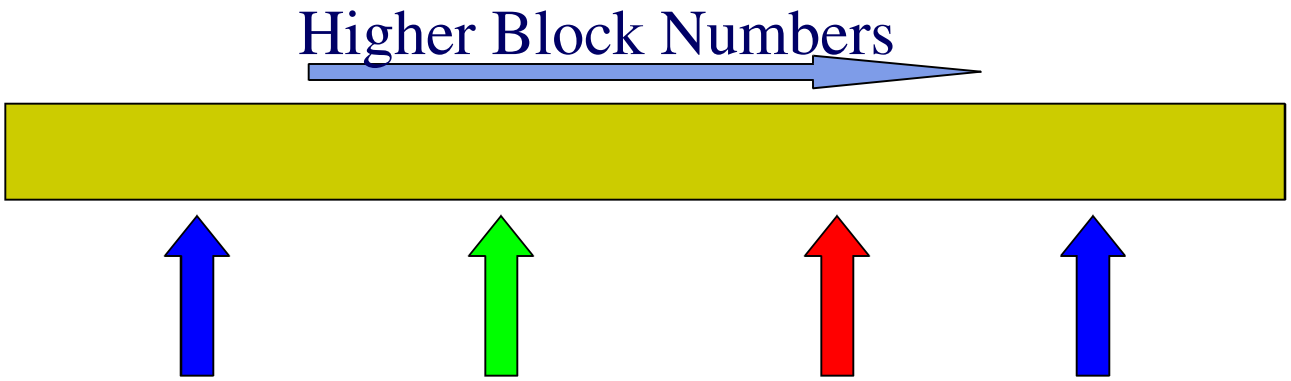

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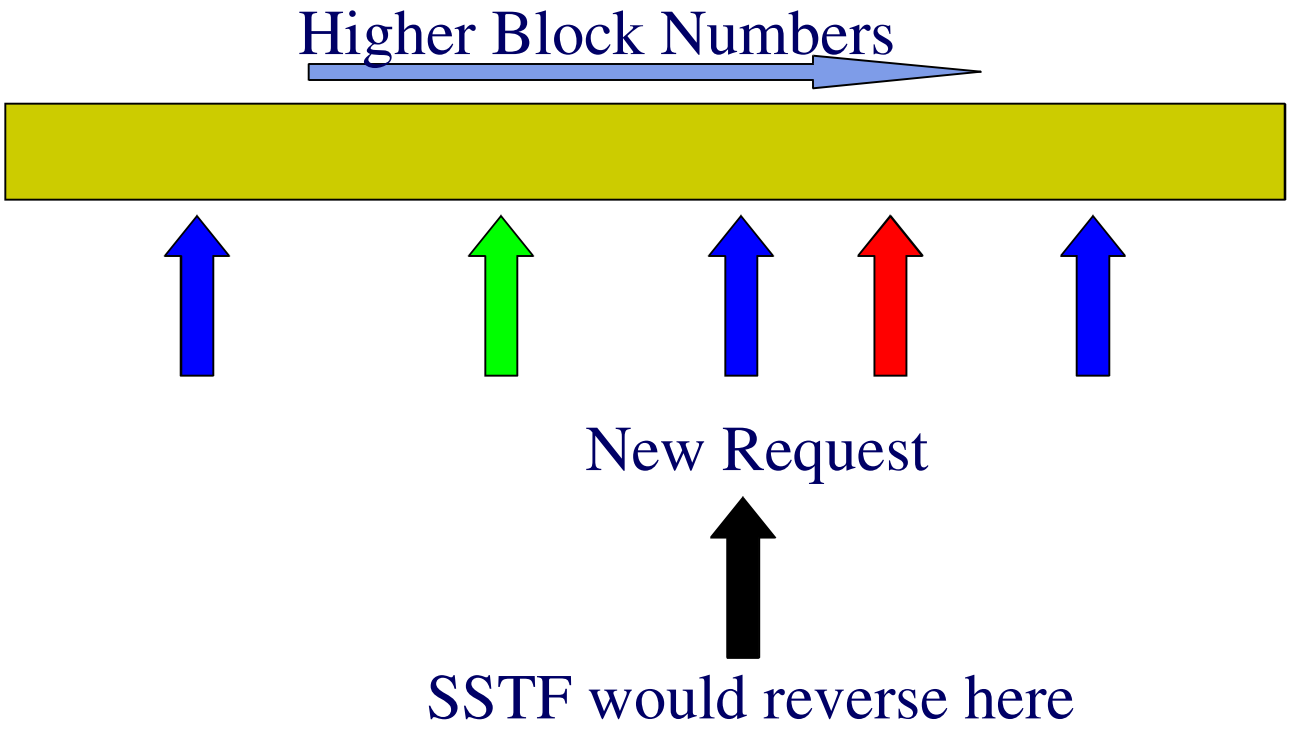

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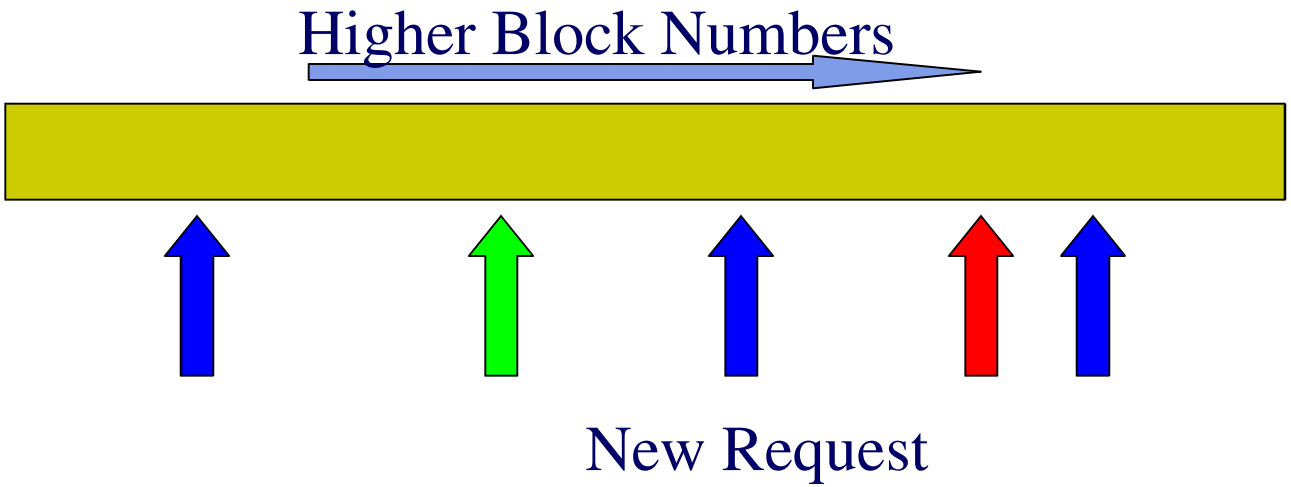

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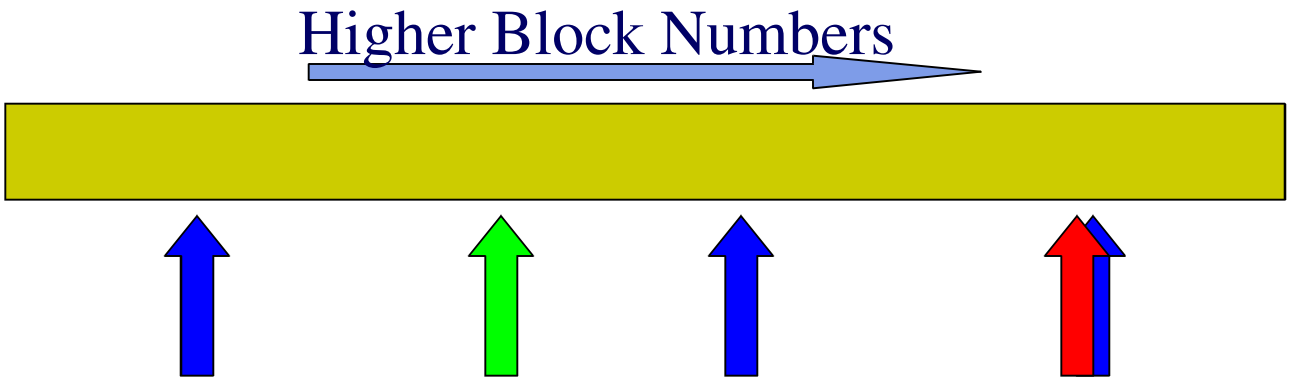
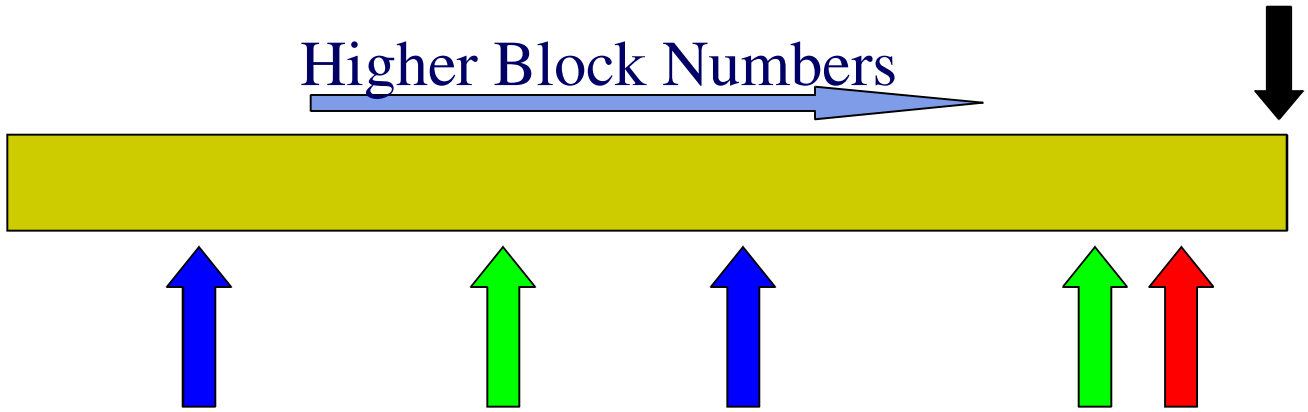


SCAN

In SCAN, we continue to the end of the disk



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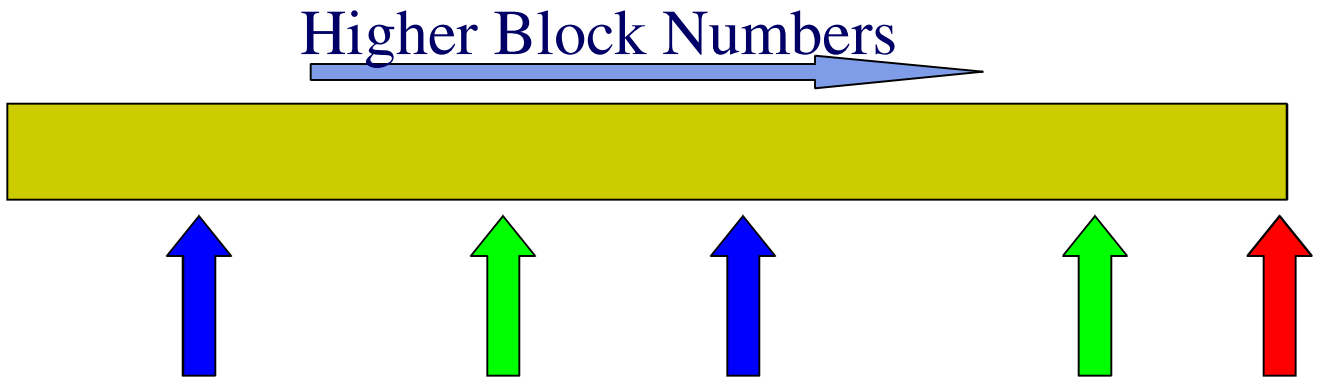

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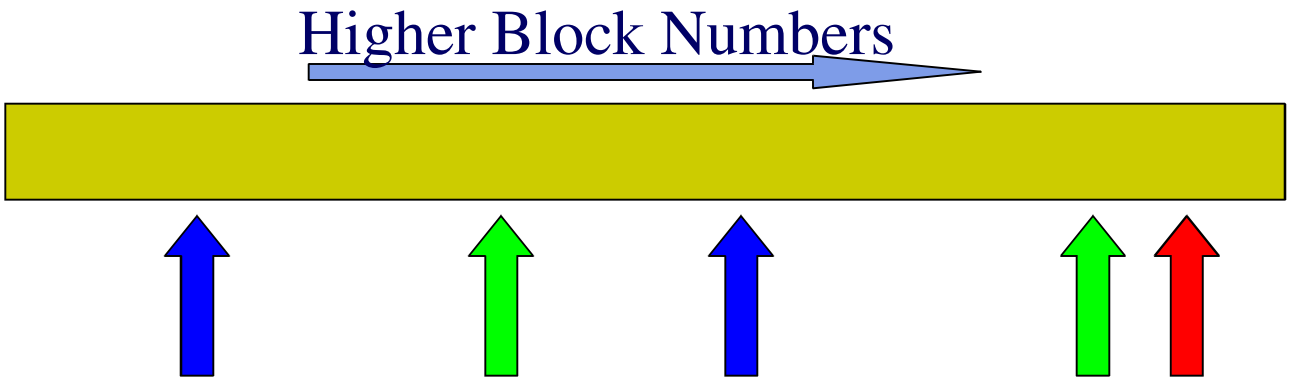

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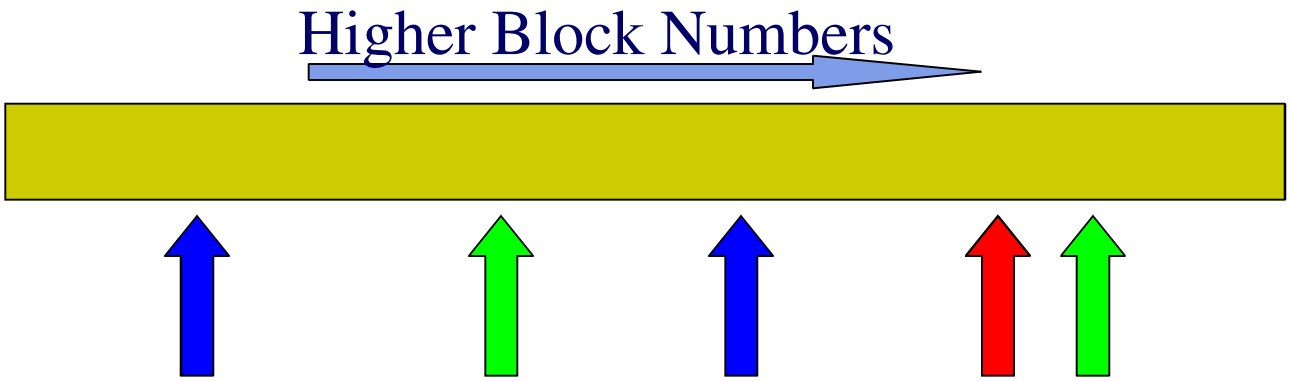

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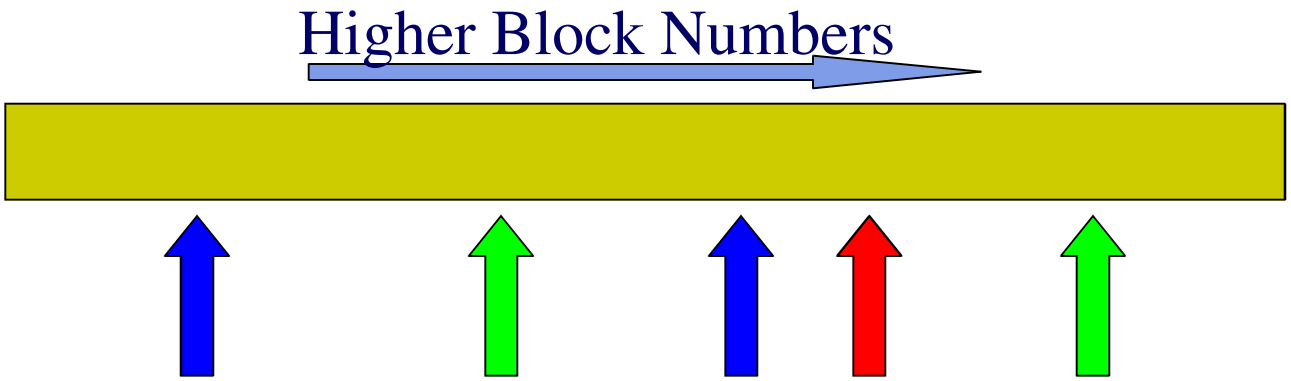

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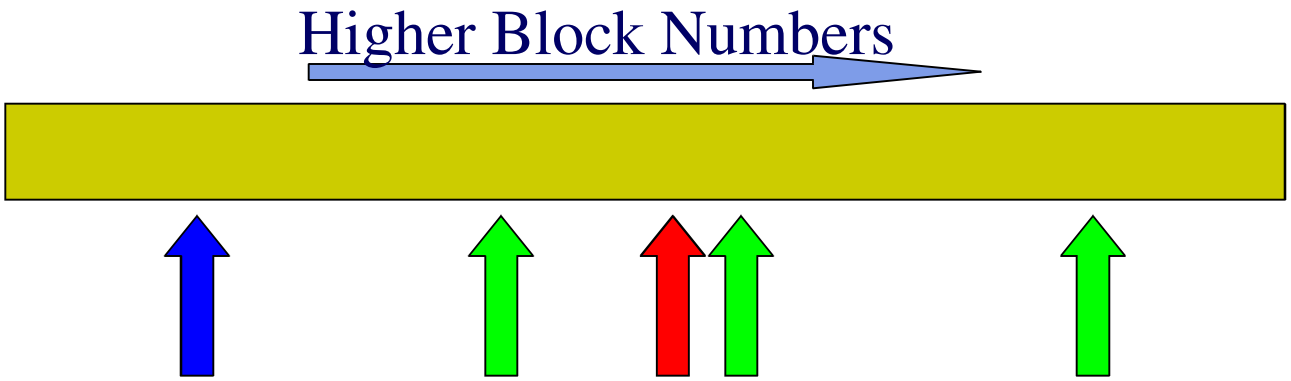

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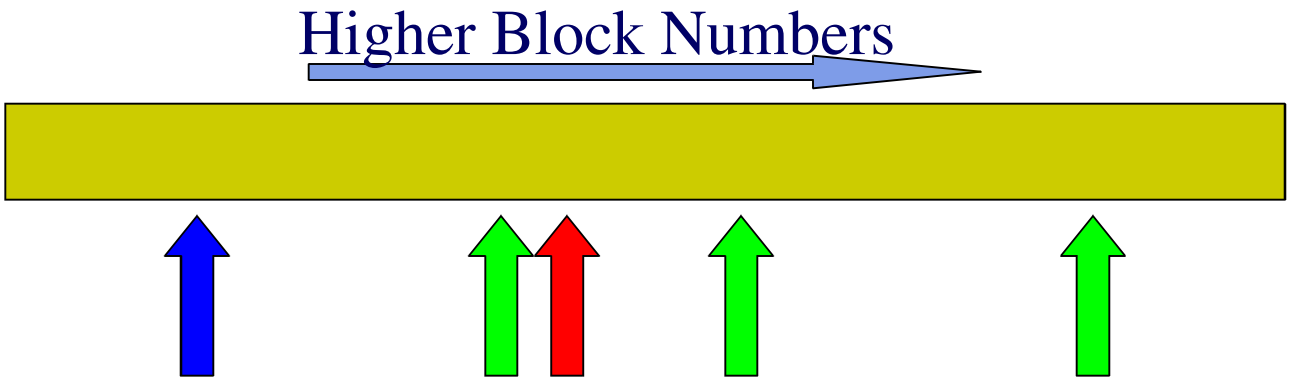

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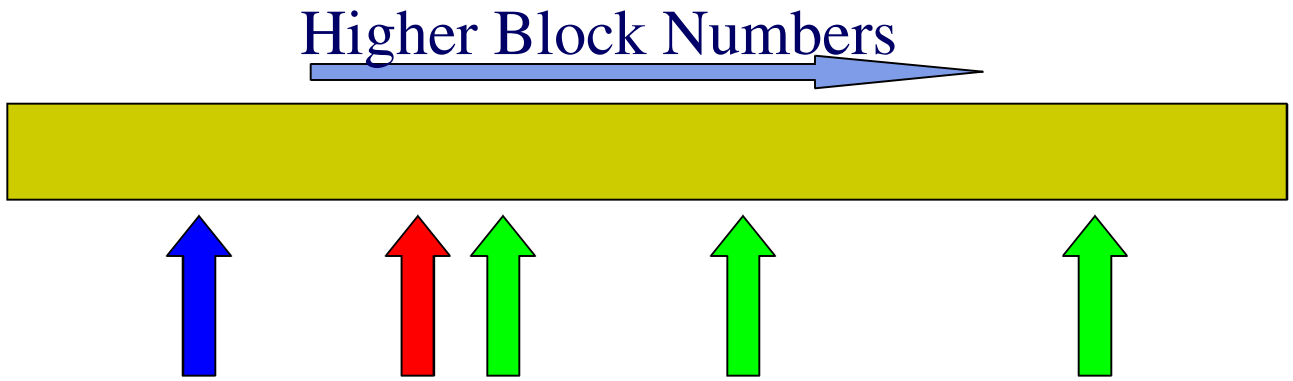

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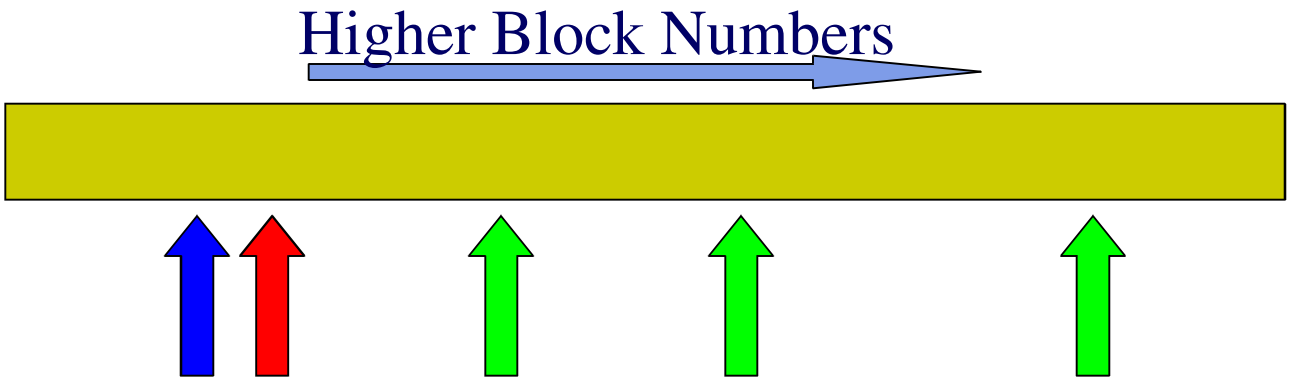

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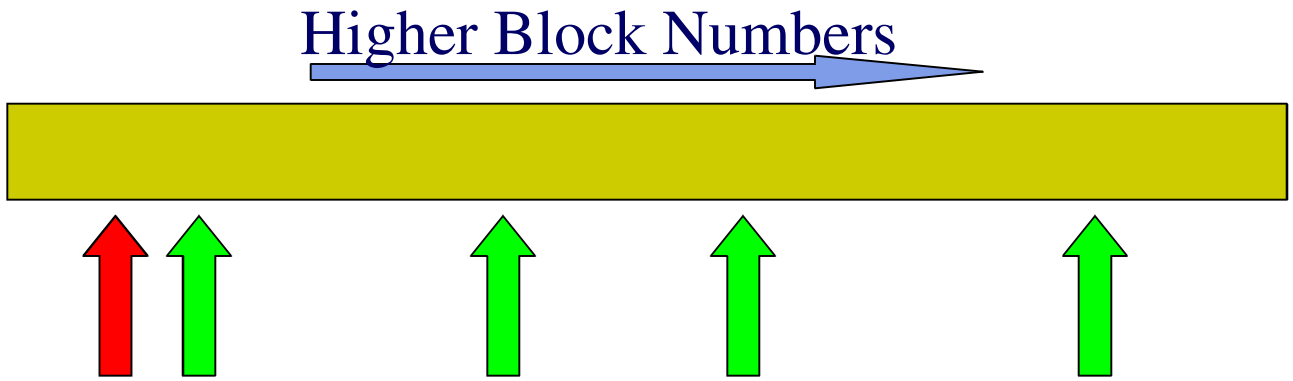

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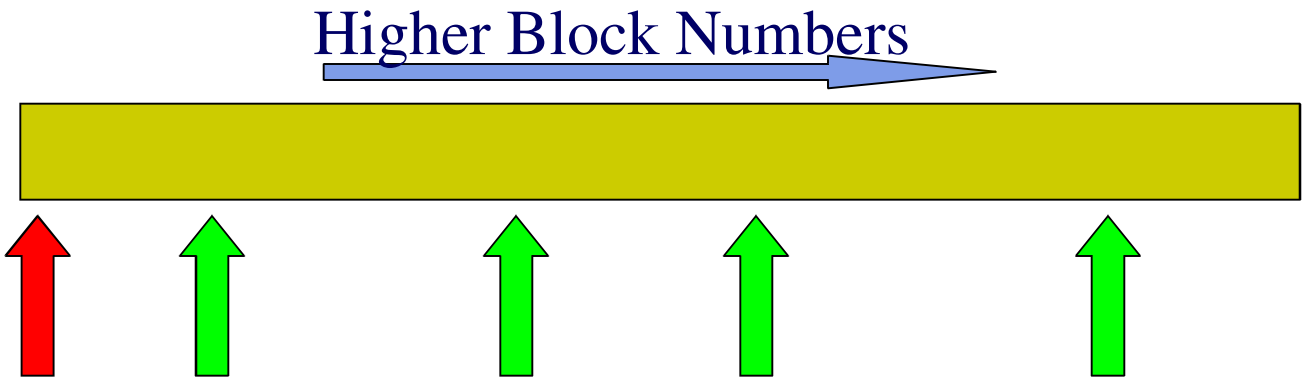


Evaluating SCAN

Mean response time

- Worse than SSTF, better than FCFS

Response time variance

- Better than SSTF

Unfair – why?

LOOK

Just like SCAN – sweep back and forth through cylinders

Don't wait for the "thud" to reverse the scan

- **Reverse when there are no requests "ahead" of the arm**

Improves mean response time, variance

Still unfair though

CSCAN - "Circular SCAN"

Send requests in ascending cylinder order

When the last cylinder is reached, seek all the way back to the first cylinder

Long seek is amortized across all accesses

- Key implementation detail
  - Seek time is a *non-linear* function of seek distance
  - One big seek is faster than N smaller seeks

Variance is improved

Fair

Still missing something though…

C-LOOK

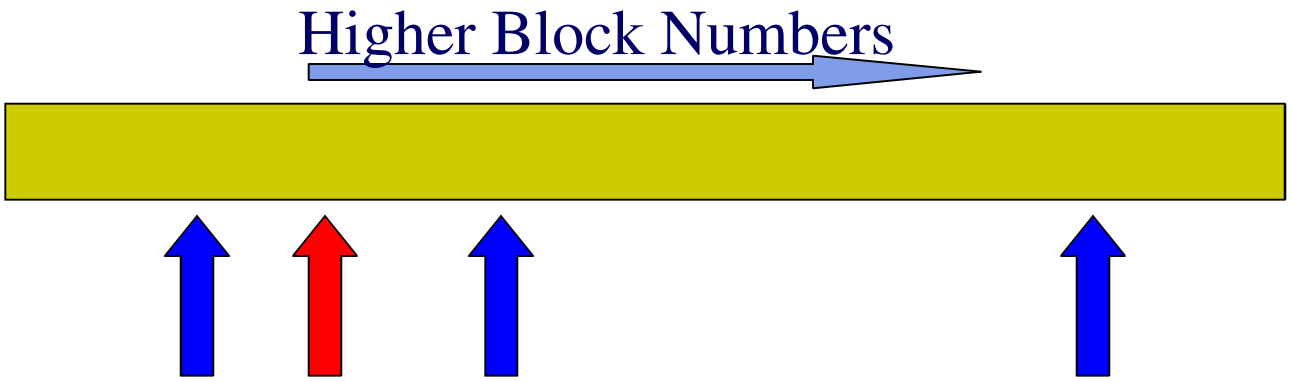
CSCAN + LOOK

Scan in one direction, as in CSCAN

If there are no more requests in current direction go back to furthest request

Very popular

C-LOOK



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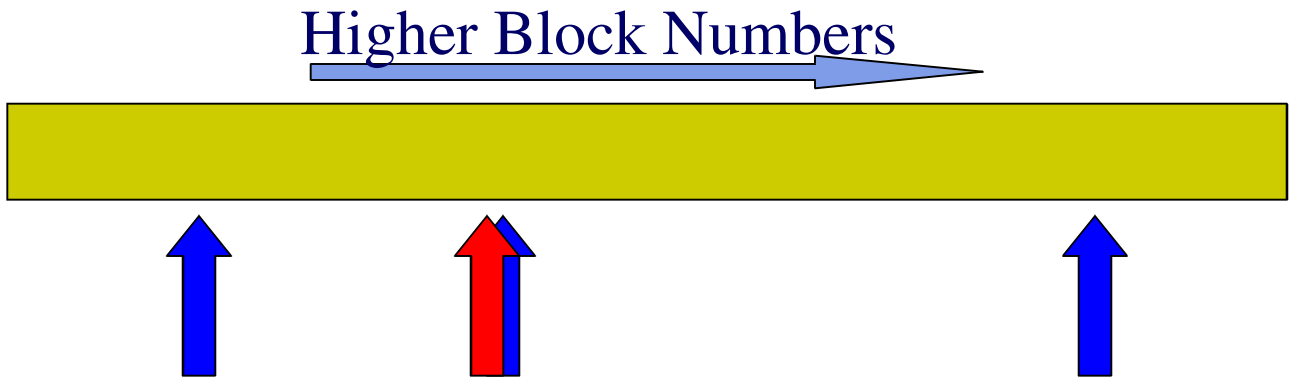

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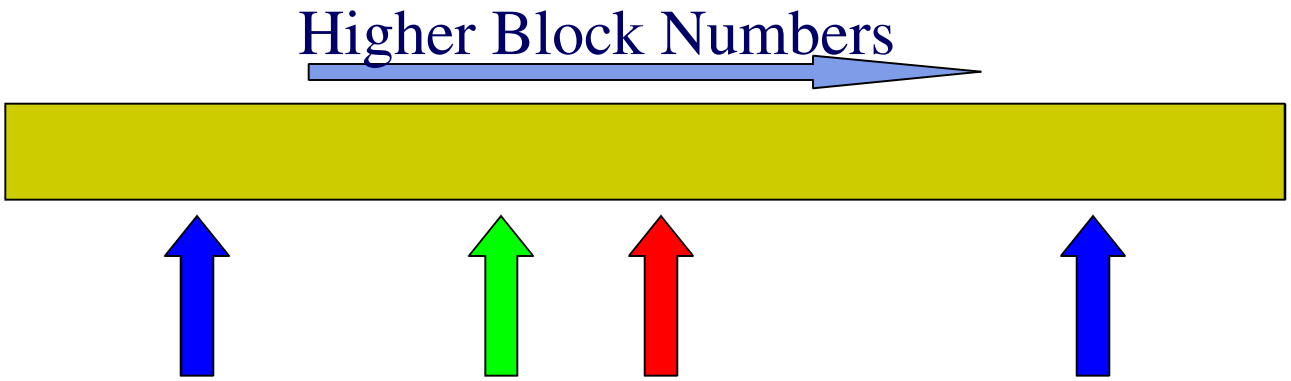

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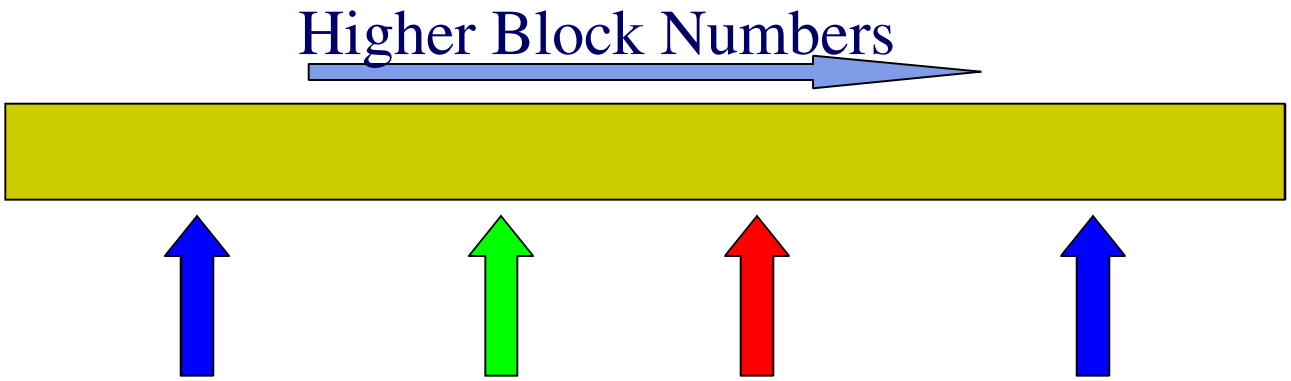

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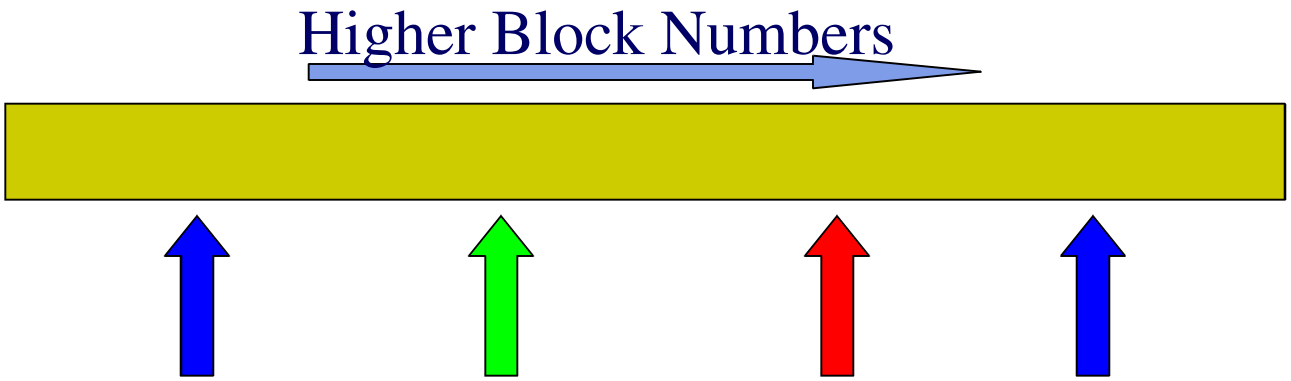
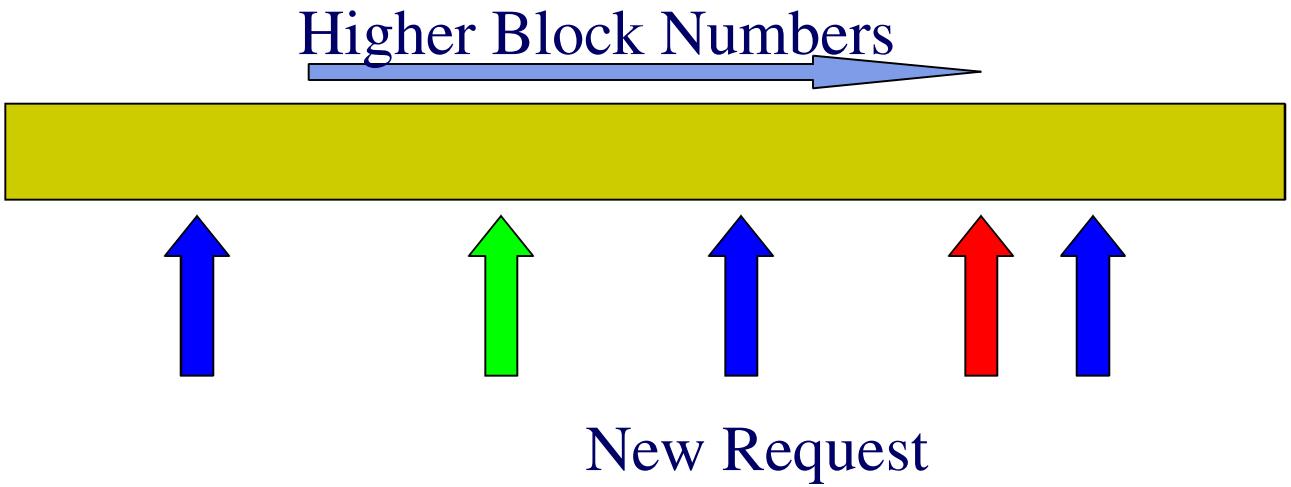
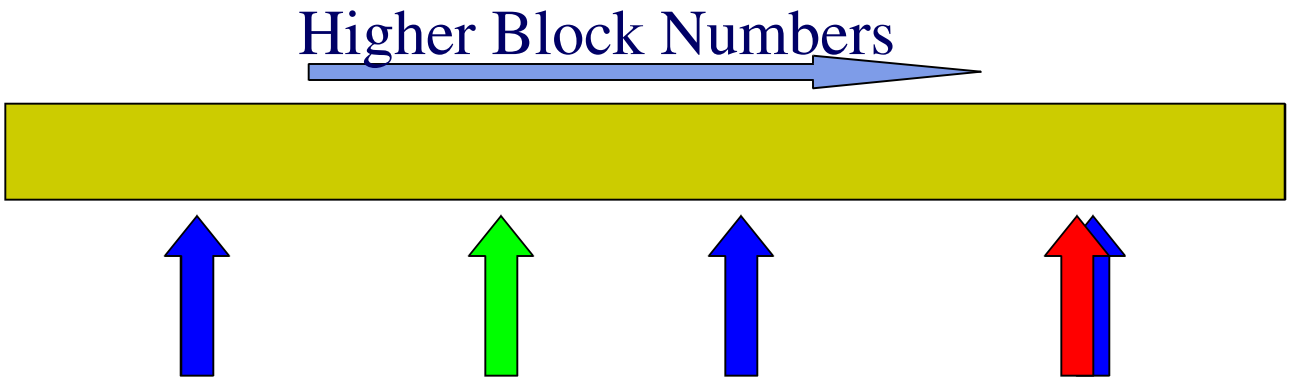


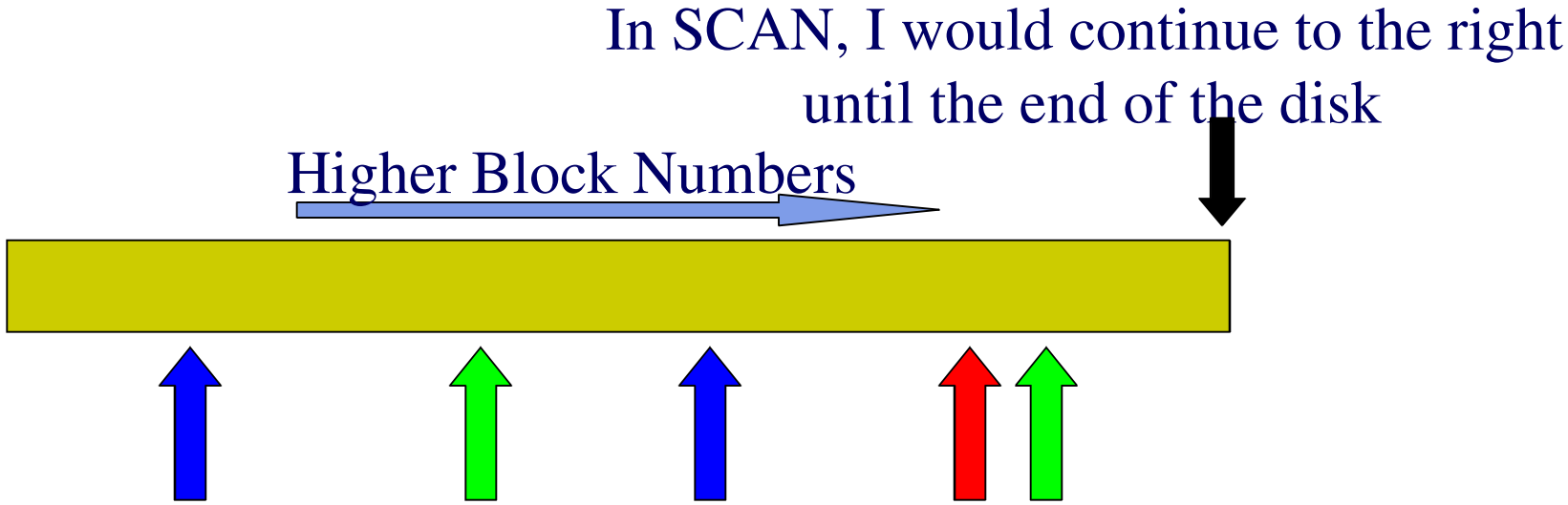
C-LOOK



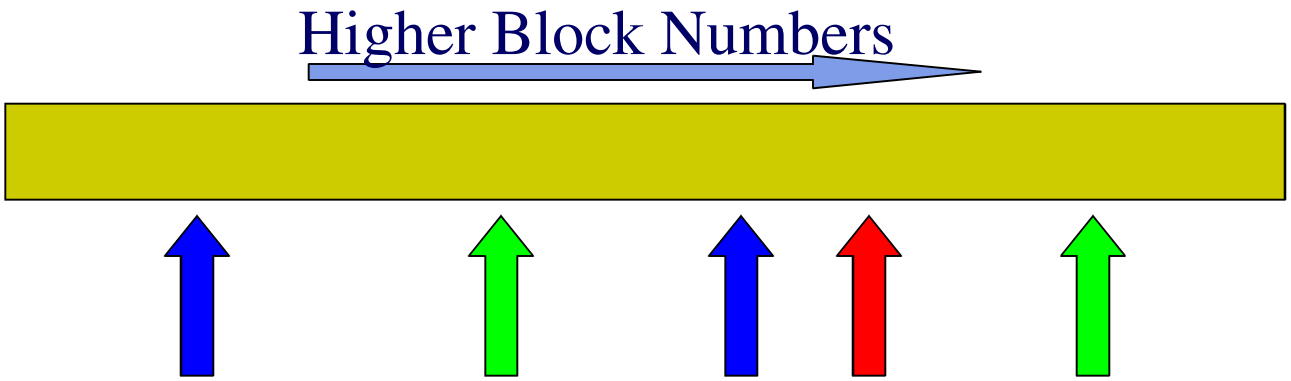
C-LOOK



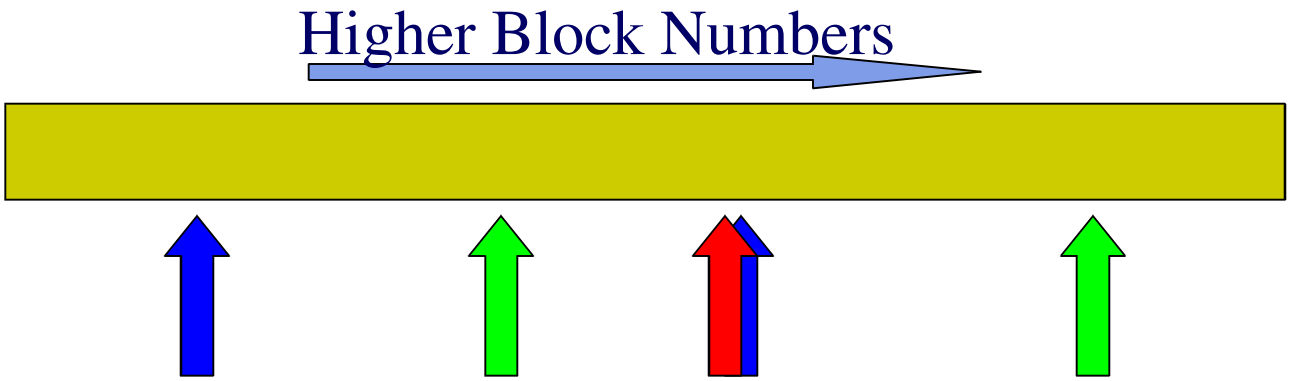
C-LOOK


In SCAN, I would continue to the right
until the end of the disk
Higher Block Numbers

C-LOOK

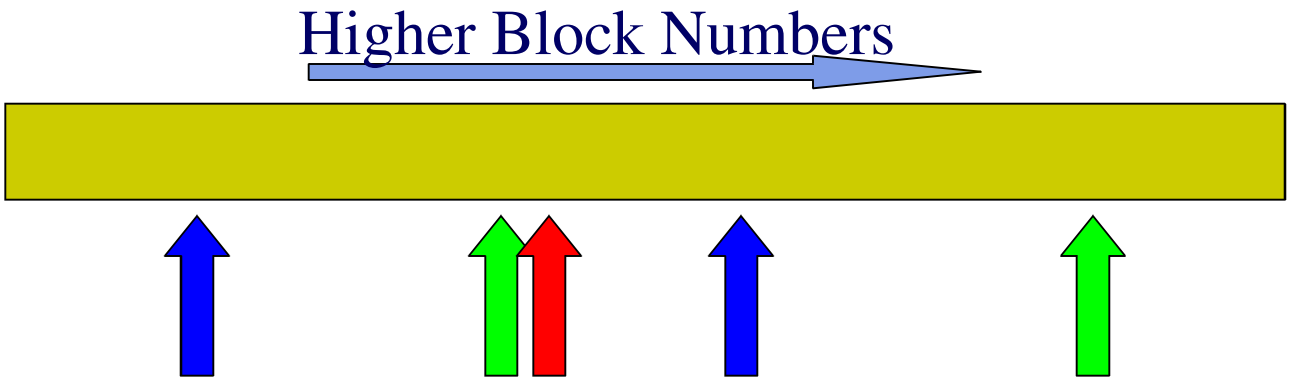


C-LOOK

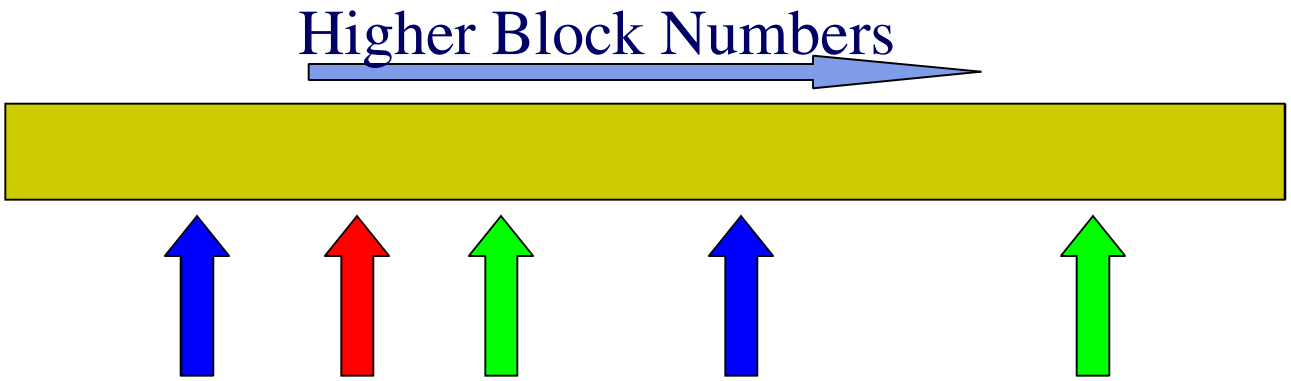


In LOOK, we would have read this request

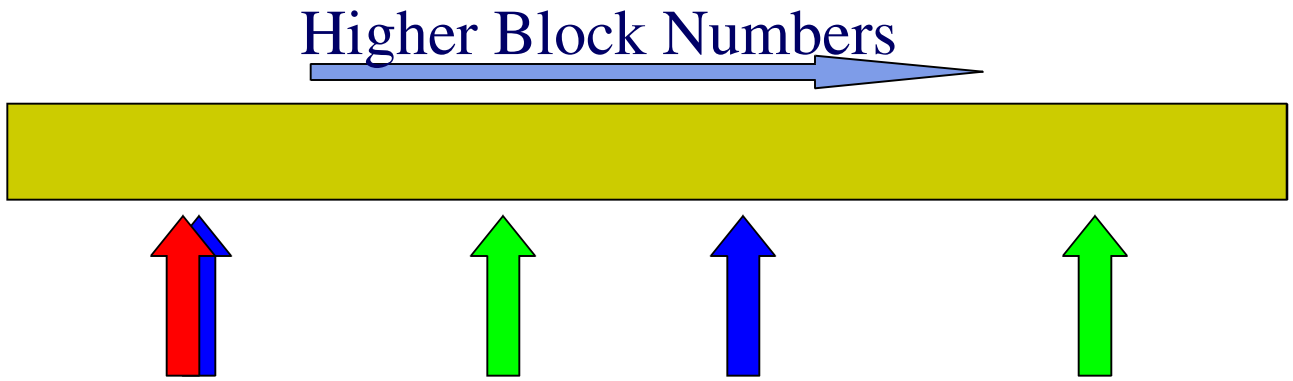
C-LOOK



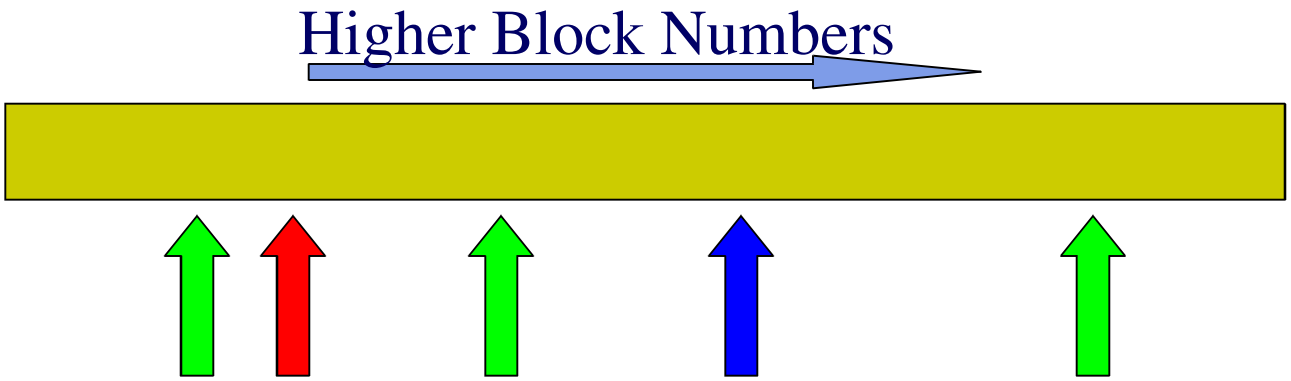
C-LOOK



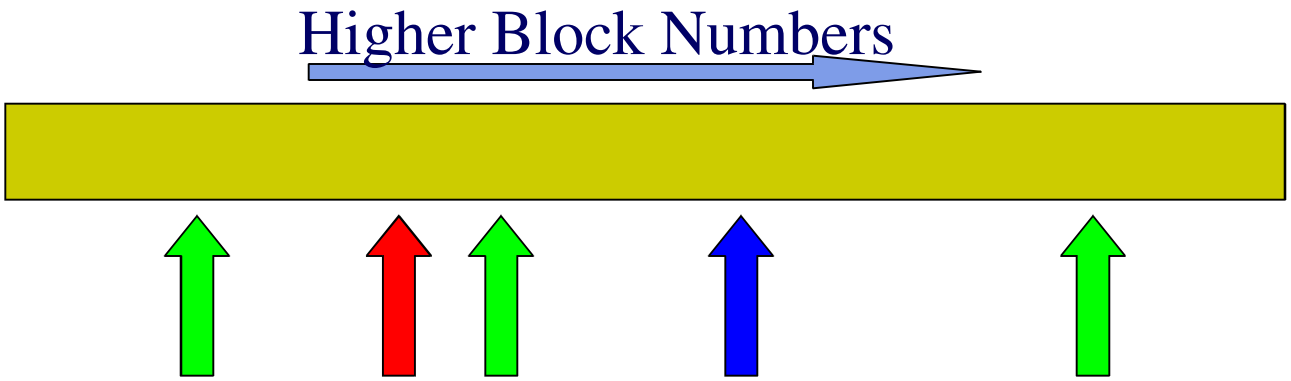
C-LOOK



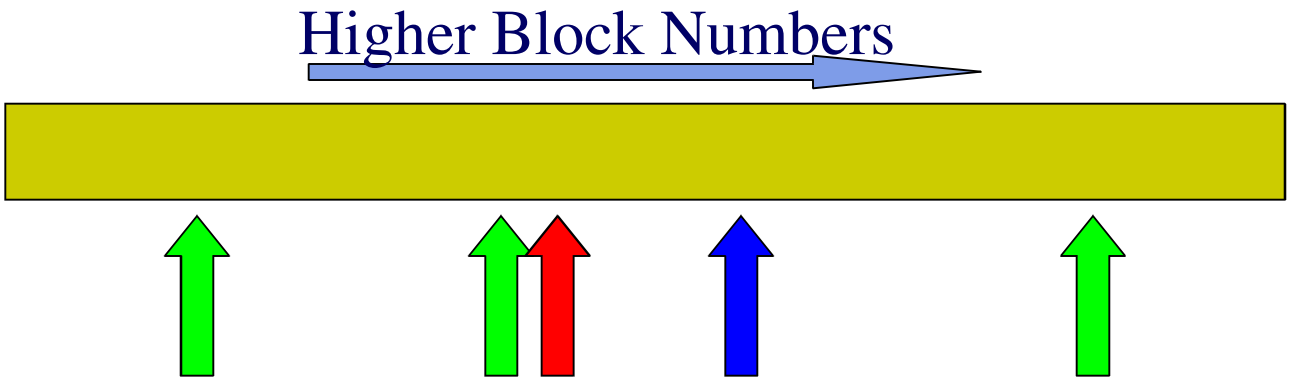
C-LOOK



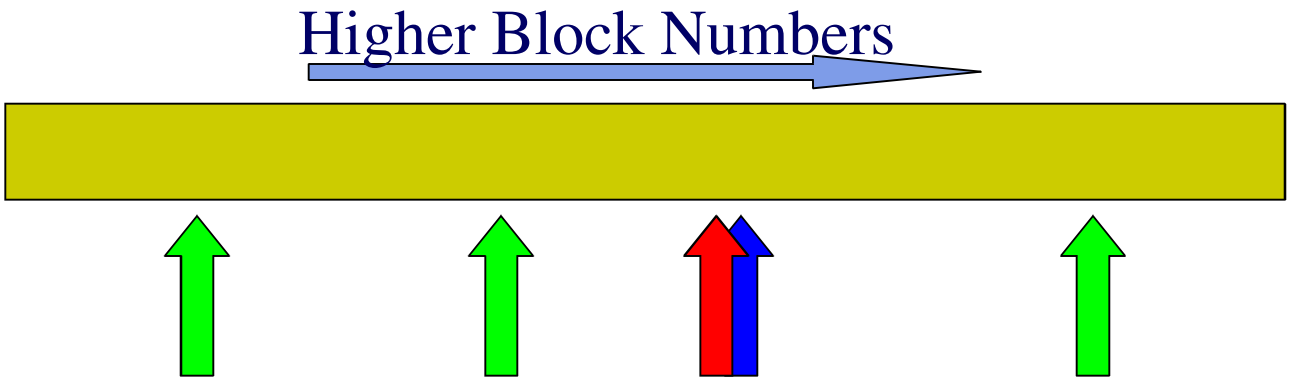
C-LOOK



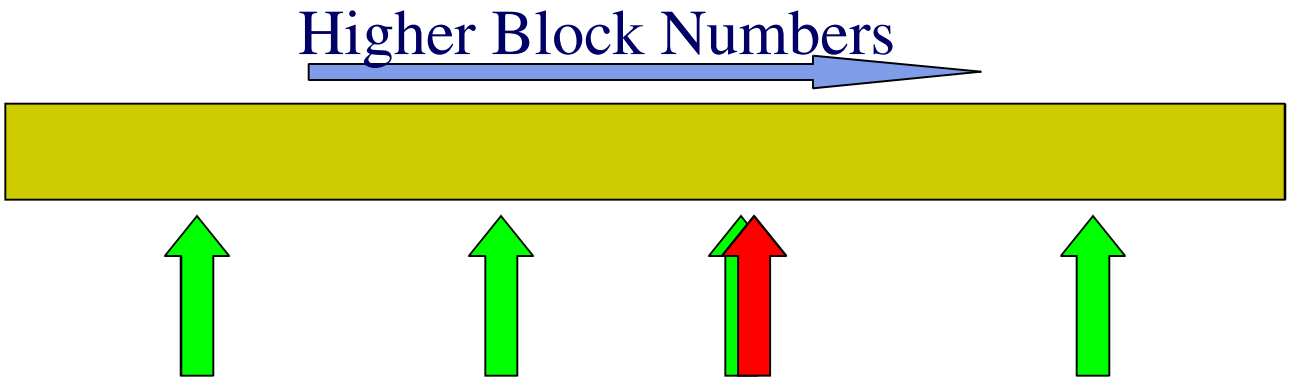
C-LOOK



C-LOOK



C-LOOK



Shortest Positioning Time First (SPTF)

Key observation

- Seek time takes a while
- But rotation time is comparable!
  - Short seeks are faster than whole-disk rotations
- What matters is *positioning* time, not seek time

SPTF is like SSTF

- Serve “temporally nearest” sector next

Challenge

- Can't estimate by subtracting sector numbers
- Must know rotation position of disk in real time!

Performs better than SSTF, but still starves requests

Weighted Shortest Positioning Time First (WSPTF)

SPTF plus fairness

Requests are "aged" to prevent starvation

- Compute "temporal distance" to each pending request
- Subtract off "age factor"
- Result: sometimes serve old request, not closest request

Various aging policies possible, many work fine

Excellent performance

Like SPTF, hard for OS to know disk status in real time

- On-disk schedulers can manage this, though...
  - Some disks (SCSI, newer IDE) accept a request queue
  - When complete, give OS both data and sector number

Lies Disks Tell

Disks re-order I/O requests

- You ask “read 37”, “read 83”, “read 2”
- Disk gives you 37, 2, 83
  - Not so bad

Disks lie about writes

- You ask “read 37”, “write 23”, “read 2”
- Disk writes 23, gives you 2, 37
  - Still not so bad
- You ask “write 23”, “write 24”, “write 1000”, “read 4-8”, ...
- Disk writes 24, 23 (!!), gives you 4, 5, 6, 7, 8
  - What if power fails before last write?
  - What if power fails between first two writes?

Lies Disks Tell

Disks lie about lies

- Special commands
  - Flush all pending writes
    - » Think “my disk is 'modern'”, think “disk barrier”
  - Disable write cache
    - » Think “please don't be quite so modern”
- Some disks ignore the special commands
  - “Flush all pending writes” ⇒ “Uh huh, sure, no problem”
  - “Disable write cache” ⇒ “Uh huh, sure, no problem”
- Result
  - Great performance on benchmarks!!!
  - Really bizarre file system corruption after power failures

Conclusions

Disks are very slow

Disks are very complicated

FCFS is a very bad idea

- C-LOOK is ok in practice
- Disks probably do something like WSPTF internally

Disks lie

- Some are vicious